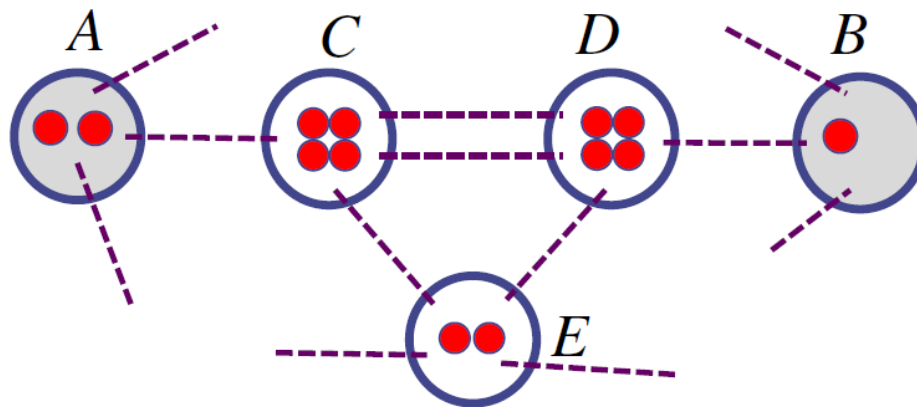


# Data Structures

## Programming Project #3

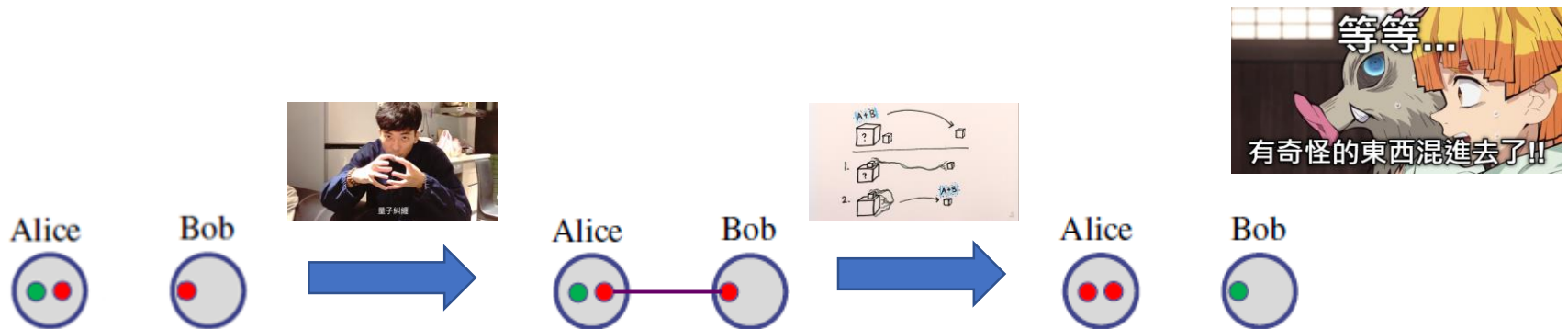
# Data Transmission in Quantum Networks

- A quantum network:
- Nodes has a limited units of **quantum memory**
- Nodes are interconnected with a limited number of **quantum channels** (e.g., optical fiber)



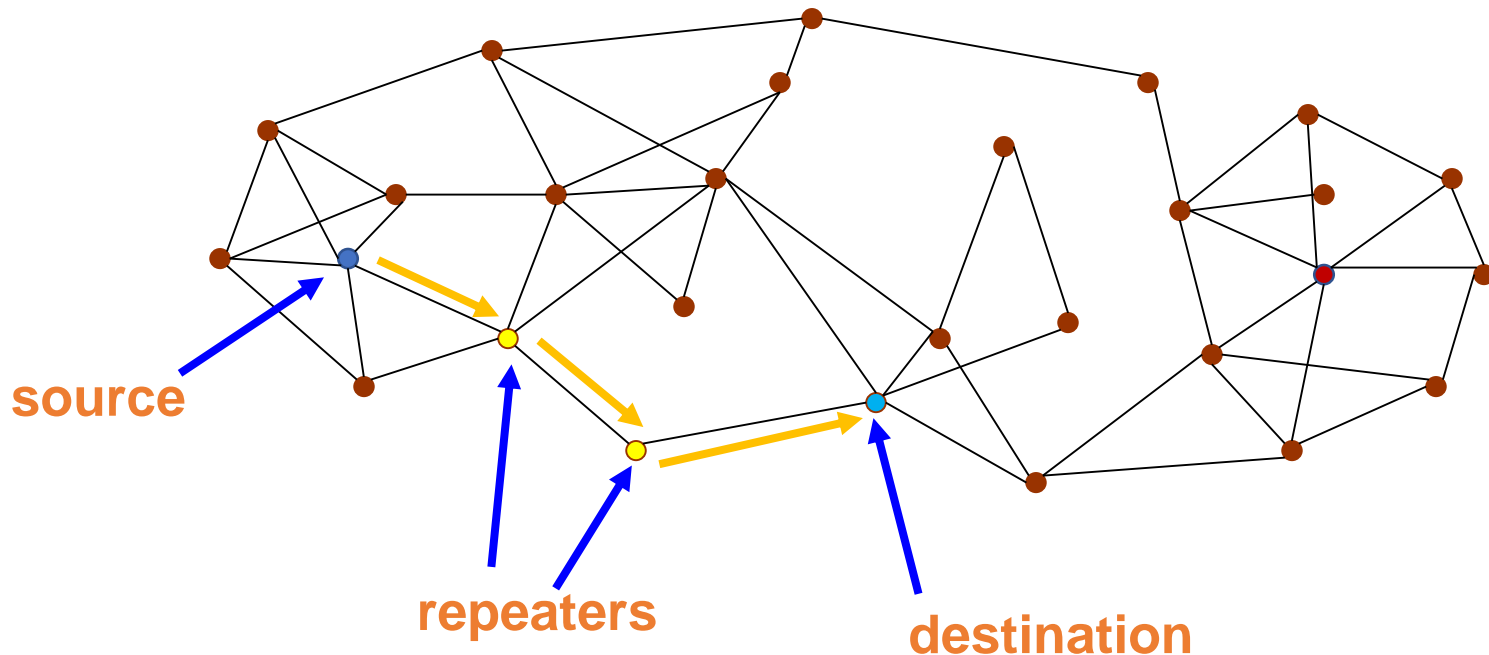
# Data Transmission in Quantum Networks

- **Entangling** (building an entangled link):  
Create an entangled pair between two nodes
- Precondition:  
Two nodes each with a quantum memory are interconnected with a quantum channel



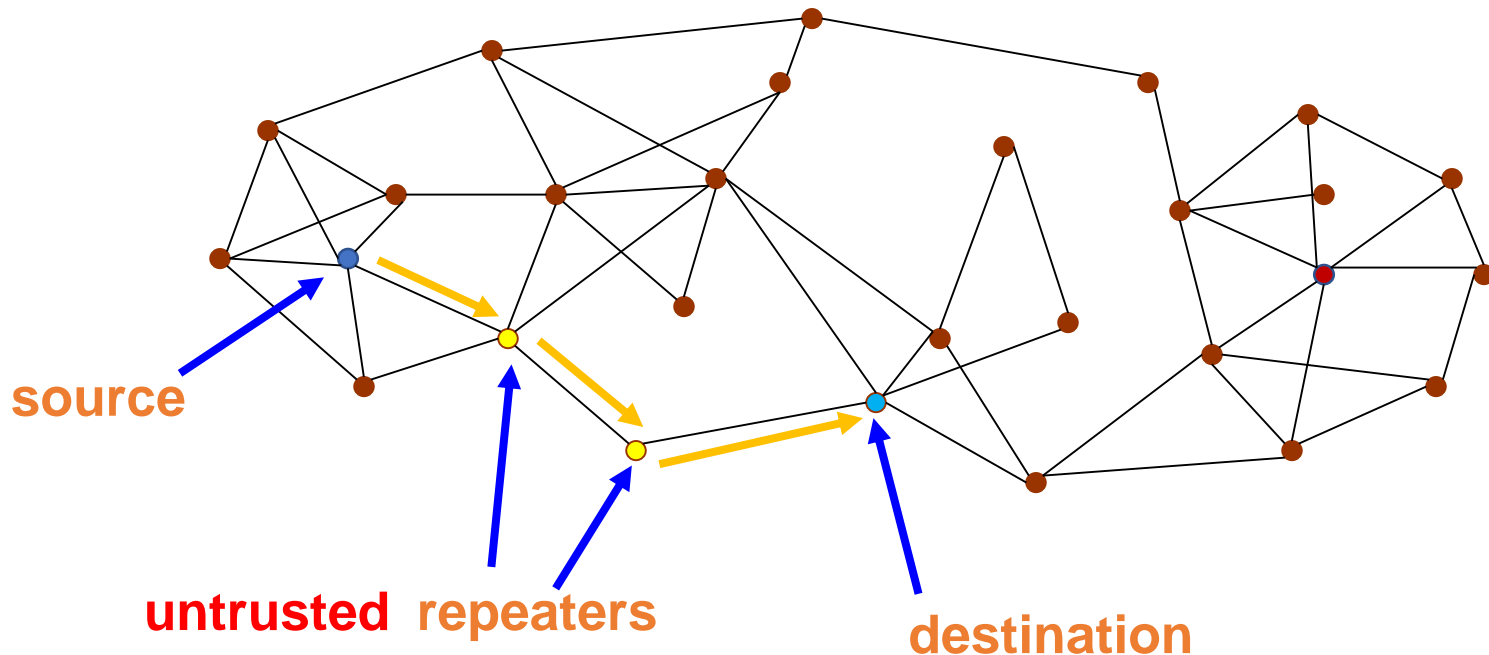
# Data Transmission in Quantum Networks

- The two nodes may be distant from each other
- **Classical networks:**  
Repeaters use **store and forward** to transmit packets from a source to a destination



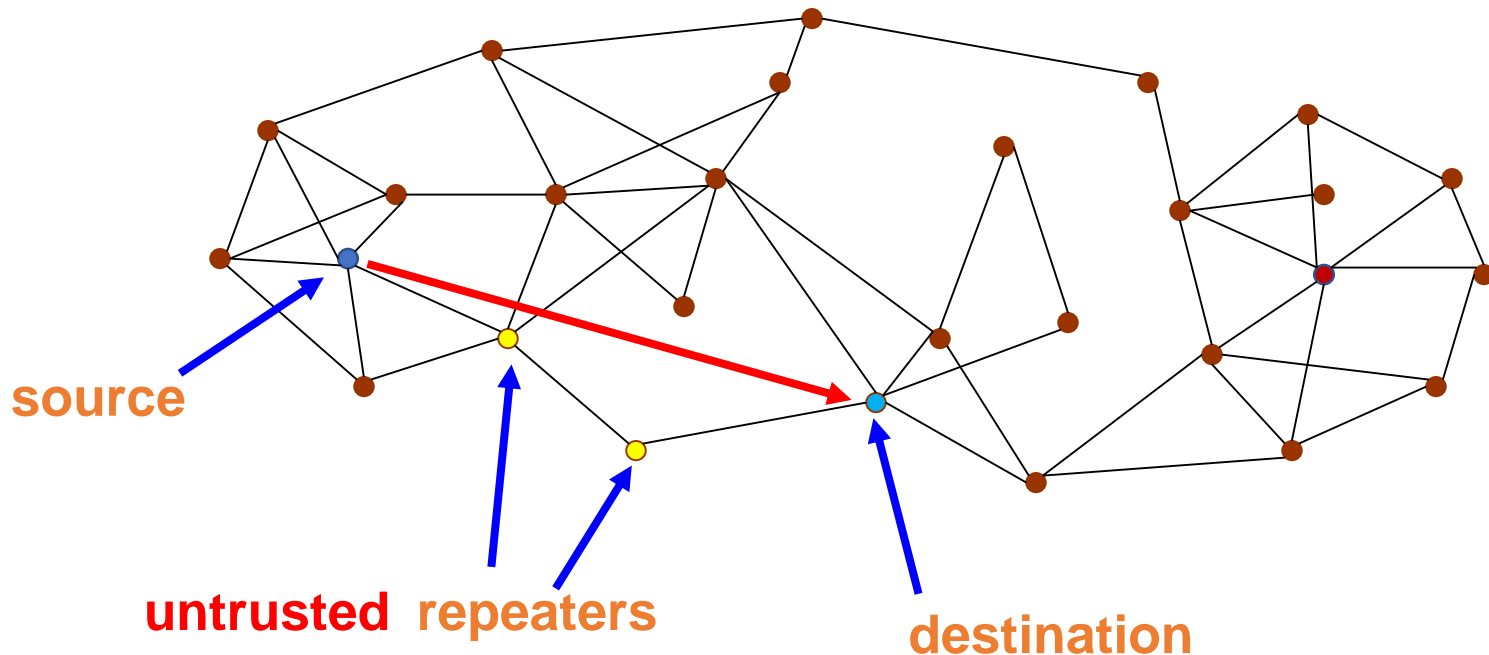
# Data Transmission in Quantum Networks

- However, the data qubit may visit **untrusted** repeaters
- It could be **destroyed, peeked at, or faked**



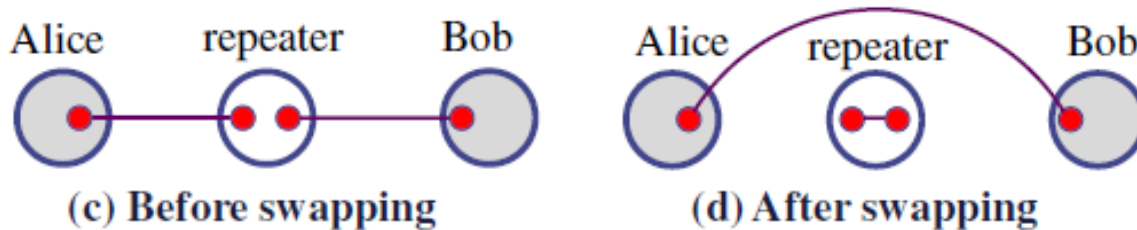
# Data Transmission in Quantum Networks

- Can we send the data qubits **without letting repeaters know**?
- Yes, via **Entanglement Swapping**



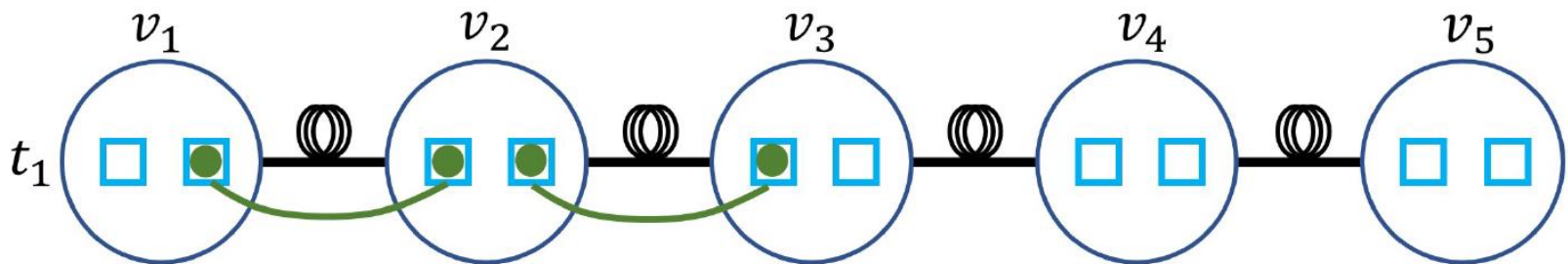
# Entanglement Swapping

- Alice has a data qubit for Bob
- **Entangling**: build the links between Alice and the repeater and between Bob and the repeater
- **Swapping**: build a long-distance entanglement



# Scheduling of Entangling and Swapping

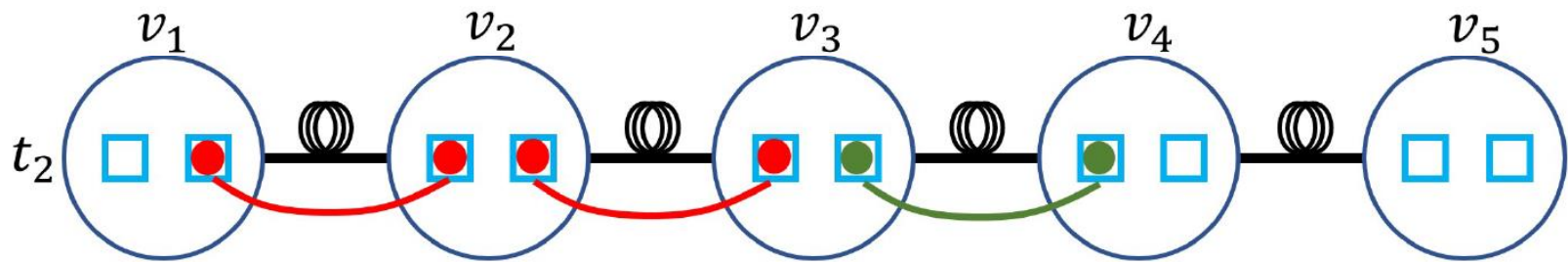
- To construct an entanglement between  $v_1$  and  $v_5$
- One possible scheduling is linear:
- Conduct swapping one hop by one hop from the source to the destination





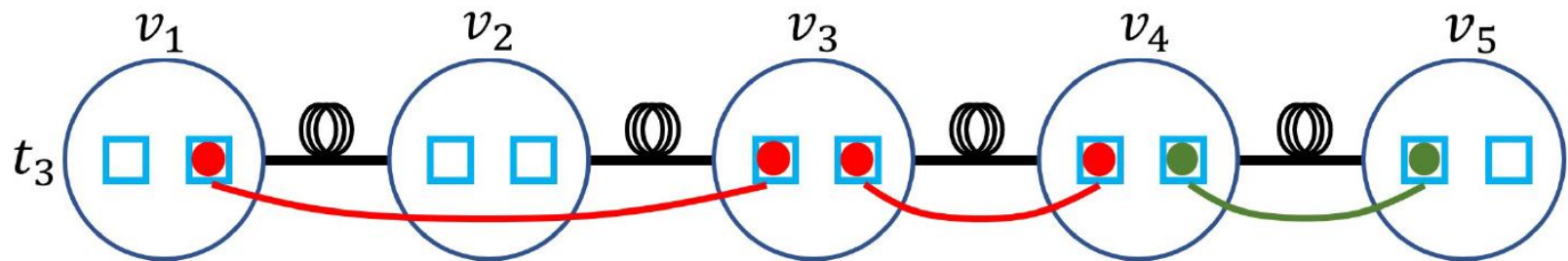
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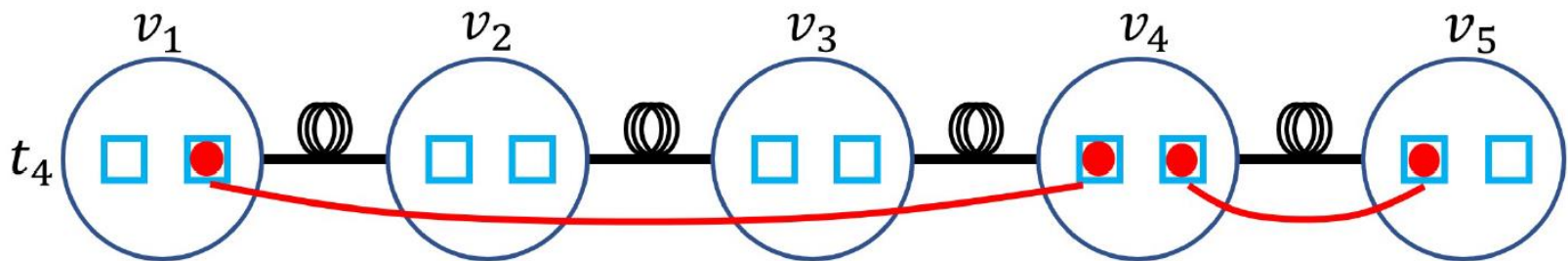
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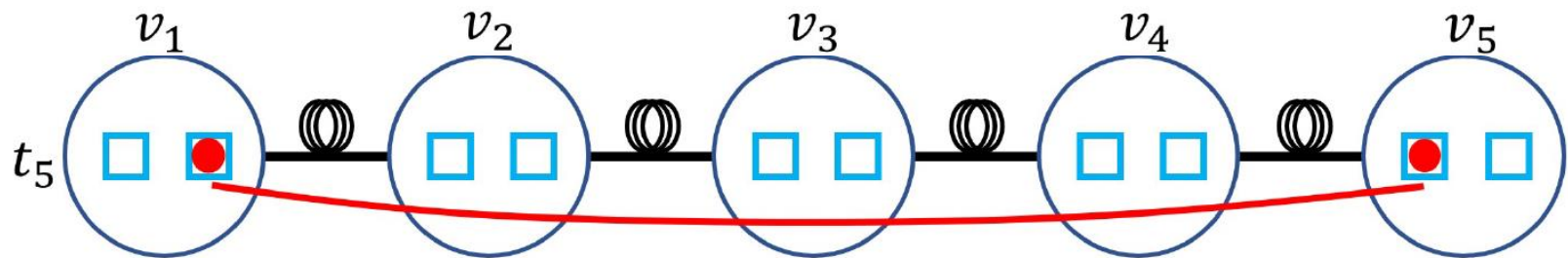
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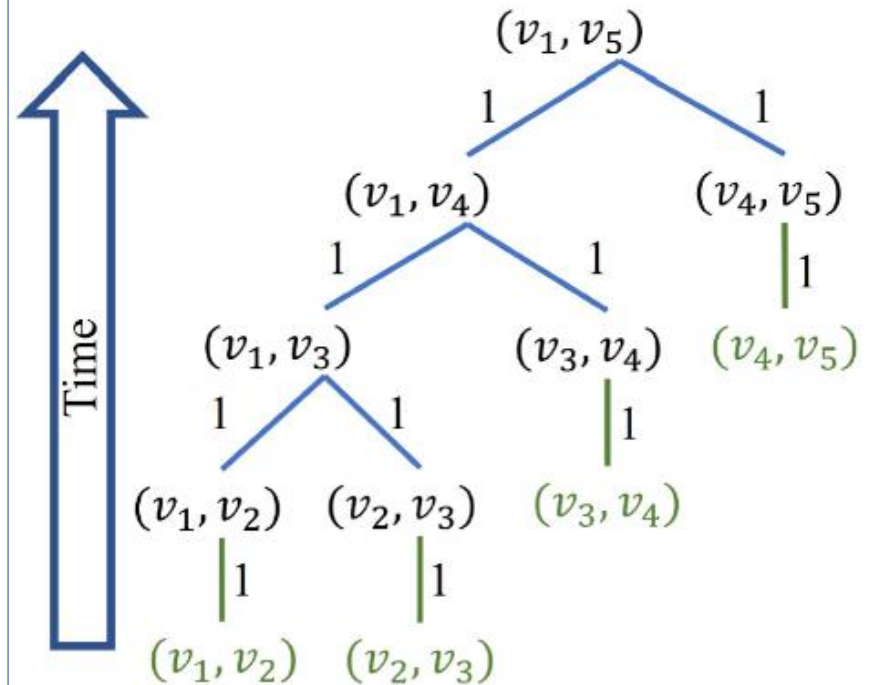
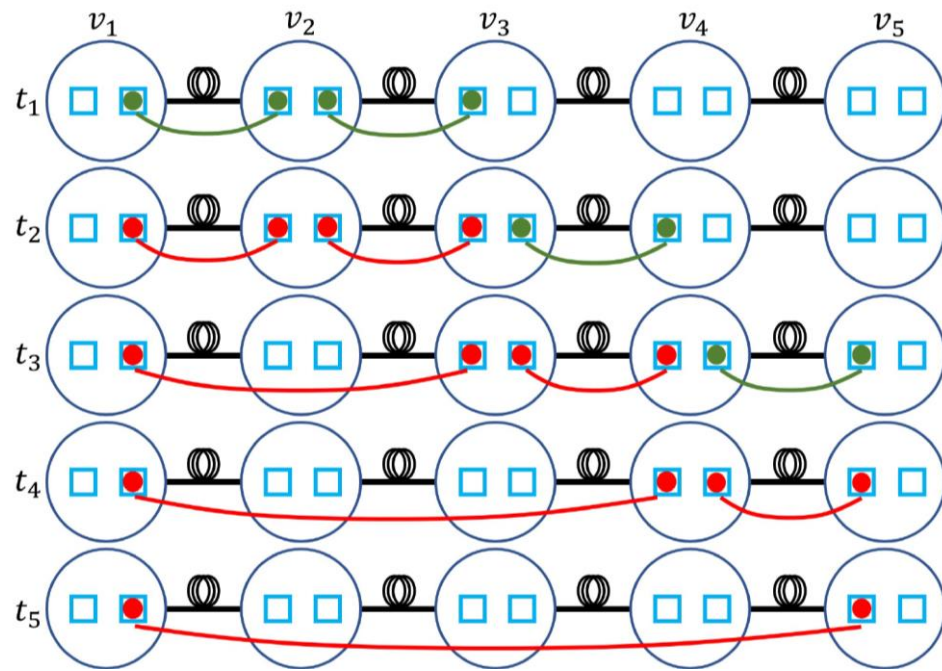
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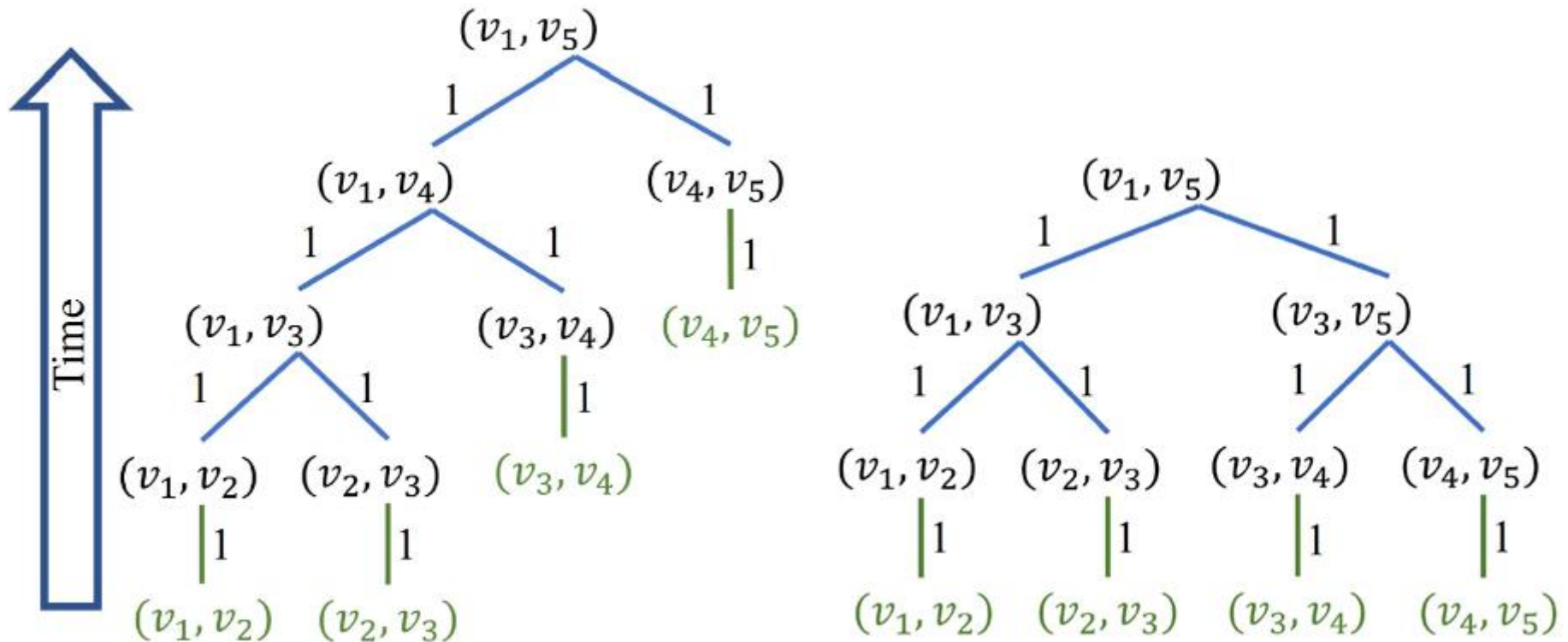
# Scheduling vs Strategy Tree (Binary Tree)

- Any scheduling can be represented by a binary tree



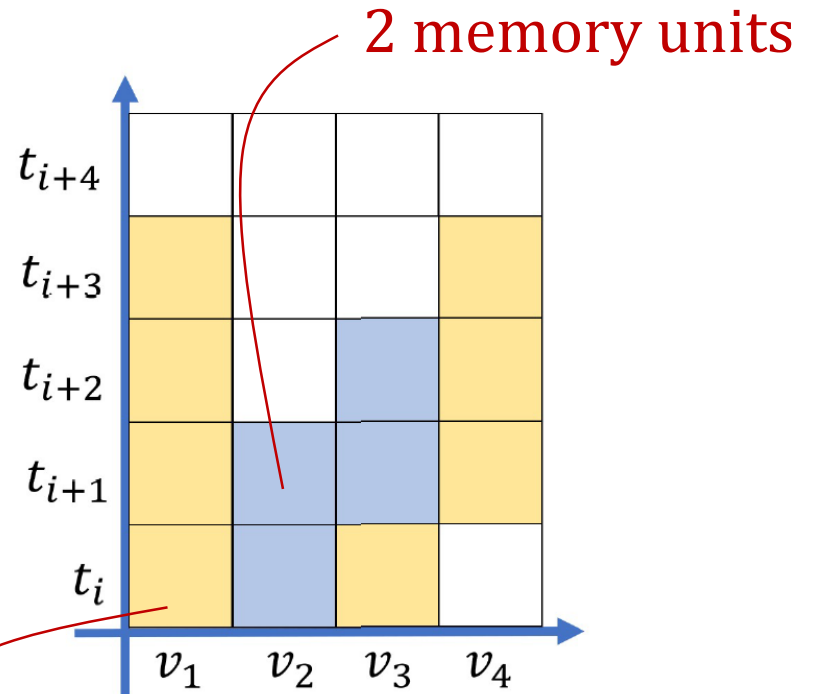
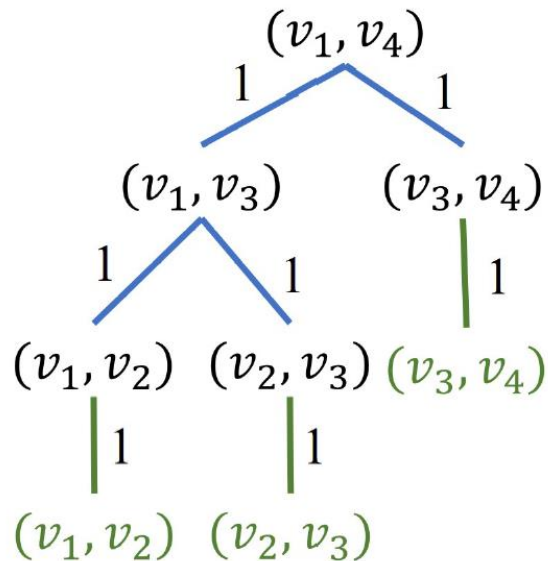
# Scheduling vs Strategy Tree (Binary Tree)

- Different schedulings lead to different binary trees



# Strategy Tree vs Numerology

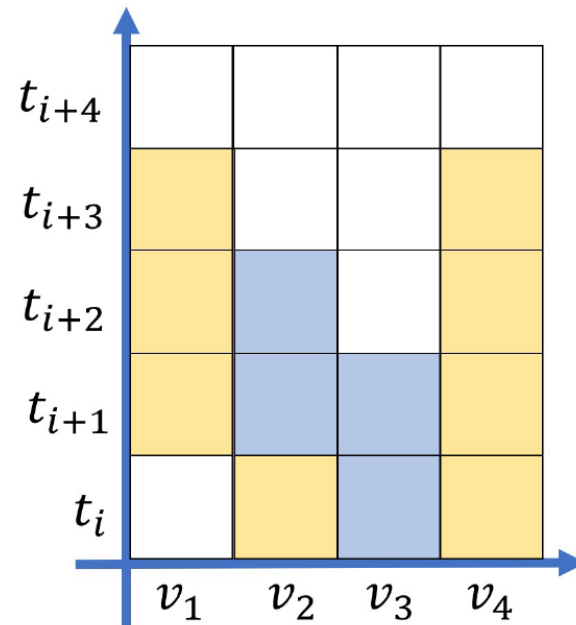
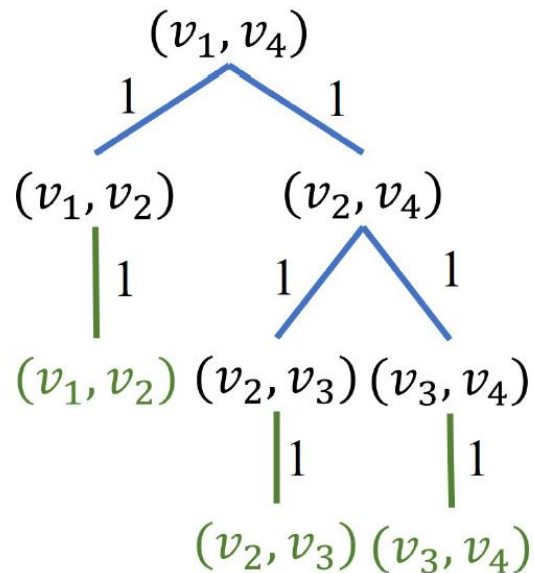
- Different trees lead to different numerologies



1 memory unit

# Strategy Tree vs Numerology

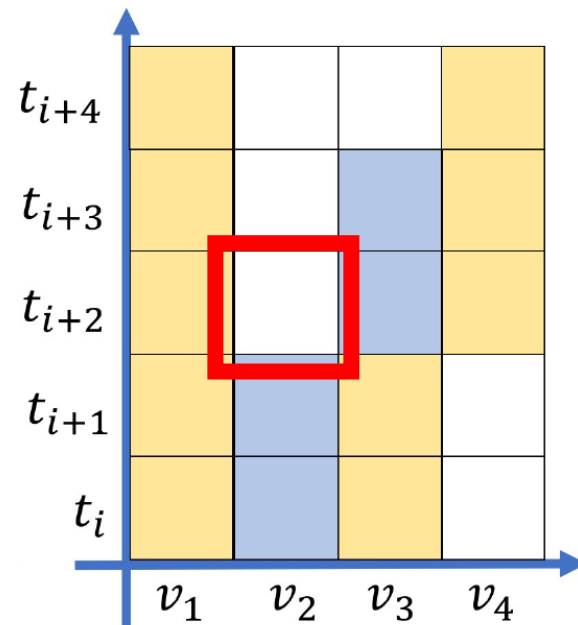
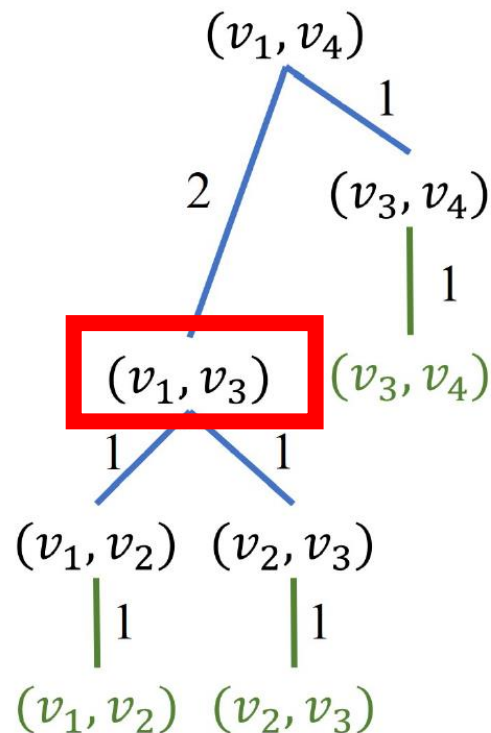
- Different trees lead to different numerologies





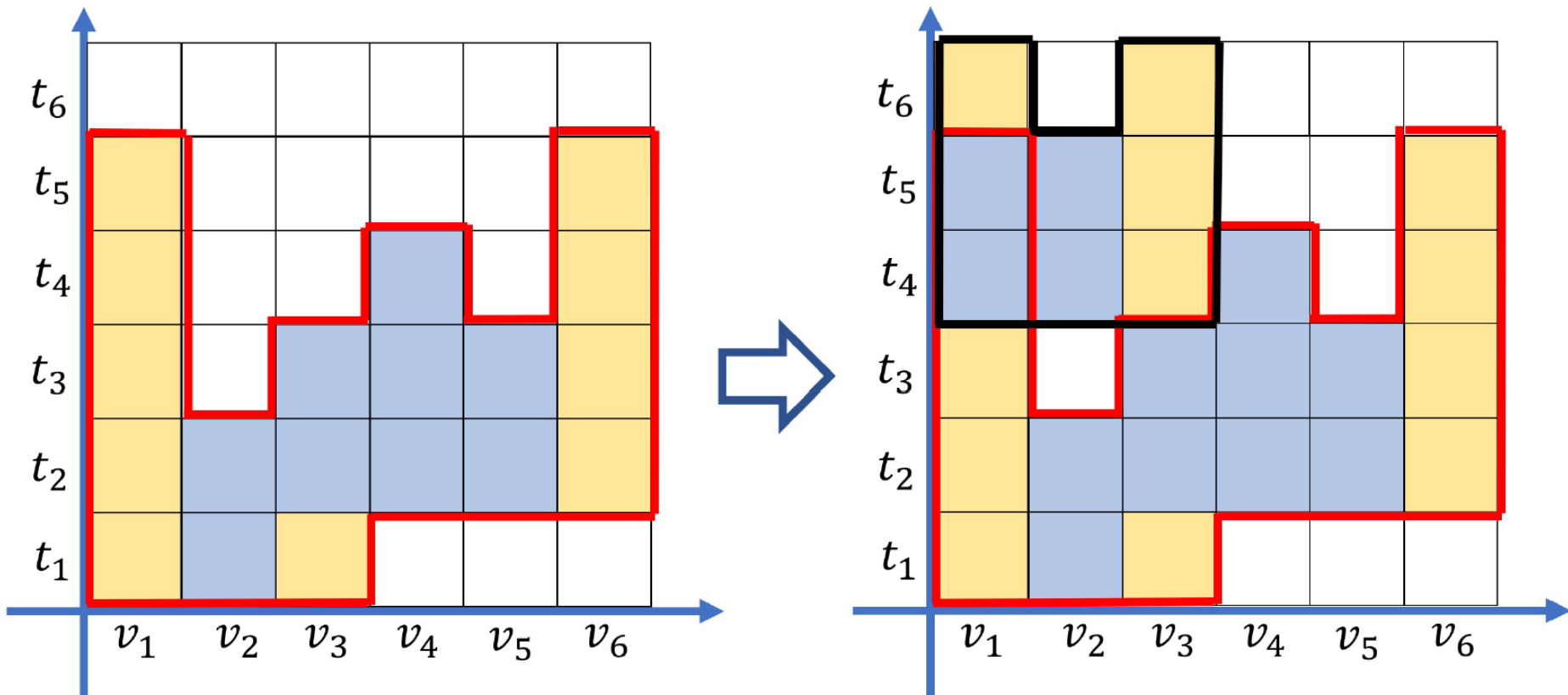
# Strategy Tree vs Numerology

- Different trees lead to different numerologies



# Quantum Resource Allocation Problem

- Multiple numerologies can overlap within  $T$  if the number of memory units is sufficient at that time



# System Model & Problem Formulation

- Given:
- A quantum network with nodes and edges
- Each node has limited memory units
- (Assume fiber channels are always sufficient)
- Multiple source-destination (SD) requests
- Number of time slots ( $T: t_1 \sim t_n$ )
- Goal: maximize # accepted requests
- Constraints:
- Find a path for each accepted SD request
- Find a numerology for each accepted SD request
- Each node's memory size cannot be violated

# Programming Project #3:

## Entanglement Routing & Numerology Selection

- Input:
  - A node-weighted network  $G = (V, E)$
  - Multiple SD requests
  - A number of time slots  $T$
- Procedure:
  - Accept or reject each SD request
  - Find a path with a numerology for accepted one
- Output:
  - The accepted SD requests
  - Their paths and numerologies
- The grade is proportional to # accepted SD requests

# The Competition

- The grade is proportional to **# accepted SD requests**
- **Basic: 60 (deadline)**
  - A baseline solution (see the following pages)
  - (Meet the coding style requirements)
- **Performance ranking** (decided after the deadline)
  - [0%, 30%) (bottom): +0
  - [30%, 50%): + 5
  - [50%, 75%): + 10
  - [75%, 85%): + 15
  - [85%, 90%): + 20
  - [90%, 95%): + 25
  - [95%, 100%] (top): + 30
- **Homework assistant** (superb deadline)
  - +10

# The Competition

- The grade
- **Basic: 60** (requests)
  - A baseline
  - (Meet the
- **Performance**
  - [0%, 30%
  - [30%, 50%
  - [50%, 75%
  - [75%, 85%
  - [85%, 90%
  - [90%, 95%
  - [95%, 100%
- **Homework assistant** (superb deadline)
  - +10



We have  
TIME LIMIT!



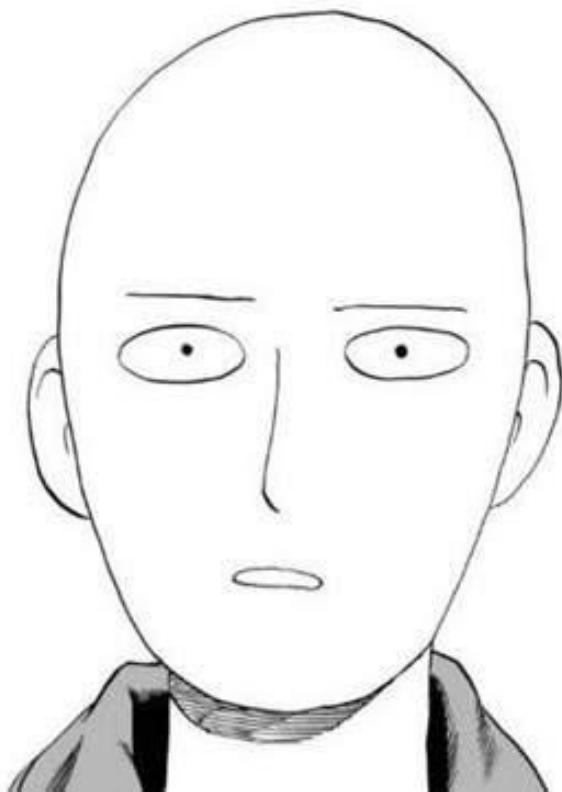
# The Competition

- The grade
- **Basic: 60** (requests
- A baseline
- (Meet the
- **Performance** (deadline)
- [0%, 30%
- [30%, 50%
- [50%, 75%
- [75%, 85%
- [85%, 90%
- [90%, 95%
- [95%, 100%
- **Homework assistant** (superb deadline)
- +10



相信你們在做完作業以後

也變強了



禿了



# The Baseline Algorithm

- Sequentially find the **minimum-hop path** for each input SD request via **BFS (on the original graph)**
- If there is a tie, select the one that visits the node with a smaller ID first
  - 10 -> 4 -> **11** -> 15 -> 12 vs 10 -> 4 -> **9** -> 17 -> 12
  - Select 10 -> 4 -> **9** -> 17 -> 12
- Only use **complete binary tree** (recall that in Chap. 5) where the bottom leaf(s) is **at time  $t_1$**  and **side length = 1**
- Examine whether the above path and numerology can be accommodated in the network
  - If yes, then accept the SD request
  - Otherwise, reject the SD request
- Repeat the above actions until all SD requests are examined

Input Sample: use scanf

Format:

#Nodes #Links #TimeSlots #Req

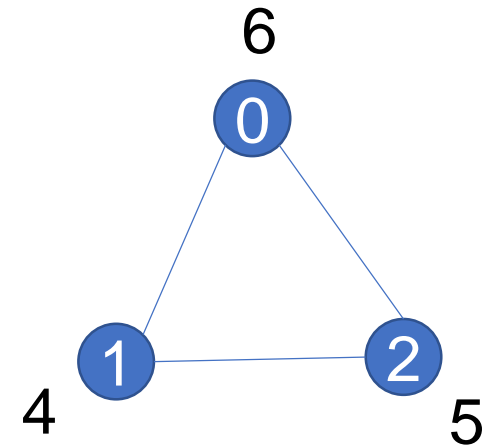
NodeID #QuantumMemories

...

LinkID LinkEnd1 LinkEnd2

...

ReqID ReqSrc ReqDst



Ex:

3 3 7 2

0 6

1 4

2 5

0 0 1

1 0 2

2 1 2

0 0 2

1 1 2

# Output Sample: use printf

Format:

#AccReq

ReqID ReqSrc Rep1 Rep2... ReqDst

Tree Structure ([detailed in the next slide](#))

...

# Output Sample: use printf

Format:

Tree Structure:  
use a postorder traversal

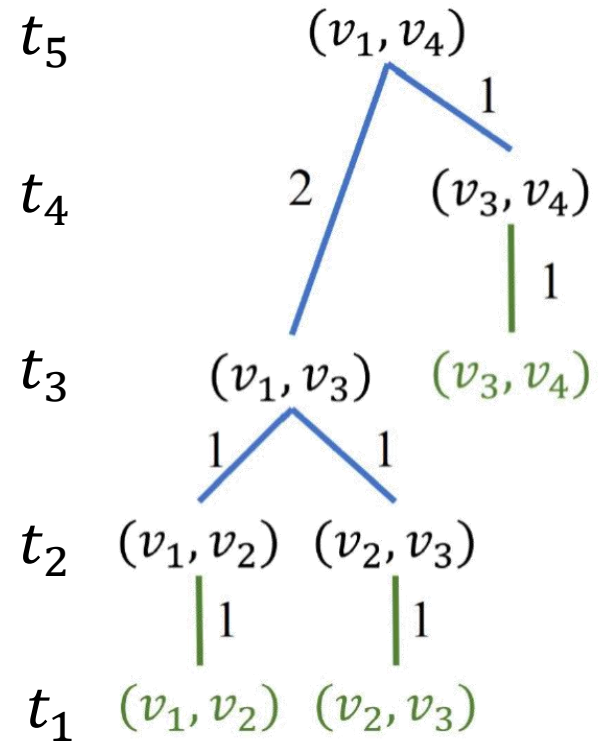
v1\_ID v2\_ID t2\_time

V2\_ID v3\_ID t2\_time

v1\_ID v3\_ID v1\_ID v2\_ID v2\_ID v3\_ID t3\_time

v3\_ID v4\_ID t4\_time

v1\_ID v4\_ID v1\_ID v3\_ID v3\_ID v4\_ID t5\_time



Output Sample: use printf

Format:

Tree Structure:  
use a postorder traversal

10 8 4

8 15 4

10 15 10 8 8 15 5

15 20 6

10 20 10 15 15 20 7

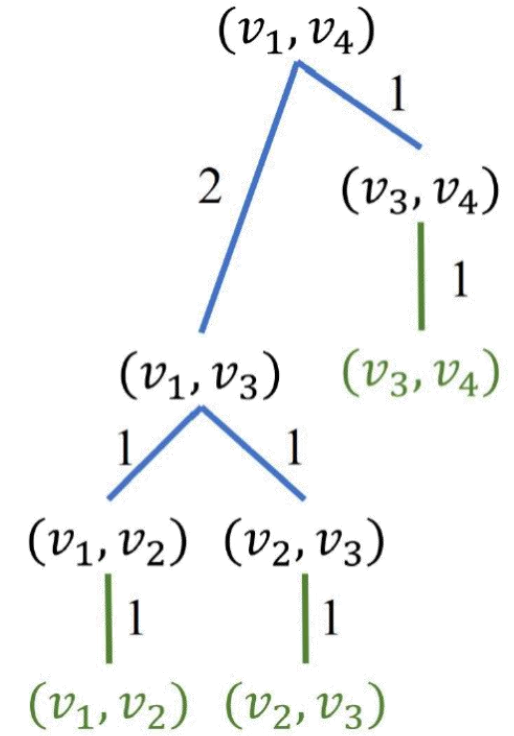
$t_5 = 7$

$t_4 = 6$

$t_3 = 5$

$t_2 = 4$

$t_1 = 3$



v1\_ID = 10

v2\_ID = 8

v3\_ID = 15

v4\_ID = 20

# Note

- Superb deadline: 12/14 Thu
- Deadline: 12/21 Thu
- Pass the test of our [online judge](#) platform
- Submit your code to [E-course2](#)
- Demonstrate your code [remotely or in person](#) with TA
- [C Source code \(i.e., only .c\)](#)
- Show a good programming style

It is worth noting that...

- The processes of entangling and swapping on the path may fail (although **we don't consider it**)
- In addition, strategy tree height also affects quality of entangled pair (although **we don't consider it**)
- For example:
  - The **success probability of entangling** is  $p = e^{-\alpha \cdot L}$ , where  $\alpha$  is a constant of the fiber material and  $L$  is the distance between the two nodes
  - The **success probability of swapping** is  $q$ , which is related to the technique