A2 Verteilte Systemsoftware Part 1 – <u>Intro "Distributed Systems"</u>

aus: Coulouris / Dollimore / Kindberg: "Distributed Systems: Concepts and Design", Addison Wesley/Pearson Internat., 5. Ausg., 2012

<u>"Distributed systems" - simple definition:</u> HW+SW components, located on global networks, which communicate & coordinate their actions only by message passing; important: common global goal!

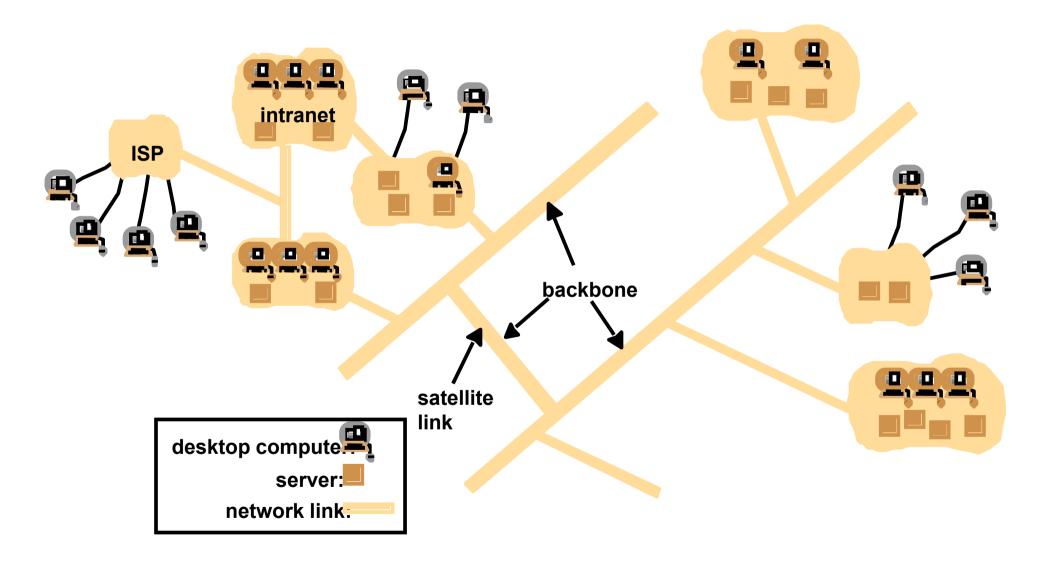
Typical characteristics:

- Concurrency: → siehe "Nebenläufigkeit & Verteilung"
- No global clock: → s. "Synchronisation" (+ Vorl. "VIS" später)
- Independent failures: → alle Kapitel + "Transparenz"-Eigenschaften

Beispiele:



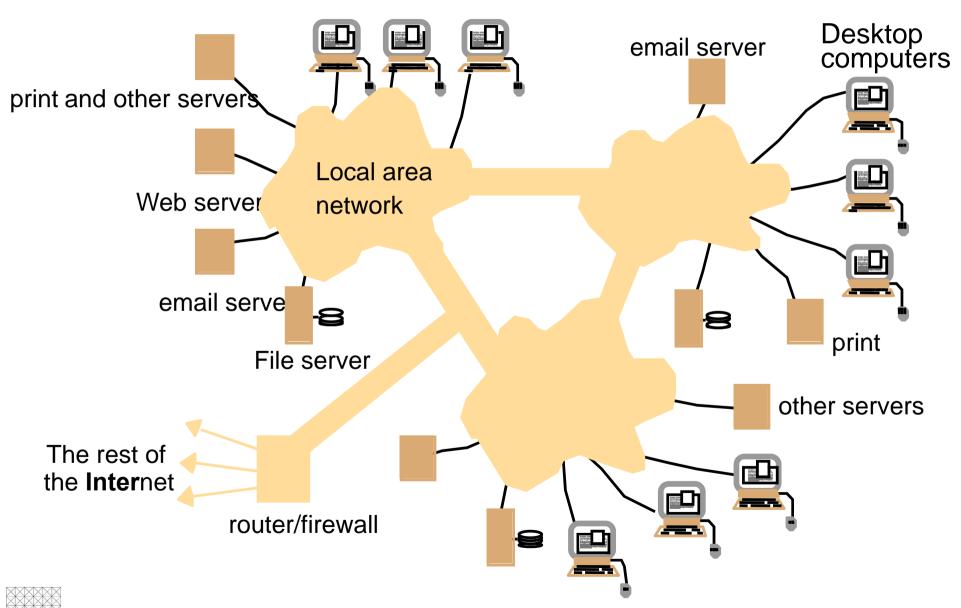
Distributed Systems: "A typical portion of the Internet"



Communication, System software (i.e. Operating System), Middleware, Security



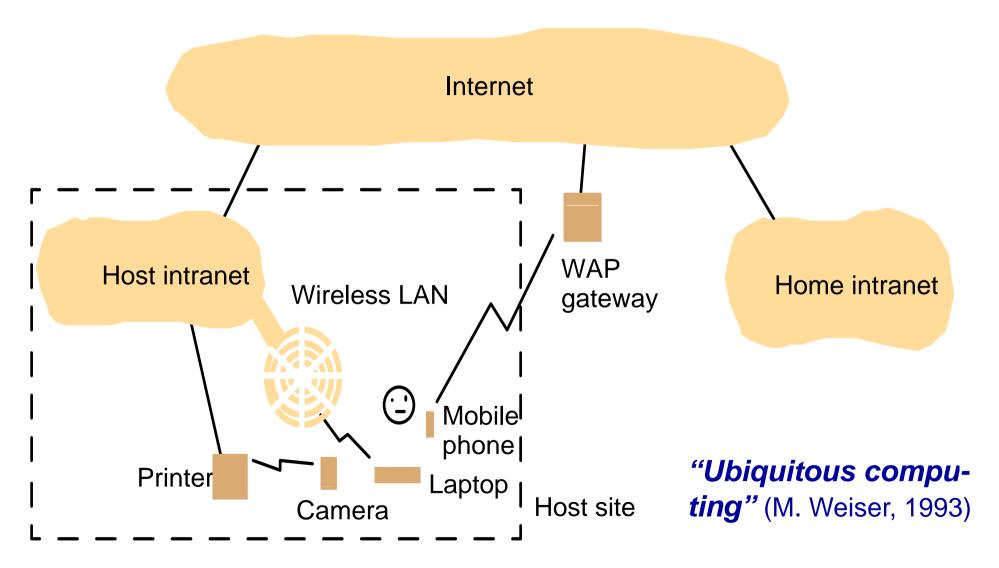
A Typical Part of an *Intra*net





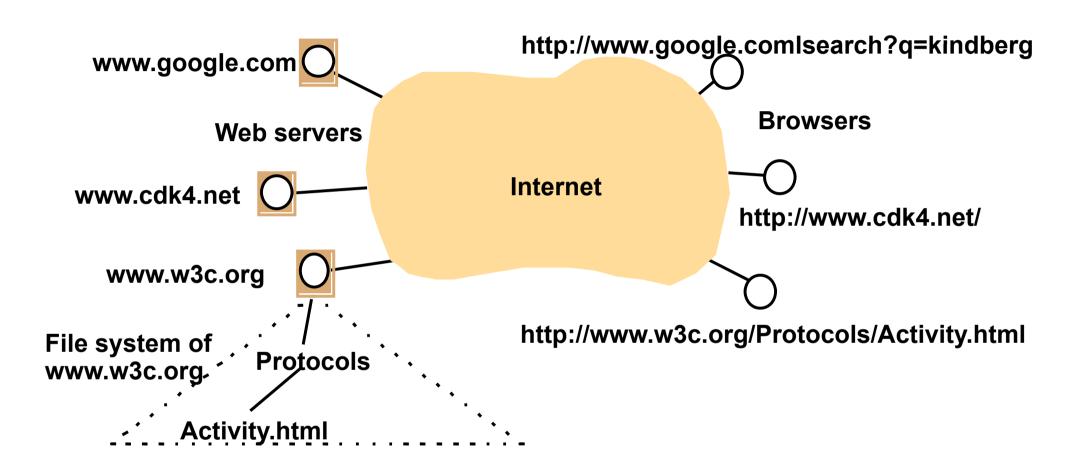
"Mobile Systems":

Portable and handheld devices in a distributed system





(Web) Server: Web servers and web browsers



Ressource sharing in the web: Servers, clients, remote service invocation, resource location (URL), interfaces (HTML), protocols (http), web services, service-oriented computing/ architecture (SOA)



Number of Computers in the Internet

Challenges:

- Heterogeneity
- Openness
- Scalability
- Security

Datum	# Computer (mit reg. IP-Addr.
1979, Dez.	188
1989, Juli	130,000
1999, Juli	56,218,000
2003, Jan.	171,638,297
 2011	über 2 Mrd. IN Benutzer

Inzwischen:

"Internet of Things"



Distributed Systems Goals: "Transparencies" (1)

Overall goal: "Distribution Transparency" – can be subdivided into:

Access transparency: enables local and remote resources to be accessed using identical operations.

Location transparency: enables resources to be accessed without knowledge of their location.

Concurrency transparency: enables several processes to operate concurrently using shared resources without interference between them.

Replication transparency: enables multiple instances of resources to be used to increase reliability and performance without knowledge of the replicas by users or application programmers.



Distributed Systems Goals: "Transparencies" (2)

Failure transparency: enables the concealment of faults, allowing users and application programs to complete their tasks despite the failure of hardware or software components.

Mobility transparency: allows the movement of resources and clients within a system without affecting the operation of users or programs.

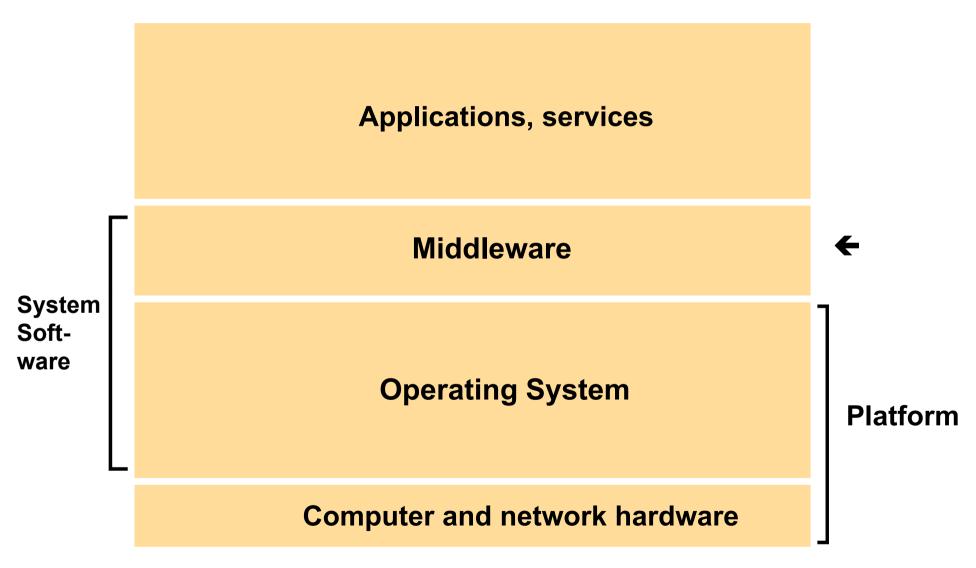
Performance transparency: allows the system to be *reconfigured* to improve performance as loads vary.

Scaling transparency: allows the system and applications to *expand* in scale without change to the system structure or the application algorithms.



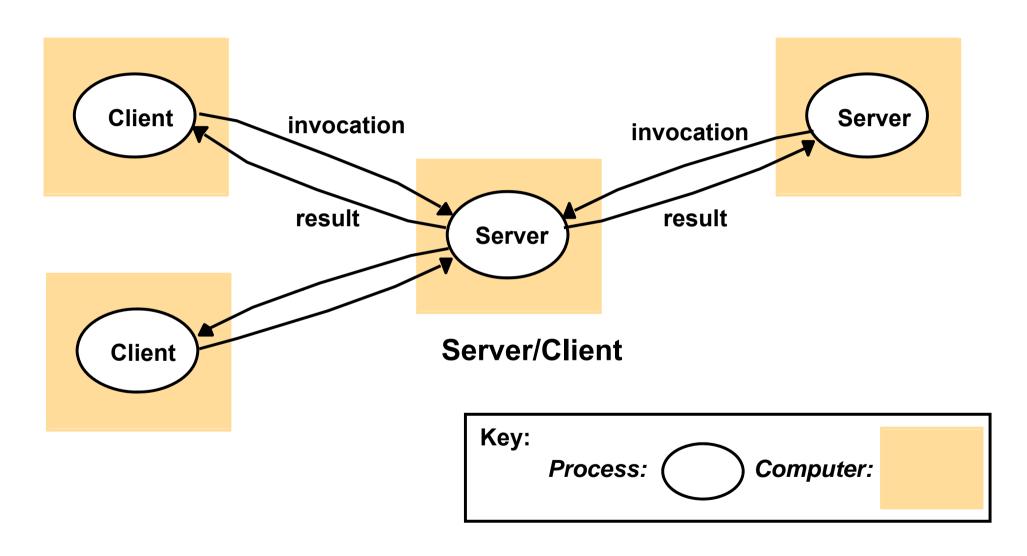
Part 2 - System Models:

Software and hardware service layers in distributed systems





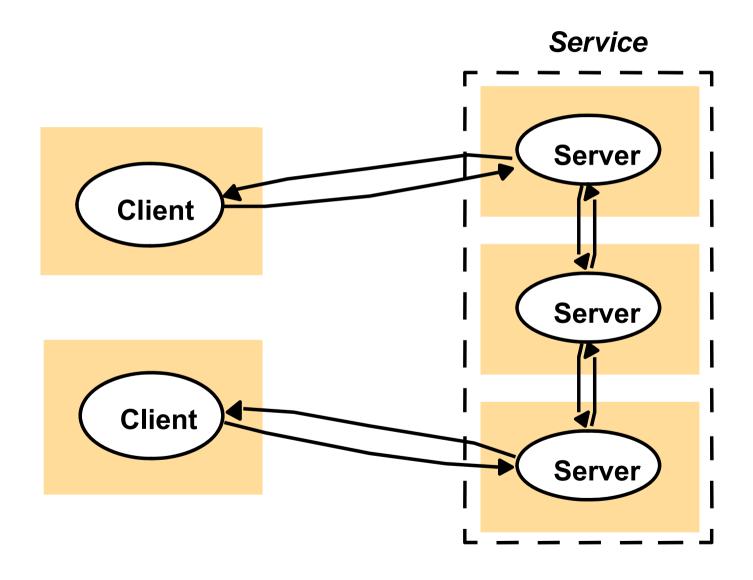
Client/Server: (a) Clients invoke individual servers





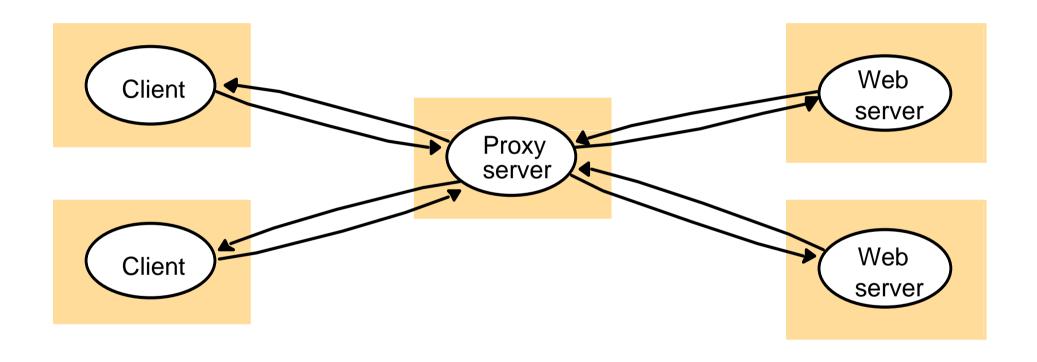
Client/Server:

(b) A service provided by multiple servers





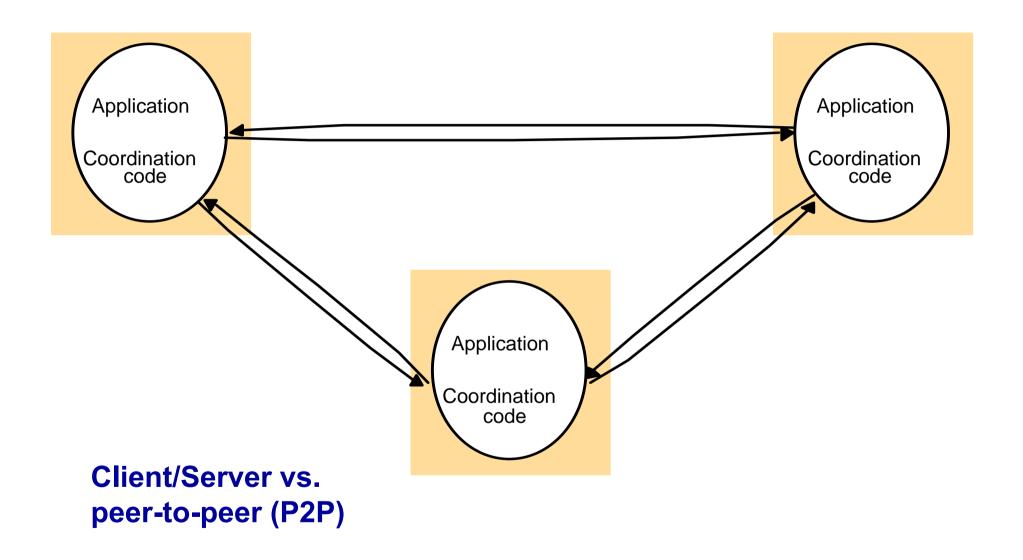
Client/Server: (c) Web proxy server



Proxy server: shared cache of web resources



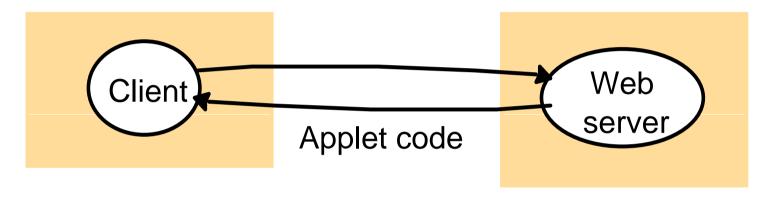
P2P: A distributed application based on peer processes



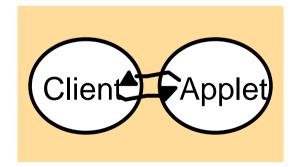


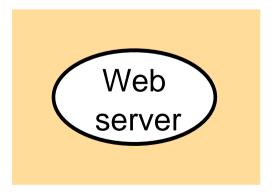
Web applications: Web applets

a) client request results in the downloading of applet code



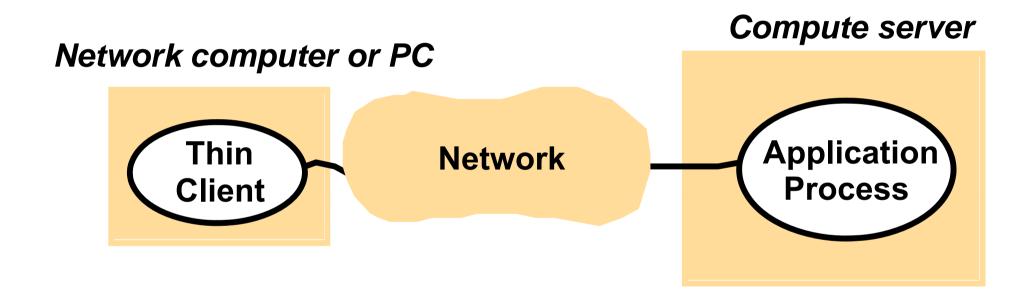
b) client interacts with the applet





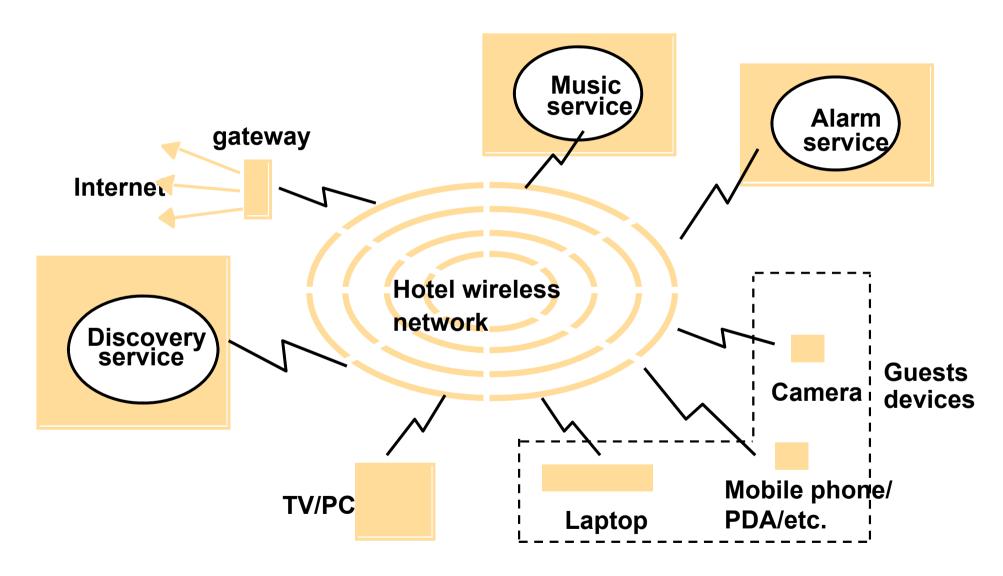


Client/Server: (d) Thin clients and compute servers





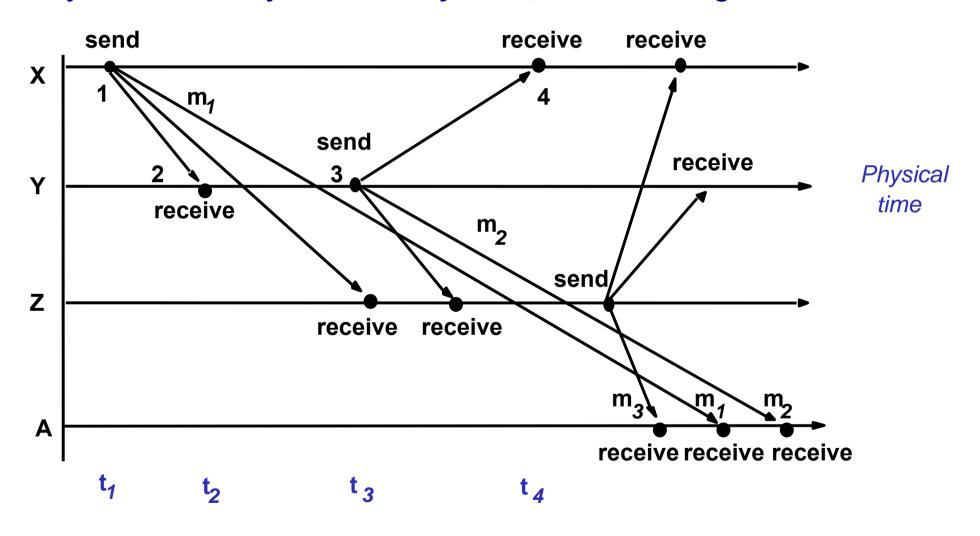
Application example: Spontaneous networking in a hotel





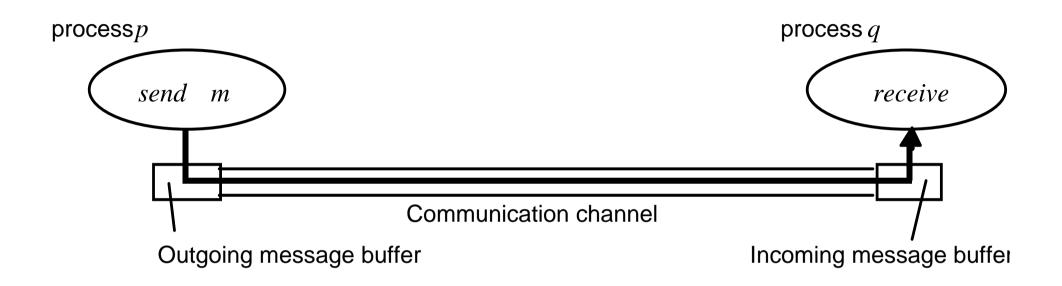
Timing models: Real-time ordering of events

Asynchronous & synchronous systems; event ordering





Failure models: Processes and channels



2 types of failures:

- timing failures: in synchronous distributed systems
- arbitrary failures: (Byzantine failure) "anything at any time"



Timing failures

Class of Failure	Affects	Description
Clock	Process	Process's local clock exceeds the bounds on its
		rate of drift from real time.
Performance	Process	Process exceeds the bounds on the interval
		between two steps.
Performance	Channel	A message's transmission takes longer than the
		stated bound.

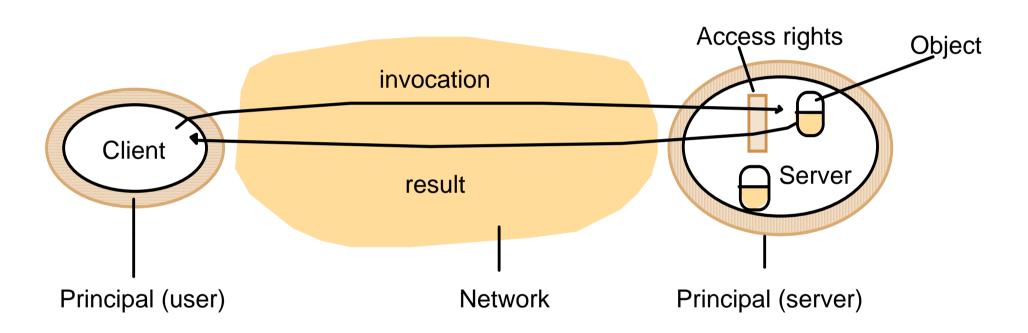


Omission and arbitrary failures

Class of failure	Affects	Description
Fail-stop	Process	Process halts and remains halted. Other processes may
		detect this state.
Crash	Process	Process halts and remains halted. Other processes may
		not be able to detect this state.
Omission	Channel	A message inserted in an outgoing message buffer never
		arrives at the other end's incoming message buffer.
Send-omission	Process	A process completes a send, but the message is not put
		in its outgoing message buffer.
Receive-omissionProcess		A message is put in a process's incoming message
		buffer, but that process does not receive it.
Arbitrary	Process or	Process/channel exhibits arbitrary behaviour: it may
(Byzantine)	channel	send/transmit arbitrary messages at arbitrary times,
		commit omissions; a process may stop or take an
		incorrect step.



Security models: Objects and principals

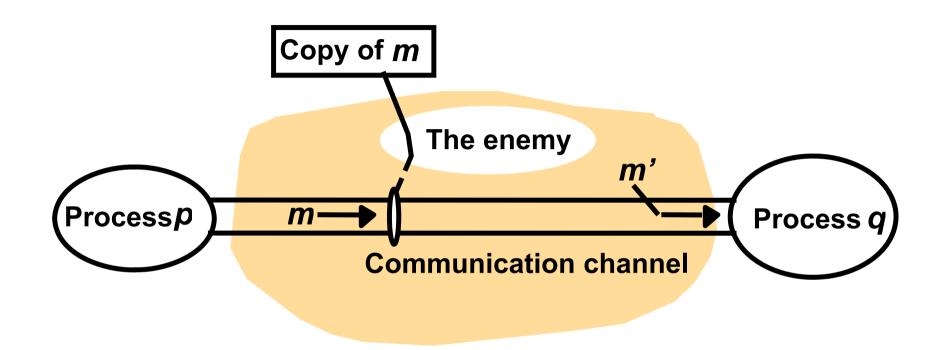


Goals:

- Protecting objects
- Securing processes and their interaction
- Clients → "Agents" (Principal-agent theory)



Security models: The enemy

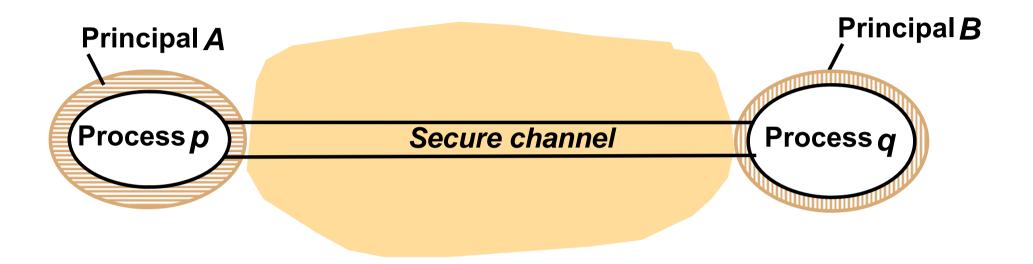


Basis security threads:

- loss of messages
- unauthorised read and/or access
- corruption
- false/stolen identities



Security models: Secure channels



Basis security <u>techniques</u>:

- cryptography & shared secrets
- authorisation & authentication
- secure channels
- → Sicherheitsaspekte: anderer Teil der Vorlesung GSS!

