OPERATION and OPTIONS Values Produced by EXPLAIN PLAN

Excerpt taken from Oracle DB's official doc (version 12.c release 1):
https://docs.oracle.com/database/121/TGSQL/tgsql_interp.htm#GUID-D7331098-AEF5-49A2-98A1-3039E2A5900C__CEGDHHFH

Operation	Option	Description
AND-EQUAL		Operation accepting multiple sets of rowids, returning the intersection of the sets, eliminating duplicates. Used for the single-column indexes access path.
BITMAP	CONVERSION	TO ROWIDS converts bitmap representations to actual rowids that you can use to access the table. FROM ROWIDS converts the rowids to a bitmap representation. COUNT returns the number of rowids if the actual values are not needed.
BITMAP	INDEX	SINGLE VALUE looks up the bitmap for a single key value in the index. RANGE SCAN retrieves bitmaps for a key value range. FULL SCAN performs a full scan of a bitmap index if there is no start or stop key.
BITMAP	MERGE	Merges several bitmaps resulting from a range scan into one bitmap.
BITMAP	MINUS	Subtracts bits of one bitmap from another. Row source is used for negated predicates. Use this option only if there are nonnegated predicates yielding a bitmap from which the subtraction can take place. An example appears in "Viewing Bitmap Indexes with EXPLAIN PLAN".
BITMAP	OR	Computes the bitwise OR of two bitmaps.
BITMAP	AND	Computes the bitwise AND of two bitmaps.
BITMAP	KEY ITERATION	Takes each row from a table row source and finds the corresponding bitmap from a bitmap index. This set of bitmaps are then merged into one bitmap in a following BITMAP MERGE operation.
CONNECT BY		Retrieves rows in hierarchical order for a query containing a CONNECT BY clause.
CONCATENATION		Operation accepting multiple sets of rows returning the union-all of the sets.
COUNT		Operation counting the number of rows selected from a table.
COUNT	STOPKEY	Count operation where the number of rows returned is limited by the ROWNUM expression in the WHERE clause.
CUBE SCAN		Uses inner joins for all cube access.
CUBE SCAN	PARTIAL OUTER	Uses an outer join for at least one dimension, and inner joins for the other dimensions.
CUBE SCAN	OUTER	Uses outer joins for all cube access.

Operation	Option	Description
DOMAIN INDEX		Retrieval of one or more rowids from a domain index. The options column contain information supplied by a user-defined domain index cost function, if any.
FILTER		Operation accepting a set of rows, eliminates some of them, and returns the rest.
FIRST ROW		Retrieval of only the first row selected by a query.
FOR UPDATE		Operation retrieving and locking the rows selected by a query containing a FOR UPDATE clause.
HASH	GROUP BY	Operation hashing a set of rows into groups for a query with a GROUP BY clause.
HASH	GROUP BY PIVOT	Operation hashing a set of rows into groups for a query with a GROUP BY clause. The PIVOT option indicates a pivot-specific optimization for the HASH GROUP BY operator.
HASH JOIN (These are join operations.)		Operation joining two sets of rows and returning the result. This join method is useful for joining large data sets of data (DSS, Batch). The join condition is an efficient way of accessing the second table. Query optimizer uses the smaller of the two tables/data sources to build a hash table on the join key in memory. Then it scans the larger table, probing the hash table to find the joined rows.
HASH JOIN	ANTI	Hash (left) antijoin
HASH JOIN	SEMI	Hash (left) semijoin
HASH JOIN	RIGHT ANTI	Hash right antijoin
HASH JOIN	RIGHT SEMI	Hash right semijoin
HASH JOIN	OUTER	Hash (left) outer join
HASH JOIN	RIGHT OUTER	Hash right outer join
INDEX (These are access methods.)	UNIQUE SCAN	Retrieval of a single rowid from an index.
INDEX	RANGE SCAN	Retrieval of one or more rowids from an index. Indexed values are scanned in ascending order.
INDEX	RANGE SCAN DESCENDING	Retrieval of one or more rowids from an index. Indexed values are scanned in descending order.
INDEX	FULL SCAN	Retrieval of all rowids from an index when there is no start or stop key. Indexed values are scanned in ascending order.
INDEX	FULL SCAN DESCENDING	Retrieval of all rowids from an index when there is no start or stop key. Indexed values are scanned in descending order.

Operation	Option	Description
INDEX	FAST FULL SCAN	Retrieval of all rowids (and column values) using multiblock reads. No sorting order can be defined. Compares to a full table scan on only the indexed columns. Only available with the cost based optimizer.
INDEX	SKIP SCAN	Retrieval of rowids from a concatenated index without using the leading column(s) in the index. Only available with the cost based optimizer.
INLIST ITERATOR		Iterates over the next operation in the plan for each value in the IN-list predicate.
INTERSECTION		Operation accepting two sets of rows and returning the intersection of the sets, eliminating duplicates.
MERGE JOIN (These are join operations.)		Operation accepting two sets of rows, each sorted by a value, combining each row from one set with the matching rows from the other, and returning the result.
MERGE JOIN	OUTER	Merge join operation to perform an outer join statement.
MERGE JOIN	ANTI	Merge antijoin.
MERGE JOIN	SEMI	Merge semijoin.
MERGE JOIN	CARTESIAN	Can result from 1 or more of the tables not having any join conditions to any other tables in the statement. Can occur even with a join and it may not be flagged as CARTESIAN in the plan.
CONNECT BY		Retrieval of rows in hierarchical order for a query containing a CONNECT BY clause.
MAT_VIEW REWITE ACCESS (These are access methods.)	FULL	Retrieval of all rows from a materialized view.
MAT VIEW REWITE ACCESS	SAMPLE	Retrieval of sampled rows from a materialized view.
MAT VIEW REWITE ACCESS	CLUSTER	Retrieval of rows from a materialized view based on a value of an indexed cluster key.
MAT VIEW REWITE ACCESS	HASH	Retrieval of rows from materialized view based on hash cluster key value.
MAT_VIEW REWITE ACCESS	BY ROWID RANGE	Retrieval of rows from a materialized view based on a rowid range.
MAT_VIEW REWITE ACCESS	SAMPLE BY ROWID RANGE	Retrieval of sampled rows from a materialized view based on a rowid range.
MAT_VIEW REWITE ACCESS	BY USER ROWID	If the materialized view rows are located using user-supplied rowids.

Operation	Option	Description
MAT_VIEW REWITE ACCESS	BY INDEX ROWID	If the materialized view is nonpartitioned and rows are located using index(es).
MAT_VIEW REWITE ACCESS	BY GLOBAL INDEX ROWID	If the materialized view is partitioned and rows are located using only global indexes.
MAT VIEW REWITE ACCESS	BY LOCAL INDEX ROWID	If the materialized view is partitioned and rows are located using one or more local indexes and possibly some global indexes. Partition Boundaries: The partition boundaries might have been computed by: A previous PARTITION step, in which case the PARTITION START and PARTITION STOP column values replicate the values present in the PARTITION step, and the PARTITION_ID contains the ID of the PARTITION step. Possible values for PARTITION_START and PARTITION_STOP are NUMBER(n), KEY, INVALID. The MAT_VIEW REWRITE ACCESS or INDEX step itself, in which case the PARTITION_ID contains the ID of the step. Possible values for PARTITION START and PARTITION STOP are NUMBER(n), KEY, ROW REMOVE LOCATION(MAT_VIEW REWRITE ACCESS only), and INVALID.
MINUS		Operation accepting two sets of rows and returning rows appearing in the first set but not in the second, eliminating duplicates.
NESTED LOOPS (These are join operations.)		Operation accepting two sets of rows, an outer set and an inner set. Oracle Database compares each row of the outer set with each row of the inner set, returning rows that satisfy a condition. This join method is useful for joining small subsets of data (OLTP). The join condition is an efficient way of accessing the second table.
NESTED LOOPS	OUTER	Nested loops operation to perform an outer join statement.
PARTITION		Iterates over the next operation in the plan for each partition in the range given by the PARTITION START and PARTITION STOP columns. PARTITION describes partition boundaries applicable to a single partitioned object (table or index) or to a set of equipartitioned objects (a partitioned table and its local indexes). The partition boundaries are provided by the values of PARTITION_START and PARTITION_STOP of the PARTITION. Refer to Table 7-1 for valid values of partition start and stop.
PARTITION	SINGLE	Access one partition.
PARTITION	ITERATOR	Access many partitions (a subset).
PARTITION	ALL	Access all partitions.
PARTITION	INLIST	Similar to iterator, but based on an IN-list predicate.
PARTITION	INVALID	Indicates that the partition set to be accessed is empty.

Operation	Option	Description
PX ITERATOR	BLOCK, CHUNK	Implements the division of an object into block or chunk ranges among a set of parallel execution servers.
PX COORDINATOR		Implements the query coordinator that controls, schedules, and executes the parallel plan below it using parallel execution servers. It also represents a serialization point, as the end of the part of the plan executed in parallel and always has a PX SEND QC operation below it.
PX PARTITION		Same semantics as the regular PARTITION operation except that it appears in a parallel plan.
PX RECEIVE		Shows the consumer/receiver parallel execution node reading repartitioned data from a send/producer (QC or parallel execution server) executing on a PX SEND node. This information was formerly displayed into the DISTRIBUTION column. See Table 7-2.
PX SEND	QC (RANDOM), HASH, RANGE	Implements the distribution method taking place between two parallel execution servers. Shows the boundary between two sets and how data is repartitioned on the send/producer side. This information was formerly displayed into the DISTRIBUTION column. See Table 7-2.
REMOTE		Retrieval of data from a remote database.
SEQUENCE		Operation involving accessing values of a sequence.
SORT	AGGREGATE	Retrieval of a single row that is the result of applying a group function to a group of selected rows.
SORT	UNIQUE	Operation sorting a set of rows to eliminate duplicates.
SORT	GROUP BY	Operation sorting a set of rows into groups for a query with a GROUP BY clause.
SORT	GROUP BY PIVOT	Operation sorting a set of rows into groups for a query with a GROUP BY clause. The PIVOT option indicates a pivot-specific optimization for the SORT GROUP BY operator.
SORT	JOIN	Operation sorting a set of rows before a merge-join.
SORT	ORDER BY	Operation sorting a set of rows for a query with an ORDER BY clause.
TABLE ACCESS (These are access methods.)	FULL	Retrieval of all rows from a table.
TABLE ACCESS	SAMPLE	Retrieval of sampled rows from a table.
TABLE ACCESS	CLUSTER	Retrieval of rows from a table based on a value of an indexed cluster key.
TABLE ACCESS	HASH	Retrieval of rows from table based on hash cluster key value.
TABLE ACCESS	BY ROWID RANGE	Retrieval of rows from a table based on a rowid range.

Operation	Option	Description
TABLE ACCESS	SAMPLE BY ROWID RANGE	Retrieval of sampled rows from a table based on a rowid range.
TABLE ACCESS	BY USER ROWID	If the table rows are located using user-supplied rowids.
TABLE ACCESS	BY INDEX ROWID	If the table is nonpartitioned and rows are located using indexes.
TABLE ACCESS	BY GLOBAL INDEX ROWID	If the table is partitioned and rows are located using only global indexes.
TABLE ACCESS	BY LOCAL INDEX ROWID	If the table is partitioned and rows are located using one or more local indexes and possibly some global indexes. Partition Boundaries: The partition boundaries might have been computed by: A previous PARTITION step, in which case the PARTITION START and PARTITION STOP column values replicate the values present in the PARTITION step, and the PARTITION_ID contains the ID of the PARTITION step. Possible values for PARTITION_START and PARTITION_STOP are NUMBER(n), KEY, INVALID. The TABLE ACCESS or INDEX step itself, in which case the PARTITION_ID contains the ID of the step. Possible values for PARTITION START and PARTITION STOP are NUMBER(n), KEY, ROW REMOVE_LOCATION (TABLE ACCESSonly), and INVALID.
TRANSPOSE		Operation evaluating a PIVOT operation by transposing the results of GROUP BY to produce the final pivoted data.
UNION		Operation accepting two sets of rows and returns the union of the sets, eliminating duplicates.
UNPIVOT		Operation that rotates data from columns into rows.
VIEW		Operation performing a view's query and then returning the resulting rows to another operation.