# How often does 2kp+1 divide $p^p-1$ ?

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May 27, 2011

### Background

This work is part of an investigation of the minimum period of the Bell numbers B(n) reduced modulo a prime p.

Many congruences have been proved for B(n) modulo p. The basic facts we need are these:

### J. Touchard's congruence [1933]

$$B(n+p) \equiv B(n) + B(n+1) \bmod p,$$

valid for any prime p and for all  $n \ge 0$ , shows that any p consecutive values of  $B(n) \mod p$  determine the sequence modulo p after that point.

It follows from this congruence that B(n) mod p must be periodic with period  $\leq p^p$ .

In 1945, G. T. Williams proved that for each prime p, the period of the Bell numbers modulo p divides

$$N_p = (p^p - 1)/(p - 1).$$

In fact the minimum period equals  $N_p$  for every prime p for which this period is known.

**Theorem 1**. The minimum period of the sequence  $\{B(n) \bmod p\}$  is  $N_p$  when p is a prime < 126 and also when p = 137, 149, 157, 163, 167 or 173.

This theorem is proved by showing that the period does not divide  $N_p/q$  for any prime divisor q of  $N_p$ .

If q divides  $N_p$  and  $N=N_p/q$ , then one can test whether the period of the Bell numbers modulo p divides N by checking whether  $B(N+i)\equiv B(i) \mod p$  for  $0 \le i \le p-1$ . The period divides N if and only if all p of these congruences hold.

A polynomial time algorithm for computing  $B(n) \mod p$  is known.

The theorem for p can be proved (or disproved) this way if we know the factorization of  $N_p$ .

It is conjectured that the minimum period of the Bell numbers modulo p equals  $N_p$  for every prime p.

The conjecture is known to be true for all primes < 126 and for a few larger primes.

We give a heuristic argument for the probability that the conjecture holds for a prime p.

The most difficult piece of this heuristic argument is determining the probability that a prime q divides  $N_p$ . This probability is studied in this talk.

# How often does 2kp + 1 divide $N_p$ as p varies?

It is well known (Euler, 1755) that when p is prime every prime factor of  $N_p$  has the form 2kp+1.

For each  $1 \le k \le 50$  and for all odd primes p < 100000, we computed the fraction of the primes q = 2kp + 1 that divide  $N_p$ .

For example, when k=5 there are 1352 primes p<100000 for which q=2kp+1 is also prime, and 129 of these q divide  $N_p$ , so the fraction is 129/1352=0.095.

This fraction is called "Prob" in the table because it approximates the probability that q divides  $N_p$ , given that p and q=2kp+1 are prime, for fixed k.

## Probability that $q = (2kp + 1) \mid N_p$

k	Prob
1	0.503
2	1.000
3	0.171
4	0.247
5	0.095
6	0.173
7	0.076
8	0.496
9	0.047
10	0.096
11	0.042
12	0.082
13	0.051
14	0.068
15	0.033
16	0.064
17	0.032
18	0.111
19	0.021
20	0.050

# Probability that $q = (2kp + 1) \mid N_p$

k	1/k	Prob	
1	1.000	0.503	
2	0.500	1.000	
3	0.333	0.171	
4	0.250	0.247	
5	0.200	0.095	
6	0.167	0.173	
7	0.143	0.076	
8	0.125	0.496	
9	0.111	0.047	
10	0.100	0.096	
11	0.091	0.042	
12	0.083	0.082	
13	0.077	0.051	
14	0.071	0.068	
15	0.067	0.033	
16	0.063	0.064	
17	0.059	0.032	
18	0.056	0.111	
19	0.053	0.021	
20	0.050	0.050	

Probability that $q = (2kp + 1) \mid N_p$						
$Odd\;k$			Even k			
k	1/(2k)	Prob	k	1/k	Prob	
1	0.500	0.503	2	0.500	1.000	
3	0.167	0.171	4	0.250	0.247	
5	0.100	0.095	6	0.167	0.173	
7	0.071	0.076	8	0.125	0.496	
9	0.056	0.047	10	0.100	0.096	
11	0.045	0.042	12	0.083	0.082	
13	0.038	0.051	14	0.071	0.068	
15	0.033	0.033	16	0.063	0.064	
17	0.029	0.032	18	0.056	0.111	
19	0.026	0.021	20	0.050	0.050	
21	0.024	0.016	22	0.045	0.054	
23	0.022	0.021	24	0.042	0.042	
25	0.020	0.021	26	0.038	0.052	
27	0.019	0.021	28	0.036	0.036	
29	0.017	0.022	30	0.033	0.031	
31	0.016	0.019	32	0.031	0.055	
49	0.010	0.014	50	0.020	0.043	

#### Obervations about the table

- 1. Prob is approximately 1/(2k) when k is odd.
- 2. Usually Prob is approximately 1/k when k is even.
- 3. Some anomalies to 2. are that Prob is about 2/k when k = 2, 18, 32 and 50.
- 4. Also, Prob is about 4/k when k = 8.
- 5. The exceptional values of k in 3. and 4. have the form  $2m^2$  for  $1 \le m \le 5$ . (These numbers also arise as the lengths of the rows in the periodic table of elements in chemistry.)

We will now explain these observations. Suppose k is a positive integer and that both p and q=2kp+1 are odd primes. Let g be a primitive root modulo q.

If  $p \equiv 1 \mod 4$  or k is even (so  $q \equiv 1 \mod 4$ ), then by the Law of Quadratic Reciprocity

$$\left(\frac{p}{q}\right) = \left(\frac{q}{p}\right) = \left(\frac{2kp+1}{p}\right) = \left(\frac{1}{p}\right) = +1,$$

so p is a quadratic residue modulo q. In this case  $g^{2s} \equiv p \mod q$  for some s. Now by Euler's criterion for power residues,  $(2kp+1) \mid (p^p-1)$  if and only if p is a (2k)-ic residue of 2kp+1, that is, if and only if  $(2k) \mid (2s)$ . It is natural to assume that  $k \mid s$  with probability 1/k because k is fixed and s is a random integer.

If  $p \equiv 3 \mod 4$  and k is odd (so  $q \equiv 3 \mod 4$ ), then

$$\left(\frac{p}{q}\right) = -\left(\frac{q}{p}\right) = -\left(\frac{2kp+1}{p}\right) = -\left(\frac{1}{p}\right) = -1,$$

so p is a quadratic nonresidue modulo q. Now  $g^{2s+1} \equiv p \mod q$  for some s. Reasoning as before,  $(2kp+1) \mid (p^p-1)$  if and only if  $(2k) \mid (2s+1)$ , which is impossible. Therefore q does not divide  $N_p$ .

Thus, if we fix k and let p run over all primes, then the probability that q=2kp+1 divides  $N_p$  is 1/k when k is even and 1/(2k) when k is odd because, when k is odd only those  $p\equiv 1 \mod 4$  (that is, half of the primes p) offer a chance for q to divide  $N_p$ .

In fact, when k = 1 and  $p \equiv 1 \mod 4$ , q always divides  $N_p$ . This theorem must have been known long ago, but we could not find it in the literature.

**Theorem 2**. If p is odd and q = 2p+1 is prime, then q divides  $N_p$  if and only if  $p \equiv 1 \mod 4$ .

*Proof.* We have just seen that q does not divide  $N_p$  when  $p \equiv 3 \mod 4$ . If  $p \equiv 1 \mod 4$ , then p is a quadratic residue modulo q, as was mentioned above, so  $p^p = p^{(q-1)/2} \equiv +1 \mod q$  by Euler's criterion. Finally, q is too large to divide p-1, so q divides  $N_p$ .

We now explain the anomalies, beginning with k = 2.

**Theorem 3**. If q = 4p + 1 is prime, then q divides  $N_p$ .

This result was an ancient problem posed and solved more than 100 years ago. Here is a modern proof.

*Proof.* Since  $q \equiv 1 \mod 4$ , there exists an integer i with  $i^2 \equiv -1 \mod q$ . Then

$$(1+i)^4 \equiv (2i)^2 \equiv -4 \equiv \frac{1}{p} \mod q.$$

Hence

$$p^p \equiv \left(\frac{1}{p}\right)^{-p} \equiv (1+i)^{-4p} \equiv (1+i)^{1-q} \equiv 1 \bmod q$$

by Fermat's theorem. Thus, q divides  $p^p - 1$ . But q = 4p + 1 is too large to divide p - 1, so q divides  $N_p$ . **Theorem 4**. Let p be an odd positive integer and m be a positive integer. If  $q = 4m^2p + 1$  is prime, then q divides  $p^{m^2p} - 1$ .

Of course, Theorem 3 is the case m=1 of Theorem 4.

Fermat's theorem says that q divides  $p^{4m^2p}-1$ .

In the case m=2, that is, k=8, we can do even better.

Fermat's theorem says that q divides  $p^{16p} - 1$ .

**Theorem 5** If q = 16p + 1 is prime, then q divides  $p^{2p} - 1$ .

*Proof.* As in the proof of the previous theorem, we have i with  $i^2 \equiv -1 \mod q$  and  $(1+i)^4 \equiv -4 \mod q$ . Therefore,  $(1+i)^8 \equiv 16 \equiv -1/p \mod q$  and so

$$p^{2p} \equiv (1+i)^{-16p} \equiv (1+i)^{1-q} \equiv 1 \mod q$$

which proves the theorem.

Thus, a prime q = 2kp + 1 = 16p + 1 divides  $(p^p - 1)(p^p + 1)$  when k = 8. Assuming that q has equal chance to divide either factor, the probability that q divides  $p^p - 1$  is 1/2.

So far, we have explained all the behavior seen in the table. Further experiments with  $q=2m^2p+1$  lead us to the following result, which generalizes Theorems 4 and 5.

The theorem lets us remove arbitrarily large powers of 2 from the exponent in certain cases.

**Theorem 6.** Suppose p, m, t are positive integers, with t a power of 2 and t > 1. Let  $k = (2m)^t/2$  and  $q = 2kp + 1 = (2m)^tp + 1$ . If q is prime, then t divides k and  $p^{kp/t} \equiv 1 \mod q$ .

When t = 2, the theorem is just Theorem 4.

When t=4, Theorem 6 says that if  $q=(2m)^4+1=16m^4+1$  is prime, then q divides  $p^{2m^4p}-1$ . Theorem 5 is the case m=1 of this statement.

When t=8, Theorem 6 says that if  $q=(2m)^8+1=256m^8+1$  is prime, then q divides  $p^{16m^8p}-1$ . The first case, m=1, of this statement is for k=128, which is beyond the end of the table.

We now apply Theorem 6. As above, let g be a primitive root modulo q and let  $a=g^{(q-1)t/k} \bmod q$ . Then  $a^j$ ,  $0 \le j < k/t$ , are all the solutions to  $x^{k/t} \equiv 1 \bmod q$ . Let  $b=p^p \bmod q$ . By the theorem,  $b^{k/t} \equiv 1 \bmod q$ , so  $b \equiv a^j \bmod q$  for some  $0 \le j < k/t$ . It is natural to assume that j=0, that is,  $q \mid N_p$ , happens with probability 1/(k/t)=t/k.

#### Summary

We have given heuristic arguments which conclude that, for fixed k, when p and q = 2kp + 1 are both prime, the probability that q divides  $N_p$  is c(k)/k, where c(k) is defined as follows:

When k is an odd positive integer, c(k) = 1/2.

When k is an even positive integer, let t be the largest power of 2 for which there exists an integer m so that  $2k = (2m)^t$ . Then c(k) = t.

We have

$$c(k) = \begin{cases} 1/2 & \text{if } k \text{ is odd,} \\ 1 & \text{if } k \text{ is even and } k \neq 2m^2, \\ O(\log k) & \text{if } k = 2m^2 \text{ for some } m. \end{cases}$$

The average value of c(k) is 3/4 because the numbers  $2m^2$  are rare.

Is the conjecture about the Bell numbers' period true? Does it always equal  $N_p$ ?

Applying the Bateman-Horn conjecture, the Prime Number Theorem, the Binomial Theorem and the divisibility results for  $N_p$ , one can show that the heuristic probability that the minimum period of the Bell numbers modulo p is  $N_p$  is

$$(1-p^{-p})^{3(\log p)/2} \approx 1 - \frac{3\log p}{2p^p},$$

exceedingly close to 1 when p is large.

Finally, we compute the expected number of primes p>x for which the conjecture fails. When x>2, this number is

$$\sum_{p>x} \frac{3\log p}{2p^p} < \sum_{p>x} p^{1-x} \le \int_x^\infty t^{1-x} dt = \frac{x^{2-x}}{x-2}.$$

By Theorem 1, the conjecture holds for all primes p < 126. Taking x = 126, the expected number of primes for which the conjecture fails is  $< 126^{-124}/124 < 10^{-262}$ . Thus, the heuristic argument predicts that the conjecture is almost certainly true.