

Model Checking of Fault-tolerant Systems

Lucas Baudin

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Model Checking Modulo Theories

Model Checking

- a state type is a list of variables: \mathbf{x}
- a state is a valuation for these variables
- a transition is a formula over the current state variables and the next state variables
(usually represented as a guard $H(\mathbf{x})$ and a (partial) assignment $V(\mathbf{x}')$)
- $(H_1(\mathbf{x}) \wedge V_1(\mathbf{x}')) \vee \dots \vee (H_n(\mathbf{x}) \wedge V_n(\mathbf{x}'))$

Modulo Theories

- the formulae can be in any theory
- example:
 - if the state type is a variable x of type int
 - transition: $x < 0 \wedge x' = x \vee x \geq 0 \wedge x' = x + 1$
 - describes a system which has a variable x which keeps increasing unless it is lower than 0.

Model Checking of a Fault-tolerant System

The Byzantine General Problem

- One general wants to give an order to $n - 1$ lieutenants.
- Some of them may be faulty (including the general).
- In the end, they have to decide on the same order (at least for all the non faulty lieutenants/general)

Pseudo-Code

% message[i][j] is the message sent by lieutenant j to i
message: array of (lieutenants * lieutenants) to message

source: the general id

% value[i] is the order that lieutenant i will follow
value: array of lieutenants to message

% good[i] is true if lieutenant i is non faulty
good: array of lieutenants to boolean

Pseudo-Code

```
propagate():  
  forall i, j do  
    if good[i]  
      message[j][i] <- message[i][source]  
  
agree():  
  forall i do  
    if good[i]  
      value[i] <- majority(message[i])  
  
propagate();  
propagate();  
agree();
```

The Byzantine General Problem

- $\exists v \forall i \text{ good}[i] \Rightarrow \text{value}[i] = v$
- How many lieutenants can be faulty? Does that work for a third of the lieutenants?
- $\#\{i \mid \text{good}[i]\} > 2N/3$

Sally

A model checker for infinite-state systems

- sri-csl.github.io/sally
- a symbolic model checker
- several engines: bmc, kind, ic3
- works with various smt solvers: mathsat, yices2, z3

Input language

- lisp-like language
- low level
- easy to parse and work with

Input language

- state type

```
(define-state-type my_state_type  
  ((x Real) (y Real))  
)
```

Input language

- state type

```
(define-state-type my_state_type  
  ((x Real) (y Real))  
)
```

- state formula

```
(define-states x_is_zero my_state_type  
  (= x 0)  
)
```

Input language

- transition: a first order formula over state variables and next state variables

```
(define-transition my_transition my_state_type
  (or
    (= next.x (+ state.x 1))
    next.x_is_zero
  )
)
```


Input language

- queries: check that a property always holds

```
(query my_system  
  ( $\geq$  x 0)  
)
```

A new old input language: Sal

Sal

- an older model checker, developed at SRI
- developed actively until 2006, minor versions until 2013
- finite state systems

Input language

- already used
- supports modules, composition

Input language

```
my_module: MODULE
BEGIN
  OUTPUT
    x: REAL,
    y: REAL
  INITIALIZATION
    x = 0;
  TRANSITION
    ...
END;
```

```
TRANSITION
  [x >= 0 -->
    x' = x + 1;
    y' IN { i: REAL | TRUE }
  []
  TRUE -->
    x' = 0;
    y' IN { i: REAL | TRUE }
  ]
```

Input language

- a lemma is translated to a Sally query
- multiple lemma
- syntax for temporal logic (not available in Sally)

```
my_context: CONTEXT =  
BEGIN  
  my_module: MODULE  
    ...  
  
  always_positive: LEMMA  
    my_module |- G(x >= 0);  
  
  wrong_lemma: LEMMA  
    my_module |- G(x > 0 -> x = 1)  
  
END
```

Parametrization

Arrays

Quantifiers

- for most examples, they can be avoided in transitions
- works only with z3
- example:

```
(forall (i Int) (select a i))
```

Counting in SMT

- $\phi(y) \wedge y = \# \{x | \psi(x)\}$
- $\psi(\cdot)$: first order formula of Presburger arithmetic, then with arrays too
- how is it solved?

State of the art

- Bradley, Manna, and Sipma (2006): a decision procedure for a fragment of arrays, with distinct theories for elements and indexes
- Alberti, Ghilardi, and Pagani (2016): a decision procedure for counting on arithmetic and arrays, via various rewriting and quantifier eliminations, mix index and elements theories, but quantify only over one element at a time
- Bjørner, Gleissenthall, and Rybalchenko (n.d.): a model checking oriented way to deal with arrays (one update at a time for every arrays)

Counting over Presburger arithmetic

- given a model, one can compute the value of $\# \{x | \psi(x)\}$
- given an ordering on the integer variables, one can compute the symbolic value of $\# \{x | \psi(x)\}$
- \Rightarrow symbolic computation of cardinalities, for the ordering of a given model

Counting over Presburger arithmetic, with an ordering

- the ordering is called an oracle: when asked whether $a > b$, it looks in the model the value of a and b and answers accordingly.
- when the cardinality is computed symbolically, it is equal to a formula which holds under some assumptions, and the oracle can say what they are

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- when the cardinality is computed symbolically, it is equal to a formula which holds under some assumptions, and the oracle can say what they are
- example: $y = \#\{x \mid 0 \leq x < z \wedge 0 \leq x < u\}$
- if the oracle says $z > u$, then y can be computed and $y = z$.
- under the assumption $z > u$, $y = z$

Counting over Presburger arithmetic, with an ordering

- compute a symbolic interval list in which the formula is satisfiable i.e. $\exists I \in V_x(\psi) \ x \in I \Leftrightarrow \psi(x)$
- if members of $V_x(\psi)$ are disjoint
$$\text{card}(V_x(\psi)) = \sum_{[v, v'] \in \#_x(\psi)} (v - v') = \#\{x | \psi(x)\}$$

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- $V_x(y < x) = \{[y, +\infty)\}$

Counting over Presburger arithmetic, with an ordering

- $V_x(\psi \wedge \phi) = V_x(\psi) \sqcap V_x(\phi)$
where $A \sqcap B$ is the set of every intersection of an interval of A
with an interval of B (of course some are empty and must be
deleted)

Counting over Presburger arithmetic, with an ordering

- $V_x(\psi \wedge \phi) = V_x(\psi) \sqcap V_x(\phi)$
where $A \sqcap B$ is the set of every intersection of an interval of A with an interval of B (of course some are empty and must be deleted)
- Intersection between two intervals: $[a, b)$ and $[c, d)$ can be computed: there is an oracle giving the ordering on the variables, so $\max(a, c)$ and $\min(b, d)$ are computable given the assumptions that the oracle makes.
Then the intersection is $[\max(a, c), \min(b, d))$ if $\max(a, c) \leq \min(b, d)$

Counting over Presburger arithmetic, with an ordering

- if $V_x(\psi) = \{[a_1, b_1), \dots, [a_n, b_n)\}$ (with $a_1 \leq b_1 \leq \dots \leq b_n$, $a_1 \neq -\infty$ and $b_n = +\infty$)
- $V_x(\neg\psi) = \{(-\infty, a_1), [b_1, a_2), \dots, [b_n, +\infty)\}$
- other cases are easy too, disjunction on $a_1 = -\infty$ and $b_n = +\infty$

Counting with multiplication

- to deal with constant multiplications, a modulo information can be added to every intervals (such as $([5, 10), = 1[3])$ are the integers x between 5 and 10 and such that $3|x - 1$)
- intersection, negation of these intervals can be done in an analog way

Future Work

- Counting over arrays
- IC3 with arrays and counting quantifiers

References

- Alberti, Francesco, Silvio Ghilardi, and Elena Pagani. 2016. "Counting Constraints in Flat Array Fragments." *CoRR* abs/1602.00458. <http://arxiv.org/abs/1602.00458>.
- Bjørner, Nikolaj, Klaus v Gleissenthall, and Andrey Rybalchenko. n.d. "Cardinalities and Universal Quantifiers for Verifying Parameterized Systems."
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