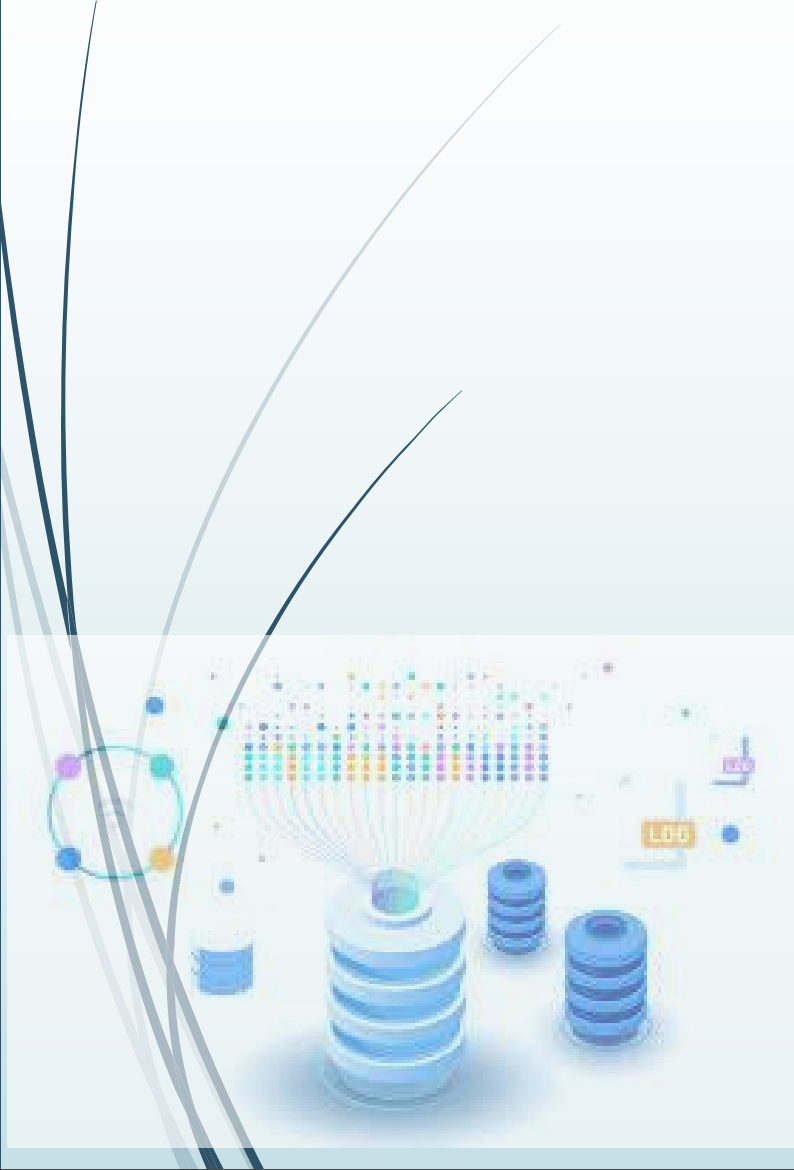


CHAPTER 2

THE RELATIONAL MODEL OF DATA



OBJECTIVES

Understand concepts of:

- ❖ What is the **relational model** and database design basing relational model.
- ❖ Conceptualize data using the relational model.
- ❖ Understand what basic **relational algebra operators** under set semantics.
- ❖ **Express queries** using relational algebra.

CONTENT

2.1 An Overview of Data Models

2.2 Basics of the Relational Model

2.3 An Algebraic Query Language

2.1 AN OVERVIEW OF DATA MODELS

- **Data model:** a collection of concepts for describing data, including 3 parts:
 - **Structure of the data**
 - Ex: arrays or objects
 - **Operations on the data**
 - Queries and modification on data
 - **Constraints on the data**
 - Limitations on the data

2.1 AN OVERVIEW OF DATA MODELS

- The **relational model**, including object-relational extensions
- The **semi-structured data model**, including XML and related standards
 - resembles **trees or graphs** rather than tables or arrays.
 - XML, a way to represent data by **hierarchically** nested tagged elements
 - Operations involve following **paths in tree** from an element to one or more of its nested sub elements, and so on
 - Constraints involve the data type of values associated with a nested tag

2.1 AN OVERVIEW OF DATA MODELS

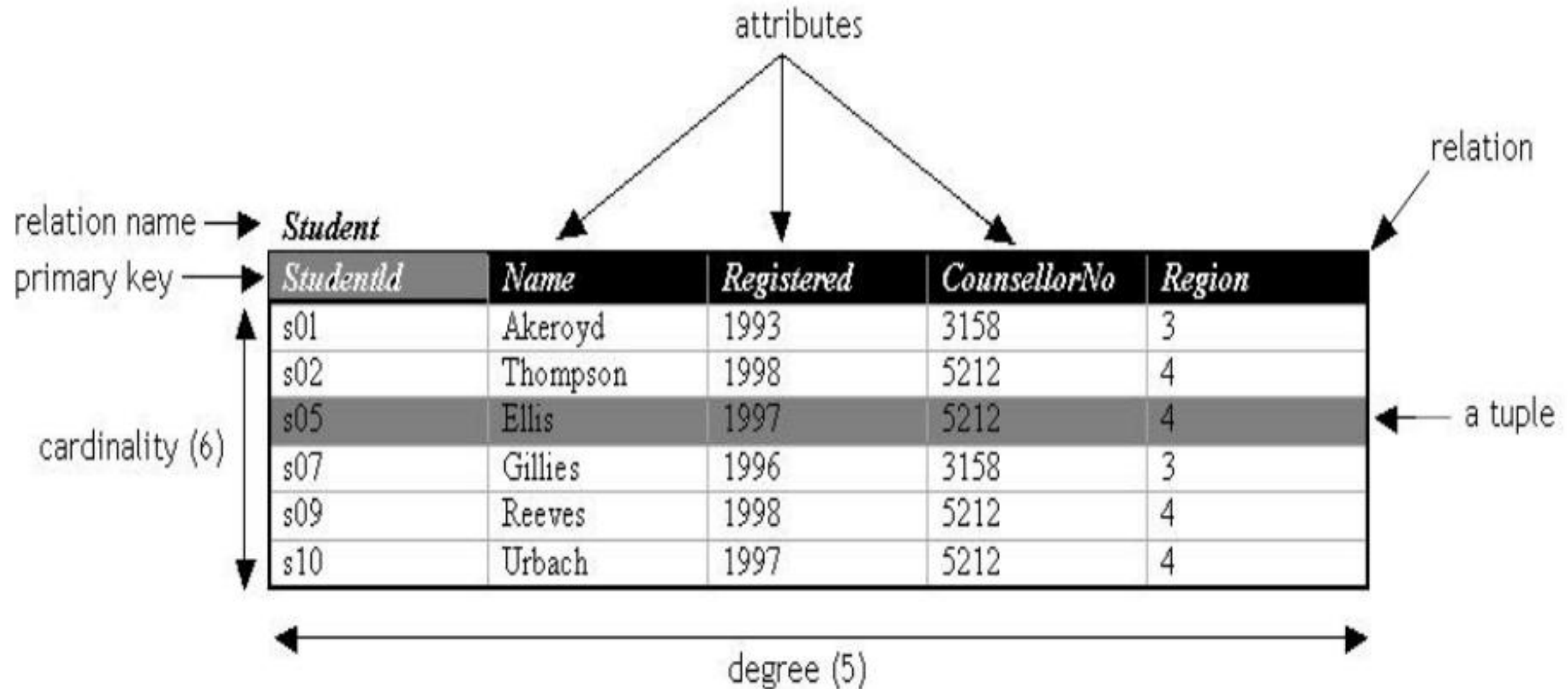
```
<?xml version="1.0"?>
<!DOCTYPE PARTS SYSTEM "parts.dtd">
<?xml-stylesheet type="text/css" href="xmlpartsstyle.css"?>
<PARTS>
  <TITLE>Computer Parts</TITLE>
  <PART>
    <ITEM>Motherboard</ITEM>
    <MANUFACTURER>ASUS</MANUFACTURER>
    <MODEL>P3B-F</MODEL>
    <COST> 123.00</COST>
  </PART>
  <PART>
    <ITEM>Video Card</ITEM>
    <MANUFACTURER>ATI</MANUFACTURER>
    <MODEL>All-in-Wonder Pro</MODEL>
    <COST> 160.00</COST>
  </PART>
  <PART>
    <ITEM>Sound Card</ITEM>
    <MANUFACTURER>Creative Labs</MANUFACTURER>
    <MODEL>Sound Blaster Live</MODEL>
    <COST> 80.00</COST>
  </PART>
  <PART>
    <ITEM>17 inch Monitor</ITEM>
    <MANUFACTURER>LG Electronics</MANUFACTURER>
    <MODEL> 995E</MODEL>
    <COST> 290.00</COST>
  </PART>
</PARTS>
```

2.2 BASICS OF THE RELATIONAL MODELS

■ Relational model

- A relation is made up from 2 parts:
 - Schema: specifies name of relation, name of attributes and domain/type of one's.
 - Ex: Student(StudentID: string, Name: string, Registered: int, CounsellorNo: int, Region: int)
 - Instance: a table with rows and columns
 - Rows ~ cardinality; columns ~ degree/arity
- **A simple thinking: a relation as a set of distinct rows or tuples**

2.2 BASICS OF THE RELATIONAL MODELS



2.2 BASICS OF THE RELATIONAL MODELS

- Database schema: a set of schemas for the relations of a database
- An example of DB schema:

- **Sailors**(sid: *integer*, sname: *string*, rating: *integer*, age: *real*)
- **Boats**(bid: *integer* , bname: *string*, color: *string*)
- **Reserves**(sid: *integer*, bid: *integer* , day: *date*)|

2.2 BASICS OF THE RELATIONAL MODELS

- Key attribute
- Non-key attribute
- Multi-valued attribute
- Derived- attribute
- Candidate key
- Primary key
- Foreign key

2.3 AN ALGEBRAIC QUERY LANGUAGE

➡ Relational Algebra

- An algebra consists of operators and atomic operands
- Relational algebra is an example of an algebra, its atomic operands are
 - Variables that stand for relations
 - Constants, which are finite relations
- Relational algebra is a set of operations on relations
- Operations operate on one or more relations to create new relation

2.3 AN ALGEBRAIC QUERY LANGUAGE

- Relational algebra fall into four classes
 - Set operations – union, intersection, difference
 - Selection and projection
 - Cartesian product and joins
 - Rename

2.3 AN ALGEBRAIC QUERY LANGUAGE

■ Set operations

Union

$$\mathbf{R} \cup \mathbf{S} = \{ t \mid t \in \mathbf{R} \vee t \in \mathbf{S} \}$$

Intersection

$$\mathbf{R} \cap \mathbf{S} = \{ t \mid t \in \mathbf{R} \wedge t \in \mathbf{S} \}$$

Difference

$$\mathbf{R} \setminus \mathbf{S} = \{ t \mid t \in \mathbf{R} \wedge t \notin \mathbf{S} \}$$

- Intersection can be expressed in terms of set difference

$$\mathbf{R} \cap \mathbf{S} = \mathbf{R} \setminus (\mathbf{R} \setminus \mathbf{S})$$

R and S must be 'type compatible'

- The same number of attributes
- The domain of corresponding attributes must be compatible

SET OPERATIONS- EXAMPLE

<i>name</i>	<i>address</i>	<i>gender</i>	<i>birthdate</i>
Carrie Fisher	123 Maple St., Hollywood	F	9/9/99
Mark Hamill	456 Oak Rd., Brentwood	M	8/8/88

Relation R

<i>name</i>	<i>address</i>	<i>gender</i>	<i>birthdate</i>
Carrie Fisher	123 Maple St., Hollywood	F	9/9/99
Harrison Ford	789 Palm Dr., Beverly Hills	M	8/8/88

Relation S

SET OPERATIONS- EXAMPLE

$R \cup S$

<i>name</i>	<i>address</i>	<i>gender</i>	<i>birthdate</i>
Carrie Fisher	123 Maple St., Hollywood	F	9/9/99
Mark Hamill	456 Oak Rd., Brentwood	M	8/8/88
Harrison Ford	789 Palm Dr., Beverly Hills	M	8/8/88

$R \cap S$

<i>name</i>	<i>address</i>	<i>gender</i>	<i>birthdate</i>
Carrie Fisher	123 Maple St., Hollywood	F	9/9/99

$R \setminus S$

<i>name</i>	<i>address</i>	<i>gender</i>	<i>birthdate</i>
Mark Hamill	456 Oak Rd., Brentwood	M	8/8/88

2.3.1. SELECTION

1. Selection

$R1 := \sigma_C (R2)$

with C illustrated conditions

Ex:

$$\begin{aligned} & \sigma_{\langle C1 \rangle} (\sigma_{\langle C2 \rangle} (R)) \\ &= \sigma_{\langle C2 \rangle} (\sigma_{\langle C1 \rangle} (R)) \\ &= \sigma_{\langle C1 \rangle \text{ AND } \langle C2 \rangle} (R) \end{aligned}$$

EX1 (Selection)

Movies table

<i>title</i>	<i>year</i>	<i>length</i>	<i>genre</i>
Gone With the Wind	1939	231	Drama
Star Wars	1977	124	Scifi
Wayne's World	1992	95	Comedy



$\sigma_{length \geq 100}(\text{Movies})$

<i>title</i>	<i>year</i>	<i>length</i>	<i>genre</i>
Gone With the Wind	1939	231	Drama
Star Wars	1977	124	Scifi

2.3.2. PROJECTION

2. Projection

$$S = \pi_{A1, A2, \dots, An} (R)$$

$A1, A2, \dots, An$ are attributes of R
 S relation schema $S(A1, A2, \dots, An)$

EX2: Projection

Movies table

<i>title</i>	<i>year</i>	<i>length</i>	<i>genre</i>
Star Wars	1977	124	Scifi
Galaxy Quest	1999	104	Comedy
Wayne's World	1992	95	Comedy

$\pi_{title, year, length}(\text{Movies})$

<i>title</i>	<i>year</i>	<i>length</i>
Star Wars	1977	124
Galaxy Quest	1999	104
Wayne's World	1992	95

$\pi_{genre}(\text{Movies})$

<i>genre</i>
Scifi
Comedy

2.3.3. CARTESIAN PRODUCT AND JOINS

a. Cartesian product $R3 := R1 \times R2$

Relation R	
A	B
1	2
3	4

Relation S		
B	C	D
2	5	6
4	7	8
9	10	11

Cartesian Product R X S				
A	R.B	S.B	C	D
1	2	2	5	6
1	2	4	7	8
1	2	9	10	11
3	4	2	5	6
3	4	4	7	8
3	4	9	10	11

2.3.3. CARTESIAN PRODUCT AND JOINS

b. theta joins $R3 := R1 \bowtie_{\langle \text{join condition} \rangle} R2$

A	B	C
1	2	3
6	7	8
9	7	8

Relation U

B	C	D
2	3	4
2	3	5
7	8	10

Relation V

A	U.B	U.C	V.B	V.C	D
1	2	3	7	8	10

Result of $U \bowtie_{A < D \text{ AND } U.B \neq V.B} V$

A	U.B	U.C	V.B	V.C	D
1	2	3	2	3	4
1	2	3	2	3	5
1	2	3	7	8	10
6	7	8	7	8	10
9	7	8	7	8	10

$U \bowtie_{A < D} V$

2.3.3. CARTESIAN PRODUCT AND JOINS

c. Natural join $R_3 := R_1 \bowtie R_2$

Relation R

A	B
1	2
3	4

Relation S

B	C	D
2	5	6
4	7	8
9	10	11

Natural Join $R \bowtie S$

A	B	C	D
1	2	5	6
3	4	7	8

Rename

The ρ operation gives a new schema to a relation

$\rho_{S(A_1, \dots, A_n)}(R)$ makes S be a relation with attributes A_1, \dots, A_n and the same tuples as R

Simplified notation: $S := R(A_1, A_2, \dots, A_n)$

Relation R

A	B
1	2
3	4

Relation S

B	C	D
2	5	6
4	7	8
9	10	11

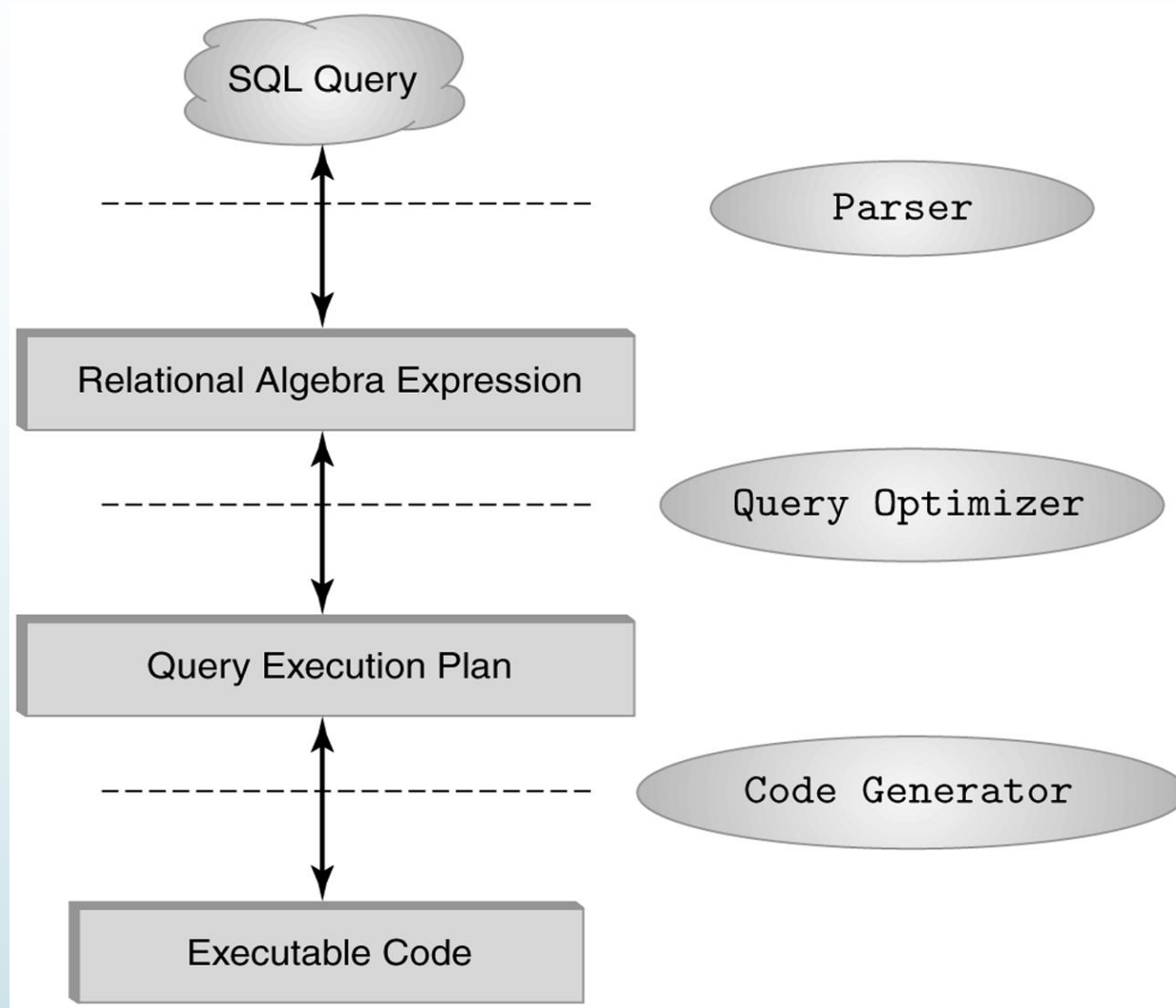
$R \times \rho_{S(X, C, D)}(S)$

A	B	X	C	D
1	2	2	5	6
1	2	4	7	8
1	2	9	10	11
3	4	2	5	6
3	4	4	7	8
3	4	9	10	11

2.3.4. RELATIONAL EXPRESSION

- How we need relational expression
- Relational algebra allows us to form expressions
- Relational expression is constructed by applying operations to the result of other operations
- Expressions can be presented as expression tree

The role of relational algebra in a DBMS

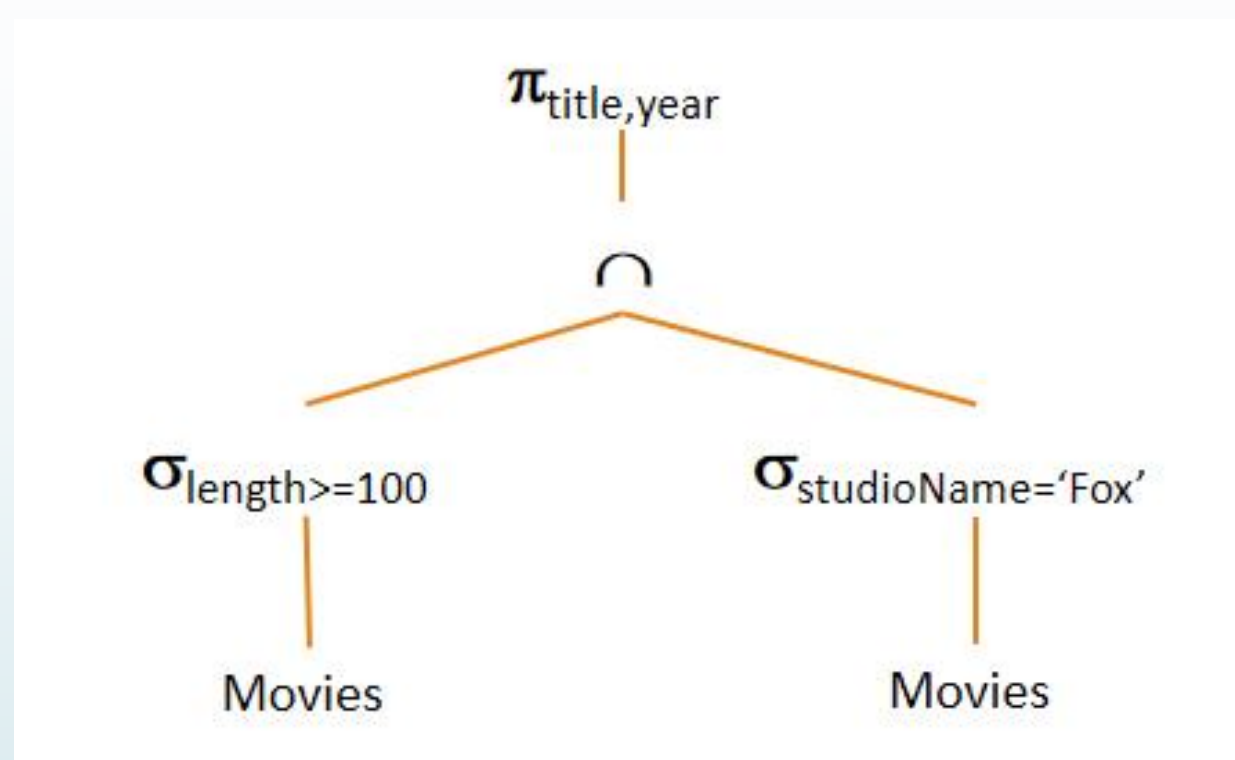


Relational Expression

Example: What are the titles and years of movies made by Fox that are at least 100 minutes long?

- (1) Select those Movies tuples that have length ≥ 100
- (2) Select those Movies tuples that have studioName='Fox'
- (3) Compute the intersection of (1) and (2)
- (4) Project the relation from (3) onto attributes title and year

Relational Expression



Expression tree for a relational algebra expression

$\pi_{\text{title, year}}(\sigma_{\text{length} \geq 100}(\mathbf{Movies}) \cap \sigma_{\text{studioName} = \text{'Fox'}}(\mathbf{Movies}))$

$\pi_{\text{title, year}}(\sigma_{\text{length} \geq 100 \text{ AND studioName} = \text{'Fox'}}(\mathbf{Movies}))$

EXERCISE

Product(maker, model, type)

PC(model, speed, ram, hd, price)

Laptop(model, speed, ram, hd, screen, price) **Printer**(model, color, type, price)

- What PC models have a speed of at least 3.00?
- Which manufacturers make laptops with a hard disk of at least 100GB?
- Find the model number and price of all products (of any type) made by manufacturer B.
- Find the model numbers of all color laser printers.
- Find those manufacturers that sell Laptops, but not PC's.