# Online Methods

#### **Last Time**

- Policy Iteration
- Value Iteration
- Does Value Iteration always converge?
- Is the optimal value function unique?

## **Guiding Questions**

- What are the differences between *online* and *offline* solutions?
- Are there solution techniques that require computation time *independent* of the state space size?

# Why Do We Need Something Else?

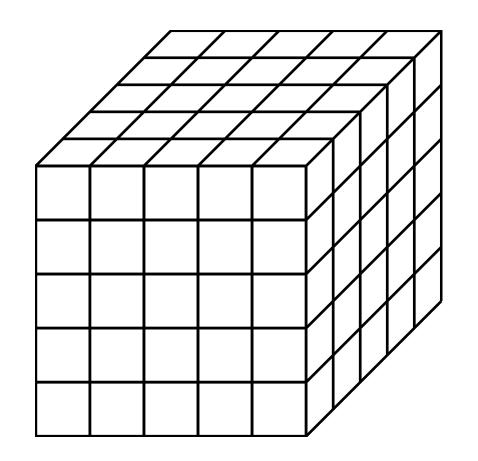
- Problems Policy and Value Iteration may struggle with?
  - Path planning across the country, or interplanetary
  - More realistic car dynamics (continuous states)
- Why are these problems hard?
  - State Space is massive (or infinite)

#### **Curse of Dimensionality**

1 dimension, 5 segments  $|\mathcal{S}|=5$ 

2 dimensions, 5 segments 
$$|\mathcal{S}|=25$$

3 dimensions, 5 segments  $|\mathcal{S}|=125$ 



n dimensions, k segments  $o |\mathcal{S}| = k^n$ 

#### Offline vs Online Solutions

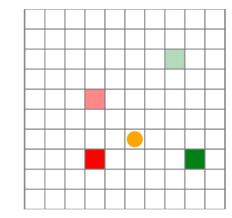
#### **Offline**

- Before Execution: find  $V^*/Q^*$
- During Execution:  $\pi^*(s) = \operatorname{argmax} Q^*(s, a)$

<b>→</b>	<b>→</b>	<b>→</b>	<b>→</b>	<b>→</b>	1	1	<b>→</b>	1	Ţ
<b>→</b>	<b>→</b>	<b>→</b>	<b>→</b>	-	1	1	<b>→</b>	1	ţ
<b>→</b>	<b>→</b>	<b>→</b>	<b>→</b>	-	1	1	t	1	ţ
<b>→</b>	t	t	<b>→</b>	-	<b>→</b>	1	1	1	Ţ
1	1	1	t	-	<b>→</b>	1	1	1	ı
Ţ	<b>→</b>	<b>→</b>	<b>→</b>	<b>→</b>	<b>→</b>	<b>→</b>	1	1	1
1	1	<b>→</b>	<b>→</b>	<b>→</b>	<b>→</b>	<b>→</b>	<b>→</b>	1	1
1	1	1	t	-	<b>→</b>	<b>→</b>	<b>→</b>	t	-
1	1	1	<b>→</b>	-	<b>→</b>	<b>→</b>	<b>→</b>	t	t
-	-	-	-	-	-	-	t	t	t

#### <u>Online</u>

- Before Execution: <nothing>
- During Execution: Consider actions and their consequences (everything)



- Why?
- Online methods are insensitive to the size of S!

#### One Step Lookahead

```
randstep(\mathcal{P}::MDP, s, a) = \mathcal{P}.TR(s, a)
function rollout (P, s, \pi, d)
     ret = 0.0
     for t in 1:d
          a = \pi(s)
          s, r = randstep(P, s, a)
          ret += \mathcal{P}.v^{(t-1)} * r
     end
     return ret
end
function (π::RolloutLookahead)(s)
     U(s) = rollout(\pi.P, s, \pi.\pi, \pi.d)
     return greedy (\pi.P, U, s).a
end
function greedy (\mathcal{P}::MDP, U, s)
     u, a = findmax(a\rightarrowlookahead(\mathcal{P}, U, s, a), \mathcal{P}.\mathcal{A})
     return (a=a, u=u)
end
```

function lookahead( $\mathcal{P}$ ::MDP, U, s, a)

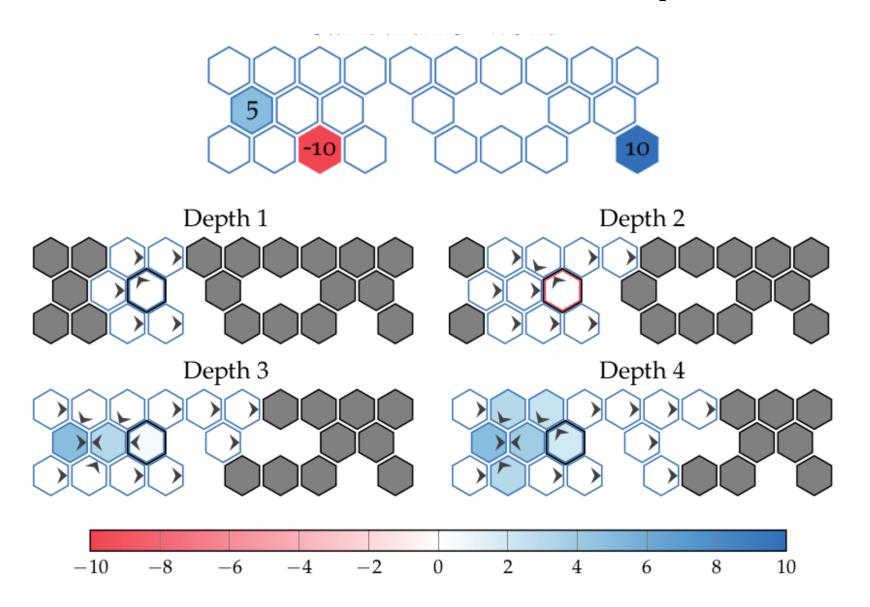
S, T, R,  $\gamma = \mathcal{P}.S$ ,  $\mathcal{P}.T$ ,  $\mathcal{P}.R$ ,  $\mathcal{P}.\gamma$ 

return R(s,a) +  $\gamma * sum(T(s,a,s')*U(s')$  for s' in S)

#### **Forward Search**

```
function forward_search(₱, s, d, U)
      if d \leq 0
           return (a=nothing, u=U(s))
      end
      best = (a=nothing, u=-Inf)
      U'(s) = forward\_search(P, s, d-1, U).u
      for a in \mathcal{P}. A
           u = lookahead(P, U', s, a)
           if u > best.u
                 best = (a=a, u=u)
           end
      end
      return best
 end
function lookahead(\mathcal{P}::MDP, U, s, a)
    S, T, R, \gamma = \mathcal{P}.S, \mathcal{P}.T, \mathcal{P}.R, \mathcal{P}.\gamma
    return R(s,a) + \gamma *sum(T(s,a,s')*U(s') for s' in S)
                   O\left((|S| \times |A|)^d\right)
```

### Forward Search depth



## **Sparse Sampling**

```
function sparse_sampling (P, s, d, m, U)
    if d \leq 0
         return (a=nothing, u=U(s))
    end
    best = (a=nothing, u=-Inf)
    for a in \mathcal{P}.\mathcal{A}
         u = 0.0
         for i in 1:m
             s', r = randstep(P, s, a)
             a', u' = sparse\_sampling(P, s', d-1, m, U)
             u += (r + \mathcal{P}.\gamma*u') / m
         end
         if u > best.u
             best = (a=a, u=u)
         end
    end
    return best
end
```

$$O\left((m|A|)^d\right)$$

$$|V^{ ext{SS}}(s) - V^*(s)| \leq \epsilon$$

m,  $\epsilon$ , and d related, but independent of |S|

#### Break

Draw the trees produced by the following algorithms for a problem with 2 actions and 3 states:

- 1. One-step lookahead with rollout
- 2. Forward search (d=2)
- 3. Sparse sampling (d=2, m=2)

#### Monte Carlo Tree Search (MCTS/UCT)

#### Keep track of:

Q(s,a): Value estimate of that state and action combo N(s,a): Number of times we visit a state and action combo

$$Q(s,a) + c\sqrt{rac{\log N(s)}{N(s,a)}} \quad Q(s,a) + crac{N(s)^{eta}}{\sqrt{N(s,a)}}$$

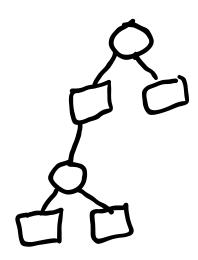
low N(s,a)/N(s) = high bonus start with  $c=2(ar{V}-\underline{V})$ , eta=1/4

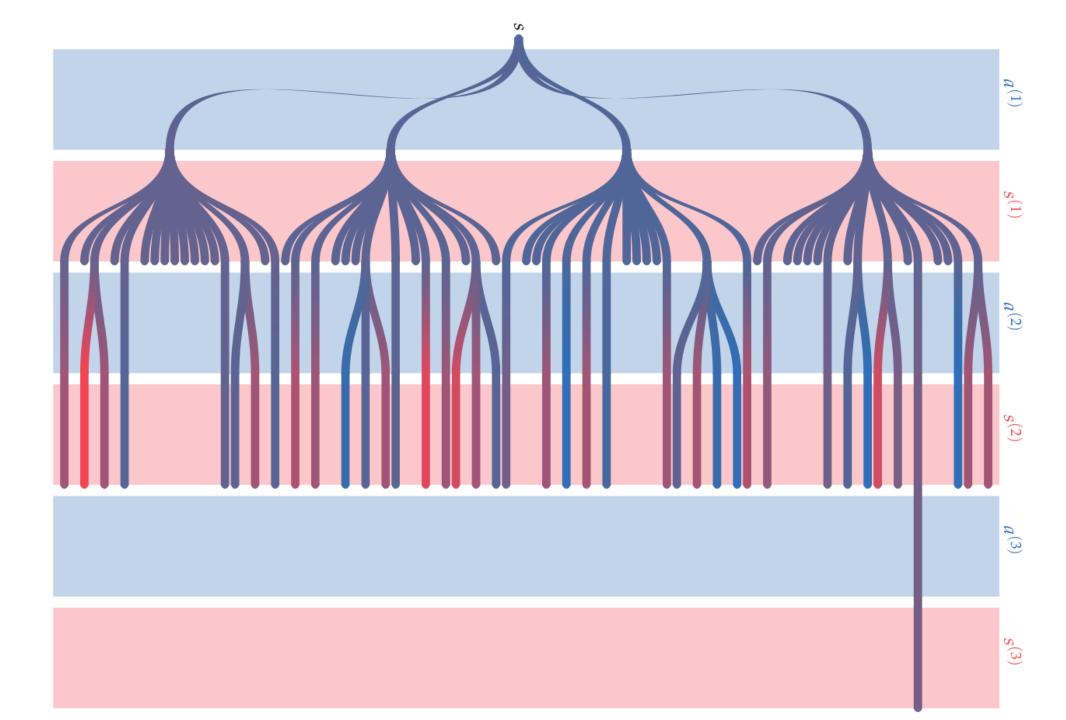
Full story can be found in https://arxiv.org/pdf/1902.05213.pdf

## Monte Carlo Tree Search (MCTS/UCT)

```
for k in 1:\pi.m
             simulate!(\pi, s)
      end
      return argmax(a \rightarrow \pi.Q[(s,a)], \pi.P.A)
end
function simulate! (\pi::MonteCarloTreeSearch, s, d=\pi.d)
     if d \le 0
            return \pi.U(s)
     end
     \mathcal{P}, N, Q, c = \pi \cdot \mathcal{P}, \pi \cdot N, \pi \cdot Q, \pi \cdot c
     \mathcal{A}, TR, \gamma = \mathcal{P} \cdot \mathcal{A}, \mathcal{P} \cdot \mathsf{TR}, \mathcal{P} \cdot \gamma
     if !haskey(N, (s, first(\Re)))
           for a in \mathcal{A}
                 N[(s,a)] = 0
                 Q[(s,a)] = 0.0
            end
           return \pi.U(s)
      a = explore(\pi, s)
      s', r = TR(s,a)
     q = r + \gamma * simulate!(\pi, s', d-1)
     N[(s,a)] += 1
     Q[(s,a)] += (q-Q[(s,a)])/N[(s,a)]
     return q
```

function (π::MonteCarloTreeSearch)(s)





## **Guiding Questions**

- What are the differences between online and offline solutions?
- Are there solution techniques that are *independent* of the state space size?

## Forward Search Sparse Sampling

(FSSS)

Paper: https://cdn.aaai.org/ojs/7689/7689-13-11219-1-2-20201228.pdf

Sparse Sampling, but only look at potentially valuable states

#### Things it keeps track of:

Q(s,a): Estimate of the value for the state action pair

U(s): Upper bound for value of state s

L(s): Lower bound for value of state s

U(s,a): Upper bound for value of state-

action

L(s,a): Lower bound for value of stateaction

## Forward Search Sparse Sampling

```
Algorithm 3 FSSS(s, d)
  if d = 1 (leaf) then
      L^d(s,a) = U^d(s,a) = R(s,a), \forall a
      L^d(s) = U^d(s) = \max_a R(s, a)
  else if n_{sd} = 0 then
      for each a \in A do
         L^d(s,a) = V_{\min}
         U^d(s,a) = V_{\text{max}}
         for C times do
            s' \sim T(s, a, \cdot)
            L^{d-1}(s') = V_{\min}
            U^{d-1}(s') = V_{\text{max}}
            K^d(s,a) = K^d(s,a) \cup \{s'\}
  a^* = \operatorname{argmax}_a U^d(s, a)
s^* = \operatorname{max}_{s' \in K^d(s, a^*)} (U^{d-1}(s') - L^{d-1}(s'))
  FSSS(s^*, d-1)
  n_{sd} = n_{sd} + 1
  L^{d}(s, a^{*}) = R(s, a^{*}) + \gamma \sum_{s' \in K^{d}(s, a^{*})} L^{d-1}(s') / C
  U^{d}(s, a^{*}) = R(s, a^{*}) + \gamma \sum_{s' \in K^{d}(s, a^{*})} U^{d-1}(s') / C
  L^d(s) = \max_a L^d(s, a)
  U^d(s) = \max_a U^d(s, a)
```

If  $L(s, a*) \ge \max_{a \ne a^*} U(s, a)$  for best action ( $a^* = \arg \max_a U(s, a)$ ): then, the node is closed because the best action is found.