vv214: Introduction: systems of linear equations, graphs, even/oddtown, Markov chains.

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UM-SJTU Joint Institute



May 16, 2019

This week

Today

- 1. Even/oddtown problems.
- 2. Card shuffling/Markov processes.
- 3. Graphs.
- Systems of linear equations: motivation with two historical problems, solutions of a SLE, geometrical interpretations of linear systems.
- 5. Gauss-Jordan elimination.
- 6. Simplest input-output Leontief models.

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- 5. Gauss-Jordan elimination.
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Next class

- 1. Matrices: coefficient, augmented, square, upper and lower triangular, identity.
- 2. Elementary transformations of a matrix.
- 3. Reduced row-echelon form of a matrix.
- 4. Rank of a matrix.



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- Club rule 1: No two clubs have exactly the same membership

$$\forall i \neq j \quad C_i \neq C_j \quad (\Rightarrow m \leq 2^n)$$

 Club rule 2: Every club must have an even number of people

$$\forall i \mid C_i \mid = \text{even}$$



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Example: Let n = 16. Then 16 residents can form $255 = 2^8 - 1$ clubs in an eventown.



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Q: What happens if we shuffle the deck multiple times?

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Definition: The adjacency matrix of the graph G with vertices v_1, \ldots, v_n is an $n \times n$ matrix $A = A_G = (a_{ij})_{n \times n}$ defined by

$$a_{ij} = \begin{cases} 1 & v_i \sim v_j \\ 0 & v_i \nsim v_j \end{cases} \quad \text{and } \quad$$

Theorem(Handshake theorem): The number of people who have made an odd number of hand shakes must be even.

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Theorem (Graph Coloring): If $deg(G) \leq d$, then G can be colored with d+1 colors.

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$$N_{fields}x + N_{fields}N_{days}y = N_{cows}N_{days}z$$



Gauss-Jordan Elimination

- 1. Proceed from equation to equation, from top to bottom.
- 2. For the *i*th equation: let x_j be the leading (pivot) variable. If x_j does not appear in the *i*th equation, swap the *i*th equation with the first equation below that does contain x_j .
- 3. Multiply the *i*th equation by the appropriate scalar so that the coefficient of the leading variable becomes 1.
- 4. Eliminate x_j from all the other equations, above and below the *i*th by subtracting suitable multiples of the *i*th equation from the others.
- 5. Now proceed to the next equation.
- 6. If an equation zero = nonzero emerges in this process, then the system fails to have solutions; the system is inconsistent.

Elementary Row Operations

Recall, that the following algebraic operations do not affect the solution(s) of systems of linear equations:

- 1. Swap equations.
- 2. Divide/multiply an equation by a nonzero scalar.
- 3. Add/substract a multiple of an equation to/from another one.
- ⇒ we can define the following **elementary row operations**:
- 1. Interchange two rows.
- 2. Multiply a row by a nonzero constant,
- 3. Add a multiple of one row to another one.

Metrodorus, Greek Anthology, 6th century A.D.

"Make me a crown weighing 60 minae, mixing gold, bronze, tin, and wrought iron. Let the gold and bronze together form two-thirds, the gold and tin together three-fourths, and the gold and iron three-fifths. Tell me how much gold, tin, bronze, and iron you must put in."

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Hint:

$$\begin{cases} x + y & = \frac{2}{3}60 \\ x + z & = \frac{3}{4}60 \\ x + t & = \frac{3}{5}60 \\ x + y & +z & +t & = 60 \end{cases}$$

Simplest Leontief input-output model

Consider an economy with several, say 2, industries. Q: What output should each of the industries in an economy produce to satisfy the total demand for all products?

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Assume that the consumer demand for their products is, respectively, 1,000 and 780, in millions of dollars per year. What outputs a and b (in millions of dollars per year) should the two industries generate to satisfy the demand?

Definition: A matrix is in reduced row-echelon form *rref* if it satisfies all of the following conditions:

- 1. If a row has nonzero entries, then the first nonzero entry is a 1, called the leading 1 in this row.
- 2. If a column contains a leading 1, then all the other entries in that column are 0.
- 3. If a row contains a leading 1, then each row above it contains a leading 1 further to the left
- \Rightarrow rows of 0's, if any, appear at the bottom of the matrix.

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Definition: The number of leading 1's in the *rref* A is called the rank of the matrix A.

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$$\Rightarrow rref A = \begin{pmatrix} \boxed{1} & 0 \\ 0 & \boxed{1} \\ 0 & 0 \end{pmatrix} \Rightarrow rank A = 2$$

$$2. A = \left(\begin{array}{ccc} 1 & 2 & 3 \\ 4 & 5 & 6 \\ 7 & 8 & 9 \end{array}\right)$$

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Questions

- 1. How many solutions may a system of linear equations have? What is a consistent system?
- 2. What is a matrix? What is a vector?
- 3. What are the elementary transformations of matrices?
- 4. When do we say that a matrix is in reduced row-echelon form?
- 5. What is the Gauss-Jordan elimination?
- 6. What is a rank of a matrix?
- 7. How is the rank of a coefficient matrix related to the number of solutions of a linear system?

The number of solutions and the rank of the coefficient matrix

Consider a linear system of equations n with m variables \Rightarrow the coefficient matrix A of the system is $A_{n \times m}$.

- 1. $rank A \leq n$, $rank A \leq m$
- 2. If rank A = n then the system is consistent.
- 3. If rank A = m then the system has at most one solution.
- 4. If rank A < m then the system either has infinitely many solutions OR inconsistent.

Remarks:

- 1. If n < m then $rank A \le n < m \Rightarrow$ infinitely many OR no solutions
- 2. If n = m and
- a. $rank A = n \Rightarrow$ there exists a unique solution
- b. $rank A < n \Rightarrow$ infinitely many OR no solutions



Matrix Algebra

1. The sum of two matrices $A_{n \times m}$ and $B_{n \times m}$ is the matrix $C_{n \times m}$ s.t.

$$c_{ij} = a_{ij} + b_{ij}, \ i = \overline{1, \ n}, j = \overline{1, \ m}.$$

- 2. The scalar product $\alpha A_{n \times m} = (\alpha a_{ij}), i = \overline{1, n}, j = \overline{1, m}$.
- 3. The product of a row-matrix $(a_1 \ a_2 \ a_3)$ and a column-matrix $\begin{pmatrix} b_1 \\ b_2 \\ b_3 \end{pmatrix}$ is

$$(a_1 \quad a_2 \quad a_3) \left(\begin{array}{c} b_1 \\ b_2 \\ b_3 \end{array} \right) = a_1b_1 + a_2b_2 + a_3b_3$$

Matrix Algebra

4. The product of a matrix $A_{n \times m}$ and a vector $\bar{x} \in \mathbb{R}^m$ is

$$A_{n\times m}\bar{x} = \begin{pmatrix} \bar{w}_1 \\ \bar{w}_2 \\ \vdots \\ \bar{w}_n \end{pmatrix} \bar{x} = (def) \begin{pmatrix} (\bar{w}_1, \bar{x}) \\ (\bar{w}_2, \bar{x}) \\ \vdots \\ (\bar{w}_n, \bar{x}) \end{pmatrix}$$

$$=(prop)(\bar{a}_1 \quad \bar{a}_2 \dots \bar{a}_m)\bar{x}=x_1\bar{a}_1+x_2\bar{a}_2+\dots+x_m\bar{a}_m$$

Next Week

- ► More practice with linear systems/rref/rank.
- ► Matrix Algebra.
- Number fields and linear spaces.
- ► Abelian groups and linear spaces. Cyclic groups.
- Linear combinations and linear dependence/independence.
- Structure of a linear space: basis, dimension.
- Lots of examples!