

OOP Midterm solution

- 1
 - a) function-to-pointer conversion, pointer conversion, qualification conversion
`void(const void*)`
 \rightarrow `void(*) (const void*)` // function-to-pointer conversion
 \rightarrow `void*` // pointer conversion
 \rightarrow `const void*` // qualification conversion
 - b) `void (*(&g())[1])()` // the name `g` may be omitted
 - c) The call `operator new[] (3*sizeof(int))` allocates more storage than the call `operator new(3*sizeof(int))`. The extra storage is used to store the number of elements in the array by a `new` expression. However, the extra storage is useless in 2), as there is no `new` expression in 2).
 - d) The function returns a reference to a local temporary location that is no longer available after the function returns.
 More precisely, the compiled code looks like

```
double const& f(int x) { const double tmp=x; return tmp; }
```
- 2
 - a) Template 2, since it is a specialist in `T*`.
 - b) Ambiguous
 Template 1 uses call-by-value, and template 2 uses call-by-reference.
 Both parameter-passing methods are equally well for the call `p(x)`.
 - c) All of 1), 2) and 3)
 A class template explicit specialization may have a different set of class members from the generic class template.
 - d) Only 1)
 A function template explicit specialization for `T = int` is obtained by substituting `int` for every occurrence of `T` in the signature of the function template.
 For example, by substituting `int` for `T` in the underlined signature

```
template<typename T> T f(T& x) { return x; }
```


 we obtain

```
template<> int f<int>(int& x) { return x; }
```


 In conclusion, a function template explicit specialization may not have a different signature format, but may have a different body, i.e. a different implementation (algorithm) of the function. For example,

```
template<> char f<char>(char& x) // same signature format
{
    // do whatever you want
}
```

2 e) Only 3)

```
template<typename T> void p(const T&) {}
```

The declaration **const T** declares the type **T** to be a constant type.

For **T = const int***, **T** is a pointer type. Thus, **const T** means the pointer type is a constant type.

Thus, the explicit specialization for **T = const int*** is

```
template<> void p(const int*const&) {}
```



3 a) The viable functions are

```
float A::f(float x);
```

```
double A::f(double x); // 2
```

```
double f<double>(double x); // 3
```

Since the ordinary function 2 is better than the instantiated function 3, the best viable function is

```
double A::f(double x);
```

b) Using the same argument as a), the compiler first resolves the call **f(7)** to a call to the instantiated function:

```
int f(int x)
```

```
{
```

```
    return numeric_limits<T>::is_integer? x: f(x); // 1
```

```
}
```

In a similar manner, the compiler then resolves the call **f(x)** in line one to be a recursive call. (Note that this recursive call will never be executed.)

Were **f(x)** replaced by **A::f(x)**, an ambiguity would occur. This is because the instantiated function is no longer a candidate, and both

```
float A::f(float x);
```

```
double A::f(double x);
```

are viable, but neither is better.

4 a) (1) **!comp(value,*it)&&!comp(*it,value)**

(2) **comp(value,*it)**

b) Given

```
binary_search(c,c+6,"Yankees",less<const char*>())
```

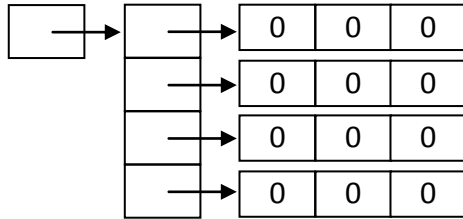
```
bool binary_search(T*,T*,const T&,Compare)
```

the compiler first determines that **T = char*** from **c** and **c+6**, but then finds out **T = const char*** from **"Yankees"**.

This ambiguity can be easily resolved by type conversion

```
const_cast<char*>("Yankees") or (char*) "Yankees"
```

5 a)



b) It still works, but the benefit of memoization disappears. To see why, observe that each time the function **c** is called, it receives a zero-initialized vector. This means two things: firstly, the function works as if there is no **cache**; and secondly, since it is as if there is no **cache**, the problem of recomputation remains.

c) See lecture note

6 (1) **m!=n && n!=0**(2) **s.push(m-1); s.push(n); m--; n--;**

or

s.push(m-1); s.push(n-1); m--;(3) **n=s.top(); s.pop(); m=s.top(); s.pop();**7 a) (1) **new int*[80]**(2) **while (!empty()) pop(); delete [] stk;**(3) **stk[++_top]=new int(n);**

or

++_top; stk[_top]=new int(n);(4) **delete stk[_top--];**

or

delete stk[_top]; --_top;(5) **return *stk[_top];**

or

return stk[_top][0];b) **bool empty(const stack* this)**

{

return this->_top==-1;

}

8 a) `template<typename T>`
`T max(T* a,int n)`
`{`
`return accumulate(a,a+n,numeric_limits<T>::min(),`
`std::max<T>);`
`}`

Note that the qualified name `std::max<T>` is necessary.

b) `char* h(char* r,char* s)`
`{`
`if (strlen(s)%2==0) itoa(atoi(r)+1,r,10);`
`return r;`
`}`

9 a) 77 2 3

b) `void* operator new(size_t sz,int* pool)`
`{`
`static int offset=0;`
`int old_offset=offset;`
`offset+=sz;`
`return reinterpret_cast<char*>(pool)+old_offset;`
`}`

or

`void* operator new(size_t sz,int* pool)`
`{`
`static char* available=reinterpret_cast<char*>(pool);`
`char* old_available=available;`
`available+=sz;`
`return old_available;`
`}`

or,

`void* operator new(size_t,int* pool)`
`{`
`static int next=0;`
`return pool+next++;`
`}`

This last version isn't recommended, for it doesn't make use of the 1st parameter.