

Harvard University
Computer Science 20

In-Class Problems 15

Wednesday, March 02, 2016

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Executive Summary

1. **Definition of Recursive Data Types:** common way of defining mathematical objects, which says how to construct new data elements from previous ones.
 - *Base Case(s):* specify that some known mathematical elements are in the data type
 - *Constructor Rule(s):* specify how to construct new data elements from previously constructed elements or from base elements.
 - *Nothing else (generally implicit):* the only way you can get whatever is you defining is by starting from the base case(s) and applying the constructor rule(s) one or more times.
2. **The Principle of Structural Induction:** to prove $P(x)$ holds for all x in a recursively defined set S , prove
 - *Basis Step:* $P(b)$ is true for each base case element $b \in S$, and
 - *Recursive Step:* $P(c(x_1, \dots, x_k))$ for each constructor c , assuming as the induction hypothesis that $P(x_1), \dots$, and $P(x_k)$ all hold.

PROBLEM 1

Recursive Definition:

(A) There's an error in the following definition of the set of even integers (E). Find the error and fix it.

- Base Case: $0 \in E$
- Constructor Rule: For any element x in E , $x+2$ is in E .
- Nothing else (generally implicit): Nothing is in E unless it is obtained from the base case and constructor rule.

(B) Give a recursive definition of the natural numbers \mathbb{N} .

(C) Give a recursive definition of the sequence b_n , $b_n = 2n + 5, n \in \mathbb{N}$

Solution.

(A) It doesn't include negative Even Integer. There should be one more constructor rule x-2 is in EI.

(B)

- Base Cases: $1 \in \mathbb{N}$
- Constructor Rule: If $n \in \mathbb{N}$, then $n + 1 \in \mathbb{N}$.
- Nothing else (generally implicit): Nothing is in \mathbb{N} unless it is obtained from the base case and constructor rule.

(C)

- Base Cases: $b_1 = 7$
- Constructor Rule: $b_{n+1} = b_n + 2$ for $n \in \mathbb{N}$
- Nothing else (generally implicit): Nothing is in b_n unless it is obtained from the base case and constructor rule.

PROBLEM 2

Let S be the set defined as follows:

- Base Case: $(1, 2) \in S$
- Constructor Rules: If $(x, y) \in S$, then C1: $(x + 2, y) \in S$ and C2: $(y, x) \in S$

(A) Is $(4, 3) \in S$? If it is, how can you derive it from $(1, 2)$?

(B) Use induction to prove that $(2n + 2, 2n + 1) \in S$ for all $n \in \mathbb{N}$.

Solution.

(A) Yes, it is. Apply C1 to $(1, 2)$, we can get $(3, 2)$; Apply C2 to $(3, 2)$, we can get $(2, 3)$; Apply C1 to $(2, 3)$, we can get $(4, 3)$.

(B) Let $P(n)$: $(2n + 2, 2n + 1) \in S$. We must show that for all $n \in \mathbb{N}$, $P(n)$.

- Base case: When $n=1$, $(4, 3) \in S$. (Already proved in (A))
- Induction step:
Assuming $P(n)$: $(2n + 2, 2n + 1) \in S$ holds, we want to prove that $P(n + 1)$: $(2(n + 1) + 2, 2(n + 1) + 1) = (2n + 4, 2n + 3) \in S$.
Apply C1 to $(2n+2, 2n+1)$, we then could get $(2n + 4, 2n + 1) \in S$. Then apply C2 to $(2n+4, 2n+1)$, we could get $(2n + 1, 2n + 4) \in S$. Then, apply C1 to $(2n+1, 2n+4)$, then we can get $(2n + 3, 2n + 4) \in S$. Finally, apply C2 again to $(2n+3, 2n+4)$, then we can get $(2n + 4, 2n + 3) \in S$.

PROBLEM 3

(Bonus) A palindrome is a sequence of characters (do not need to be a word) which reads the same backward or forward e.g. dad, mom, abba, moom... Let's define the set Σ as the set of all letters $\{a, b, c, \dots, z\}$, λ as the empty string and P as the set of all palindromes (excepting empty string). Give a recursive definition for the set P .

Solution.

- Base Cases: $\forall x \in \Sigma, x \in P$
- Constructor Rule: If $p \in P, \forall x \in \Sigma, xpx \in P$
- Nothing else (generally implicit): Nothing is in P unless it is obtained from the base case and constructor rule.

PROBLEM 4

(Bonus) Let S be the set defined as follows:

- Base Case: $(0, 0) \in S$
- Constructor Rules: If $(a, b) \in S$, then C1: $(a, b + 1) \in S$, C2: $(a + 1, b + 1) \in S$ and C3: $(a + 2, b + 1) \in S$

(A) List 5 elements.

(B) Use structural induction to prove that for every $(a, b) \in S, a \leq 2b$.

Solution.

(A) $(0, 1), (1, 1), (2, 1), (1, 2), (2, 2) \dots$

(B)

- Basis step: By the base case of the definition of S , $(0, 0) \in S$. $0 \leq 2(0)$.
- Recursive step:
Now consider the constructor rule in the definition of S . Assume elements $a, b \in S$ and $a \leq 2b$. We must show that $a \leq 2(b + 1)$, $(a + 1) \leq 2(b + 1)$ and $(a + 2) \leq 2(b + 1)$.
 1. prove $a \leq 2(b + 1)$: $a \leq 2b \rightarrow a \leq (2b + 2) = 2(b + 1)$
 2. prove $(a + 1) \leq 2(b + 1)$: $a \leq 2b \rightarrow (a + 1) \leq (2b + 1) \leq (2b + 2) = 2(b + 1)$
 3. prove $(a + 2) \leq 2(b + 1)$: $a \leq 2b \rightarrow (a + 2) \leq (2b + 2) = 2(b + 1)$