

1 Master theorem Examples - part 2

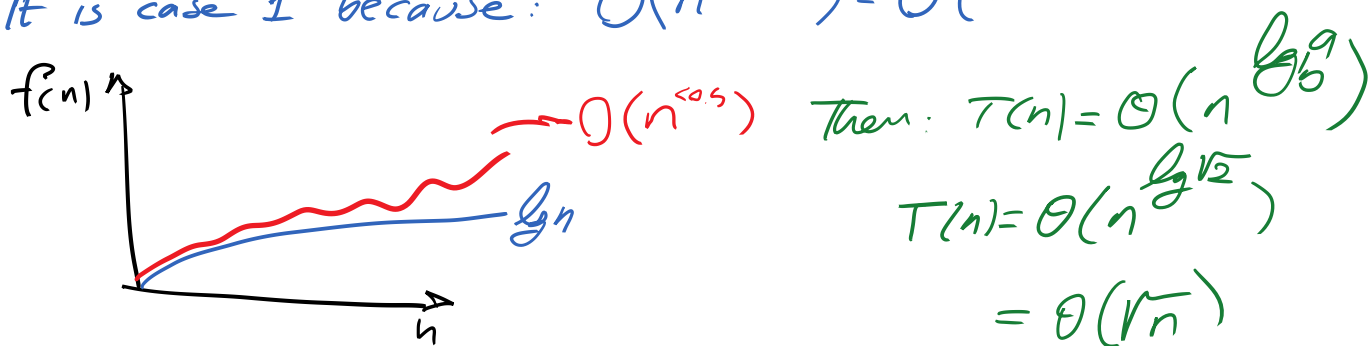
⑪ $T(n) = \sqrt{2}T(n/2) + \log n$:

$a = \sqrt{2}$

$b = 2$

$f(n) = \log n$

It is case 1 because: $O(n^{\frac{\log a}{\log b}}) = O(n^{\frac{\log \sqrt{2}}{\log 2}}) = O(n^{\frac{1}{2} - \epsilon}) \rightarrow < 0.5$



⑬ $T(n) = 3T(n/3) + \sqrt{n}$

$a = 3$

$b = 3$

$f(n) = \sqrt{n}$

It is case 1: $O(n^{\frac{\log 3}{\log 3} - \epsilon}) = O(n^{1 - \epsilon})$ for small $\epsilon > 0 \rightarrow \sqrt{n} \leq O(n^{1 - \epsilon})$

Then: $T(n) = O(n^{\frac{\log 3}{\log 3}}) = O(n)$

⑭ $T(n) = 4T(n/2) + cn$

$a = 4$

$b = 2$

$f(n) = cn$

It is case 2 because: $O(n^{\frac{\log 4}{\log 2}}) = O(n^2)$

for all $0 < \epsilon < 2$, $n \leq O(n^{2 - \epsilon})$

Then $T(n) = O(n^2)$

$$15 \quad T(n) = 3T\left(\frac{n}{4}\right) + n \lg n$$

$$a = 3$$

$$b = 4$$

$$f(n) = n \lg n$$

It is not case 1 because: $O\left(n^{\frac{3-\epsilon}{4}}\right) < 1$
 Then $n \lg n > n^{\frac{3-\epsilon}{4}}$

It is not case 2: $\Theta\left(n^{\frac{3}{4}} \lg^k n\right) \Rightarrow$
 we cannot find a range for k to sandwich $f(n)$

for $k=0 \rightarrow n^{\frac{3}{4}} (\lg n)^0 = n^{\alpha} < n$
 $\alpha < 1$

$$k=1 \rightarrow n^{\alpha} \lg n$$

It might be case 3: $\Omega\left(n^{\frac{3+\epsilon}{4}}\right) \rightarrow n \lg n > \Omega\left(n^{\beta}\right)$
 $f(n) \gg \Omega\left(n^{\frac{3+\epsilon}{4}}\right)$

But we have to examine regularity condition before applying the rule in case 3.

$$3\left(\frac{n}{4} \lg \frac{n}{4}\right) \leq c n \lg n$$

$$\frac{3n}{4} \lg\left(\frac{n}{4}\right) \leq c n \lg n$$

$$\frac{3}{4} (\lg n - \lg 4) \leq c \lg n$$

for large $n \rightarrow \frac{3}{4} \lg n \leq c \lg n \rightarrow \frac{3}{4} \leq c < 1$

Then: $T(n) = \Theta(n \lg n)$

$$(6) T(n) = 3T(n/3) + \frac{n}{2}$$

$$a=3$$

$$b=3$$

$$f(n) = \frac{n}{2}$$

It is not case 1 because: $\frac{n}{2} > O(n \lg^{3-\epsilon} n)$ \leftarrow linear

It is case 2: $\theta(n \lg^3 n) = \theta(n \lg^k n)$

for $k=0 \rightarrow f(n) \gg \theta(n)$

for $k=1 \rightarrow f(n) < \theta(n \lg n)$

Then: $T(n) = \theta(n \lg^3 n) = \theta(n \lg n)$

$$(7) T(n) = 6T(n/3) + n^2 \lg n$$

$$a=6$$

$$b=3$$

$$f(n) = n^2 \lg n$$

It is case 3 because: $n^2 \lg n \gg \Omega(n \lg^{6+\epsilon} n)$ \leftarrow for $0 < \epsilon < 3$

We have to examine regularity condition:

$$6\left(\frac{n}{3}\right)^2 \lg \frac{n}{3} \leq c n^2 \lg n$$

$$\frac{6}{9} n^2 \lg \frac{n}{3} \leq c n^2 \lg n$$

\hookrightarrow for large $n \rightarrow \lg \frac{n}{3} \approx \lg n$

$$\frac{6}{9} < c < 1$$

Then: $T(n) = \theta(n^2 \lg n)$

$$(19) T(n) = 64T(n/8) - n^2 \lg n$$

$$a = 64$$

$$b = 8$$

$$f(n) = \underline{-n^2 \lg n}$$

$f(n)$ is negative.

Master theorem cannot be applied.

$f(n)$ must be positive.

$$(20) T(n) = 7T(n/3) + n^2$$

$$a = 7$$

$$b = 3$$

$$f(n) = n^2$$

It is case 3, because $n^2 > \Omega(n^{\lg_{\frac{7-\epsilon}{3}}}) \rightarrow \epsilon < 2$ for all $0 < \epsilon < 2$

But we have to examine regularity condition:

$$7\left(\frac{n}{3}\right)^2 \leq cn^2 \rightarrow \frac{7}{9}n^2 \leq cn^2 \rightarrow \frac{7}{9} \leq c < 1$$

Then: $\boxed{T(n) = \Theta(n^2)}$

$$(21) T(n) = 4T(n/2) + \lg n$$

$$a = 4$$

$$b = 2$$

$$f(n) = \lg n$$

It is case 1, $\lg n \leq O(n^{\lg_{\frac{4-\epsilon}{2}}}) \rightarrow$ for $0 < \epsilon < 2$

Then $T(n) = \Theta(n^{\lg_{\frac{4}{2}}}) = \Theta(n^2)$

