

### Question 1a: Paired Sum

```
public static boolean pairSum(int[] arr, int sum) {  
    for (int i = 0; i < arr.length; i++) {  
        for (int j = i + 1; j < arr.length; j++) {  
            if (arr[i] + arr[j] == sum) {  
                System.out.println("Found pair: " +  
                    arr[i] + " + " + arr[j] +  
                    " = " + sum);  
                return true;  
            }  
        }  
    }  
    return false;  
}
```

This method would, in worse case scenario, run in time complexity  $O(n^2)$ , because:

- The outer for loop runs  $n$  times.
- The inner for loop runs  $\sum_{i=0}^n n - i - 1$  times in worst case (pair is never found)
  - This series evaluates to some factor of  $n^2$

The  $n^2$  dominates the  $n$  with large array sizes, so the final time complexity is  $O(n^2)$ .

### Question 1b: Tripled Sum

```
public static boolean tripletSum(int[] arr, int sum) {  
    for (int i = 0; i < arr.length; i++) {  
        for (int j = i + 1; j < arr.length; j++) {  
            for (int k = i + 2; k < arr.length; k++) {  
                if (arr[i] + arr[j] + arr[k] == sum) {  
                    System.out.println("Found triplet: " + arr[i] +  
                        " + " + arr[j] + " + " +  
                        arr[k] + " = " + sum);  
                    return true;  
                }  
            }  
        }  
    }  
    return false;  
}
```

This method would, in worse case scenario, run in time complexity  $O(n^3)$ , because:

- We know from the previous example that the middle loop runs some factor of  $n^2$  times in worst case
- Without doing an obscene amount of series calculations, we can logically suggest that the inner loop adds another “layer” to the iteration, and dominates the smaller  $n$  terms, resulting in  $O(n^3)$  runtime.

- I believe the inner loop runs  $\sum_{i=0}^n (n - i - 1 + \sum_{i=0}^n n - i - 2)$  times in worst case scenario

## Question 1c: Matrix Multiplication

```
for (int i = 0; i < n; i++) {  
    for (int j = 0; j < n; j++) {  
        double sum = 0;  
        for (int k = 0; k < n; k++) {  
            sum += a[i][k] * b[k][j];  
        }  
        c[i][j] = sum;  
    }  
}
```

This method would run in **time** complexity  $O(n^3)$ , because:

- Outer for loop runs  $n$  times
- Middle for loop runs  $n * n = n^2$  times
- Inner for loop runs  $n * n * n = n^3$  times

$n^3$  dominates smaller  $n$  terms, so  $O(n^3)$  runtime.

This method would run in **space** complexity  $O(n^2)$ , because:

- The method accesses three  $n * n$  arrays: the 2 matrices to be multiplied and the product matrix
- $3 * n * n = 3n^2 \Rightarrow O(n^2)$