

in drivers/tty/n_tty.c

```

1. struct tty_ldisc_ops tty_ldisc_N_TTY = {
2.     .magic          = TTY_LDISC_MAGIC,
3.     .name           = "n_tty",
4.     .open           = n_tty_open,
5.     .close          = n_tty_close,
6.     .flush_buffer   = n_tty_flush_buffer,
7.     .chars_in_buffer = n_tty_chars_in_buffer,
8.     .read           = n_tty_read,
9.     .write          = n_tty_write,
10.    .ioctl           = n_tty_ioctl,
11.    .set_termios     = n_tty_set_termios,
12.    .poll            = n_tty_poll,
13.    .receive_buf     = n_tty_receive_buf,
14.    .write_wakeup    = n_tty_write_wakeup,
15.    .fasync          = n_tty_fasync,
16.    .receive_buf2    = n_tty_receive_buf2,
17. };
18.
19.
20. static void n_tty_receive_buf(struct tty_struct *tty, const unsigned char *cp,
21.                             char *fp, int count)
22. {
23.     n_tty_receive_buf_common(tty, cp, fp, count, 0);
24. }
25.
26. static int n_tty_receive_buf2(struct tty_struct *tty, const unsigned char *cp,
27.                              char *fp, int count)
28. {
29.     return n_tty_receive_buf_common(tty, cp, fp, count, 1);
30. }
31.
32. static int
33. n_tty_receive_buf_common(struct tty_struct *tty, const unsigned char *cp,
34.                         char *fp, int count, int flow)
35. {
36.     struct n_tty_data *ldata = tty->disc_data;
37.     int room, n, rcvd = 0;
38.
39.     down_read(&tty->termios_rwsem);
40.
41.     while (1) {
42.         room = receive_room(tty);           ❶
43.         n = min(count, room);
44.         if (!n) {
45.             if (flow && !room)               ❷
46.                 ldata->no_room = 1;
47.             break;
48.         }
49.         __receive_buf(tty, cp, fp, n);      ❸
50.         cp += n;
51.         if (fp)
52.             fp += n;
53.         count -= n;

```

```

54.         rcvd += n;
55.     }
56.
57.     tty->receive_room = room;
58.     n_tty_check_throttle(tty);           ④
59.     up_read(&tty->termios_rwsem);
60.
61.     return rcvd;
62. }

```

unsigned char *cp中是data,而char *fp是每个char对应的flag

①

得到n_tty line discipline接收buffer (ldata->read_buf)的空闲空间

②

n = 0, 三种可能, 一是ldata->read_buf[]还有空闲空间, 但已经接收完count的data; 二是虽然接收的data还未到count, 但ldata->read_buf[]已经满了; 三是接收完所有的count的data, 并且ldata->read_buf[]也满了。

这里如果是n_tty_receive_buf(), 即flow = 0, 则直接break处while loop; 而对于n_tty_receive_buf2(),

即flow = 1, 那么在ldata->read_buf[]已经满了的情况下, ldata->no_room = 1.

对ldata->no_room的access

```

1. static void n_tty_set_room(struct tty_struct *tty)
2. {
3.     struct n_tty_data *ldata = tty->disc_data;
4.
5.     /* Did this open up the receive buffer? We may need to flip */
6.     if (unlikely(ldata->no_room) && receive_room(tty)) {
7.         ldata->no_room = 0;
8.
9.         WARN_RATELIMIT(tty->port->itty == NULL,
10.            "scheduling with invalid itty\n");
11.         /* see if ldisc has been killed - if so, this means that
12.          * even though the ldisc has been halted and ->buf.work
13.          * cancelled, ->buf.work is about to be rescheduled
14.          */
15.         WARN_RATELIMIT(test_bit(TTY_LDISC_HALTED, &tty->flags),
16.            "scheduling buffer work for halted ldisc\n");
17.         queue_work(system_unbound_wq, &tty->port->buf.work);
18.     }
19. }

```

即在n_tty_receive_buf()接收data的情况下，ldata->no_room总是被初始化为0，没有机会变1，那么n_tty_set_room()等于是empty function。而在n_tty_receive_buf2()接收data的情况下，

if (unlikely(ldata->no_room) && receive_room(tty))

即在n_tty_receive_buf2()中ldata->read_buf[]已经满，而此时，在invoke n_tty_set_room()时ldata->read_buf[]又有空闲空间了，那么除了ldata->no_room = 0外，运行->buf.work?

③

根据tty的mode来填充n_tty line discipline的buffer,ldata->read_buf[]

④

这里是一种line discipline与low-level driver间的flow control。如果line discipline接收

buffer空闲空间太小了，则要让low-level driver控制一下传入data的速度。

```
1. static int receive_room(struct tty_struct *tty)
2. {
3.     struct n_tty_data *ldata = tty->disc_data;
4.     int left;
5.
6.     if (I_PARMRK(tty)) { ①
7.         /* Multiply read_cnt by 3, since each byte might take up to
8.          * three times as many spaces when PARMRK is set (depending on
9.          * its flags, e.g. parity error). */
10.        left = N_TTY_BUF_SIZE - read_cnt(ldata) * 3 - 1;
11.    } else ②
12.        left = N_TTY_BUF_SIZE - read_cnt(ldata) - 1;
13.
14.    /*
15.     * If we are doing input canonicalization, and there are no
16.     * pending newlines, let characters through without limit, so
17.     * that erase characters will be handled. Other excess
18.     * characters will be beeped.
19.     */
20.    if (left <= 0) ③
21.        left = ldata->icanon && ldata->canon_head == ldata->read_tail;
22.
23.    return left;
24. }
```

①

if (I_PARMRK(tty))

判断接收的data是否有校验错，这里要多预留空间，没找到详细的资料解释 × 3 的原因。

②

返回空闲空间

③

?

```

1. static void __receive_buf(struct tty_struct *tty, const unsigned char *cp,
2.                          char *fp, int count)
3. {
4.     struct n_tty_data *ldata = tty->disc_data;
5.     bool preops = I_ISTRIP(tty) || (I_IUCLC(tty) && L_IEXTEN(tty));
6.
7.     if (ldata->real_raw)
8.         n_tty_receive_buf_real_raw(tty, cp, fp, count);
9.     else if (ldata->raw || (L_EXTPROC(tty) && !preops))
10.        n_tty_receive_buf_raw(tty, cp, fp, count);
11.     else if (tty->closing && !L_EXTPROC(tty))
12.        n_tty_receive_buf_closing(tty, cp, fp, count);
13.     else {
14.         if (ldata->lnext) {
15.             char flag = TTY_NORMAL;
16.
17.             if (fp)
18.                 flag = *fp++;
19.             n_tty_receive_char_lnext(tty, *cp++, flag);
20.             count--;
21.         }
22.
23.         if (!preops && !I_PARMRK(tty))
24.             n_tty_receive_buf_fast(tty, cp, fp, count);
25.         else
26.             n_tty_receive_buf_standard(tty, cp, fp, count);
27.
28.         flush_echoes(tty);
29.         if (tty->ops->flush_chars)
30.             tty->ops->flush_chars(tty);
31.     }
32.
33.     if ((!ldata->icanon && (read_cnt(ldata) >= ldata->minimum_to_wake)) ||
34.        L_EXTPROC(tty)) {
35.         kill_fasync(&tty->fasync, SIGIO, POLL_IN);
36.         if (waitqueue_active(&tty->read_wait))
37.             wake_up_interruptible_poll(&tty->read_wait, POLLIN);
38.     }
39. }

```

对raw mode的处理比较简单，因为实质上line discipline不需要处理什么，而对canon mode则由于没有详细处理的资料，所以只能看懂个大概。

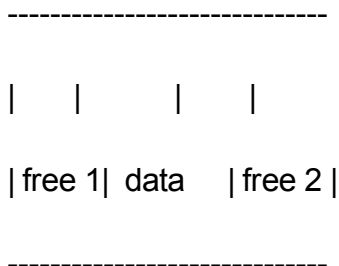
```

1. static void
2. n_tty_receive_buf_real_raw(struct tty_struct *tty, const unsigned char *cp,
3.                           char *fp, int count)
4. {
5.     struct n_tty_data *ldata = tty->disc_data;
6.     size_t n, head;
7.
8.     head = ldata->read_head & (N_TTY_BUF_SIZE - 1);           (A)
9.     n = N_TTY_BUF_SIZE - max(read_cnt(ldata), head);
10.    n = min_t(size_t, count, n);
11.    memcpy(read_buf_addr(ldata, head), cp, n);
12.    ldata->read_head += n;
13.    cp += n;
14.    count -= n;
15.
16.    head = ldata->read_head & (N_TTY_BUF_SIZE - 1);           (B)
17.    n = N_TTY_BUF_SIZE - max(read_cnt(ldata), head);
18.    n = min_t(size_t, count, n);
19.    memcpy(read_buf_addr(ldata, head), cp, n);
20.    ldata->read_head += n;
21. }

```

n_tty line discipline对接收到的data不做任何处理，直接放入ldata->read_buf[]。

这里分(A)，(B)两部分copy是源于ldata->read_buf[]实际上是个环形buffer,为了处理如下情况。



(A)先填充free 2 buffer，(B)填充free 1 buffer.

另外由于是raw mode(完全raw),所以这里fp所指向的flags buffer是没意义的。

```

1.  static void
2.  n_tty_receive_buf_raw(struct tty_struct *tty, const unsigned char *cp,
3.                        char *fp, int count)
4.  {
5.      struct n_tty_data *ldata = tty->disc_data;
6.      char flag = TTY_NORMAL;
7.
8.      while (count--) {
9.          if (fp)
10.             flag = *fp++; (A)
11.          if (likely(flag == TTY_NORMAL)) (B)
12.             put_tty_queue(*cp++, ldata); (C)
13.          else
14.             n_tty_receive_char_flagged(tty, *cp++, flag); (D)
15.      }
16.  }

```

这里的raw不太彻底，还是要处理一些特殊字符，所以要判断字符的flag。

(A)

获得当前处理character的flag。TTY_NORMAL表示不是特殊处理的字符。

(B)

normal character

(C)

static inline void put_tty_queue(unsigned char c, struct n_tty_data *ldata)

```

{
    *read_buf_addr(ldata, ldata->read_head) = c;

    ldata->read_head++;
}

```


指示放入line discipline buffer即可

(D)

不是normal character，则要处理一下

```
1. static void
2. n_tty_receive_char_flagged(struct tty_struct *tty, unsigned char c, char flag)
3. {
4.     char buf[64];
5.
6.     switch (flag) {
7.     case TTY_BREAK:
8.         n_tty_receive_break(tty);
9.         break;
10.    case TTY_PARITY:
11.    case TTY_FRAME:
12.        n_tty_receive_parity_error(tty, c);
13.        break;
14.    case TTY_OVERRUN:
15.        n_tty_receive_overrun(tty);
16.        break;
17.    default:
18.        printk(KERN_ERR "%s: unknown flag %d\n",
19.               tty_name(tty, buf), flag);
20.        break;
21.    }
22. }
```

由于资料缺乏，所以只看懂了个大概。