

Harvard University
Computer Science 20

In-Class Problems 10

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Executive Summary

1. Propositions and predicates

- A proposition P is like a boolean variable. Its value is either “true” or “false.”
- A predicate $P(x)$ is like a boolean-valued function. It may have the value “true” for some values of the x and the value “false” for others.
- A predicate can have more than one argument; e.g. $\text{Enrolled}(x, y)$ might mean “student x is enrolled in course y .”
- In principle it is important to know the (possibly infinite) set of values D that x and y might assume (the domain of the function).

2. Quantifiers

- The existential quantifier: $\exists x P(x)$ or $\exists x.P(x)$ or $\exists x \text{ s.t. } P(x)$ means “there exists at least one x in the domain D such that $P(x)$ is true.”
- The universal quantifier: $\forall x P(x)$ or $\forall x.P(x)$ means “for every x in the domain D , $P(x)$ is true.”

3. Multiple quantifiers

- $\exists x P(x)$ and $\forall x P(x)$ are both propositions, subject to the rules of logic that you already know.
- $\exists x P(x, y)$ and $\forall x P(x, y)$ are both predicates of the form $Q(y)$, subject to the rules of quantificational logic that you are learning.
- $\exists x \exists y P(x, y)$ and $\forall x \forall y P(x, y)$ are both propositions. The order of the quantifiers is irrelevant.
- $\exists x \forall y P(x, y)$ and $\forall y \exists x P(x, y)$ are both propositions, but they are different! The order of the quantifiers is important.

4. Negation and quantifiers

- $\neg(\exists x.P(x)) \leftrightarrow \forall x.(\neg P(x)).$
- $\neg(\forall x.P(x)) \leftrightarrow \exists x.(\neg P(x)).$

PROBLEM 1

Recall from Problem Set 1: Let A be the set of your pigeons, and let B be the set of pigeonholes in which they live. The *Generalized Pigeonhole Principle* states that for a natural number k , if $|A| > k|B|$, then there is a pigeonhole in which more than k pigeons live. Restate the GPP using quantifiers. Then negate the GPP (as you would for a proof by contradiction).

Solution.

PROBLEM 2

The domain of discourse is the set of all finite-length binary strings. The predicates $\text{Sub}(x, y)$ (meaning x is a substring of y) and $\text{Pre}(x, y)$ (meaning x is a prefix of y) are available.

- (A) Write an expression that means x consists of alternating 0s and 1s, e.g. 01010 or 101010.
- (B) Write two different expressions that mean x consists of one or more 1s, and no 0s. Important caveat: neither of your expressions may contain “0”.
- (C) (BONUS) Write two additional expressions for (c) under the same constraint.

Solution.

PROBLEM 3

(BONUS) We define a committee to be a subset of senators $S = \{s_1, s_2, \dots, s_n\}$. The predicate $M(s, C)$ means “Senator s is a member of committee C .” Rewrite the following in terms of predicate logic. You may use “=” and “ \in ” in your expressions.

- (A) Every committee has at least two senators serving on it.
- (B) No two senators serve on more than one committee together.

Solution.

PROBLEM 4

The domain of discourse is the set of integers. Let $S(x, y, z)$ mean that “ z is the sum of x and y .”

- (A) Write a formula that means x is an even integer.
- (B) Write a formula that symbolizes the commutative property for addition ($x + y = y + x$) of integers.
- (C) Write a formula that symbolizes the associative law for addition of integers:
 $x + (y + z) = (x + y) + z$.

Solution.