

A-star

This is a specific case of the [Best-first search](#) that uses the following evaluation function:

$$f(n) = g(n) + h(n)$$

Where:

- $g(n)$ is the cost to reach n .
- $h(n)$ = estimated movement cost to move from n to the final destination. [Heuristic function](#).

So that $f(n)$ is the expected cost of the most convenient walk from the initial state to the goal state that covers node n .

At each step we choose the node from the frontier with minimal value of $f(n)$.

Whether A* is cost-optimal depends on certain properties of the heuristic:

- Admissibility
- Consistency

Admissibility and consistency

An heuristic function is said to be admissible if it never overestimates the cost of reaching the goal.

Let $H(n)$ be the actual path from n to G .

$h(n)$ must never surpass $H(n)$.

An admissible heuristic is one that never overestimates the cost to reach a goal.

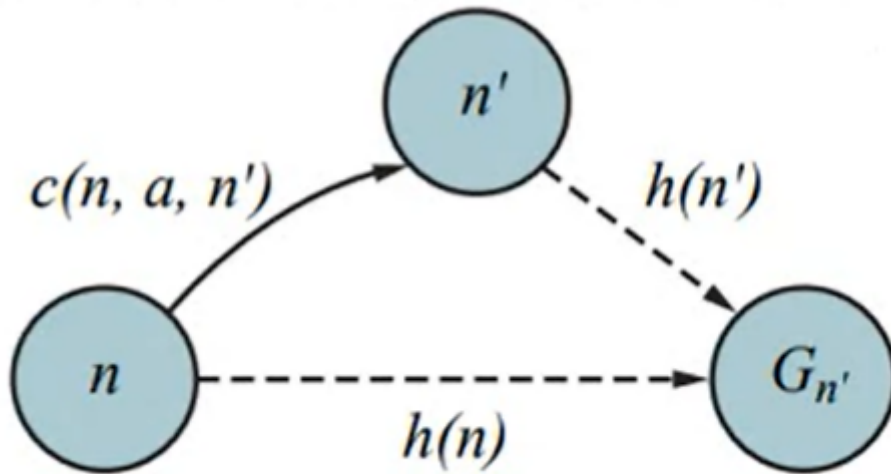
So an admissible heuristic must be "optimistic".

Admissibility: $h(n)$ is \leq to the minimum path cost from n to the goal

Consistency:

For every node n and its successor n' of n generated by action a ,

$$h(n) \leq c(n, a, n') + h(n')$$



Every consistent heuristic is admissible but not viceversa.

With a consistent heuristic, the first time we reach a state, it is on the optimal path.

If there is a cost-optimal path on which $h(n)$ is admissible for all nodes on the path, that path will be found no matter what.

If the optimal solution has cost C , the second best has cost C_2 , and $h(n)$ overestimates by never more than $C_2 - C$, then jjuhghhjb