

## PL Midterm

104 points in total

- 1 So far, we have already mentioned two design considerations that make Java suitable for use in the Internet. What are the two design considerations? (4%)
- 2 The statement "C++ is a compiled language." is problematic. What is the correct wording? (4%)
- 3 What are the pros and cons of dynamic scoping? (4%)
- 4 What are the pros and cons of dynamic type binding? (4%)

- 5 Consider the Javascript function below.

```
function f(n)
{
    var r="1"
    for (var i="2";i<=n;i++) { r*=i; alert(typeof i); }
    return r
}
```

- a) What are the messages yielded by **alert** during the evaluation of **f(3)**? (4%)
- b) What is the value of **f(10)**? of **f("10")**? (4%)

N.B.  $9! = 362880$

- 6 The following three files constitute a C++ program.

### File 1.cpp

```
#include <iostream>
using namespace std;
extern int x;
int g();
static int f() { return x; }
int main()
{
    cout << f() << g();
}
```

### File 2.cpp

```
int x=1;
int f()
{
    static int x=2; return x;
}
```

### File 3.cpp

```
struct A { static const int x=3; }
int f();
int g() { return A::x+f(); }
```

C++ prefers replacing one of the 3 occurrences of **static** by other language construct. Which one and how? (4%)

- 7 The following Fortran 95 function contains two syntax errors.

Figure out the errors. (Hint: Indicate which line contains what error. Use the line numbers.)

Also, figure out the phrase (i.e. lexical, syntax, or semantic analyzer) of the compiler that detects each error. (4%)

```

function fact(n)                ! 1
integer :: n,fact              ! 2
if (n==0) then                  ! 3
    fact = 1                     ! 4
else                           ! 5
    fact = n*fact(n-1)          ! 6
end                             ! 7
end                             ! 8

```

- 8 Show the output of the following Fortran 95 program. (4%)

```

program test
integer, dimension(2,2,2) :: a
read *, a ! 1,2
print *, a(1,::)
print *, a(2,::)
end program test

```

- 9 Consider the following Perl program with two blanks  $\alpha$  and  $\beta$ .

```

$x=0;
sub sub2
{
    sub sub3 {  $\alpha$  $x=2; sub2(); }    # local function
    print $x;                        # free variable
    $c++; if ($c<2) { sub3; }
}
sub sub1 {  $\beta$  $x=1; sub2; }
$c=0;
sub1;

```

Show the output of the program if (8%)

- a)  $\alpha = \text{my}$ ,  $\beta = \text{my}$
- b)  $\alpha = \text{my}$ ,  $\beta = \text{local}$
- c)  $\alpha = \text{local}$ ,  $\beta = \text{my}$
- d)  $\alpha = \text{local}$ ,  $\beta = \text{local}$

- 10 a) The following Perl subroutine is meant to multiply the elements of the array passed to it. Fill in the blank to complete the subroutine. (2%)

```
sub f
{
    my $r = 1;
    for (____) { $r *= $_; }
    return $r;
}
@a=(1..5);
print f @a; # output 120
```

- b) Redo a), but this time assumes that the **for** loop is written incompletely as (2%)

```
for (my $i=0;____;$i++) { $r *= $_[$i]; }
```

- 11 a) The following Scheme function solves the Towers of Hanoi problem. It moves  $n$  disks from peg  $a$  to peg  $b$  with the help of the auxiliary peg  $c$ . What is the value of the call `(hanoi 2 'a 'b 'c)?` (4%)

```
(define hanoi
  (lambda (n a b c)
    (if (= n 0)
        '()
        (append (hanoi (- n 1) a c b)
                  (cons `(. ,a . ,b) (hanoi (- n 1) c b a))))))
```

- b) Let's turn the function of part a) into an equivalent function in accumulator-passing style.

```
(define hanoiAPS
  (lambda (n a b c acc)
    (if (= n 0) acc (hanoiAPS (- n 1) a c b _____))))
```

Fill in the blank to complete the function.

Also, what should be the initial value of the parameter **acc**? (4%)

- c) Let's turn the function of part a) into an equivalent function in continuation-passing style.

```
(define hanoiCPS
  (lambda (n a b c con)
    (if (= n 0) (con '()) (hanoiCPS (- n 1) a c b _____))))
```

Fill in the blank to complete the function.

Also, what should be the initial value of the parameter **con**? (4%)

12 Given

```
(define x '(a (b c) d))
(define y (cons (list x) x))
```

- a) Draw a diagram showing the internal list structures bound to **x** and **y**. (4%)  
 b) What is the value of each expression below? (4%)

```
1 y
2 (caddr y)
```

13 a) What are the **best-case** space and time complexities of function **isort** below? (4%)

```
(define isort
  (lambda (xs)
    (if (null? xs)
        '()
        (let insert ((x (car xs)) (ys (isort (cdr xs))))
          (cond ((null? ys) (list x))
                ((<= x (car ys)) (cons x ys))
                (else (cons (car ys) (insert x (cdr ys))))))))))
```

- b) Redo a) for imperative-style insertion sort. (4%)

14 a) Infer the type of the recursive function **fix** defined by (4%)

$\text{fix} = \lambda f.f(\text{fix } f)$

N.B. Using ML's notation, the function **fix** is defined as

```
val fix = fn f => f (fix f);
```

- b) Explain why the **Y** combinator is untypable in ML. (4%)

N.B. Recall that  $Y = \lambda f.(\lambda x.f(xx))(\lambda x.f(xx))$

15 In this problem, we shall use the abbreviated  $\lambda$ -expression, say,  $\lambda abcde.exp$  to denote the curried function  $\lambda a \lambda b \lambda c \lambda d \lambda e.exp$ .

Now, let  $\mathcal{E} = \underbrace{\lambda abcdefghijklmnopqrstuvwxyz}_{26 \text{ English letters (Letter r at the end)}} \underbrace{r(\text{thisisafixedpointcombinator})}_{27 \text{ letters}}$

- a) Prove that  $\$ = \underbrace{\mathcal{E} \mathcal{E} \mathcal{E} \mathcal{E} \mathcal{E} \mathcal{E} \mathcal{E} \mathcal{E} \mathcal{E} \mathcal{E} \mathcal{E} \mathcal{E} \mathcal{E} \mathcal{E} \mathcal{E} \mathcal{E} \mathcal{E} \mathcal{E} \mathcal{E} \mathcal{E} \mathcal{E} \mathcal{E} \mathcal{E} \mathcal{E} \mathcal{E} \mathcal{E}}_{26 \mathcal{E}'s}$  is a fixed-point combinator with lazy

evaluation. (4%)

- b) Rephrase  $\$$  as a fixed-point combinator with eager evaluation. (Proof isn't needed.) (4%)

- c) Use  $\$$  to obtain the solution to **fix** of the equation

$\text{fix} = \lambda f.f(\text{fix } f)$

Note: It suffices to use  $\lambda$ -notation. No Scheme code is needed. (4%)

- 16 a) Compile the following  $\lambda$ -expression into code consisting of S, K, and I (4%)

$(\lambda x.x\ 2\ 3) +$

- b) Reduce the compiled code of part a) with lazy evaluation (4%)

**Hint**

The definitions of S, K, and I as well as the compilation algorithm are given below.

$S = \lambda x.\lambda y.\lambda z.x\ z\ (y\ z)$

$K = \lambda x.\lambda y.x$

$I = \lambda x.x = S\ K\ K$

compile  $x \Rightarrow x$ , if  $x$  is a variable, built-in constant, or built-in function

compile  $(e1\ e2) \Rightarrow \text{compile } e1\ (\text{compile } e2)$

compile  $\lambda x.e \Rightarrow \text{abstract } x\ (\text{compile } e)$

abstract  $x\ x \Rightarrow I$

abstract  $x\ y \Rightarrow K\ y$ , if  $y$  is a variable ( $\neq x$ ), built-in constant or built-in function

abstract  $x\ (e1\ e2) \Rightarrow S\ (\text{abstract } x\ e1)\ (\text{abstract } x\ e2)$