

Tree Structured Traffic Coordination with Decentralized Optimization

Xiaotian Fang

Abstract—This paper proposes a distributed model predictive control algorithm for tree structured traffic coordination problem by using two dimensional models for autonomous vehicles. With the definition of conflict zone, A complete strategy is proposed to model the structure of the vehicles' network to a tree topology. The shape of vehicles is also be considered in this paper, and the collision avoidance constraints are formulated by enclosing each vehicle in an ellipse. Then the closed-loop control scheme is constructed based on a model predictive method. And a decentralized optimization method over tree graph, called multi-sweep method, is used to solve the traffic coordination problem in each MPC loop. Two numerical cases under two realistic scenarios are introduced to illustrate the performance of the proposed method.

I. INTRODUCTION

Autonomous vehicles have received much attention in recent years. Interest in coordinating vehicles such as unmanned ground vehicles (UGVs) [1], unmanned underwater vehicles (UUVs) [2], and unmanned air vehicles (UAVs) [3] has been on the rise in control system community. Many methods [4], [5], [6], [7] for coordinating self-driving cars at traffic intersections have been proposed in the past few years. For example, in [7], a cooperative management system has been proposed to coordinate cars at traffic intersections. A connected communication network has been designed for coordination of vehicles in [5]. Moreover, a receding horizon based the cooperative approach has been developed in [8].

The coordination of autonomous vehicles can be formulated as a structured optimal control problem. Concerning various scenarios, different methods has been proposed to model collision avoidance constraints. Based on double integrator dynamics, a spatial method has been proposed in [9], which enforces a safety distance between two vehicles. Moreover, in [1] a collision avoidance controller has been developed, which allows only one car on the intersection at a time. In [10], hybrid system theory has been used to address the traffic control problem under a verifiable safety condition.

MPC controllers for traffic coordination have been developed in [6], [8]. Due to the fast dynamic of autonomous vehicles, a real time iteration based strategy was designed in [11] for obstacle avoidance constraints. If we consider the connection structure of vehicles as a multi-agent network, if a vehicle is plugged or unplugged from the network, this results in a change of the number of total differential states in the optimal control problem. In [12], a closed-loop control scheme with recursive feasibility guarantee for time-varying networks has been proposed.

Modern wireless networks can provide a reliable and fast channel for information exchange between multiple agents.

Many researchers have proposed methods that exploit the network structure of interconnected vehicles to solve multi-vehicle coordination problem in a distributed manner. By distributing the sensitivity evaluation in sequential quadratic programming (SQP) [13], a primal decomposition based SQP method [14] was proposed to solve the optimal control problem. The follow-up asynchronous version [15] reduces the communication overhead. While this method is distributed, the corresponding QP is solved by one car, which is chosen as the central unit. A fully parallelizable MPC based approach was proposed in [9]. Another class of distributed optimization methods are based on the alternating direction method of multipliers (ADMM) [16]. In [17], ADMM has been illustrated to be applicable for nonconvex optimization problem. An ADMM based nonlinear MPC approach proposed in [18] aims to coordinating multiple vehicles at intersections. Moreover, in [19] ALADIN has been used to solve the traffic optimal control problem in a distributed way by negotiating a time schedule between the vehicles.

Meanwhile, a huge amount of research focus on distributed optimization algorithms over networks with tree topology, which arise in many applications such as traditional optimal control problems [20] or receding horizon control problems [21], where linear trees occur, scenario multi-stage MPC problems [22], [23], [24], which have less trivial tree structures, or radial power grid networks that possess non-trivial tree structure too [25], [26].

This paper is mainly based on the work of [19], [12] for coordination of vehicles at traffic intersections, and [27] for decentralized optimization method over tree graphs. Modeling the vehicles in two dimensions to enclose each vehicle by a two-dimensional ellipsoid and to enforce a minimal safety distance between them for collision avoidance. Conflict region and complete regulations are defined to help model the connections between vehicles to tree topology. The closed-loop scheme is constructed by model predictive control method while the resulting online optimization problem with tree topology is solved with multi-sweep method.

The rest of this paper is organized as follows. Section II presents the model of the vehicle and the frame of local and global coordinates. Section III introduces a tractable formulation of the collision avoidance constraint. Section IV summarizes the optimal control problem and its reformulation to tree-topology problem, and meanwhile briefly introduces the solving approach, multi-sweep method. Section V presents the closed-loop MPC implement details. Two realistic case studies are presented in Section VI. At last, Section VII concludes this paper.

II. SYSTEM MODEL

In this paper, we consider a traffic system with some vehicles under several realistic scenarios. In order to formulate the coordination problem, we have following Assumptions:

- A1** The vehicles follow fixed, predetermined paths.
- A2** Vehicle states are measurable without uncertainty.
- A3** Vehicles in different paths only need to consider each other inside the conflict regions.
- A4** Vehicle paths are straight outside the conflict region.
- A5** All vehicles are equipped with V2V communication devices.
- A6** The MPC solutions at time k are available to all vehicles at time $k + 1$.

A. Vehicle Kinematics

Let N_v denote the number of vehicles and $\mathcal{V} := \{1, \dots, N_v\}$. According to assumption A1, the position of vehicle i , which moves along a fixed path, can be parametrized by only one scalar s_i . Let $x_{i,k} = [s_{i,k}; v_{i,k}]$ denote the state of each vehicle $i \in \mathcal{V}$ at time $k \in \mathbb{Z}_0^K = \{t_0, t_1, \dots, t_K\}$, where $s_{i,k}$ and $v_{i,k}$ respectively represent the scalar longitudinal position and velocity. And control variable $u_{i,k}$ represents the scalar longitudinal acceleration. Then, each vehicle is modeled as a discrete time double integrator

$$x_{i,k+1} = A \cdot x_{i,k} + B \cdot u_{i,k}, \quad (1)$$

where $A = [1, h; 0, 1]$, $B = [0; h^2/2]$, and $h = T/K$ denotes one time interval.

Meanwhile, considering the traffic regulations and the parameters of real vehicles. Each vehicle is supposed to follow the state and control constraints below:

$$v_i^{min} \leq v_{i,k} \leq v_i^{max}, \quad (2)$$

$$u_i^{min} \leq u_{i,k} \leq u_i^{max}, \quad (3)$$

In the following, we use shorthand $h(x_i, u_i) \leq 0$ to collect (2) and (3) for all $k \in \mathbb{Z}_0^K$.

B. Local and Global Coordinate Frames

The position of vehicles in 2-D can be given by function $p_i : \mathbb{R} \rightarrow \mathbb{R}^2$,

$$p_i(s_{i,k}) = (p_{i,x}(s_{i,k}), p_{i,y}(s_{i,k})).$$

For the given predetermined path, refer to [28], we use n_p -th order B-splines to represent it,

$$\begin{bmatrix} p_{i,x}(s_{i,k}) \\ p_{i,y}(s_{i,k}) \end{bmatrix} = \begin{bmatrix} \sum_{l=0}^{n_p} \alpha^{x,l} B_n^{x,l}(s_{i,k}) \\ \sum_{l=0}^{n_p} \alpha^{y,l} B_n^{y,l}(s_{i,k}) \end{bmatrix}. \quad (4)$$

Similarly, the navigation angle $\theta(s_{i,k})$ can be also represented as functions of $s_{i,k}$. In the following, we use notation $p_{i,k} := p(s_{i,k})$ and $\theta_{i,k} := \theta(s_{i,k})$.

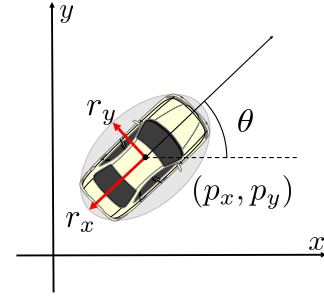


Fig. 1: Elliptical vehicle model.

C. Vehicle Shape

In order to analyze the distance between two vehicle in 2-D plane, we model vehicle as an ellipse as shown in Figure 1. We denote the shape ellipse by $\mathcal{E}(p_{i,k}, \mathcal{Q}(\theta_{i,k}))$ with center $p_{i,k}$ and shape matrix

$$\mathcal{Q}(\theta_{i,k}) = \begin{bmatrix} r_x^2 \cos^2 \theta_{i,k} + r_y^2 \sin^2 \theta_{i,k}, & 0.5 \sin 2\theta_{i,k} (r_x^2 - r_y^2) \\ 0.5 \sin 2\theta_{i,k} (r_x^2 - r_y^2), & r_x^2 \sin^2 \theta_{i,k} + r_y^2 \cos^2 \theta_{i,k} \end{bmatrix}.$$

Here, both the center of the ellipse and the shape matrix are parameterized over $s_{i,k}$. And (r_x, r_y) denotes the length of semi-major radius of ellipses.

III. COLLISION AVOIDANCE

A. Conflict Region

For vehicle i , we define the conflict region

$$\mathcal{C}_i := [s_i^{in}, s_i^{out}],$$

Where s_i^{in} and s_i^{out} denote the path coordinates of conflict region's entry and exit, respectively. Before vehicles enter the conflict region, let assumption A3 hold, vehicles in different paths move independently such that we only need consider rear-end collision. Once vehicles has been inside the conflict region, the possibility of side collision between vehicles in different lanes occurs. Notice that the choosing principles of conflict region might have some difference in different scenarios, we need to analyse specific scenarios to acquire a suitable conflict region. Some articles could be referred here [29] [28]. And in the case study section, we will define conflict region specifically in two different scenarios.

B. Rear-end Collision Avoidance

According to the assumption A3, we only need to consider the rear-end collision outside conflict region. Therefore, we denote by a_i the vehicle which is ahead of vehicle i in same path. When $s_{i,k} \notin \mathcal{C}_i$, vehicle i only need to avoid rear-end collision with vehicle a_i . Since the path is straight before the intersection, as assumed in A4, the distance between two vehicles can be calculated by

$$\delta_{i,a_i,k} = s_{i,k} - s_{a_i,k} - r_{i,x} - r_{a_i,x}.$$

and the collision avoidance constraints can be given by

$$\delta_{i,a_i,k} \leq \varepsilon_{out}, \quad (5)$$

where $\varepsilon_{out} > 0$ is the safe distance for vehicles outside conflict region.

C. Side Collision Avoidance

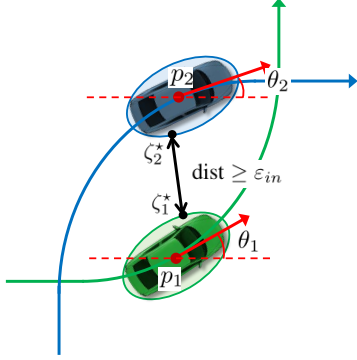


Fig. 2: collision avoidance inside conflict region

According to the definition of conflict region, vehicles start to consider the side collision when they enter the conflict region. In this phase, not only the rear-end collision at the same path but also the side collision between different lanes need to be considered. Here, we denote by \mathcal{V}_i^b the set of vehicles, except vehicle i , inside conflict region. Therefore, in order to avoid collision inside conflict region, we can enforce a safe distance between two vehicle ellipses inside the conflict region as shown in Figure 2. For vehicle i , the collision avoidance constraint can be then formulated as

$$\text{dist}(\mathcal{E}(p_{i,k}, \mathcal{Q}_i(\theta_{i,k})), \mathcal{E}(p_{j,k}, \mathcal{Q}_j(\theta_{j,k}))) \geq \varepsilon_{in} \quad (6)$$

for all $j \in \mathcal{V}_i^b$. Here, ε_{in} denote the safety distance between vehicles inside the conflict region. However, constraint (6) is not applicable in numerical algorithm design such that we introduce the following technical result.

Theorem 1 For any given two ellipsoids $\mathcal{E}(q_1, Q_1)$ and $\mathcal{E}(q_2, Q_2)$ in two dimensions, the inequality

$$\text{dist}(\mathcal{E}(q_1, Q_1), \mathcal{E}(q_2, Q_2)) \geq d_{12}$$

holds for a given distance d_{12} if and only if there exist a $\delta = (\delta_1, \delta_2) \in \mathbb{R}^2$ such that

$$\begin{cases} \|\zeta_1(\delta) - \zeta_2(\delta)\|_2 \geq d_{12}, \\ (\zeta_1(\delta) - q_1)^\top Q_1^{-1}(\zeta_1(\delta) - q_1) = 1, \\ (\zeta_2(\delta) - q_2)^\top Q_2^{-1}(\zeta_2(\delta) - q_2) = 1, \\ \delta \geq 0, \end{cases} \quad (7)$$

where the functions $\zeta_1(\cdot)$ and $\zeta_2(\cdot)$ are given by

$$\begin{aligned} \zeta_1(\delta) &= [(I + \delta_2 Q_2^{-1}(I + \delta_1 Q_1^{-1}) - I)^{-1} \\ &\quad [\delta_1(I + \delta_2 Q_2^{-1})Q_1^{-1}q_1 + \delta_2 Q_2^{-1}q_2]], \\ \zeta_2(\delta) &= [(I + \delta_1 Q_1^{-1})(I + \delta_2 Q_2^{-1}) - I]^{-1} \\ &\quad [\delta_2(I + \delta_1 Q_1^{-1})Q_2^{-1}q_2 + \delta_1 Q_1^{-1}q_1] \\ &\quad \forall \delta \geq 0. \end{aligned} \quad (8)$$

Proof. The square of distance between two ellipsoids is given by the optimal value of the following optimization problem

$$\begin{aligned} \min_{\xi_1, \xi_2} \quad & \|\xi_1 - \xi_2\|_2^2 \\ \text{s.t.} \quad & \begin{cases} 1 \geq (\xi_1 - q_1)^\top Q_1^{-1}(\xi_1 - q_1) & | \kappa_1 \\ 1 \geq (\xi_2 - q_2)^\top Q_2^{-1}(\xi_2 - q_2) & | \kappa_2 \end{cases} \end{aligned} \quad (9)$$

We denote the Lagrangian of (9) by

$$\begin{aligned} L(\xi, \kappa) &= \|\xi_1 - \xi_2\|_2^2 + \kappa_1(\xi_1 - q_1)^\top Q_1^{-1}(\xi_1 - q_1) \\ &\quad - \kappa_1 + \kappa_2(\xi_2 - q_2)^\top Q_2^{-1}(\xi_2 - q_2) - \kappa_2. \end{aligned}$$

Since (9) is a strongly convex QCQP problem, whose feasible set has a non-empty interior, there is no duality gap. By expanding the first order optimality condition of (9), we find that

$$\begin{aligned} (\xi_1^* - \xi_2^*) + \kappa_1^* Q_1^{-1}(\xi_1^* - q_1) &= 0, \\ (\xi_2^* - \xi_1^*) + \kappa_2^* Q_2^{-1}(\xi_2^* - q_2) &= 0, \end{aligned} \quad (10)$$

where (ξ^*, κ^*) is the primal and dual solution of (9). By simplifying (10), we obtain (8). According to the definition of the distance between two ellipsoids, substituting

$$\xi_1^* = \zeta_1(\delta) \quad \text{and} \quad \xi_2^* = \zeta_2(\delta)$$

yields the first inequality in (7). Moreover, the last inequality in (7) is given by the dual feasibility condition. \square

According to Theorem 1, constraint (6) can be replaced by (8), which is summarized as

$$\mathcal{G}(\delta_{i,j,k}, p_{i,k}, p_{j,k}, \theta_{i,k}, \theta_{j,k}, \varepsilon_{in}) \leq 0. \quad (11)$$

with

$$\mathcal{G} := \begin{pmatrix} \varepsilon_{in} - \|\zeta_1(\delta_{i,j,k}) - \zeta_2(\delta_{i,j,k})\|_2, -\delta_{i,j,k} \\ (\zeta_1(\delta_{i,j,k}) - p_{i,k})^\top Q^{-1}(\theta_{i,k})(\zeta_1(\delta_{i,j,k}) - p_{i,k}) - 1 \\ 1 - (\zeta_1(\delta_{i,j,k}) - p_{i,k})^\top Q^{-1}(\theta_{i,k})(\zeta_1(\delta_{i,j,k}) - p_{i,k}) \\ (\zeta_2(\delta_{i,j,k}) - p_{j,k})^\top Q^{-1}(\theta_{j,k})(\zeta_2(\delta_{i,j,k}) - p_{j,k}) - 1 \\ 1 - (\zeta_2(\delta_{i,j,k}) - p_{j,k})^\top Q^{-1}(\theta_{j,k})(\zeta_2(\delta_{i,j,k}) - p_{j,k}) \end{pmatrix}.$$

IV. PROBLEM FORMULATION

A. Centralized Problem Formulation

Firstly, we construct the traffic coordination problem as a centralized problem,

$$\min_{x, u} \sum_{i \in \mathcal{V}} J_i(x_i, u_i) \quad (12a)$$

$$\text{s.t.} \quad \forall i \in \mathcal{V}$$

$$x_{i,k+1} = Ax_{i,k} + Bu_{i,k}, \quad k \in \mathbb{Z}_0^{K-1} \quad (12b)$$

$$x_{i,0} = \hat{x}_i, \quad h_i(x_i, u_i) \leq 0, \quad (12c)$$

$$\begin{cases} \text{if } \hat{x}_i \notin \mathcal{C}_i, \quad \forall k \in \mathbb{Z}_0^K \\ \delta_{i,a_i,k} \leq \varepsilon_{out}, \end{cases} \quad (12d)$$

$$\begin{cases} \text{if } \hat{x}_i \in \mathcal{C}_i, \quad \forall j \in \mathcal{V}_i^b, \quad \forall k \in \mathbb{Z}_0^K \\ \mathcal{G}(\delta_{i,j,k}, p_{i,k}, p_{j,k}, \theta_{i,k}, \theta_{j,k}, \varepsilon_{in}) \leq 0, \end{cases} \quad (12e)$$

where \hat{x}_i denotes the initial state of vehicle i . with $x_i = (x_{i,0}, \dots, x_{i,K})$, $u_i = (u_{i,0}, \dots, u_{i,K-1})$, the objective function is given by

$$J_i(x_i, u_i) = \|x_i - x_{ref}\|_Q^2 + \|u_i\|_R^2. \quad (13)$$

where $R \geq 0$ and $Q \geq 0$ are block diagonal weighting matrices of appropriate dimensions.

B. Decentralized Tree-Topology Problem Formulation

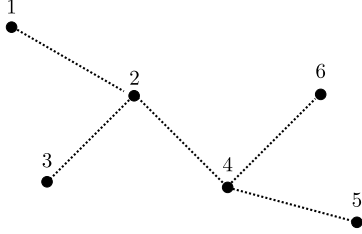


Fig. 3: Example for a tree graph with $N = 6$ nodes.

Let $(\mathcal{N}, \mathcal{E})$ denote a graph with node set $\mathcal{N} = \{1, \dots, N\}$ and edge set $\mathcal{E} \subseteq \mathcal{N} \times \mathcal{N}$, and $\mathcal{N}_i = \{j \in \mathcal{N} \mid (i, j) \in \mathcal{E}\}$ denotes the set of neighbors of vehicle i . To reduce duplication, we enumerate tree graph $(\mathcal{N}, \mathcal{E})$ to graph $(\mathcal{N}, \mathcal{E}^+)$, with

$$\mathcal{E}^+ = \{(i, j) \in \mathcal{E} \mid i < j\}.$$

Refer to [27], a standard form of tree-structured optimization problems is like following

$$V^* = \min_x \sum_{i \in \mathcal{N}} F_i(x_i) \quad \text{s.t.} \quad \begin{cases} \forall (i, j) \in \mathcal{E}^+, \\ S_{i,j}x_i = S_{j,i}x_j. \end{cases} \quad (14)$$

Here, the functions $F_i : \mathbb{R}^{n_i} \rightarrow \mathbb{R}$ denote objective functions, $S_{i,j} \in \mathbb{R}^{n_i, j \times n_i}$ and $S_{j,i} \in \mathbb{R}^{n_i, j \times n_j}$ given connectivity matrices.

To reformulate the centralized problem, we construct new state variable

$$x_i = (x_i, x_{j,j \in \mathcal{N}_i}, u_i)$$

and use short hand $d(x_i) \geq \varepsilon$ to collect the collision avoidance constraints (12e) and (12d). Then, the origin problem can be written as following format

$$V^* = \min_x \sum_{i \in \mathcal{N}} J_i(x_i) \quad \text{s.t.} \quad \begin{cases} \forall (i, j) \in \mathcal{E}^+, \\ S_{i,j}x_i = S_{j,i}x_j, \\ d(x_i) \geq \varepsilon. \end{cases} \quad (15)$$

Remark 1 To make adding constraints into problems feasible, we must make sure that the locally active set of adding constraint is fixed, and solution of each sub problem has regular KKT. Traffic coordination problems in this paper satisfy the two conditions above.

V. DISTRIBUTED MPC SCHEME

In this section, we propose a distributed MPC scheme outlined in Algorithm 1.

Algorithm 1 Distributed MPC for Traffic Coordination

Online:

- 1) Each vehicle measures \hat{x}_i .
- 2) Each vehicle communicate with its neighbors to update a_i, \mathcal{V}_i^b .
- 3) Each vehicle constructs x_i , objective function $J_i(x_i)$, connectivity matrix $S_{i,j}$ and constraint $d(x_i)$.
- 4) Use multi-sweep method to solve (15).
- 5) Each vehicle applies the optimal control $u_{i,0}^*$ during the first one sampling time.

Once vehicles receives the new measurement, we need to update the a_i and \mathcal{V}_i^b based on current states \hat{x}_i . Besides, if there exists vehicles plugged into the coordination system as a new neighbors, then we expand vehicles set \mathcal{V} . Once a_i, \mathcal{V}_i^b and x_i are updated, Problem (15) is solved by a multi-sweep method, a decentralized optimization algorithm over tree topology. Here, we don't elaborate the details about how this method works in a distributed manner, more details can be found in [27]. After that, the first step of optimal solution $u_{i,0}^*$ is applied.

VI. CASE STUDY

This section illustrates the detailed defination of conflict region and the performance of proposed closed-loop scheme for two scenarios. Both cases were implemented with CasADi v3.4.5 [30] in MATLAB 2018a. The closed-loop setting and parameters are listed in Table I.

TABLE I: Parameter values

r_x	2 m	r_y	1 m
ε_{out}	1 m	ε_{in}	1 m
\underline{u}_i	$-20 \frac{\text{m}}{\text{s}^2}$	\bar{u}_i	$20 \frac{\text{m}}{\text{s}^2}$
\bar{v}_i	$15 \frac{\text{m}}{\text{s}}$	$v_{i,ref}$	$8 \frac{\text{m}}{\text{s}}$
Q	$\text{diag}(5, 5)$	R	$\text{diag}(1, 1)$

A. Intersection scenario

1) *Conflict region*: As shown in Figure 4, once vehicles has been inside the intersection, side collision between vehicles in different lanes possibly occurs. In practice, it's too late for vehicles to start to consider side collision when it drive into intersection due to the braking distance not zero. Therefore, we make conflict region to expand intersection region with a buffer, as yellow region shown in Figure 4. In order to clearly describe this case, every enter or exit point is given an index here. The length of buffer for a given vehicle is denoted by

$$l_b = d_b + \bar{r}_x \quad (16)$$

where d_b is the minimum braking distance under the max speed, \bar{r}_x denote the longest semi-axis of vehicle ellipses.

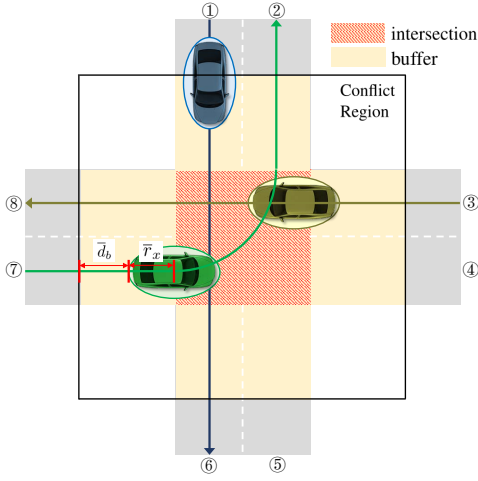


Fig. 4: Conflict Region illustrations

2) *Numerical result:* We make seven vehicles driving in this scenario, set the side length of intersection to $L_s = 6m$, time interval $h = 0.04s$ and total steps $K = 60$. For each vehicle i , in the mid point of entrance and exit of intersection, the scalar position $s_i = 0$. Here, all initial velocities and accelerations are set to $v_{i,ref}$ and 0. Table II presents the initial position, entrance index and exit index of all vehicles.

TABLE II: Vehicles' paths and initial states

i	enter index	exit index	$s_{i,0}/m$
1	3	8	-4
2	3	2	-11
3	7	2	-6
4	5	8	-10
5	7	4	-15
6	5	2	-16
7	1	6	-9

Figure 5 shows part of the trajectories for vehicles in this intersection scenarios. We choose the vehicle which is closet to the center of intersection as the root of tree graph. The evolution of tree structure is also shown in Figure 5.

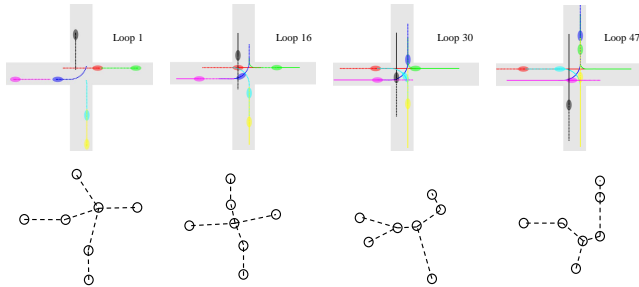


Fig. 5: Evolution of tree structure

Figure 6 shows the acceleration and velocity trajectories for vehicles. From the picture, we notice that some vehicles have quite big accelerations and velocities changed a lot when they need to avoid collision with other vehicles. When vehicles pass the intersection and the collision avoidance

constraints are no longer active, the accelerations move to zero and the velocities approach to reference velocity.

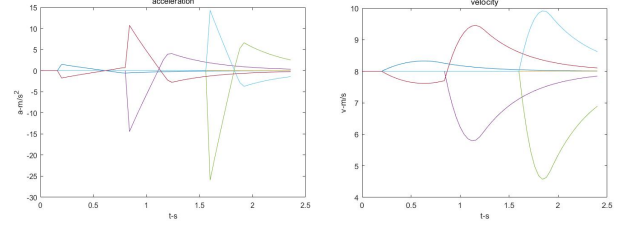


Fig. 6: Acceleration and velocity trajectories for vehicles

B. Lanes merge scenario

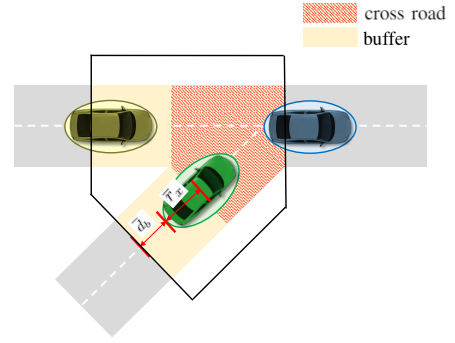


Fig. 7: Conflict Region illustrations

1) *Conflict region:* As shown in Figure 7, the definition of conflict region in merge scenario is similar to that in intersection scenario. One thing different is that there is no merge region on the merged road, because all the vehicles on the right side of the cross road have the same direction.

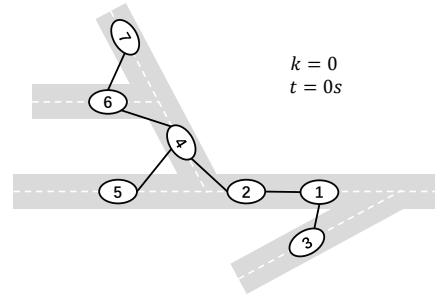


Fig. 8: Initial structure of merge scenario

2) *Numerical result:* A merge scenario is constructed as Figure 8, all vehicles will travel from left to right and eventually merge into one lane. We set the time interval $h = 0.05s$ and total steps $K = 20$. The scalar position of three merge points are respectively $2m$, $-10m$ and $-14m$, from right to left, and the angles are $\pi/6, \pi/3, \pi/3$ respectively. In order to Magnify the effect of collision avoidance, We enlarge the range of the state and control variables, $\underline{v}_i = 30m/s^2$

and $[u_i, \bar{u}_i] = [-60, 60]m/s^2$. Table III presents the initial positions and velocities of all vehicles.

TABLE III: Vehicles' initial states

i	$s_{i,0}/m$	$v_{i,0}/m \cdot s^{-2}$
1	-2	5
2	-7	8
3	-5	8
4	-11	4
5	-13	7
6	-15	3
7	-17	9

According to the Figure 10 trajectories of vehicles and Figure 9 evolution of tree structure in this merge scenario, movements of all vehicles are reasonable and interpretable. First, we focus on the behaviours of car 4, 5, 6 around time $t = 0.15s$. Before that time, the car 4 did not pass the local merge point, so there is no connection between car 5 and 6, their scalar positions s_i are very close. After that time, car 4 passed the local merge point, tree structure changed and connection arose between car 5 and 6. Therefore, car 5 speed up, car 6 speed down, and car 4, 7 also had corresponding movement because they need to keep distance with their neighbors. Similarly, the movements of car 1-4 around time $t = 0.65s$ can also be explained in the same way.

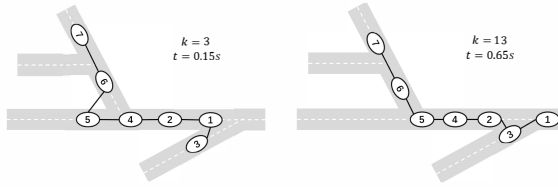


Fig. 9: Evolution of tree structures

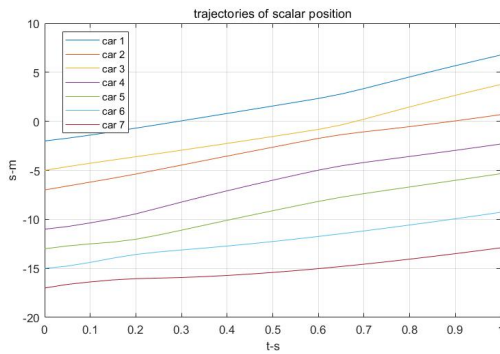


Fig. 10: Trajectories of vehicles' positions

VII. CONCLUSION

This paper has proposed a distributed closed-loop control algorithm for autonomous vehicles coordination in two-dimensional plane. Vehicles' dynamic are modeled by double integrator, the shape of vehicles is modeled as ellipse, and the predetermined paths are presented using B-splines. In order

to construct the closed-loop scheme, we use model predictive control method. Furthermore, we model the connections between vehicles as tree topology, and use a decentralized algorithm over tree graph, multi-sweep method, to solve the coordination problem online in each loop. Two cases show feasibility and good performance of this method. However, the obvious drawback of the proposed approach is the efficiency. Even it is a decentralized algorithm, each agent takes a unrealistically long time to calculate the optimal solution. Future research will focus on the improvement of vehicles and traffic systems modeling, as well as developing suitable distributed algorithms.

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