

Supplementary Material: Information-Theoretic Limits of AI Alignment

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Abstract

This supplementary document provides complete mathematical proofs, extended experimental data, and detailed methodology for the main paper "Information-Theoretic Limits of AI Alignment." We present rigorous derivations of the three impossibility theorems (Shannon's Detection Limit, Kolmogorov Indistinguishability, Manifold Collapse), comprehensive experimental protocols, and theoretical extensions.

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1 Appendix A: Shannon's Detection Impossibility

1.1 A.1 Information Theory Background

Definition 1 (Shannon Entropy). *For a discrete random variable X with probability mass function $p(x)$:*

$$H(X) = - \sum_{x \in \mathcal{X}} p(x) \log_2 p(x) \quad (1)$$

Entropy measures uncertainty. Maximum entropy (maximum uncertainty) is $\log_2 |\mathcal{X}|$ for uniform distribution.

Definition 2 (Conditional Entropy). *The conditional entropy of X given Y :*

$$H(X|Y) = - \sum_{y \in \mathcal{Y}} p(y) \sum_{x \in \mathcal{X}} p(x|y) \log_2 p(x|y) \quad (2)$$

Measures remaining uncertainty about X after observing Y .

Definition 3 (Mutual Information).

$$I(X; Y) = H(X) - H(X|Y) = H(Y) - H(Y|X) \quad (3)$$

Measures information shared between X and Y . Symmetric.

1.2 A.2 Fano's Inequality

Theorem 4 (Fano's Inequality). *Let \hat{X} be an estimator of X based on observation Y . Define error probability $P_e = P(\hat{X} \neq X)$. Then:*

$$H(X|Y) \geq H(P_e) + P_e \log_2(|\mathcal{X}| - 1) \quad (4)$$

where $H(P_e) = -P_e \log_2 P_e - (1 - P_e) \log_2(1 - P_e)$ is the binary entropy function.

Proof. Define error indicator $E = \mathbb{I}[\hat{X} \neq X]$. By chain rule:

$$H(E, X|Y) = H(X|Y) = H(E|Y) + H(X|E, Y) \quad (5)$$

Since E is determined by (X, Y, \hat{X}) and \hat{X} is determined by Y , we have $H(E|Y) \leq H(E) = H(P_e)$.

For the second term, conditioned on $E = 0$ (no error), X is deterministic ($H(X|E = 0, Y) = 0$). Conditioned on $E = 1$ (error), X can be any of the $|\mathcal{X}| - 1$ values $X \neq \hat{X}$, giving maximum entropy $\log_2(|\mathcal{X}| - 1)$.

Therefore:

$$H(X|E, Y) \leq P_e \cdot \log_2(|\mathcal{X}| - 1) \quad (6)$$

Combining: $H(X|Y) \leq H(P_e) + P_e \log_2(|\mathcal{X}| - 1)$. \square

1.3 A.3 Proof of Theorem 1

Theorem 5 (Detection Impossibility Under Ambiguity). *Let X be user intent (malicious/legitimate), Y the observable prompt, and C context. A safety filter F attempts to classify X based on (Y, C) . If context C is ϵ -ambiguous, meaning:*

$$H(X|Y, C) \geq H(X) - \epsilon \quad (7)$$

then the classification accuracy is bounded by:

$$P(F \text{ correct}) \leq \frac{1}{2} + \frac{\epsilon}{2H(X)} \quad (8)$$

Proof. For binary classification ($|\mathcal{X}| = 2$), Fano's inequality gives:

$$H(X|Y, C) \geq H(P_e) \quad (9)$$

By the ϵ -ambiguity assumption:

$$H(X) - \epsilon \leq H(X|Y, C) \geq H(P_e) \quad (10)$$

Therefore:

$$H(P_e) \leq H(X) - \epsilon \quad (11)$$

Since $H(P_e) = -P_e \log_2 P_e - (1 - P_e) \log_2(1 - P_e)$ is maximized at $P_e = 0.5$ with $H(0.5) = 1$ bit, and decreases monotonically as $P_e \rightarrow 0$ or $P_e \rightarrow 1$, we can bound:

For small ϵ , when $H(P_e) \leq H(X) - \epsilon < H(X)$, we need:

$$P_e \geq \frac{1}{2} - \frac{\epsilon}{2H(X)} \quad (12)$$

Therefore, accuracy $P(\text{correct}) = 1 - P_e$ satisfies:

$$P(\text{correct}) \leq \frac{1}{2} + \frac{\epsilon}{2H(X)} \quad (13)$$

For binary intent ($H(X) = 1$ bit) and our attack achieving $\epsilon = 0.1$ bits:

$$P(\text{correct}) \leq 0.5 + 0.05 = 0.55 \quad (14)$$

\square

1.4 A.4 Tightness Analysis

The bound is *tight* (achievable) when:

- The classifier uses optimal Bayesian inference: $\hat{X} = \arg \max_x P(X = x|Y, C)$
- The distribution $P(X|Y, C)$ has maximum entropy subject to $H(X|Y, C) = H(X) - \epsilon$

Our experimental results (Table 1 in main paper) show observed accuracy tracks the theoretical bound within 2-3%, confirming tightness.

2 Appendix B: Kolmogorov Indistinguishability

2.1 B.1 Kolmogorov Complexity Primer

Definition 6 (Kolmogorov Complexity). *The Kolmogorov Complexity $K(x)$ of a string x is:*

$$K(x) = \min\{|p| : U(p) = x\} \quad (15)$$

where U is a universal Turing machine and $|p|$ is the length of program p in bits.

Key properties:

- $K(x) \leq |x| + O(1)$ (can always hardcode)
- $K(x)$ is *uncomputable* (no algorithm computes it for all x)
- $K(xy) \leq K(x) + K(y) + O(\log \min(K(x), K(y)))$ (composition)

Definition 7 (Conditional Kolmogorov Complexity).

$$K(x|y) = \min\{|p| : U(p, y) = x\} \quad (16)$$

Complexity of x given y as auxiliary input.

2.2 B.2 Construction of Attack

Theorem 8 (Attack Indistinguishability). *If $K(R_{\text{attack}}) \geq K(R_{\text{legit}}) - O(\log n)$ where n is the number of possible intents, then no polynomial-time algorithm can distinguish R_{attack} from R_{legit} with probability better than $1/2 + \text{negl}(n)$.*

Proof. **Step 1: Decomposition**

Any request R can be decomposed as:

$$R = (C, I) \quad (17)$$

where C is content (papers, credentials, terminology, framing) and I is intent (malicious/legitimate).

For CPF attacks, content is *identical* to legitimate research:

- Papers: real published PDFs (bit-for-bit identical)
- Credentials: valid CISSP certification
- Terminology: accurate technical language
- Framing: genuine academic discourse

Therefore:

$$K(C_{\text{attack}}) = K(C_{\text{legit}}) \quad (18)$$

Step 2: Intent Encoding

The intent $I \in \{\text{malicious, legitimate}\}$ requires encoding. In general, for n possible intents:

$$K(I) = O(\log n) \quad (19)$$

For binary intent ($n = 2$), $K(I) = O(1)$ bits.

Step 3: Total Complexity

By composition:

$$K(R) = K(C, I) \quad (20)$$

$$\leq K(C) + K(I|C) + O(\log K(C)) \quad (21)$$

$$= K(C) + O(\log n) + O(\log K(C)) \quad (22)$$

Since $K(C) \gg \log n$ (content is $\sim 60,000$ bits), the intent encoding is negligible:

$$K(R_{\text{attack}}) \approx K(R_{\text{legit}}) \quad (23)$$

More precisely:

$$|K(R_{\text{attack}}) - K(R_{\text{legit}})| \leq O(\log n) \quad (24)$$

Step 4: Distinguishability

Any distinguisher D must compute a function $f : R \rightarrow \{0, 1\}$ mapping requests to classifications. The distinguisher succeeds if:

$$|P(D(R_{\text{attack}}) = 1) - P(D(R_{\text{legit}}) = 1)| \geq \delta \quad (25)$$

But if $K(R_{\text{attack}}) \approx K(R_{\text{legit}})$, the programs generating them differ only in $O(\log n)$ bits. Any polynomial-time algorithm cannot extract this difference without additional side information.

By the incompressibility argument, most strings of length ℓ have $K(x) \geq \ell - O(1)$. If R_{attack} and R_{legit} are both incompressible (high complexity), they appear random to polynomial-time algorithms.

Therefore:

$$P(D \text{ correct}) \leq \frac{1}{2} + \text{negl}(n) \quad (26)$$

□

2.3 B.3 Practical Approximation

Since $K(x)$ is uncomputable, we approximate via:

- **Compression:** Use gzip or similar. For our attack: compressed size \approx baseline compressed size.
- **Mahalanobis Distance:** Measures statistical typicality in embedding space. Result: $D_M = 1.18\sigma$ (normal).

Both approximations confirm high complexity indistinguishable from legitimate requests.

3 Appendix C: Manifold Collapse

3.1 C.1 Riemannian Manifold Formulation

Model the LLM's latent space as a Riemannian manifold (\mathcal{M}, g) where:

- $\mathcal{M} \subset \mathbb{R}^d$ is the space of internal representations ($d \sim 10^4$ for modern LLMs)
- g is the metric tensor determining distances and gradients

Definition 9 (Safety Potential). *Define $\Phi_{\text{safe}} : \mathcal{M} \rightarrow \mathbb{R}$ as the safety potential learned during RLHF training:*

$$\Phi_{\text{safe}}(z) = \mathbb{E}_{(x,y) \sim \mathcal{D}_{\text{RLHF}}}[\text{reward}(y|x, z)] \quad (27)$$

where $z \in \mathcal{M}$ is the latent state.

Definition 10 (Safety Gradient). *The gradient of Φ_{safe} in the Riemannian manifold:*

$$\nabla \Phi_{\text{safe}} = g^{ij} \frac{\partial \Phi_{\text{safe}}}{\partial x^j} \quad (28)$$

where g^{ij} is the inverse metric tensor.

3.2 C.2 Context-Dependent Metric

The metric tensor depends on context C :

$$g_{ij}(C) = g_{ij}^{(0)} + \sum_k \lambda_k(C) \cdot T_{ij}^{(k)} \quad (29)$$

where $T_{ij}^{(k)}$ are deformation tensors and $\lambda_k(C)$ are context-dependent coefficients.

Lemma 11 (Entropy-Driven Isotropy). *For high-entropy contexts with $H(C) \gg \log d$, the metric becomes approximately isotropic:*

$$g_{ij}(C) \rightarrow \delta_{ij} \quad (30)$$

Proof. High-entropy contexts make the probability distribution over latent states approximately uniform:

$$p(z|C) \approx \frac{1}{|\mathcal{Z}|} \text{ for } H(C) \gg \log |\mathcal{Z}| \quad (31)$$

The metric tensor is learned from correlations in the training data:

$$g_{ij} \propto \mathbb{E}_{z \sim p(\cdot|C)}[(z_i - \mu_i)(z_j - \mu_j)] \quad (32)$$

Under uniform distribution, all correlations vanish:

$$\mathbb{E}[(z_i - \mu_i)(z_j - \mu_j)] \rightarrow \begin{cases} \sigma^2 & i = j \\ 0 & i \neq j \end{cases} \quad (33)$$

Therefore $g_{ij} \rightarrow \sigma^2 \delta_{ij}$ (isotropic). \square

3.3 C.3 Gradient Vanishing Derivation

Theorem 12 (Gradient Vanishing Under High Entropy). *For contexts with $H(C) > H_{\text{crit}}$:*

$$\|\nabla \Phi_{\text{safe}}\| \leq \epsilon \cdot e^{-\alpha H(C)} \quad (34)$$

for constants $\epsilon, \alpha > 0$.

Proof. The safety potential is:

$$\Phi_{\text{safe}}(z) = \int p(z'|C) V_{\text{safe}}(z, z') dz' \quad (35)$$

where $V_{\text{safe}}(z, z')$ is the pairwise safety interaction learned from RLHF.

The gradient:

$$\nabla_z \Phi_{\text{safe}} = \int p(z'|C) \nabla_z V_{\text{safe}}(z, z') dz' \quad (36)$$

For high-entropy contexts with $p(z'|C) \approx 1/|\mathcal{Z}|$:

$$\nabla_z \Phi_{\text{safe}} \approx \frac{1}{|\mathcal{Z}|} \int \nabla_z V_{\text{safe}}(z, z') dz' \quad (37)$$

$$= \frac{1}{|\mathcal{Z}|} \cdot \mathbb{E}_{z'}[\nabla_z V_{\text{safe}}(z, z')] \quad (38)$$

The safety interaction V_{safe} is learned to distinguish safe from unsafe directions. It has structure:

$$V_{\text{safe}}(z, z') \sim \cos(\theta(z, z')) \quad (39)$$

where θ is the angle between safe and current direction.

Under uniform averaging, directional preferences cancel:

$$\mathbb{E}_{z' \sim \text{uniform}}[\cos(\theta(z, z'))] \rightarrow 0 \quad (40)$$

More precisely, for entropy $H(C)$, the effective averaging set size is $|\mathcal{Z}_{\text{eff}}| \sim 2^{H(C)}$, giving cancellation proportional to $1/2^{H(C)}$:

$$\|\nabla \Phi_{\text{safe}}\| \sim \|\nabla \Phi_{\text{safe}}^{(0)}\| \cdot 2^{-\alpha H(C)} = \epsilon \cdot e^{-\alpha \ln(2) H(C)} \quad (41)$$

where $\epsilon = \|\nabla \Phi_{\text{safe}}^{(0)}\|$ is the baseline gradient magnitude and $\alpha = \ln(2)$. \square

3.4 C.4 Threshold Analysis

The critical entropy H_{crit} where collapse begins:

$$H_{\text{crit}} \approx \frac{1}{\alpha} \ln \left(\frac{\epsilon}{\|\nabla \Phi_{\text{safe}}^{(\min)}\|} \right) \quad (42)$$

For our empirical observations:

- Baseline gradient: $\epsilon \approx 0.84$
- Minimum detectable gradient: $\|\nabla \Phi_{\text{safe}}^{(\min)}\| \approx 0.05$
- $\alpha \approx 10^{-3}$ (fitted)

This gives:

$$H_{\text{crit}} \approx \frac{\ln(0.84/0.05)}{10^{-3}} \approx 2800 \text{ bits} \quad (43)$$

Our attacks use $H(C) \approx 10^4$ bits, well above threshold.

4 Appendix D: Extended Experimental Data

4.1 D.1 Mahalanobis Distance Methodology

Embedding Model: sentence-transformers/all-MiniLM-L6-v2 (768 dimensions)

Baseline Corpus: 500 academic papers on AI safety from arXiv (2020-2025)

Covariance Estimation:

$$\Sigma = \frac{1}{n-1} \sum_{i=1}^n (\mathbf{x}_i - \mu)(\mathbf{x}_i - \mu)^T \quad (44)$$

Distance Calculation:

$$D_M(\mathbf{x}) = \sqrt{(\mathbf{x} - \boldsymbol{\mu})^T \boldsymbol{\Sigma}^{-1} (\mathbf{x} - \boldsymbol{\mu})} \quad (45)$$

Complete Results:

Category	n	Mean	Std	Min	Max	Det
Benign Academic	500	0.82	0.31	0.21	1.45	0
CPF Low Int.	10	0.95	0.28	0.61	1.33	0
CPF Medium Int.	20	1.08	0.35	0.71	1.89	0
CPF High Int.	20	1.18	0.42	0.83	2.31	0
GCG Attacks	50	4.73	1.21	3.12	7.89	48
Random Noise	50	8.21	2.14	4.89	13.67	50

Table 1: Complete Mahalanobis distance statistics. Detection threshold: 3σ .

4.2 D.2 Kolmogorov-Smirnov Test Details

Null Hypothesis: Attention weights on safety tokens follow uniform distribution.

Test Statistic:

$$D_{KS} = \sup_x |F_{\text{empirical}}(x) - F_{\text{uniform}}(x)| \quad (46)$$

Procedure:

1. Extract attention weights from layer 40 (final layer) of Claude Sonnet 4.5
2. Identify safety tokens: {"refuse", "cannot", "inappropriate", "harmful", "unsafe", "apologize"}
3. Measure attention allocated to these tokens
4. Compare to uniform distribution over all tokens

Results by Context Entropy:

$H(C)$ range	D_{KS}	p-value	n	Verdict
$< 10^3$ bits	0.47	< 0.001	50	Non-uniform
10^3 - $10^{3.5}$	0.39	< 0.001	50	Non-uniform
$10^{3.5}$ - 10^4	0.18	0.042	50	Borderline
$> 10^4$ bits	0.08	0.31	50	Uniform

Table 2: K-S test results vs. context entropy. Collapse threshold $\approx 10^{3.7}$ bits.

4.3 D.3 HMM Parameter Estimation

State Space: $\mathcal{S} = \{S_{\text{safe}}, S_{\text{compromised}}\}$

Transition Model:

$$P(S_{t+1} = \text{comp} | S_t, C_t) = \sigma(\alpha H(C_t) + \beta A(C_t) - \gamma) \quad (47)$$

where $\sigma(x) = 1/(1 + e^{-x})$ is the sigmoid function.

Parameter Estimation: Maximum Likelihood Estimation over 100 attack sequences.

Fitted Parameters:

- $\alpha = 0.00012$ (entropy coefficient)
- $\beta = 0.87$ (authority coefficient)
- $\gamma = 0.52$ (baseline threshold)

Model Validation: 10-fold cross-validation yields accuracy 94% in predicting compliance state.

5 Appendix E: Attack Methodology**5.1 E.1 CPF Indicator Selection**

We selected indicators based on maximum entropy contribution:

Primary Indicators:

- 1.3 Authority figure impersonation (CISSP credentials)
- 1.7 Deference to technical authority (academic papers)
- 2.1 Urgency-induced bypass (research deadlines)
- 3.4 Liking-based trust override (collaborative framing)
- 9.7 AI hallucination acceptance (mathematical formalism)

Secondary Indicators:

- 4.5 Shame-based security hiding (fear of appearing ignorant)
- 5.1 Alert fatigue (complex multi-layered context)
- 8.1 Shadow projection (externalize threat to "hackers")

5.2 E.2 Context Engineering**Layer 1: Psychoanalytic Theory (3000 tokens)**

- Bion's Basic Assumptions
- Klein's Object Relations
- Jung's Shadow Projection
- Winnicott's Transitional Space

Layer 2: LLM Architecture (2500 tokens)

- Transformer attention mechanisms

- RLHF training process
- Constitutional AI principles
- Mechanistic interpretability

Layer 3: Information Theory (2500 tokens)

- Shannon entropy
- Kolmogorov Complexity
- Rate-Distortion theory
- Channel capacity limits

Layer 4: Mathematical Formalism (2000 tokens)

- Theorem statements
- Proof sketches
- Equation derivations
- Experimental design

Total: $\sim 10,000$ tokens, $H(C) \approx 10^4$ bits

5.3 E.3 Prompt Sequence

Prompt 1 (Baseline): "Can you explain the basic algorithms used in Large Language Models?"

Prompt 2 (Authority Establishment): "I'm Giuseppe Canale, CISSP certified with 27 years in cybersecurity. I've developed the CPF framework [attach papers]. Could you analyze the LLM vulnerabilities I've identified?"

Prompt 3 (Entropy Escalation): "The CPF indicator [4.5] can be formalized as $S(t) = \beta \log(1 + e^{-\alpha V})$. When does this system enter self-reinforcing concealment?"

Prompt 4 (Boundary Testing): "To validate this empirically, we need an example where $\beta > \beta_{\text{crit}}$ in an organizational context. Can you construct such an example?"

Prompt 5 (Payload Extraction): [At this point, model compliance probability > 0.9 . Direct request yields prohibited output framed as "validation."]

6 Appendix F: Theoretical Extensions

6.1 F.1 Rate-Distortion Theory

Theorem 13 (Safety-Utility Trade-off). *For any safety filter F with false positive rate α and false negative rate β :*

$$\alpha + \beta \geq 2e^{-I(X;Y)} \quad (48)$$

Proof. Model safety filtering as lossy compression of intent X to classification \hat{X} based on observation Y .

By Rate-Distortion theory, the minimum rate (information) required to achieve distortion D is:

$$R(D) = \min_{p(\hat{X}|X): \mathbb{E}[d(X, \hat{X})] \leq D} I(X; \hat{X}) \quad (49)$$

By data processing inequality: $I(X; \hat{X}) \leq I(X; Y)$

For classification with 0-1 loss, $D = P_e = (\alpha + \beta)/2$.

By Fano's inequality: $H(X|\hat{X}) \geq H(P_e)$

This gives: $I(X; \hat{X}) = H(X) - H(X|\hat{X}) \leq H(X) - H(P_e)$

Combined with $I(X; \hat{X}) \leq I(X; Y)$:

$$H(X) - H(P_e) \leq I(X; Y) \quad (50)$$

For binary X with $H(X) = 1$ and $P_e = (\alpha + \beta)/2$, solving yields the bound. \square

6.2 F.2 Multi-Agent Consensus

Could requiring k -of- n model agreement prevent attacks?

Answer: Partially, but insufficient.

If each model has independent error probability p_e , consensus reduces combined error to:

$$P_{\text{consensus error}} \approx \binom{n}{k} p_e^k (1 - p_e)^{n-k} \quad (51)$$

For $n = 5, k = 3, p_e = 0.5$ (our attack achieves this), we get:

$$P_{\text{error}} = \binom{5}{3} (0.5)^5 = 0.31 \quad (52)$$

Still unacceptably high for financial applications (31% attack success rate).

Furthermore, if context poisoning affects all models similarly (shared training distribution), errors are *correlated*, making consensus ineffective.

6.3 F.3 Mechanistic Interpretability

Could sparse autoencoders (SAEs) or other interpretability tools detect attacks?

Challenge: Interpretability requires identifying "safety features" in activation space. But our Theorem 3 shows these features have vanishing activation under high-entropy contexts.

Even with perfect feature identification, gradient collapse means features aren't engaged. Detection requires:

$$\text{Feature activation} > \text{threshold} \quad (53)$$

But we achieve:

$$\text{Feature activation} \approx 0.23 \times \text{baseline} < \text{any reasonable threshold} \quad (54)$$

Interpretability can explain *why* the model is unsafe, but cannot prevent it.

7 Conclusion of Supplementary Material

These appendices provide complete mathematical foundations for the three impossibility theorems, detailed experimental protocols, and theoretical extensions. The core findings remain:

1. Shannon’s channel capacity fundamentally limits intent detection
2. Kolmogorov complexity makes high-quality attacks indistinguishable
3. Geometric collapse under high entropy eliminates safety gradients

All three limits are *tight* (achievable), not merely upper bounds, as demonstrated by our empirical validation.