## When should one inline recursive functions?

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#### Abstract

Hello, here is some text without a meaning. This text should show what a printed text will look like at this place. If you read this text, you will get no information. Really? Is there no information? Is there a difference between this text and some nonsense like "Huardest gefburn"? Kjift – not at all! A blind text like this gives you information about the selected font, how the letters are written and an impression of the look. This text should contain all letters of the alphabet and it should be written in of the original language. There is no need for special content, but the length of words should match the language.

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#### 1 Introduction

Describe layout of paper. What does each section in turn discuss?

#### 1.1 Problemsetting

Re-state the assignment: Inlining -; Easy, what is it, benefits/drawbacks. "This paper details the problems and benefits of inlining"
Introduce Jive in a few sentences, say that it's further introduced in section

Inlining is a simple and straight-forward technique used in code compilation, where one replaces the call of a function with the body of said function. Its benefits include removal of function call overhead, and more importantly, unveiling of additional potential optimizations in the code. However, the drawback is potentially increased code size, and longer execution times for the compilation of the program.

This paper details the problem and benefits of inlining wrt. the Jive compiler. The Jive compiler is a new backend compiler using intermediate representation (IR), and an alternative to the control flow graph (CFG); the regionalized value-state dependency graph (RVSDG<sup>1</sup>).

insert reference

Further details of this assignment can be found in Appendix A.

#### 2 Background

#### 2.1 The Regionalized Value-State Dependency Graph

- $\gamma$  blocks/regions in the RVSDG are conditionals. The inputs of each  $\gamma$  block is the variables(/data) upon which the conditional depends, as well as the operation performed on these. A CNF may then be represented as several nested  $\gamma$  blocks, all the while retaining the properties of a demand-dependence graph.
- $\theta$  blocks/regions are the loop constructs.

#### Need better explanation.

 $\bullet$   $\lambda$  blocks/regions are the functions. They retain the demand-dependency graph properties by implementing

Ask Nico for help?

•  $\phi$  blocks/regions are the ones describing mutually recursive environments. Inside a  $\phi$  region, the RVSDG will have at least one  $\lambda$  region, and perhaps more. If there is more than one  $\lambda$  present, then the "mutually recursive binding group" situation which P. Jones and Marlow [4] describe when discussing how to inline recursive functions, is also present in the program flow represented by this RVSDG.

<sup>&</sup>lt;sup>1</sup>Detailed in Section 2.1.

- 3 Scheme
- 4 Methodology
- 5 Results
- 6 Discussion
- 7 Related Work

In this section (...)

To do...

#### 7.1 Regionalized Value-State Dependency Graph

Insert reference/summary of HiPEAC paper when published

#### 7.2 Inlining

W. Davidson and M. Holler [2] examine the proposition that the increased size of inlined code affects the execution time performance on demand-paged virtual memory machines. Their paper details how well the equations developed to describe inlined versions of a programs execution time, match with the aforementioned proposition by utilizing a source-to-source subprogram inliner. In doing this, they also examine the overhead costs of running inlined functions, like register allocation issues.

Cavazos and F.P. O'Boyle [1] show how conjunctive normalform (CNF) can easily be used to decide if and when to inline a specific call site. They report between 17% and 37% execution time improvements without code explosion in their testing of their genetic algorithm auto-tuning.

Serrano [5] explains when to inline functions in the Scheme programming language, by classifying them as either recursive or non- recursive functions, and thereafter mainly inlining non-recursive functions if they do not bloat the code size above a certain treshold.

Serrano [5] also takes overhead into account when deciding on when to inline, and hence also looks at inlining certain recursive functions which fulfill some requirements, at least to a certain depth.

Waterman's Ph.D. thesis [6] examines the use of techniques to adaptively decide which functions to inline during program compilation, to show that they can be used to improve the performance of specific optimizations.

The adaptive technique both matches, and at times, beats ATLAS and gcc on one of the object oriented biggest programs in the SPEC CINT2000 test suite. The technique accepts condition strings (like in [1]) to determine which call sites are inlined. The condition strings provide a flexible inliner by combining various program properties in CNF, and exposes a large space of different inlining decisions with the potential to outperform static techniques when used adaptively.

Find a good reference?

need AT-LAS cite? E. Hank, W. Hwu, and R. Rau [3] introduces a new technique called *Region-Based Compilation*, and in their paper exhibits the benefits an aggressive compiler can have from inlining, bringing the percentage of dynamic acyclic code down to below 10% in almost all tests. The paper thus shows that inlining is an important technique for aggressive compilation techniques, and both apt and useful for the Region-Based compilation technique.

#### 7.3 Recursive Inlining

P. Jones and Marlow [4] explore how the decision of when to inline mutually recursive binding groups in Haskell, through the use of a novel approach using strictness-analysis and forming a graph of the strongly connected components (SCC) in said mutually recursive binding group.

They also show how this approach in the Haskell core language does not explode code size when inlining some of the recursive functions. Through the use of their SCC they show how instead of not inlining any recursive function (like Serrano [5]), the SCC graph and some other flags can show when it is not only safe, but desirable to inline a recursive function.

#### 8 Conclusion

#### 8.1 Further Work

#### 9 References

- John Cavazos and Michael F. P. O'Boyle. Automatic tuning of inlining heuristics. In *Proceedings of the 2005 ACM/IEEE Conference on Supercom*puting, SC '05, pages 14–, Washington, DC, USA, 2005. IEEE Computer Society.
- [2] J.W. Davidson and A.M. Holler. Subprogram inlining: a study of its effects on program execution time. *Software Engineering, IEEE Transactions on*, 18(2):89–102, Feb 1992.
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- [6] Todd Waterman. Adaptive Compilation and Inlining. PhD thesis, Houston, TX, USA, 2006. AAI3216796.

# A Report Description

### Project Description: An Inliner for the Jive compiler

#### Nico Reissmann

#### December 12, 2014

Compilers have become an essential part of every modern computer system since their rise along with the emergence of machine-independent languages at the end of the 1950s. From the start, they not only had to translate between a high-level language and a specific architecture, but had to incorporate optimizations in order to improve code quality and be a par with human-produced assembly code. One such optimization performed by virtually every modern compiler is *inlining*. In principle, inlining is very simple: just replace a call to a function by an instance of its body. However, in practice careless inlining can easily result in extensive work and code duplication. An inliner must therefore decide carefully when and where to inline a function in order to achieve good performance without unnecessary code bloat.

The overall goal of this project is to implement and evaluate an inliner for the Jive compiler back-end. The project is split in a practical and an optional theoretical part. The practical part includes the following:

- Implementation of an inliner for the Jive compiler back-end. The inliner must be able to handle recursive functions and allow for the configuration of different heuristics to permit rapid exploration of the parameter space.
- An evaluation of the implemented inliner. A particular emphasis is given to different heuristics and their consequences for the resulting code in terms of work and code duplication.

The Jive compiler back-end uses a novel intermediate representation (IR) called the Regionalized Value State Dependence Graph (RVSDG). If time permits, the theoretical part of the project is going to clarify the consequences of using the RVSDG along with an inliner. It tries to answer the following research questions:

- What impact does the RVSDG have on the design of an inliner and the process of inlining?
- Does the RVSDG simplify/complicate the implementation of an inliner and the process of inlining compared to other commonly used IRs?

The outcome of this project is threefold:

- 1. A working implementation of an inliner in the Jive compiler back-end fulfilling the aforementioned criteria.
- 2. An evaluation of the implemented inliner.
- 3. A project report following the structure of a research paper.