

01 LOGIC

01_04 De Morgan's Laws and Their Applications



Augustus De Morgan

De Morgan's Laws

$$\sim(p \wedge q) \Leftrightarrow \sim p \vee \sim q,$$

$$\sim(p \vee q) \Leftrightarrow \sim p \wedge \sim q.$$

They can be proved respectively by the truth tables below.

| p | q | $p \wedge q$ | $\sim(p \wedge q)$ | $\sim p$ | $\sim q$ | $\sim p \vee \sim q$ | $\sim(p \wedge q) \Leftrightarrow \sim p \vee \sim q$ |
|---|---|--------------|--------------------|----------|----------|----------------------|---|
| T | T | T | F | F | F | F | T |
| T | F | F | T | F | T | T | T |
| F | T | F | T | T | F | T | T |
| F | F | F | T | T | T | T | T |

| p | q | $p \vee q$ | $\sim(p \vee q)$ | $\sim p$ | $\sim q$ | $\sim p \wedge \sim q$ | $\sim(p \vee q) \Leftrightarrow \sim p \wedge \sim q$ |
|---|---|------------|------------------|----------|----------|------------------------|---|
| T | T | T | F | F | F | F | T |
| T | F | T | F | F | T | F | T |
| F | T | T | F | T | F | F | T |
| F | F | F | T | T | T | T | T |

One application of De Morgan's Laws

Find the negation of the following propositions.

(1) Today is Sunday and $2 + 3 = 5$.

(2) It is raining or $6 > 8$.

[Solution of (1)]

p: Today is Sunday,

q: $2 + 3 = 5$.

By De Morgan's laws,

$$\sim(\text{Today is Sunday and } 2 + 3 = 5) = \sim(p \wedge q) \Leftrightarrow$$

$$\sim p \vee \sim q = \text{Today is not Sunday or } 2 + 3 \neq 5.$$

[Solution of (2)]

p: It is raining,

q: $6 > 8$.

By De Morgan's laws,

$$\sim(\text{It is raining or } 6 > 8) = \sim(p \vee q) \Leftrightarrow$$

$$\sim p \wedge \sim q = \text{It is not raining and } 6 \leq 8.$$

One application of De Morgan's Laws

Show that $\sim(p \rightarrow q)$ and $p \wedge \sim q$ are logically equivalent.

[Proof] Recall $p \rightarrow q \Leftrightarrow \sim p \vee q$. Thus, by De Morgan's

$$\text{Laws, } \sim(p \rightarrow q) \Leftrightarrow \sim(\sim p \vee q) \Leftrightarrow \sim(\sim p) \wedge \sim q \Leftrightarrow$$

$$p \wedge \sim q.$$

One application of De Morgan's Laws

Show that $(p \wedge q) \rightarrow p$ is a tautology.

[Proof] Recall $r \rightarrow s \Leftrightarrow \sim r \vee s$. Thus, by De Morgan's

$$\text{Laws, } (p \wedge q) \rightarrow p \Leftrightarrow \sim(p \wedge q) \vee p \Leftrightarrow (\sim p \vee \sim q) \vee p$$

$$\begin{aligned} &\Leftrightarrow \sim p \vee (\sim q \vee p) \Leftrightarrow \sim p \vee (p \vee \sim q) \Leftrightarrow (\sim p \vee p) \vee \sim q \\ &\Leftrightarrow T \vee \sim q \Leftrightarrow T. \end{aligned}$$

Similarly, we can show that $(p \wedge q) \rightarrow q$ is a tautology.

One application of De Morgan's Laws

Show that $p \rightarrow p \vee q$ is a tautology.

[Proof] Recall $r \rightarrow s \Leftrightarrow \sim r \vee s$. Thus, by De Morgan's

$$\begin{aligned} &\text{Laws, } p \rightarrow p \vee q \Leftrightarrow \sim p \vee (p \vee q) \Leftrightarrow (\sim p \vee p) \vee q \\ &\Leftrightarrow T \vee q \Leftrightarrow T. \end{aligned}$$

Similarly, we can show that $q \rightarrow p \vee q$ is a tautology.

More information on Augustus De Morgan can be found at

https://en.wikipedia.org/wiki/Augustus_De_Morgan