Getting Started In Haskell

Wayne Chang (wayne@wycd.net)

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Slides and Exercises

With Git installed, you can run this command to download the slides and exercises.

\$ git clone https://github.com/wyc/haskell-intro.git

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These slides have coverage of:

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glhf!

Introduction

Haskell is statically-typed, functional, immutable, and lazy.

Statically-Typed

• Types are checked at compile time, as opposed to runtime.

Functional

• "Code is data."

Immutable

• Values do not get overwritten.

Lazy

• Evaluation is deferred for as long as possible.

Statically-Typed

Types are checked at compile time, as opposed to runtime.

```
int add(int x, int y) {
          return x + y;
}

int main() {
          add(1, 2);
          add(1, "two");
          return 0;
}

$ gcc -o s s.c
s.c: In function 'main':
s.c:9:20: warning: passing argument 2 of 'add'
    makes integer from pointer without a cast
          add(1, "two");
```

Functional

"Code is data."

Functions are passed around as easily as any other values, getting "first-class" treatment. This style allows us operate on whole data structures as opposed to just pieces of them at a time.

```
$ node
> [1, 2, 3, 4, 5]
    .map((elem) => elem * 2)
    .reduce((acc, elem) => acc * elem, 1)
3840
```

Immutable

Values do not get overwritten, or else you get yelled at.

```
$ node
> const a = 5
undefined
> a = 6
TypeError: Assignment to constant variable.
```

Immutable

What's the difference between Caching and Memoization?

- Caching is hoping that the result happens to be in a fast store, and then fetching it if it isn't.
- Memoization is remembering the mapping of specific inputs to their results for a function.

A function can only be memoized if it is *referentially transparent*, otherwise known as *immutable* or *pure*; the same inputs must always produce the same result.

Lazy

Evaluation is deferred for as long as possible.

"Don't start cleaning the apartment until the doorbell rings, and then only clean the parts that they can see."

Goodbye to Pretense

You don't need category theory, abstract algebra, or graduate-level algorithms to use or understand Haskell.

Hello, functional programming

To to remove the first three elements:

OOP-like:

```
[1,2,3,4,5].drop(3) // [4,5]
Functional:
drop 3 [1,2,3,4,5] -- [4,5]
```

In functional programming, we focus on the functions! When you focus on the objects themselves, you're "object-oriented".

Functional programming

Focus on functions.

- We ask "what can this function operate on?" and not "what can this object do?"
- What if we didn't have countless POJOs, and instead some core functions that we can use over and over again?

Composability is key!

```
Prelude> map (drop 1) [[1, 2, 3], [4, 5, 6], [7, 8, 9]] [[2,3],[5,6],[8,9]]
```

Getting Started (1/5)

- 1. Install Stack and the Glasgow Haskell Compiler (GHC) for GNU/Linux, Mac OSX, or Windows from https://docs.haskellstack.org/en/stable/README/
- 2. Setup Stack.
 - \$ stack setup
- 3. Write hello.hs:

```
main :: IO ()
main = do
putStrLn "Hello, World!"
```

Getting Started (2/5)

```
4. Compile it and run!
```

```
$ stack ghc hello.hs
$ ./hello
Hello, World!
```

5. Alternatively, you can use runhaskell:

```
$ stack runhaskell hello.hs
Hello, World!
```

Getting Started (3/5)

7. Write a function to use in GHC's interactive environment, ghci. myfuncs.hs:

```
salesTax :: Double -> Double
salesTax price = price * 0.089
```

8. Enter the ghci shell and load your source file.

```
$ stack ghci
GHCi, version 8.0.2: http://www.haskell.org/ghc/
Prelude> :load myfuncs
[1 of 1] Compiling Main ( myfuncs.hs, interpreted )
Ok, modules loaded: Main.
*Main> salesTax 100
8.9
```

Getting Started (4/5)

7. Importing modules:

```
-- Haskell
import Data.Char
import Data.Char (ord)
import qualified Data.Char as C
```

```
ord 'a' -- 97
C.ord 'a' -- 97
# Python
import data.char
import ord from data.char
import data.char as C

ord('a') # 97
C.ord('a') # 97
```

Getting Started (5/5)

8. Try importing from ghci.

```
$ stack ghci
GHCi, version 8.0.2: http://www.haskell.org/ghc/
Prelude> import Data.Char (ord)
Prelude Data.Char> ord 'a'
97
```

Basic Types

Types start with a Capital letter, variables (including functions and type variables) start with lowercase.

- Int
 - Fixed-Precision Whole Number
- Integer
 - Unbounded Whole Number
- Float
 - Single-Precision Floating
- Double
 - Double-Precision Floating
- Bool
 - True or False
- Char
 - Single character, such as 'a'

Lists (1/2)

• List

- For example, list of Char looks like [Char]
- The type resembles generic lists in Java such as List<T> (instantiated to List<Character>), but that's where the similarities end.
- Lisp-style construction (Horribly inefficient for, e.g., binary data? Yes.
 More on this later.) Illustration in JSON:

```
{
    'value': 5,
    'next': {
        'value': 4,
        'next': {
            value: 3,
            'next': null
        }
}
```

Lists (2/2)

- List (continued...)
 - Literal notation: [1,2,3]
 - Syntactic sugar for: 1:2:3:[], where : prepends a single element to a list, and [] is an empty list:

```
Prelude> 1:[2,3] [1,2,3]
```

- List items must all be of the same type. Heterogenous lists are possible using advanced language features, but they are usually an anti-pattern. When have you ever been happy to see List<Object> in Java?

Operations on Lists (1/4)

• Take first element with head:

```
Prelude> head [1,2,3]
1
Prelude> head []
*** Exception: Prelude.head: empty list
```

• Take *n*th element using the index (!!) operator:

```
Prelude> [10, 11, 12] !! 0
10
Prelude> [10, 11, 12] !! 2
```

```
Prelude> [10, 11, 12] !! 5
*** Exception: Prelude.!!: index too large
```

Operations on Lists (2/4)

• Check for the presence of a value using elem:

```
Prelude> elem 99 [42, 52, 62]
False
Prelude> elem 42 [42, 52, 62]
```

• Reverse a list with reverse:

```
Prelude> reverse [1,2,3,4] [4,3,2,1]
```

Operations on Lists (3/4)

• Add two lists together with ++:

```
Prelude> [1,2] ++ [3,4]
[1,2,3,4]
Prelude> "type" ++ "writer"
"typewriter"
```

• Grab the first n elements with take:

```
Prelude> take 2 [1,2,3,4,5]
[1,2]
```

Operations on Lists (4/4)

• Mapping a function across a list with map:

```
Prelude> let addFive x = x + 5
Prelude> map addFive [1,2,3]
[6,7,8]
```

• Reducing a list with foldl:

```
Prelude> foldl (\acc elem -> acc + elem) 0 [1,2,3]
6
```

Tuples (1/2)

- Tuple
 - Store several values into a single value. They can be of any type, but you must know how many values there will be.
 - For example,
 Prelude> let point = (3,4)
 Prelude> fst point
 3
 Prelude> snd point
 4
 Prelude> let point3d = (3,4,5)
 Prelude> let location = ("Germany", (51.1657, 10.4515))
 Prelude> fst location
 "Germany"
 Prelude> snd location

Tuples (2/2)

(51.1657, 10.4515)

- - Tuples are of different types when they vary in size or constituent types:

```
Prelude> (5, 5) :: (Int, Int)
Prelude> (5, 5, 5) :: (Int, Int, Int)
Prelude> (5, "Five") :: (Int, String)
```

Reading Types (1/5)

:: can generally be read as "has type", and it is a more of a declaration than an observation.

```
-- 5 ``has type'' Int

Prelude> 5 :: Int
-- 5 ``has type'' Integer

Prelude> 5 :: Integer
-- 5 ``has type'' String...uh oh, we can't do that!
```

Reading Types (2/5)

The function arrow or function operator -> denotes a function that takes an argument of the type on the left and returns a value of the type on the right.

```
-- addFive ``has type'' function that takes
-- an Int and returns an Int
addFive :: Int -> Int
addFive x = x + 5

-- mkAdder ``has type`` function that takes
-- an Int and returns a (function that takes
-- an Int and returns an Int)
mkAdder :: Int -> (Int -> Int)
mkAdder x = (\y -> y + x)
addFive = (mkAdder 5)
```

Reading Types (3/5)

More practice reading types.

```
import Data.Char (toLower)
-- toLower ``has type'' function that takes a Char
-- and returns a Char
toLower :: Char -> Char
-- lowerString ``has type`` function that takes
-- a list of characters and returns a list of
-- characters
lowerString :: [Char] -> [Char]
lowerString = map toLower
Prelude> map toLower "ABCDE"
"abcde"
```

```
Prelude> lowerString "ABCDE"
"abcde"
```

Reading Types (4/5)

You can use :t in GHCi to ask about an expression's type.

```
Prelude> :t "Hello, World!"
"Hello, World!" :: [Char]
-- String and [Char] are the same type

Prelude> let addFive x = x + 5
Prelude> :t addFive
addFive :: Num a => a -> a
Prelude> addFive 1.2
6.2
```

Reading Types (5/5)

```
Locking down the type:
```

Currying (1/3)

Currying is the process of transforming a function that takes multiple arguments into a function that takes just a single argument and returns another function if any arguments are still needed.

```
-- Uncurried
add (x, y) = x + y
add (5, 2)
```

Curried forms are preferred because they allow for partial application:

```
-- Curried
add x y = x + y
addFive = add 5
addFive 2
```

Currying (2/3)

Curried 1, 2, and 3 are equivalent.

```
-- Uncurried

add :: (Int, Int) -> Int

add = \(x, y) -> x + y

-- Curried 1

add :: Int -> Int -> Int

add x y = x + y

-- Curried 2

add :: Int -> (Int -> Int)

add x = (\y -> x + y)

-- Curried 3

add :: (Int -> (Int -> Int))

add = (\x -> (\y -> x + y))
```

Currying (3/3)

The arrow operator is right-associative. Each value we apply returns a value conforming to "the rest" of the type signature.

```
a -> a -> a -> a
a -> a -> (a -> a)
a -> (a -> (a -> a))
```

Function application is left-associative due to currying. It "cancels out" because we apply values one at a time.

```
fabc (fa)bc
```

```
((f a) b) c
```

Reading the Type of map (1/3)

```
Prelude > import Data.Char (toLower)

Prelude Data.Char > :t map
map :: (a -> b) -> [a] -> [b]

-- `a' and `b' are known as type variables,
-- because they can represent any type. They
-- are uncapitalized to distinguish them
-- from specific types such as `Int'.

Prelude Data.Char > :t toLower
toLower :: Char -> Char

Prelude Data.Char > :t map toLower
map toLower :: [Char] -> [Char]
```

Reading the Type of map (2/3)

Hindley–Milner type inference!

```
1. Pass toLower to map.
```

```
map :: (a -> b ) -> [a] -> [b] toLower :: (Char -> Char)
```

2. See if there could be a fit.

3. There it is! Substitute let a = Char, b = Char in

```
(Char -> Char) -> [Char] -> [Char] (Char -> Char)
```

4. "Subtract," and then our result is

```
[Char] -> [Char]
map toLower :: [Char] -> [Char]
```

Reading the Type of map (3/3)

We've just done partial function application!

```
map :: (a -> b) -> [a] -> [b]
lowerString = (map toLower :: [Char] -> [Char])
```

lowerString is the resulting function of a partially-applied map.

 \mathtt{map} has a and b so we can change the type of the list.

```
Prelude> :t even
even :: Int -> Bool
Prelude> :t [1,2,3]
[1,2,3] :: [Int]

Prelude> map even [1,2,3]
[False,True,False]
Prelude> :t [False,True,False]
[False,True,False] :: [Bool]
```

Reading Typeclasses (1/5)

Num is a typeclass that defines some common numeric operations. Think interfaces in Java or Go, except that typeclasses apply to types.

```
Prelude> let addFive x = x + 5
Prelude> :t addFive
addFive :: Num a => a -> a
Prelude> addFive 1.2
6.2
```

Reading Typeclasses (2/5)

List function examples with type variables:

```
fst :: (a,b) -> a
head :: [a] -> a

take :: Int -> [a] -> [a]

zip :: [a] -> [b] -> [(a,b)]
id :: a -> a
```

Reading Typeclasses (3/5)

Abbreviated examples:

```
sum :: Num a => [a] -> a
sum = foldl (+) 0

-- Num - Numeric Types
(+) :: Num a => a -> a -> a
-- Eq - Equality Types
(==) :: Eq a => a -> a -> Bool
-- Ord - Ordered Types
(<) :: Ord a => a -> a -> Bool
```

Reading Type classes (4/5)

Nearly-complete definition from the GHC source code:

```
-- | Basic numeric class.
class Num a where
   (+), (-), (*) :: a -> a -> a
   -- / Unary negation.
   negate
            :: a -> a
   -- / Absolute value.
                     :: a -> a
   -- | Sign of a number.
                     :: a -> a
   -- | Conversion from an 'Integer'.
   fromInteger :: Integer -> a
   х - у
                     = x + negate y
                     = 0 - x
   negate x
```

Reading Typeclasses (5/5)

Typeclasses are useful because they allow us to define functions that can work across a variety of types sharing common properties.

```
class Show a where
    show :: a -> String
    -- ...more function requirements

instance Show Int where
    show = showSignedInt' -- Some low-level function
```

```
prefixedShow :: (Show a) => String -> a -> String
prefixedShow prefix x = prefix ++ ": " ++ show x

prefixedShow "MyValue" 42 -- "MyValue: 42"
prefixedShow "MyPet" "cat" -- "MyPet: \"cat\""
prefixedShow "MyPoint" (4,5) -- "MyPoint: (4,5)"
```

Infixed Functions (1/3)

Functions can be put in between arguments like this:

```
Prelude> 5 + 5
10
They are of type a -> b -> c.
Prelude> :t (+)
(+) :: Num a => a -> a -> a
```

When surrounded by parentheses, they are considered in *prefix notation*, meaning that they are called like any other non-infixed functions.

```
Prelude> (+) 5 5 10
```

Infixed Functions (2/3)

Likewise, non-infix functions matching type $a \rightarrow b \rightarrow c$ may be used as infix functions.

```
Prelude> mod 12 5
2
Prelude> 12 `mod` 5
2
```

This is all just syntactic sugar.

Infixed Functions (3/3)

In other languages such as C, infix functions are built in. In Haskell, we can define our own.

```
infixr 0 ~ -- right associative
-- a ~ b ~ c becomes a ~ (b ~ c)

infixl 0 ~ -- left associative
-- a ~ b ~ c becomes (a ~ b) ~ c

infix 0 ~ -- non-associative
-- a ~ b ~ c becomes (a ~ b) ~ c OR a ~ (b ~ c)
```

They have associativity and precendence (0 the weakest, 9 the strongest).

Anonymous Functions

Anonymous functions help us not think about names. They can be defined as follows:

```
add :: Int -> Int -> Int
add = \x y -> x + y

addFive :: Int -> Int
addFive = \x -> add x 5
```

They are expressions, and can be used as inputs to functions that take functions as arguments or assigned to variables.

```
Prelude> map (x \rightarrow x + 5) [1,2,3] [6,7,8]
Prelude> let double = (x \rightarrow x * 2) in double 5
```

Evaluation Order

When in doubt, use parentheses.

```
Prelude> 5 - 3 * 2
-1
Prelude> (5 - 3) * 2
```

You can then look up the associativity and precendence per operator to reason about what the execution would be.

Symbols Related to Evaluation Order

You will encouner two symbols that help dictate order of evaluation, \$ is known as *application*, and . is known as *function composition*.

• (\$) calls the function which is its left-hand argument on the value which is its right-hand argument.

```
map addFive (map addFive [1,2,3])
-- Same as
map addFive $ map addFive [1,2,3]
```

• (.) composes the function which is its left-hand argument on the function which is its right-hand argument.

```
addFive (double 5)

-- Same as
(addFive . double) 5

(Source: ellisbben, StackOverflow)
```

PITFALL: Mismatched Arguments (1/3)

Watch out for what gets passed in as arguments!

PITFALL: Mismatch Arguments (2/3)

```
-- ...continued
<interactive>:3:13: error:
• Couldn't match expected type
    '[Integer]'
    with actual type
        '(a0 -> b0) -> [a0] -> [b0]'
• Probable cause: 'map' is applied to too few arguments
    In the second argument of 'map', namely 'map'
    In the expression:
        map addFive map addFive [1, 2, 3]
    In an equation for 'it':
        it = map addFive map addFive [1, 2, 3]
```

PITFALL: Mismatched Arguments (3/3)

Pattern Matching

Pattern matching is an extremely useful technique that is a staple in functional programming languages. It allows us to "search" for the data we want from an expression. We will cover:

- Pattern matching constants, lists, and tuples
- Using pattern matching in
 - Function Definitions (we've already been using them!)
 - Guards
 - where
 - let

```
- case...of
```

Pattern Matching in Haskell is top to bottom, left to right.

Pattern Matching: In JavaScript

Here is an example of code with and without destructuring syntax in ES6:

While destructuring isn't pattern matching, it has a similar intent, and there are even proposals to add true pattern matching to the language (see references).

Pattern Matching: Matching Constants (1/2)

```
Somewhere in test.hs:

guess :: Int -> Bool
guess 5 = True

(guess 5) -- True
(guess 10) -- BOOM!
-- *** Exception: Non-exhaustive
-- patterns in function guess

GHC can warn us about this before it happens!

$ ghc -Wall test.hs
test.hs:83:1: warning: [-Wincomplete-patterns]
Pattern match(es) are non-exhaustive
In an equation for 'guess':
Patterns not matched:
p where p is not one of {5}
```

Pattern Matching: Matching Constants (2/2)

Fixed, in test.hs:
guess :: Int -> Bool
guess 5 = True
guess _ = False

- Underscore (_) means "throw this value out, we're not going to use it."
- Patterns are tried sequentially until there's a match.
- If no match could be found, then an exception is thrown. Therefore, it's good practice to have catch-all patterns at the end.

Pattern Matching: Matching Lists

Lists can be matched in a way that grants you their first few elements, which is very useful in recursive functions. The form is (x:xs).

```
map :: (a \rightarrow b) \rightarrow [a] \rightarrow [b]

map f (x:xs) = (f x):(map f xs)

map f [] = [] -- Cover all the cases
```

If necessary, we can also specify a handle for the whole list:

```
addListLengthToFirstElement :: [Int] -> Int
addListLengthToFirstElement lst@(x:_) = length lst + x
addListLengthToFirstElement [] = error "empty list"
```

PITFALL: Matching Lists

Be wary of these syntax fails that can lead to confusing compiler errors:

```
sum [x:xs] = x + sum xs -- NO
sum x:xs = x + sum xs -- NO
sum (x:xs) = x + sum xs -- ok
```

Pattern Matching: Matching Tuples

We can pull the values right out of tuples using pattern matching.

```
dist :: (Double, Double) -> (Double, Double) -> Double
dist (x, y) (x', y') =
    sqrt $ (x - x') ** 2 + (y - y') ** 2
Prelude> dist (0.0, 0.0) (1.0, 1.0)
1.4142135623730951

Even with nesting:
showIdAndFullName :: (Int, (String, String)) -> String
showIdAndFullName (id', (firstName, lastName)) =
    show id' ++ ": " ++ firstName ++ " " ++ lastName
Prelude> showIdAndFullName (5, ("Rebecca", "Jones"))
"5: Rebecca Jones"
```

Pattern Not-Matching: Guards

Guards are ways to test whether some properties are true or false. They don't actually do pattern matching. However, they can be used with the where clause to do pattern matching. They are evaluated sequentially.

```
scovilleComment :: Int -> String
scovilleComment score
   | score < 0 = "Huh?"
   | score <= 300 = "What spice?"
   | score <= 700 = "Eh...mild." -- Frank's
   | score <= 3500 = "A bit of kick!" -- Tabasco
   | score <= 10000 = "Mmm, spicy." -- Orange Kush
   | score <= 1500000 = "WOW THAT'S HOT" -- Pyro Diablo
   | otherwise = "OMG"</pre>
```

PITFALL: Guards

Do not put an equal sign after the function name and variables:

```
-- ok
guess x | x == 5 = "You got it!"
| otherwise = "Wrong!"
```

Pattern Matching: where

where can be used to help things read nicer and avoid repetition.

```
totalCompensation :: Int -> Int -> Int -> Int -> Int
totalCompensation base yearsEmployed rank performance =
  base + bonus
where rankMultiplier = (yearsEmployed / rank) `mod` 2
  bonus = rankMultiplier * performance * 10000
```

We can see right away that totalCompensation simply returns base + bonus, but if we want more detail, then we can dig in to the where clause. They can also be used to pattern match:

```
showPoint :: (Double, Double) -> String
showPoint p = "In " ++ show p ++ " " ++ x' ++ "," ++ y'
    where (x, y) = p
        x' = "x=" ++ show x
        y' = "y=" ++ show y

Prelude> showPoint (1,2)
"In (1,2) x=1, y=2"
```

Pattern Matching: let

let can be used to pattern match in the form of:

```
let <bindings> in <expression>
```

let itself is an expression and therefore produces a value, which is an important difference from where.

```
quadratic :: Double -> Double -> Double -> Int -> Double
quadratic a b c sign =
  let sign' = if sign >= 0 then 1.0 else -1.0
    rootTerm = sqrt $ (b**2) - 4*a*c
    numerator = (- b) + (sign') * rootTerm
    denominator = 2 * a
  in
    numerator / denominator
```

Pattern Matching: case...of

Case expressions can be thought as guards that do pattern matching instead of property testing. Like let, they are expressions.

if...then...else

if...then...else constructs in Haskell are better thought of as *if expressions* than *if statements*. They evaluate to a value depending on what comes after if. They cannot be used to pattern match.

```
addIfEven :: Int -> Int -> Int
addIfEven x y = x + if even y then y else 0
```

PITFALL: Sometimes spacing matters

```
zeroIfEven :: Int -> Int
zeroIfEven x = -- NO
    if even x then
        0
    else
        x
zeroIfEven x = -- ok
    if even x
-- ^ Code Block 1
        then 0
```

```
else x
-- ^ Code Block 2
```

Sum Types (1/7)

If we have three types of people: employees, managers, and clients, how do we cleanly represent them? Let's say we're using Go.

```
type PersonType string

const (
    EmployeePersonType PersonType = "Employee"
    ManagerPersonType PersonType = "Manager"
    ClientPersonType PersonType = "Client"
)
```

Sum Types (2/7)

```
type PersonType string
```

The primary benefit is that the compiler gives us some type checks for PersonType.

Sum Types (3/7)

```
type PersonType string
Disadvantages:
```

- Strings are expensive to pass around
- Strings are OVERKILL. c^{strlen} possibilities!
- Someone can just decide to invent a new type via cast: PersonType("Daughter")
- No exhaustiveness checks

We could consider using enum-like iota, but the last 3 points would still exist.

Sum Types (4/7)

```
type PersonType string
What if we forget one of the checks?
          case EmployeePersonType:
               return false
Runtime error!
$ go run test.go
panic: Unknown PersonType: Employee
goroutine 1 [running]:
panic(0x48c5c0, 0xc42000a320)
Sum Types (5/7)
Enter sum types!
-- Person is a sum type with zero-argument value
-- constructors of Employee, Manager, and Client.
data PersonType = Employee | Manager | Client
greet :: PersonType -> String
greet pt = case pt of
    Employee -> "G'day, employee!"
    Manager -> "Hello, manager."
main :: IO ()
main = do
   putStrLn $ greet Client
Sum Types (6/7)
```

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Run it:

Sum Types (7/7)

Fixed:

```
greet :: PersonType -> String
greet pt = case pt of
    Employee -> "G'day, employee!"
    Manager -> "Hello, manager."
    Client -> "Hey!"

$ ghc -Wall test.hs
$ ./test
Hey!
```

Product Types (1/2)

Sum types can be combined with product types to store additional data:

The data can be accessed using pattern matching, as seen with time.

Product Types (2/2)

Running our latest example:

```
main :: IO ()
main = do
    currentTime <- getCurrentTime
    putStrLn $ apptStatus (At currentTime)

$ ghc -Wall test.hs
$ ./test
Arrive by 2017-05-08 10:04:09.234585398 UTC</pre>
```

Algebraic Data Types

Sum and product types together make up the algebraic data types (ADTs) in Haskell. These are not to be confused with abstract data types. There are also generic algebraic data types (GADTs), but those are out of scope for this presentation.

```
-- Sum Types are combined with '|',
-- behaving like an "OR" operator

data MySumType = A | B
-- Product Types are combined with ' ',
-- behaving like an "AND" operator

data MyProductType = C D
```

Namespaces and Constructors (1/12)

Haskell has two namespaces:

- One for values.
- One for types.

Constructors:

- A data/value constructor is a "function" that takes 0 or more values and gives you back a new value.
- A type constructor is a "function" that takes 0 or more types and gives you back a new type.

(Source: kqr, StackOverflow)

Namespaces and Constructors (2/12)

The data operator constructs a new type.

```
data Color = Red | Green | Blue
Red :: Color
Green :: Color
Blue :: Color

data Color = RGB Int Int Int
RGB :: Int -> Int -> Color
```

For all practical purposes you can just think of data constructors as constants belonging to a type.

Namespaces and Constructors (3/12)

Same-named constructors

Namespaces and Constructors (4/12)

PITFALL: Constructor Naming

Namespaces and Constructors (5/12)

Pattern matching is the primary way we access the data wrapped in value constructors.

```
data Person = Person String Int
showNameAndAge :: Person -> String
showNameAndAge (Person name age) =
    name ++ " age=" ++ show age

Prelude> showNameAndAge (Person "Fred" 2)
"Fred age=2"
```

Namespace and Constructors (6/12)

Next example:

```
-- Cats have a certain number of lives
data Creature = Dog String | Cat String Int
showCreature :: Creature -> String
showCreature c = case c of
```

```
(Dog name) -> name
  (Cat name lives) -> name ++ " lives=" ++ show lives

Prelude> showCreature (Dog "Fido")
"Fido"

Prelude> showCreature (Cat "Mittens" 9)
"Mittens lives-left=9"
```

Namespaces and Constructors (7/12)

Another pattern matching example for constructed values:

```
-- Function from the test-taking belt of New Jersey
data Grade = A | B | C | D | F
data TestResult = Complete Grade | Missed
showTestResult :: TestResult -> String
showTestResult Missed = "The test was missed."
showTestResult (Complete g) = case g of
A -> "Top score!"
B -> "Horrible score, but you'll live!"
otherwise -> "Why even bother?"

Prelude> showTestResult (Complete A)
"Top score!"
```

Namespaces and Constructors (8/12)

Here is an example of a recursive type definition:

```
data List a = Nil | Cons a (List a)
```

• List is a unary type constructor, meaning it takes one type argument. This can be anything from Int, to Char, to [Char], etc.

```
List :: * -> *
```

• Nil is a nullary value constructor, taking no value arguments.

```
Nil :: List a
```

• Cons is a binary value constructor, taking two value arguments.

```
Cons :: a -> List a -> List a
```

Namespaces and Constructors (9/12)

Here is an example of a recursive type definition:

```
data List a = Nil | Cons a (List a)
```

Kinds

List can be viewed as an even higher-order function with the kind of * -> *, in the same fashion as Int -> Int.

Kinds are higher than types.

* represents a type argument, and in this case, the $type\ variable\ {\tt a},$ which contains a type.

Namespaces and Constructors (10/12)

You can ask GHCi to print out some kinds:

```
-- Int is a nullary type constructor

Prelude> :k Int
Int :: *

-- Int is a unary type constructor

Prelude> :k []
[] :: * -> *

-- We feed [] :: (* -> *) an Int :: *

Prelude> :k [Int]
[Int] :: *
```

PITFALL: Value Constructor Argument Mismatch

Value constructors only understand nullary type constructors (of kind *). Watch out for argument mismatches. In this case, the RHS List does not get passed the type argument a. Instead, List gets passed to Cons as its second argument and a gets passed as its third:

```
-- NO
data List a = Nil | Cons a List a
```

Instead, wrap the type constructor and its type argument in parentheses.

```
-- ok
data List a = Nil | Cons a (List a)
--
-- 1 2
```

Namespaces and Constructors (11/12)

Let's re-examine our recursive type definition:

```
data List a = Nil | Cons a (List a)
```

Value constructors explained (again):

• Nil is a nullary value constructor: it takes no values to make a new constructed value. This is just like a previous example of three of them:

```
data = Employee | Manager | Client
```

Cons is a binary value constructor. It takes a first argument of type a and a second argument of type List a. This is okay because List is of kind *
-> *, so when it is passed the type variable a in the expression List a, the resulting kind is *, just like Int.

Namespaces and Constructors (12/12)

Our List in action:

```
(Cons 5 Nil)
                             :: Num a => List a
(Cons 5 (Cons 4 Nil))
                           :: Num a => List a
(Cons 5 (Cons 4 (Cons 3 Nil))) :: Num a => List a
car
                   :: List a -> a
                   = x
car (Cons x _)
car Nil
                   = error "car called on Nil List"
-- Compare to the GHC source code:
-- | Extract the first element of a list...
                      :: [a] -> a
head
head (x:_)
head []
                       = badHead
```

Records (1/2)

Records are a way to generate accessor functions for a constructed value.

Without records:

```
data Person = Person String Int deriving Show
Prelude> :t Person "Fred" 5
Person "Fred" 5 :: Person
Prelude> let (Person _ age) = (Person "Fred" 5) in age
5
With records:
data Person = Person { name :: String, age :: Int } deriving Show
Prelude> Person {name="Fred", age=5}
Person {name = "Fred", age = 5}
Prelude> age Person {name="Fred", age=5}
```

Records (2/2)

Trouble in paradise. There is a serious namespace issue:

Google "haskell record problem" to learn more about the workarounds.

Typeclasses Revisited (1/6)

- Type classes were created to express "ad-hoc polymorphism," AKA, over-loaded functions.
- They turned out to be nicer in many ways than those in Java or C++.
- I stole some great examples from Philip Wadler's talk at Microsoft. He implemented typeclasses for Haskell and generics for Java.

Typeclasses Revisited (2/6)

We want a generic function max so that we don't have to define it over and over again for different types. Let's write it naively in C++ using templates.

```
#include <iostream>
template<class T>
T max(T x, T y) {
   return x < y ? y : x;
}</pre>
```

Typeclasses Revisited (3/6)

```
#include <iostream>

template<class T>
T max(T x, T y) {
    return x < y ? y : x;
}

void main() {
    // 2
    std::cout << max<int>(1, 2) << std::endl; // 2
    // 'b' (on most machines)
    std::cout << max<char>('a', 'b')) << std::endl;
    const char* a = "zzz"; const char* b = "aaa";
    // ???
    std::cout << max<const char*>(a, b)) << std::endl;
    // It's "aaa"
}</pre>
```

Typeclasses Revisited (4/6)

Instead of C++ templates, let's try it using Haskell typeclasses.

```
class Ord a where
    (<) :: a -> a -> Bool

instance Ord Int where
    (<) = primitiveLessInt

instance Ord Char where</pre>
```

```
(<) = primitiveLessChar

max :: Ord a => a -> a -> a
max x y = if x < y then y else x</pre>
```

Typeclasses Revisited (5/6)

List of Char? No problem. Just add:

Typeclasses Revisited (6/6)

But we can go deeper. String is just [Char]. Char already implements Ord. What if we abstracted this?

```
-- Instead of this tautology
Ord Char => [Char]
-- Let's say
Ord a => [a]
instance Ord a => Ord [a] where
[] < [] = False
[] < y:ys = True
x:xs < [] = False
x:xs < y:ys | x < y = True
| y < x = False
| otherwise = xs < ys
```

Now we can handle comparisons with any level of list nesting.

The Maybe Type

```
data Maybe a = Just a | Nothing
Just (5 :: Int) :: Maybe Int
```

```
Nothing :: Maybe a

addMaybe :: Int -> Maybe Int -> Int
addMaybe (x (Just y)) = x + y
addMaybe (x Nothing) = x

addMaybe 5 (Just 3) -- 8
addMaybe 5 (Nothing) -- 5
```

Functors Introduction (1/2)

Functors are essentially types that define fmap, as follows:

```
class Functor f where
  fmap :: (a -> b) -> f a -> f b
```

fmap takes a function that maps from a to b. It lifts the function so that it can operate on the *argument* of a value constructor f. That is, fmap gives us a function that can take f a and return f b!

Functors Introduction (2/2)

Additionally, functors must also satisfy some rules known as the functor laws. Namely,

```
-- functor identity law
fmap id = id
-- functor composition law
fmap (f . g) = fmap f . fmap g
```

We don't have time to go over these in detail, but the motivation behind supporting these laws is that they're necessary to reduce complexity by defining specific but still "small" behavior.

The Maybe Functor

The Maybe functor is defined as follows:

```
instance Functor Maybe where
  fmap f (Just x) = Just (f x)
  fmap f Nothing = Nothing
```

Usage example:

```
Prelude> fmap (+1) $ Just 5

Just 6

Prelude> fmap (+1) $ Nothing

Nothing

Prelude> fmap (\x -> x^x) $ Just 10

Just 10000000000

Prelude> fmap (\x -> x^x) $ Nothing

Nothing
```

Interview Questions: Select Permutations (1/2)

```
-- Given a string and list of case-switching characters,
-- return all versions of the string that can be formed
-- by toggling the case-switching characters.
-- Inputs will be lowercase.
-- "test", ['s'] -> ["test", "test"]
-- "test", ['e', 's'] -> ["test", "test", "test", "test", "]
import Data.Char (toLower, toUpper)

data Token = Normal Char | Switching Char deriving Show

tokenize :: String -> [Char] -> [Token]
tokenize str lst = map toToken str
where toToken x = if elem x lst
then Switching x
else Normal x
```

Interview Questions: Select Permutations (2/2)

```
main = do
   putStrLn $ show $ generate "test" ['s', 'e', 't']
-- Output:
-- ["test", "Test", "Te
```

Interview Questions: Check Leading Ones (1/3)

Interview Questions: Check Leading Ones (2/3)

Interview Questions: Check Leading Ones (3/3)

```
-- Determines if the bytestring bs has
-- at least n leading ones
hasLeadingOnes :: Int -> BS.ByteString -> Bool
hasLeadingOnes n bs = hasLeadingOnes' n bs 0

hasLeadingOnes' :: Int -> BS.ByteString -> Int -> Bool
hasLeadingOnes' n bs count = case BS.uncons bs of
Nothing -> count >= n
Just (b, bs') -> hasLeadingOnes' n bs' newCount
where currentCount = countByteLeadingOnes b
newCount = currentCount + newCount
```

Where to go from here

- Haskell Book by Chris Allen http://haskellbook.com/
- Learn You a Haskell by Miran Lipovača, http://learnyouahaskell.com
- Gentle Introduction To Haskell, version 98 by Paul Hudak, John Peterson, and Joseph Fasel, https://www.haskell.org/tutorial/
- Real World Haskell by Bryan O'Sullivan, http://book.realworldhaskell.org/
- An emporium of resources on the Haskell wiki: https://wiki.haskell.org/ Learning_Haskell

Video learning materials

- Functional Programming Fundamentals by Dr. Erik Meijer (Video Lectures) https://channel9.msdn.com/Series/C9-Lectures-Erik-Meijer-Functional-Programming-Fundamentals/
- Haskell Course Taught by Philip Wadler, https://www.youtube.com/watch?v=AOl2y5uW0mA&list=PLtRG9GLtNcHBv4cuh2w1cz5VsgY6adoc3
- Adventure with Types in Haskell by Simon Peyton Jones, https://www.youtube.com/watch?v=6COvD8oynmI
- Haskell is Not For Production and Other Tales by Katie Miller, https://www.youtube.com/watch?v=mlTO510zO78

Resources to learn about monads (1/2)

There are so many monad tutorials online, so I'll only list a few that I felt really helped me. Your learning style might be different, so please check out many of them! However, here's some ground advice that could really help your venture:

- Know how to define typeclasses
- Be comfortable with type and value constructors and how they interact with functions and typeclasses
- Learn functors, applicative functors, and monoids first. LYAH will go over these beforehand.
- If you just need to *use* (as opposed to construct) a monad, you don't need to fully understand it.

Resources to learn about monads (2/2)

- Read "What is a Monad?" asked by Peter Mortensen on StackOverflow for a variety of good answers, https://stackoverflow.com/questions/44965/what-is-a-monad
- Start with Learn You a Haskell by Miran Lipovača, http://learnyouahaskell.com
- My personal favorite explanation after somewhat grasping the concept was You Could Have Invented Monads! by sigfpe. The practice problems were immensely helpful in consolidating my knowledge, http://blog.sigfpe.com/ 2006/08/you-could-have-invented-monads-and.html
- There are so many monad tutorials that Haskell Wiki gave it a whole section called "Monad tutorials timeline" with dozens of them from over the years, https://wiki.haskell.org/Monad_tutorials_timeline

References (1/2)

These information sources were heavily leaned on for amazing examples, explanations, and more found in these slides.

- Learn You a Haskell, Miran Lipovača, http://learnyouahaskell.com
- Glasgow Haskell Compiler, https://www.haskell.org/ghc/
- "Haskell Type vs Data Constructor", kqr on StackOverflow, https://stackoverflow.com/questions/18204308/haskell-type-vs-data-constructor

• "Haskell function composition (.) and function application (\$) idioms: correct use", ellisbben on StackOverflow, https://stackoverflow.com/questions/3030675/haskell-function-composition-and-function-application-idioms-correct-us

References (2/2)

(continued)

- "Faith, Evolution, and Programming Languages: from Haskell to Java", Philip Wadler, https://www.youtube.com/watch?v=NZeDRs6snm0
- "How to Learn Haskell in Less Than 5 Years", Chris Allen, https://www.youtube.com/watch?v=Bg9ccYzMbxc
- "Gentle Introduction To Haskell, version 98", Paul Hudak, John Peterson, Joseph Fasel, https://www.haskell.org/tutorial/
- "Algebraic data type", Haskell Wiki, https://wiki.haskell.org/Algebraic_data_type
- ECMAScript Pattern Matching Proposal, https://github.com/tc39/proposal-pattern-matching

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