

CIT 593 Homework 12: File I/O and Dynamic Memory

Due Date: Tuesday 12/1 @11:59pm via canvas upload ONLY

Note: Because LC-4's C doesn't have support for malloc & free, we will once again be using our school server eniac and its C-compiler clang. The book will discuss malloc() and free() in **Chapter 19** and it discusses file I/O in **Chapter 18**. While these more mature functions may operate differently internally on different OS's, they all behave the same way across the C-language. So the book is definitely a good source for learning about file I/O & dynamic memory!

Assigned Problems (to be done individually, NOT group work):

For this assignment you will need to work & compile your C-programs on the university's server: eniac.

1) OVERVIEW: Implementing the PennSim "loader" using Linked Lists

Overview: The goal of this HW is for you to write a program that can open and read in an object file created by PennSim, parse it, and load it into a linked list that will represent the LC4's program and data memories. In the last HW you worked with linked lists, now you will adapt that knowledge to a new linked list format and also work with files in this assignment.

INPUT FILE FORMAT:

The following is the format for the binary .OBJ files created by PennSim from your .ASM files. It represents the contents of memory (both program and data) for your assembled LC-4 Assembly programs. In a .OBJ file, there are 3 basic sections indicated by 3 header "types" = CODE, DATA, SYMBOL.

- **Code:** 3-word header (xCADE, <address>, <n>), n-word body comprising the instructions. This corresponds to the .CODE directive in assembly.
- **Data:** 3-word header (xDADA, <address>, <n>), n-word body comprising the initial data values. This corresponds to the .DATA directive in assembly.
- **Symbol:** 3-word header (xC3B7, <address>, <n>), n-character body comprising the symbol string. Note, each character in the file is 1 byte, not 2. There is no null terminator. Each symbol is its own section. These are generated when you create labels (such as "END") in assembly.

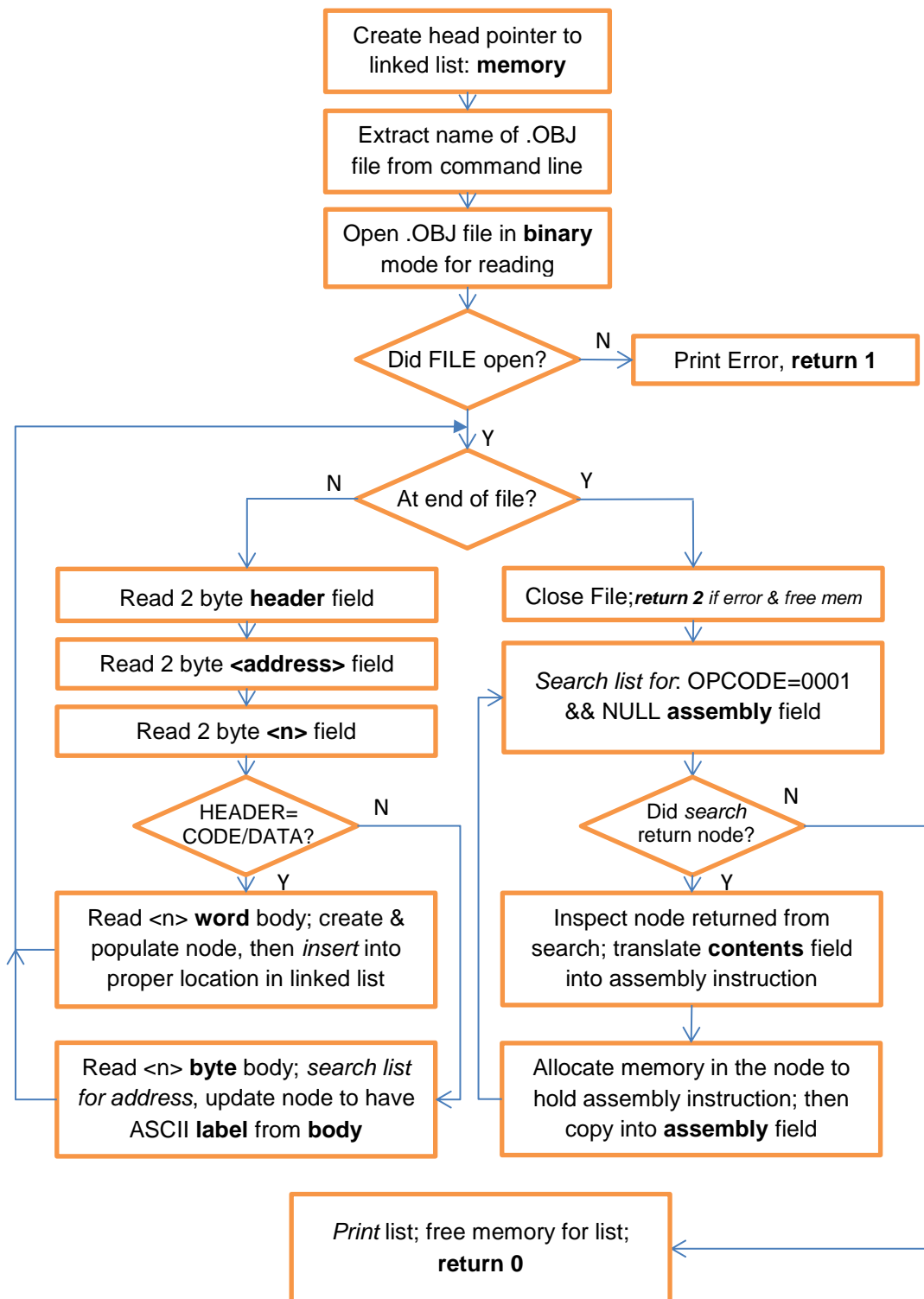
LINKED LIST NODE STRUCTURE:

The following structure (notice it has 5 **fields**) will be referenced in the problem description:

```
struct row_of_memory {
    short unsigned int address ;
    char * label ;
    short unsigned int contents ;
    char * assembly ;
    struct row_of_memory *next ;
} ;
```

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FLOW CHART: Overview of Program Operation



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IMPLEMENTATION DETAILS:

A helper file is available on canvas with this HW assignment, download and extract its contents to follow this discussion. The first files to view in the helper file are **lc4_memory.h** and **lc4_memory.c**. In these files you will notice the structure that represents a **row_of_memory** as referenced above (see the section: LINKED_LIST_NODE_STRUCTURE above for the node's layout). You will also see several helper functions that serve to manage a linked list of "rows_of_memory" nodes. Your job will be to implement these simple linked list helper functions using your knowledge from the last HW assignment.

Next, you will modify the file called: **lc4.c**. It serves as the "main" for the entire program. The head of the linked list must be stored in main(), you will see in the provided lc4.c file a pointer named: **memory** will do just that. Main() will then extract the name of the .OBJ file the user has passed in when they ran your program from the argv[] parameter passed in from the user. Upon parsing that, it will call lc4_loader.c's open_file() and hold a pointer to the open file. It will then ask call lc4_loader.c's parse_file() to interpret the .OBJ file the user wishes to have your program process. Lastly it will reverse assemble the file, print the linked list, and finally delete it when the program ends. These functions are described in greater detail below. The order of the function calls and their purpose is shown in comments in the lc4.c file that you will implement as part of this assignment.

Once you have properly implemented lc4.c and have it accept input from the command line, a user should be able to run your program as follows:

```
./lc4 my_file.obj
```

Where "**my_file.obj**" can be replaced with any file name the user desires as long as it is a valid .OBJ file that was created by PennSim. If no file is passed in, your program should generate an error telling the user what went wrong!

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Most of the work of your program will take place in the file: called: **lc4_loader.c**. In this file, you will implement a function: **open_file()** to take in the name of the file the user of your program has specified on the command line. If the file exists, the function should return a handle to that open file, otherwise a NULL should be returned.

Also in **lc4_loader.c**, you will implement a second function: **parse_file()** that will read in and parse the contents of the open .OBJ file as well as populate the linked_list as it reads the .OBJ file. The format of the .OBJ input file was discussed in lecture, but its layout is reprinted above (see section: INPUT_FILE_FORMAT). As shown in the flowchart above, have the function read in the 3-word header from the file. You'll notice that all of the LC4 .OBJ file headers consist of 3 fields: header type, <address>, <n>. As you read in the first header in the file, store the address field and the <n> filed into local variables. Then determine the type of header you have read in: CODE/DATA/SYMBOL.

If you have read in a CODE header in the .OBJ file, from the file format for a .OBJ file, you'll recall the body of the CODE section is <n>-words long. As an example, see the hex listed below, this is a sample CODE section, notice the field we should correlate with **n=0x000C**, or decimal: 12. This indicates that the next 12-words in the .OBJ file are in fact 12 LC-4 instructions. Recall each instruction in LC4 is 1 word long.

```
CA DE 00 00 00 0C 90 00 D1 40 92 00 94 0A 25 00 0C 0C 66 00 48
01 72 00 10 21 14 BF 0F F8
```

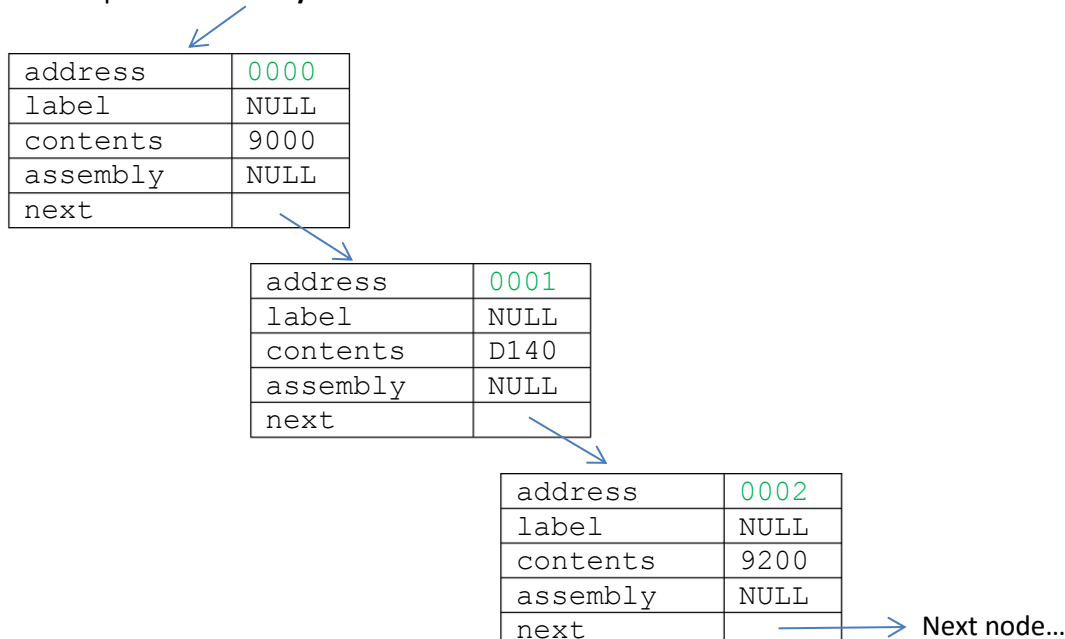
From the example above, we see that the first LC-4 instruction in the 12-word body is: 9000. (that happens to be a CONST assembly instruction if you convert to binary). Allocate memory for a new node in your linked list to correspond to the first instruction (the section above: **LINKED LIST NODE STRUCTURE**, declares a structure that will serve as a blue-print for all your linked list nodes called: "row_of_memory"). As it is the first instruction in the body, and the address has been listed as 0000, you would populate the row_of_memory structure as follows.

address	0000
label	NULL
contents	9000
assembly	NULL
next	NULL

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In a loop, read in the remaining instructions from the .OBJ file; allocate memory for a corresponding **row_of_memory** node for each instruction. As you create each **row_of_memory** add these nodes to your linked list (you should use the functions you've created in `lc4_memory.c` to help you with this). For the first 3 instructions listed in the sample above, your linked list would look like this:

Header pointer: **memory**



The procedure for reading in the DATA sections would be identical to reading in the CODE sections. These would become part of the same linked list, as we remember PROGRAM and DATA are all in one “memory” on the LC-4, they just have different addresses.

For the following SYMBOL header/body:

C3 B7 00 00 00 04 49 4E 49 54

The address field is: **0x0000**. The symbol field itself is: **0x0004** bytes long. The next 4 bytes: 49 4E 49 54 are ASCII for: **INIT**. This means that the label for address: 0000 is **INIT**. Your program must search the linked list: **memory**, find the appropriate address that this label is referring to and populate the “label” field for the node. Note: the field: **<n>** tells us exactly how much memory to `malloc()` to hold the string, however you must add a byte to hold the **NULL**. 5 bytes in the case of: **INIT**. For the example above, the node: 0000 in your linked list, would be updated as follows:

address	0000
label	INIT
contents	9000
assembly	NULL
next	

Once you have read the entire file; created and added the corresponding nodes to your linked list, close the file and return to `main()`. If you encounter an error in closing the file, before exiting, print an error, but also `free()` all the memory associated with the linked list prior to exiting the program.

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2) Implementing a Reverse Assembler

In a new file: **lc4_disassembler.c**: write a third function that will take as input the populated “memory” linked list that contains the .OBJ’s contents and translate the hex representation of some instructions into their assembly equivalent. You will need to reference the LC4’s ISA to author this function. To simplify this problem a little, you **DO NOT** need to translate every single HEX instruction into its assembly equivalent. Only translate instructions with the OPCODE: 0001 (ADD REG, MUL, SUB, DIV, ADD IMM)

As shown in the flowchart, this function will call your linked list’s “search_by_opcode()” helper function. Your search_by_opcode() function should take as input an OPCODE and return the first node in the linked list that matches the OPCODE passed in, but also has a NULL assembly field. When/if a node in your linked list is returned, you’ll need to examine the “contents” field of the node and translate the instruction into its assembly equivalent. Once you have translated the contents field into its ASCII Assembly equivalent, allocate memory for and store this as string in the “assembly” field of the node. Repeat this process until all the nodes in the linked list with an OPCODE=0001 have their assembly fields properly translated.

As an example, the figure below shows a node on your list that has been “found” and returned when the search_by_opcode() function was called. From the contents field, we can see that the HEX code: 128B is 0001 001 010 001 011 in binary. From the ISA, we realize the sub-opcode reveals that this is actually a MULTIPLY instruction. We can then generate the string **MUL R1, R2, R3** and store it back in the node in the assembly field. For this work, I strongly encourage you to investigate the **switch()** statement in C (any good book on C will help you understand how this works and why it is more practical than multiple if/else/else/else statements). *I also remind you that you must allocate memory strings before calling strcpy()!*

NODE BEFORE UPDATE

address	0009
label	NULL
contents	128B
assembly	NULL
next	

NODE AFTER UPDATE

address	0009
label	NULL
contents	128B
assembly	MUL R1, R2, R3
next	

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Putting it all together:

As you may have realized `main()` should do only 3 things: 1) create and hold the pointer to your memory linked list. 2) Call the parsing function in `lc4_loader.c`. 3) Call the disassembling function in `lc4_disassembler.c`. One last thing to do in `main()` is to call a function to print the contents of your linked list to the screen. Call the `print_list()` function in `lc4_memory.c`; you will need to implement the printing helper function to display the contents of your `lc4`'s memory list like this:

<label>	<address>	<contents>	<assembly>
INIT	0000	9000	
	0001	D140	
	0002	9200	
	...		
	0009	128B	MUL R1, R2, R3

(and so on...)

Several things to note: There can be multiple CODE/DATA/SYMBOL sections in one .OBJ file. If there is more than one CODE section in a file, there is no guarantee that they are in order in terms of the address. In the file shown above, the CODE section starting at address 0000, came before the CODE section starting at address: 0010; there is no guarantee that this will always happen, your code must be able to handle that variation. Also, SYMBOL sections can come before CODE sections! What all of this means is that before one creates/allocates memory for a new node in the memory list, one should “search” the list to make certain it does not already exist. If it exists, update it, if not, create it and add it to the list!

Prior to exiting your program, you must properly “free” any memory that you allocated. We will be using a memory checking program known as `valgrind` to ensure your code properly releases all memory allocated on the heap! Simply run your program: `lc4` as follows:

```
valgrind lc4
```

Valgrind should report 0 errors AND there should be no memory leaks prior to submission.

Also note: If your code doesn't compile or even run, you will lose most of the points of this assignment!

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STRUCTURING YOUR CODE:

In the file: **CIT593_HW13_Helper_Files.zip** you'll find the following files that you must use:

lc4.c	- must contain your main() function.
lc4_memory.c	- must contain your linked list helper functions.
lc4_memory.h	- must contain the declaration of your <u>row of memory</u> structure & linked list helper functions
lc4_loader.c	- must contain your .OBJ parsing function.
lc4_disassembler.c	- must contain your disassembling function.
Makefile	- must contain an "all" and "clean" targets

To compile the program type:

```
make all
```

The **Makefile** included automates the compiling process for you. *It has 1 error*. It produces a file called: **a.out** when you run make all. **Update the Makefile to produce a file called: lc4 instead of a.out**. You must edit the Makefile to make this correction.

3) EXTRA CREDIT: A complete reverse assembler:

Finish the disassembler to translate any/all instructions in the ISA. Have your program print the linked list to the screen still, but also create a new output file: <users_input>.asm. In that file it should contain only the assembly program that you disassembled. If it works, I should be able to load it into PennSim, assemble it, and reproduce the identical .OBJ file that your .ASM file was derived from! Don't forget to add in the directives (.CODE, .DATA)...the ultimate test of your program will be getting it to assemble using PennSim!

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Directions on how to submit your work:

- Create a single zip file called: **LAST_FIRST_HW12.zip**
- The zip file should contain the named in this assignment.
- There should not be any sub-directories within your zip file.

As an example, I would turn in the following SINGLE zip file:

FARMER_THOMAS_HW12.zip

This single zip file would contain 8 files only:

```
lc4.c
lc4_loader.h
lc4_loader.c
lc4_memory.h
lc4_memory.c
lc4_disassembler.h
lc4_disassembler.c
Makefile
```

You will then upload ONLY 1 file to canvas: **FARMER_THOMAS_HW12.zip**

- **DO NOT TURN IN ANY files created by the compiler, .obj, or PennSim.jar**
- **Make certain that you submit the latest version of your code.**
- **Submitting using any other compression type (.RAR, TAR, GZIP) will be rejected.**

Paper/Email submissions will not be accepted for this assignment.