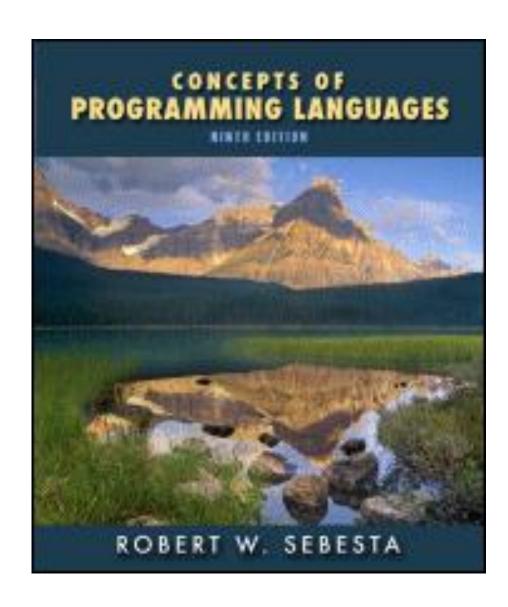
Chapter 7

Expressions and Assignment Statements



Ch07 – Expressions and Assignment Statements

- 7.1 Introduction*
- 7.2 Arithmetic Expressions
- 7.3 Overloaded Operators*
- 7.4 Type Conversions*
- 7.5 Relational and Boolean Expressions*
- 7.6 Short-Circuit Evaluation*
- 7.7 Assignment statements*
- 7.8 Mixed-mode Assignment*

- Side effects (7.2.2)
 - Assignment and I/O expressions yield values and side effects

	value	Side effect
x = 7	7	Store 7 in x's location
<pre>printf("Snoopy")</pre>	6	Snoopy is displayed

 In a boarder sense, assignment and I/O statements also yield side effects, e.g.

```
int f() { static int x = 0; x++; return x; }
```

The call f() yields a side effect, due to the assignment statement x++.

Such a side effect is called a functional side effect.

- Referential transparency and opaque
 - Referential transparency
 Contexts do not affect the meanings of expressions
 - Referential opaque
 Contexts do affect the meanings of expressions
 - Side effects make a language referential opaque.
 - For example, f() yields different values in different places cout << f();
 cout << f();
 cout << f();
 In a referential opaque language, exp==exp may not be true, e.g f()==f()

Con for side effects

- Hard to read
- Hard to prove recall { x=5 } y = x++ { x=6 }
- Prohibit optimization
 exp+exp cannot be optimized to 2*exp, if exp yields side
 effects.
- Make parallel processing difficult
 Suppose exp1 reads from x and exp2 writes to x, what will be the value of exp1+exp2, if exp1 and exp2 are evaluated in parallel?
- Affect portability
 due to side effects + unspecified operand evaluation order

- Argument (operand) evaluation order
 - exp1+exp2*exp3

Every language specifies operator evaluation order, e.g.

- * has a higher precedence than +.
- But, what is the evaluation order of the three operands exp1, exp2, and exp3?
- f(exp1,exp2)
 What is the evaluation order of the two arguments exp1 and exp2?
- Does a language have to specified argument (or operand) evaluation order?

- Side effects and argument evaluation order
 - Without side effects
 - Argument evaluation order is immaterial
 - Purely functional languages won't specify argument evaluation order, e.g. Haskell.
 - With side effects
 - Argument evaluation order is important
 - Some specify argument evaluation order, e.g. Java I-2-r
 - Most languages don't, e.g. C, C++, Scheme
 (N.B. In C/C++, ?:, &&, ||, and , have required order.)
- Pro for specifying argument evaluation order
 - Increase portability

- Pro for not specifying argument evaluation order
 - Increase efficiency leave a room for the compiler writer
 - Optimization
 With required order, e+e cannot be optimized to 2*e
 - Parallel processing
 With required order, exp1+exp2 cannot be evaluated in parallel
 - Intelligent compilation
 With required order, 5/n must be evaluated in n = 0; ... n*(5/n) ...
 But, a compiler may choose not to evaluate 5/n at all, for the reason that 0×anything=0

Example
 Given
 int x=3;
 int f(int y) { x++; return x+y; }
 What might be the value of the expression x*f(x)+f(x)?

Operand evaluation order	x*f(x)+f(x)
1 2 3	3*7+9=30
1 3 2	3*9+7=34
2 1 3	4*7+9=37
2 3 1	4*9+7=43
3 1 2	5*7+9=44
3 2 1	5*9+7=52

Example (cont'd)

Instead of fetching the value of x from memory thrice, some compilers generate code to fetch the value of x once and store it in a register. That is, the compiled code reads as

```
reg = x;
reg*f(reg)+f(reg)
```

Since the value 3 stored in the register doesn't change, the compiled code has the same effect as

$$3*f(3)+f(3)$$

Finally, depending on the order of evaluating the two calls, we have

Example (cont'd)

Operand evaluation order	3*f(3)+f(3)
1 2	3*7+8=29
2 1	3*8+7=31

Lesson

Never write code that depends on argument evaluation order.