## Changing Linux Page Replacement Algorithm

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#### Overview

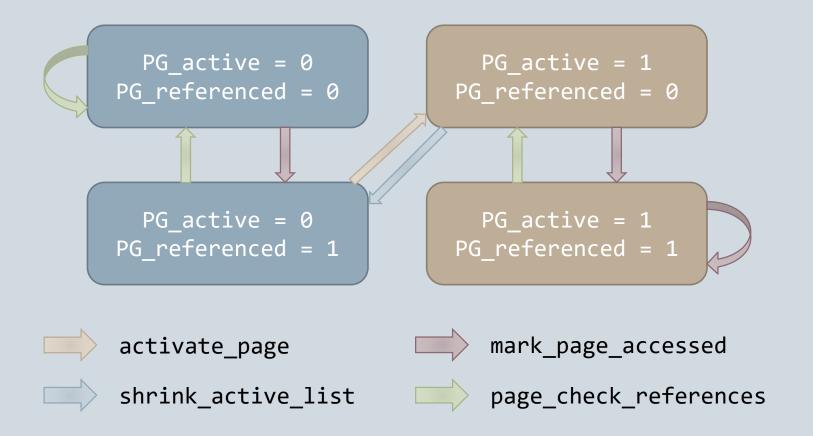
- Study of the Original Algorithm
- Changing to a new LRU Algorithm
- Changing to an LFU Algorithm

# Study of the Original Algorithm

#### Basic Ideas

- Reclaim 'unused' pages only.
- Associate a counter storing the age of the page with each page in RAM.

## State Diagram

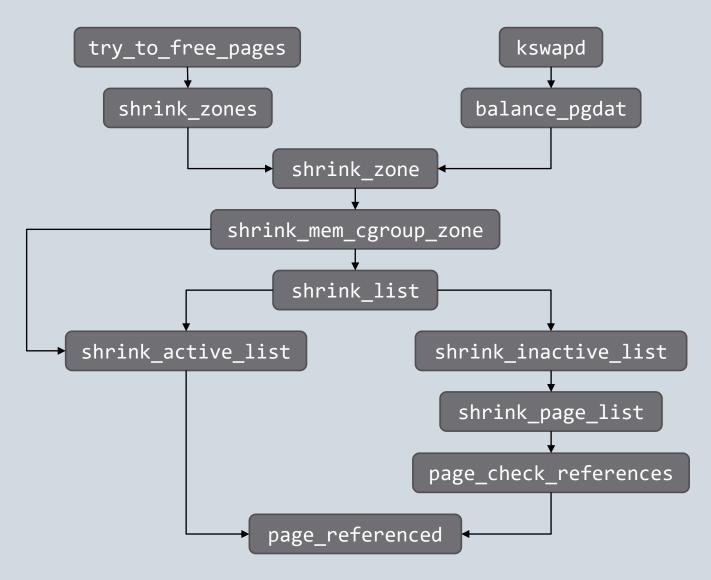


## Condition of Reclaiming

- Low on memory: The kernel has difficulty allocating new pages.
- Periodic: Kernel threads check periodically whether
   free pages are above a certain threshold.

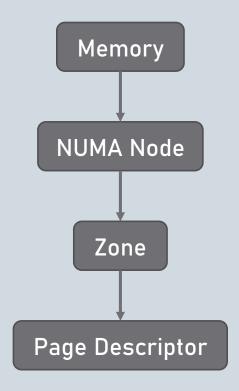
#### LOW ON MEMORY RECLAIMING

#### PERIODIC RECLAIMING



#### Memory-Related Structs

- NUMA: The time required to access different memory locations is different.
- Zone: Accounting for hardware constraints of the type of data stored in pages.



# Changing to a New LRU Algorithm

#### Scope of Work

- The algorithm involved just select the candidates to be replaced.
- We just need to change how the status of the pages are transitioned.

## Comparison of Two Algorithms

Task	Original	New LRU	
Renew page	If PG_referenced is 0, set it. Otherwise clear it and add to active list.	Right shift PG_referenced and add by an offset.	
Age page		• • —	

Condition	Original	New LRU
Add to active list	PG_referenced is 1 and the page is referenced recently.	PG_referenced is above some threshold.
Add to inactive list	PG_referenced is 0 and the page is not referenced recently.	PG_referenced is below some threshold.

#### **Parameters**

- •The offset to add to PG\_referenced when page is accessed 20.
- The threshold of moving pages between lists  $2^T$ .
- $\bullet O T$  may affect list sizes.

## **Test Results**

New LRU Algorithm

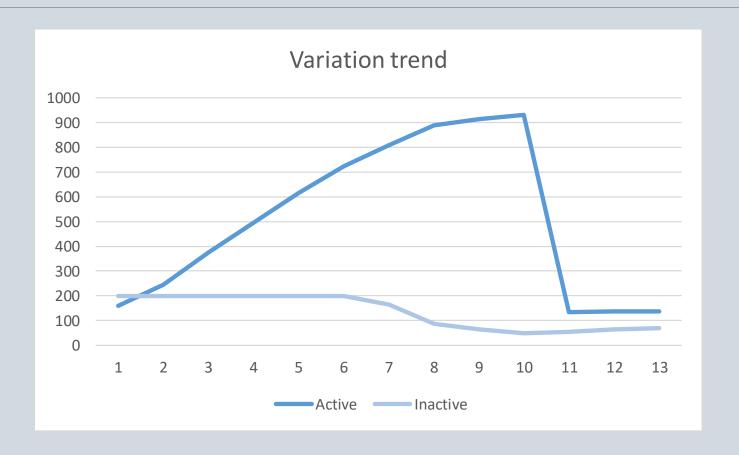
#### Test the Algorithm

- Write a program to occupy as much memory as possible.
- Access /proc/meminfo periodically to get the sizes
   of active and inactive lists.

#### **Notice**

- •kswapd invokes lowmemkiller, which kills processes to reclaim memory space. The test process is then terminated.
- List sizes change differently for mmap and malloc.

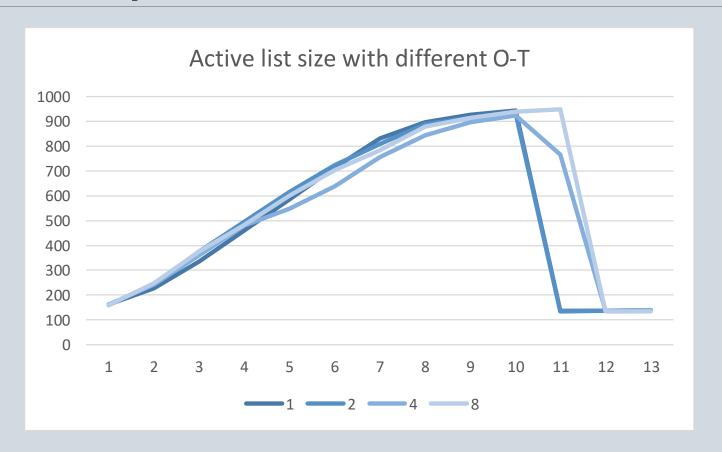
#### **Trend**



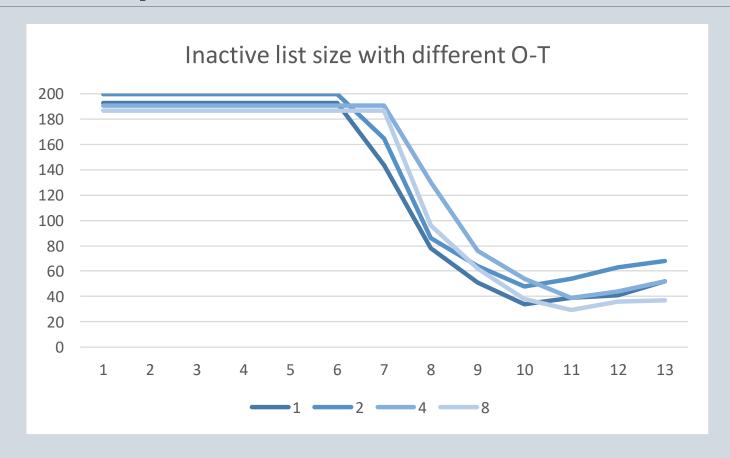
## **Trend**

Time	Trend	Explanation	
1-6s	Active list continues to grow. Inactive list remains.	Once a page is accessed, it is marked active.	
6-10s	Active list grows more slowly. Inactive list begins to shrink.	Some pages are reused.	
10-11s	Active list drops dramatically.	Some processes, including the test program, are killed.	

## The Impact of O-T



## The Impact of O-T



#### Observation

- Hard to find correlation between sizes of LRU lists and parameters.
- The difference in list sizes is more of randomness than some inherent mechanism.

## To an LFU Algorithm

With a few modifications to the LRU algorithm.

#### Differences from LRU

- LRU: A page is immediately added to active list when it's accessed.
- LFU: A page is added to active list only when it's accessed frequently enough.

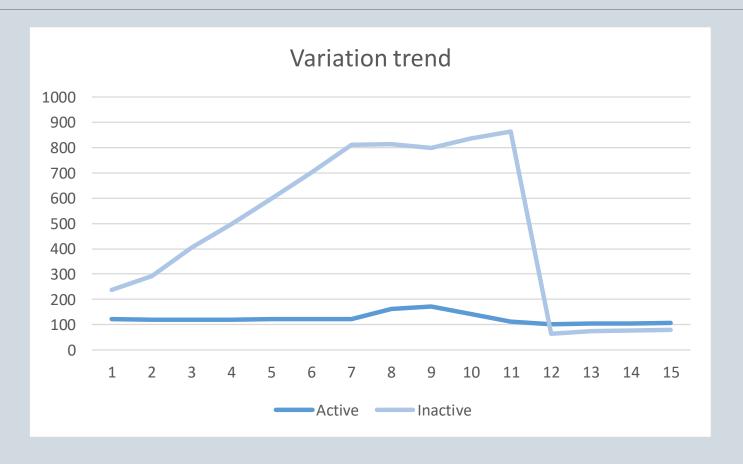
#### Implementation Details

- When a page is accessed, left shift PG\_referenced and add by 1.
- •Maximal PG\_referenced: necessary.

## **Test Results**

LFU Algorithm

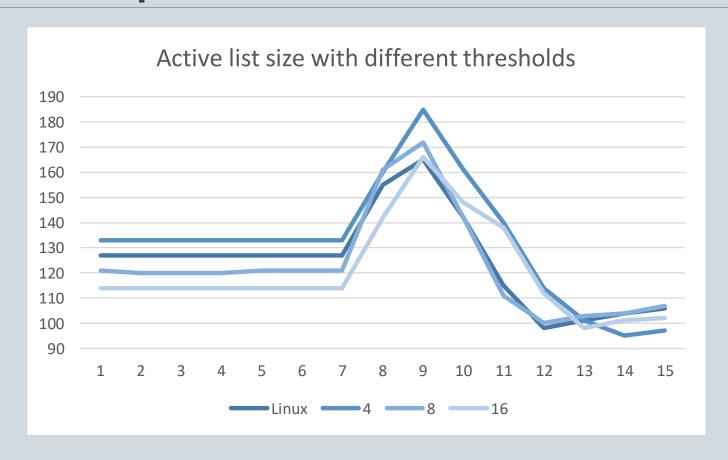
#### **Trend**



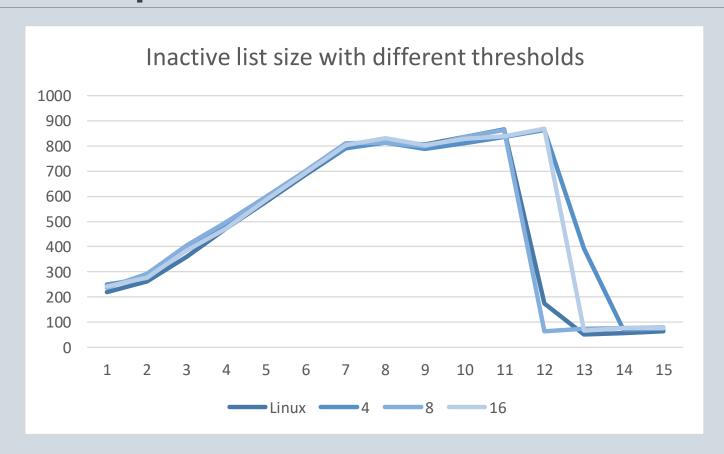
## **Trend**

Time	Trend	Explanation	
1-6s	Inactive list continues to grow. Active list remains.	Test program access every page only once.	
7-9s	Inactive list grows more slowly. Active list begins to grow.	Some pages are reused.	
9-11s	Inactive list grows, while active list shrinks.	Some pages are aged in the replacement algorithm.	
11-12s	Inactive list drops dramatically.	Some processes, including the test program, are killed.	

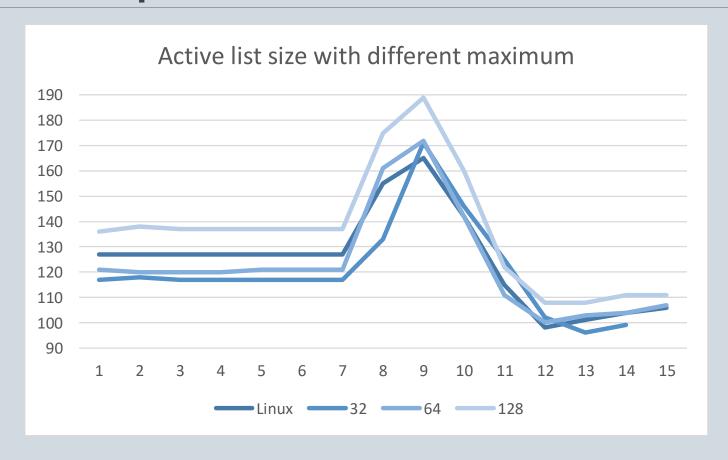
## The Impact of Threshold



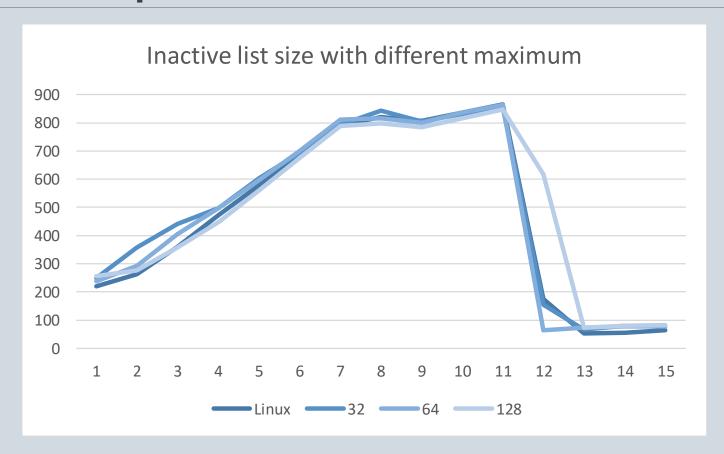
## The Impact of Threshold



## The Impact of Maximum



## The Impact of Maximum



#### Observation

- •The smaller the threshold, the larger the active list.
- •The higher the maximum, the larger the active list.
- •The size of active list varies more greatly in the LFU algorithm, compared with the original.
- •The parameters seem to have little impact on the size of inactive list.

#### **Choice of Parameters**

Criterion	Threshold	Maximum
Similar to the original	8	64
Smaller active list size	16 or higher	Threshold × 2

#### Reference

 Understanding the Linux Kernel. Daniel Bovet and Marco Cesati. O'Reilly Media.

## The End

**Thanks**