

Value Iteration Convergence

Review

Review

- How do we reason about the **future consequences** of actions in an MDP?

Review

- How do we reason about the **future consequences** of actions in an MDP?
- What are the basic **algorithms for solving MDPs**?

Guiding Questions

Guiding Questions

- Does value iteration always converge?
- Is the value function unique?
- Can there be multiple optimal policies?
- Is there always a deterministic optimal policy?

Value Iteration: The Bellman Operator

Value Iteration: The Bellman Operator

Algorithm: Value Iteration

while $\|V - V'\|_\infty > \epsilon$

$V \leftarrow V'$

$V' \leftarrow B[V]$

return V'

$$|S| = 3 \quad V = \begin{bmatrix} 0.4 \\ 0.9 \\ 1.2 \end{bmatrix}$$

Value Iteration: The Bellman Operator

Algorithm: Value Iteration

while $\|V - V'\|_\infty > \epsilon$

$V \leftarrow V'$

$V' \leftarrow B[V]$

return V'

$$B[V](s) = \max_{a \in A} (R(s, a) + \gamma E[V(s')])$$

Value Iteration Convergence

Value Iteration Convergence

Theorem 1: Let $\{V_1, \dots, V_\infty\}$ be a sequence of value functions for a discrete MDP generated by the recurrence $V_{k+1} = B[V_k]$. If $\gamma < 1$, then $\lim_{k \rightarrow \infty} V_k = V^*$.

Metrics

Metrics

Definition: Let M be a set. A *metric* on M is a function $d : M \times M \rightarrow [0, \infty)$ which satisfies the following three conditions for all $x, y, z \in M$:

Metrics

Definition: Let M be a set. A *metric* on M is a function $d : M \times M \rightarrow [0, \infty)$ which satisfies the following three conditions for all $x, y, z \in M$:

1. $d(x, y) = 0$ if and only if $x = y$

Metrics

Definition: Let M be a set. A *metric* on M is a function $d : M \times M \rightarrow [0, \infty)$ which satisfies the following three conditions for all $x, y, z \in M$:

1. $d(x, y) = 0$ if and only if $x = y$
2. $d(x, y) = d(y, x)$

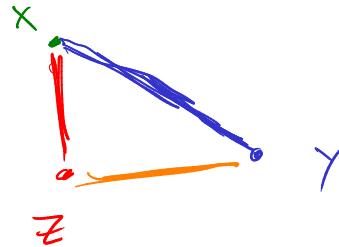
Metrics

Definition: Let M be a set. A *metric* on M is a function $d : M \times M \rightarrow [0, \infty)$ which satisfies the following three conditions for all $x, y, z \in M$:

1. $d(x, y) = 0$ if and only if $x = y$

2. $d(x, y) = d(y, x)$

3. $d(\underline{x}, \underline{y}) \leq d(\underline{x}, \underline{z}) + d(\underline{z}, \underline{y})$



Contraction Mappings

Contraction Mappings

Definition: A *contraction mapping* on metric space (M, d) is a function $f : M \rightarrow M$ satisfying

$$d(f(x), f(y)) \leq \alpha d(x, y)$$

for some α , $0 \leq \alpha \leq 1$ and all x and y in M .

Contraction Mappings

Definition: A *contraction mapping* on metric space (M, d) is a function $f : M \rightarrow M$ satisfying

$$d(f(x), f(y)) \leq \alpha d(x, y)$$

for some α , $0 \leq \alpha \leq 1$ and all x and y in M .

Definition: x^* is said to be a *fixed point* of f if $f(x^*) = x^*$.

Contraction Mappings

Definition: A *contraction mapping* on metric space (M, d) is a function $f : M \rightarrow M$ satisfying

$$d(f(x), f(y)) \leq \alpha d(x, y)$$

for some α , $0 \leq \alpha \leq 1$ and all x and y in M .

Definition: x^* is said to be a *fixed point* of f if $f(x^*) = x^*$.

$$M = \mathbb{R}^2 \quad f(x) = \begin{bmatrix} \frac{x_2}{2} + 1 \\ \frac{x_1}{2} + \frac{1}{2} \end{bmatrix}$$

Script: contraction_mapping.jl

Banach's Theorem

.

Banach's Theorem

Theorem (Banach): If f is a contraction mapping on metric space (M, d) , then

1. f has a single, unique fixed point x^* .
2. If $\{x_k\}$ is a sequence defined by $x_{k+1} = f(x_k)$, then $\lim_{k \rightarrow \infty} x_k = x^*$.

Max Norm

Max Norm

Lemma 1: $(\mathbb{R}^{|S|}, \|\cdot\|_\infty)$ is a metric space.

Max Norm

Lemma 1: $(\mathbb{R}^{|S|}, \|\cdot\|_\infty)$ is a metric space.

Definition: Let M be a set. A *metric* on M is a function $d : M \times M \rightarrow [0, \infty)$ which satisfies the following three conditions for all $x, y, z \in M$:

1. $d(x, y) = 0$ if and only if $x = y$
2. $d(x, y) = d(y, x)$
3. $d(x, y) \leq d(x, z) + d(z, y)$

Max Norm

Lemma 1: $(\mathbb{R}^{|S|}, \|\cdot\|_\infty)$ is a metric space.

Definition: Let M be a set. A *metric* on M is a function $d : M \times M \rightarrow [0, \infty)$ which satisfies the following three conditions for all $x, y, z \in M$:

Proof:

1. $d(x, y) = 0$ if and only if $x = y$
2. $d(x, y) = d(y, x)$
3. $d(x, y) \leq d(x, z) + d(z, y)$

Max Norm

Lemma 1: $(\mathbb{R}^{|S|}, \|\cdot\|_\infty)$ is a metric space.

Definition: Let M be a set. A *metric* on M is a function $d : M \times M \rightarrow [0, \infty)$ which satisfies the following three conditions for all $x, y, z \in M$:

1. $d(x, y) = 0$ if and only if $x = y$
2. $d(x, y) = d(y, x)$
3. $d(x, y) \leq d(x, z) + d(z, y)$

Proof:

Note: $\|x - y\|_\infty = \max_i |x_i - y_i|$

Max Norm

Lemma 1: $(\mathbb{R}^{|S|}, \|\cdot\|_\infty)$ is a metric space.

Definition: Let M be a set. A *metric* on M is a function $d : M \times M \rightarrow [0, \infty)$ which satisfies the following three conditions for all $x, y, z \in M$:

1. $d(x, y) = 0$ if and only if $x = y$
2. $d(x, y) = d(y, x)$
3. $d(x, y) \leq d(x, z) + d(z, y)$

Proof: Note: $\|x - y\|_\infty = \max_i |x_i - y_i|$

$$1. \max |x - y| = 0 \text{ iff } x_i = y_i \quad \forall i$$

Max Norm

Lemma 1: $(\mathbb{R}^{|S|}, \|\cdot\|_\infty)$ is a metric space.

Definition: Let M be a set. A *metric* on M is a function $d : M \times M \rightarrow [0, \infty)$ which satisfies the following three conditions for all $x, y, z \in M$:

1. $d(x, y) = 0$ if and only if $x = y$
2. $d(x, y) = d(y, x)$
3. $d(x, y) \leq d(x, z) + d(z, y)$

Proof: Note: $\|x - y\|_\infty = \max_i |x_i - y_i|$

$$1. \max |x - y| = 0 \text{ iff } x_i = y_i \quad \forall i$$

$$2. |x - y| = |-(x - y)| = |y - x|$$
$$\therefore \max |x - y| = \max |y - x|$$

Max Norm

Lemma 1: $(\mathbb{R}^{|S|}, \|\cdot\|_\infty)$ is a metric space.

Definition: Let M be a set. A *metric* on M is a function $d : M \times M \rightarrow [0, \infty)$ which satisfies the following three conditions for all $x, y, z \in M$:

1. $d(x, y) = 0$ if and only if $x = y$
2. $d(x, y) = d(y, x)$
3. $d(x, y) \leq d(x, z) + d(z, y)$

Proof: Note: $\|x - y\|_\infty = \max_i |x_i - y_i|$

$$1. \max |x - y| = 0 \text{ iff } x_i = y_i \quad \forall i$$

$$2. |x - y| = |-(x - y)| = |y - x|$$
$$\therefore \max |x - y| = \max |y - x|$$

$$3. \max |x - z| = \max |x - y + y - z|$$

Max Norm

Lemma 1: $(\mathbb{R}^{|S|}, \|\cdot\|_\infty)$ is a metric space.

Definition: Let M be a set. A *metric* on M is a function $d : M \times M \rightarrow [0, \infty)$ which satisfies the following three conditions for all $x, y, z \in M$:

1. $d(x, y) = 0$ if and only if $x = y$
2. $d(x, y) = d(y, x)$
3. $d(x, y) \leq d(x, z) + d(z, y)$

Proof: Note: $\|x - y\|_\infty = \max_i |x_i - y_i|$

$$1. \max |x - y| = 0 \text{ iff } x_i = y_i \quad \forall i$$

$$2. |x - y| = |-(x - y)| = |y - x|$$
$$\therefore \max |x - y| = \max |y - x|$$

$$3. \max |x - z| = \max |x - y + y - z|$$
$$\leq \max(|x - y| + |y - z|)$$

Max Norm

Lemma 1: $(\mathbb{R}^{|S|}, \|\cdot\|_\infty)$ is a metric space.

Definition: Let M be a set. A *metric* on M is a function $d : M \times M \rightarrow [0, \infty)$ which satisfies the following three conditions for all $x, y, z \in M$:

1. $d(x, y) = 0$ if and only if $x = y$
2. $d(x, y) = d(y, x)$
3. $d(x, y) \leq d(x, z) + d(z, y)$

Proof: Note: $\|x - y\|_\infty = \max_i |x_i - y_i|$

$$1. \max |x - y| = 0 \text{ iff } x_i = y_i \quad \forall i$$

$$2. |x - y| = |-(x - y)| = |y - x|$$
$$\therefore \max |x - y| = \max |y - x|$$

$$3. \max |x - z| = \max |x - y + y - z|$$
$$\leq \max(|x - y| + |y - z|)$$
$$\leq \max |x - y| + \max |y - z|$$

Bellman Operator Contraction

Bellman Operator Contraction

Lemma 2: B is a γ contraction mapping on $(\mathbb{R}^{|S|}, \|\cdot\|_\infty)$.

Bellman Operator Contraction

Lemma 2: B is a γ contraction mapping on $(\mathbb{R}^{|S|}, \|\cdot\|_\infty)$.

Proof

Bellman Operator Contraction

Lemma 2: B is a γ contraction mapping on $(\mathbb{R}^{|S|}, \|\cdot\|_\infty)$.

Proof

$$\|B[V_1] - B[V_2]\|_\infty = \max_{s \in S} |B[V_1](s) - B[V_2](s)|$$

Bellman Operator Contraction

Lemma 2: B is a γ contraction mapping on $(\mathbb{R}^{|S|}, \|\cdot\|_\infty)$.

Proof

$$\begin{aligned}\|B[V_1] - B[V_2]\|_\infty &= \max_{s \in S} |B[V_1](s) - B[V_2](s)| \\ &= \max_{s \in S} \left| \max_{a \in A} \left(R(s, a) + \gamma \sum_{s' \in S} T(s'|s, a) V_1(s') \right) - \max_{a \in A} \left(R(s, a) + \gamma \sum_{s' \in S} T(s'|s, a) V_2(s') \right) \right|\end{aligned}$$

Bellman Operator Contraction

Lemma 2: B is a γ contraction mapping on $(\mathbb{R}^{|S|}, \|\cdot\|_\infty)$.

Proof

$$\begin{aligned}\|B[V_1] - B[V_2]\|_\infty &= \max_{s \in S} |B[V_1](s) - B[V_2](s)| \\ &= \max_{s \in S} \left| \max_{a \in A} \left(R(s, a) + \gamma \sum_{s' \in S} T(s'|s, a) V_1(s') \right) - \max_{a \in A} \left(R(s, a) + \gamma \sum_{s' \in S} T(s'|s, a) V_2(s') \right) \right| \\ &\leq \max_{s \in S} \left| \max_{a \in A} \left(R(s, a) + \gamma \sum_{s' \in S} T(s'|s, a) V_1(s') - R(s, a) - \gamma \sum_{s' \in S} T(s'|s, a) V_2(s') \right) \right|\end{aligned}$$

Bellman Operator Contraction

Lemma 2: B is a γ contraction mapping on $(\mathbb{R}^{|S|}, \|\cdot\|_\infty)$.

Proof

$$\begin{aligned}\|B[V_1] - B[V_2]\|_\infty &= \max_{s \in S} |B[V_1](s) - B[V_2](s)| \\ &= \max_{s \in S} \left| \max_{a \in A} \left(R(s, a) + \gamma \sum_{s' \in S} T(s'|s, a) V_1(s') \right) - \max_{a \in A} \left(R(s, a) + \gamma \sum_{s' \in S} T(s'|s, a) V_2(s') \right) \right| \\ &\leq \max_{s \in S} \left| \max_{a \in A} \left(R(s, a) + \gamma \sum_{s' \in S} T(s'|s, a) V_1(s') - R(s, a) - \gamma \sum_{s' \in S} T(s'|s, a) V_2(s') \right) \right| \\ &\qquad\qquad\qquad | \max(x) | \leq \max |x|\end{aligned}$$

Bellman Operator Contraction

Lemma 2: B is a γ contraction mapping on $(\mathbb{R}^{|S|}, \|\cdot\|_\infty)$.

Proof

$$\begin{aligned} \|B[V_1] - B[V_2]\|_\infty &= \max_{s \in S} |B[V_1](s) - B[V_2](s)| \\ &= \max_{s \in S} \left| \max_{a \in A} \left(R(s, a) + \gamma \sum_{s' \in S} T(s'|s, a) V_1(s') \right) - \max_{a \in A} \left(R(s, a) + \gamma \sum_{s' \in S} T(s'|s, a) V_2(s') \right) \right| \\ &\leq \max_{s \in S} \left| \max_{a \in A} \left(R(s, a) + \gamma \sum_{s' \in S} T(s'|s, a) V_1(s') - R(s, a) - \gamma \sum_{s' \in S} T(s'|s, a) V_2(s') \right) \right| \\ &\leq \max_{s \in S, a \in A} \left| \gamma \sum_{s' \in S} T(s'|s, a) (V_1(s') - V_2(s')) \right| \qquad \qquad \qquad |\max(x)| \leq \max |x| \end{aligned}$$

Bellman Operator Contraction

Lemma 2: B is a γ contraction mapping on $(\mathbb{R}^{|S|}, \|\cdot\|_\infty)$.

Proof

$$\begin{aligned}\|B[V_1] - B[V_2]\|_\infty &= \max_{s \in S} |B[V_1](s) - B[V_2](s)| \\&= \max_{s \in S} \left| \max_{a \in A} \left(R(s, a) + \gamma \sum_{s' \in S} T(s'|s, a) V_1(s') \right) - \max_{a \in A} \left(R(s, a) + \gamma \sum_{s' \in S} T(s'|s, a) V_2(s') \right) \right| \\&\leq \max_{s \in S} \left| \max_{a \in A} \left(R(s, a) + \gamma \sum_{s' \in S} T(s'|s, a) V_1(s') - R(s, a) - \gamma \sum_{s' \in S} T(s'|s, a) V_2(s') \right) \right| \\&\leq \max_{s \in S, a \in A} \left| \gamma \sum_{s' \in S} T(s'|s, a) (V_1(s') - V_2(s')) \right| \quad \left| \max(x) \right| \leq \max |x| \\&\leq \max_{s \in S, a \in A} \gamma \sum_{s' \in S} T(s'|s, a) |V_1(s') - V_2(s')|\end{aligned}$$

Bellman Operator Contraction

Lemma 2: B is a γ contraction mapping on $(\mathbb{R}^{|S|}, \|\cdot\|_\infty)$.

Proof

$$\begin{aligned} \|B[V_1] - B[V_2]\|_\infty &= \max_{s \in S} |B[V_1](s) - B[V_2](s)| \\ &= \max_{s \in S} \left| \max_{a \in A} \left(R(s, a) + \gamma \sum_{s' \in S} T(s'|s, a) V_1(s') \right) - \max_{a \in A} \left(R(s, a) + \gamma \sum_{s' \in S} T(s'|s, a) V_2(s') \right) \right| \\ &\leq \max_{s \in S} \left| \max_{a \in A} \left(R(s, a) + \gamma \sum_{s' \in S} T(s'|s, a) V_1(s') - R(s, a) - \gamma \sum_{s' \in S} T(s'|s, a) V_2(s') \right) \right| \\ &\leq \max_{s \in S, a \in A} \left| \gamma \sum_{s' \in S} T(s'|s, a) (V_1(s') - V_2(s')) \right| \quad | \max(x) | \leq \max |x| \\ &\leq \max_{s \in S, a \in A} \gamma \sum_{s' \in S} T(s'|s, a) |V_1(s') - V_2(s')| \\ &\leq \max_{s \in S, a \in A} \gamma \sum_{s' \in S} T(s'|s, a) \|V_1 - V_2\|_\infty \end{aligned}$$

Bellman Operator Contraction

Lemma 2: B is a γ contraction mapping on $(\mathbb{R}^{|S|}, \|\cdot\|_\infty)$.

Proof

$$\begin{aligned} \|B[V_1] - B[V_2]\|_\infty &= \max_{s \in S} |B[V_1](s) - B[V_2](s)| \\ &= \max_{s \in S} \left| \max_{a \in A} \left(R(s, a) + \gamma \sum_{s' \in S} T(s'|s, a) V_1(s') \right) - \max_{a \in A} \left(R(s, a) + \gamma \sum_{s' \in S} T(s'|s, a) V_2(s') \right) \right| \\ &\leq \max_{s \in S} \left| \max_{a \in A} \left(R(s, a) + \gamma \sum_{s' \in S} T(s'|s, a) V_1(s') - R(s, a) - \gamma \sum_{s' \in S} T(s'|s, a) V_2(s') \right) \right| \\ &\leq \max_{s \in S, a \in A} \left| \gamma \sum_{s' \in S} T(s'|s, a) (V_1(s') - V_2(s')) \right| \quad | \max(x) | \leq \max |x| \\ &\leq \max_{s \in S, a \in A} \gamma \sum_{s' \in S} T(s'|s, a) |V_1(s') - V_2(s')| \\ &\leq \max_{s \in S, a \in A} \gamma \sum_{s' \in S} T(s'|s, a) \|V_1 - V_2\|_\infty \\ &= \gamma \|V_1 - V_2\|_\infty \max_{s \in S, a \in A} \sum_{s' \in S} T(s'|s, a) \end{aligned}$$

Bellman Operator Contraction

Lemma 2: B is a γ contraction mapping on $(\mathbb{R}^{|S|}, \|\cdot\|_\infty)$.

Proof

$$\begin{aligned} \|B[V_1] - B[V_2]\|_\infty &= \max_{s \in S} |B[V_1](s) - B[V_2](s)| \\ &= \max_{s \in S} \left| \max_{a \in A} \left(R(s, a) + \gamma \sum_{s' \in S} T(s'|s, a) V_1(s') \right) - \max_{a \in A} \left(R(s, a) + \gamma \sum_{s' \in S} T(s'|s, a) V_2(s') \right) \right| \\ &\leq \max_{s \in S} \left| \max_{a \in A} \left(R(s, a) + \gamma \sum_{s' \in S} T(s'|s, a) V_1(s') - R(s, a) - \gamma \sum_{s' \in S} T(s'|s, a) V_2(s') \right) \right| \\ &\leq \max_{s \in S, a \in A} \left| \gamma \sum_{s' \in S} T(s'|s, a) (V_1(s') - V_2(s')) \right| \quad | \max(x) | \leq \max |x| \\ &\leq \max_{s \in S, a \in A} \gamma \sum_{s' \in S} T(s'|s, a) |V_1(s') - V_2(s')| \\ &\leq \max_{s \in S, a \in A} \gamma \sum_{s' \in S} T(s'|s, a) \|V_1 - V_2\|_\infty \\ &= \gamma \|V_1 - V_2\|_\infty \max_{s \in S, a \in A} \sum_{s' \in S} T(s'|s, a) \\ &= \gamma \|V_1 - V_2\|_\infty \end{aligned}$$

Value Iteration Convergence

Value Iteration Convergence

Theorem 1: Let $\{V_1, \dots, V_\infty\}$ be a sequence of value functions for a discrete MDP generated by the recurrence $V_{k+1} = B[V_k]$. If $\gamma < 1$, then $\lim_{k \rightarrow \infty} V_k = V^*$.

Value Iteration Convergence

Theorem 1: Let $\{V_1, \dots, V_\infty\}$ be a sequence of value functions for a discrete MDP generated by the recurrence $V_{k+1} = B[V_k]$. If $\gamma < 1$, then $\lim_{k \rightarrow \infty} V_k = V^*$.

Lemma 2: B is a γ contraction mapping on $(\mathbb{R}^{|S|}, \|\cdot\|_\infty)$.

Value Iteration Convergence

Theorem 1: Let $\{V_1, \dots, V_\infty\}$ be a sequence of value functions for a discrete MDP generated by the recurrence $V_{k+1} = B[V_k]$. If $\gamma < 1$, then $\lim_{k \rightarrow \infty} V_k = V^*$.

Lemma 2: B is a γ contraction mapping on $(\mathbb{R}^{|S|}, \|\cdot\|_\infty)$.

Theorem (Banach): If f is a contraction mapping on metric space (M, d) , then

1. f has a single, unique fixed point x^* .
2. If $\{x_k\}$ is a sequence defined by $x_{k+1} = f(x_k)$, then $\lim_{k \rightarrow \infty} x_k = x^*$.

Value Iteration Convergence

Theorem 1: Let $\{V_1, \dots, V_\infty\}$ be a sequence of value functions for a discrete MDP generated by the recurrence $V_{k+1} = B[V_k]$. If $\gamma < 1$, then $\lim_{k \rightarrow \infty} V_k = V^*$.

Proof:

Lemma 2: B is a γ contraction mapping on $(\mathbb{R}^{|S|}, \|\cdot\|_\infty)$.

Theorem (Banach): If f is a contraction mapping on metric space (M, d) , then

1. f has a single, unique fixed point x^* .
2. If $\{x_k\}$ is a sequence defined by $x_{k+1} = f(x_k)$, then $\lim_{k \rightarrow \infty} x_k = x^*$.

Value Iteration Convergence

Theorem 1: Let $\{V_1, \dots, V_\infty\}$ be a sequence of value functions for a discrete MDP generated by the recurrence $V_{k+1} = B[V_k]$. If $\gamma < 1$, then $\lim_{k \rightarrow \infty} V_k = V^*$.

Proof:

Lemma 2: B is a γ contraction mapping on $(\mathbb{R}^{|S|}, \|\cdot\|_\infty)$.

Theorem (Banach): If f is a contraction mapping on metric space (M, d) , then

1. f has a single, unique fixed point x^* .
2. If $\{x_k\}$ is a sequence defined by $x_{k+1} = f(x_k)$, then $\lim_{k \rightarrow \infty} x_k = x^*$.

By Lemma 2 and Banach's theorem (part 2), repeated application of the Bellman operator always has a fixed point limit, \hat{V} .

Value Iteration Convergence

Theorem 1: Let $\{V_1, \dots, V_\infty\}$ be a sequence of value functions for a discrete MDP generated by the recurrence $V_{k+1} = B[V_k]$. If $\gamma < 1$, then $\lim_{k \rightarrow \infty} V_k = V^*$.

Proof:

Lemma 2: B is a γ contraction mapping on $(\mathbb{R}^{|S|}, \|\cdot\|_\infty)$.

Theorem (Banach): If f is a contraction mapping on metric space (M, d) , then

1. f has a single, unique fixed point x^* .
2. If $\{x_k\}$ is a sequence defined by $x_{k+1} = f(x_k)$, then $\lim_{k \rightarrow \infty} x_k = x^*$.

By Lemma 2 and Banach's theorem (part 2), repeated application of the Bellman operator always has a fixed point limit, \hat{V} .

By Banach's theorem (part 1), $\hat{V} = B[\hat{V}]$. Since \hat{V} satisfies Bellman's equation, it is optimal and $\hat{V} = V^*$.

Does Policy Iteration Converge?

Does Policy Iteration Converge?

Theorem: Policy iteration converges to an optimal policy for a finite MDP in finite time.

Does Policy Iteration Converge?

Theorem: Policy iteration converges to an optimal policy for a finite MDP in finite time.

Proof (sketch):

Does Policy Iteration Converge?

Theorem: Policy iteration converges to an optimal policy for a finite MDP in finite time.

Proof (sketch):

1. The policy will either improve or stay the same at each iteration

Does Policy Iteration Converge?

Theorem: Policy iteration converges to an optimal policy for a finite MDP in finite time.

Proof (sketch):

1. The policy will either improve or stay the same at each iteration
2. The policy will stay the same if and only if $V^\pi = V^*$

Does Policy Iteration Converge?

Theorem: Policy iteration converges to an optimal policy for a finite MDP in finite time.

Proof (sketch):

1. The policy will either improve or stay the same at each iteration
2. The policy will stay the same if and only if $V^\pi = V^*$
3. There are a finite number of possible policies

Does Policy Iteration Converge?

Theorem: Policy iteration converges to an optimal policy for a finite MDP in finite time.

Proof (sketch):

1. The policy will either improve or stay the same at each iteration
2. The policy will stay the same if and only if $V^\pi = V^*$
3. There are a finite number of possible policies
4. By (1), (2), and (3), the policy will improve until it finds the optimal policy, and it will always find the optimal policy.

**Is there always a deterministic
optimal policy?**

Guiding Questions

Guiding Questions

- Does value iteration always converge?
- Is the value function unique?
- Can there be multiple optimal policies?
- Is there always a deterministic optimal policy?

Break

Conservation MDP

