Metaprogramming with Macros

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Macros

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Macros

Macros realize the notion of textual abstraction.

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Macros

Macros realize the notion of textual abstraction.

Textual abstraction:

- Recognize pieces of text that match a specification
- ▶ Replace them according to a procedure

Example

```
(let (x 42) (print x))
```

Example

```
(let (x 42) (print x))
```

```
((lambda (x) (print x)) 42)
```

Step 1. Recognize pieces of text

```
(let (x 42) (print x))
(defmacro let args
```

```
((lambda (x) (print x)) 42)
```

Step 1. Recognize pieces of text

```
(let (x 42) (print x))
(defmacro let args
((lambda (x) (print x)) 42)
```

Step 2. Replace them according to a procedure

```
(let (x 42) (print x))
(defmacro let args
  (cons
   (cons 'lambda
         (cons (list (caar args))
               (cdr args)))
   (cdar args)))
((lambda (x) (print x)) 42)
```

Step 2. Replace them according to a procedure

```
(let (x 42) (print x))
(defmacro let args
  (cons
   (cons 'lambda
         (cons (list (caar args))
               (cdr args)))
   (cdar args)))
((lambda (x) (print x)) 42)
```

The essence of macros

- ▶ Recognize pieces of text that match a specification
- ▶ Replace them according to a procedure

Why macros?

- ► Deeply embedded DSLs (database access, testing)
- Optimization (programmable inlining, fusion)
- Analysis (integrated proof-checker)
- Effects (effect containment and propagation)
- **.**..

Today's talk

Macrology is vast:

- Notation
- ► Variable capture
- ► Typechecking meta-programs
- Syntax extensibility
- **.**..

Surveyed papers are versatile as well.

Today's talk

Going into all the details would be a genuine pleasure.

But instead let me tell you a story.

Outline

The prelude of macros

The tale of bindings

The trilogy of tongues

The vision of the days to come

Anaphoric if

```
(aif (calculate)
  (print it)
  (error "does not compute"))
```

Anaphoric if

```
(aif (calculate)
  (print it)
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```

```
(let* ((temp (calculate))
        (it temp))
(if temp
        (print it)
        (error "does not compute")))
```

The aif macro

```
(aif (calculate)
  (print it)
  (error "does not compute"))
(defmacro aif args
```

```
(let* ((temp (calculate))
          (it temp))
  (if temp
          (print it)
          (error "does not compute")))
```

Start with a notation

```
(aif (calculate)
 (print it)
  (error "does not compute"))
(defmacro aif args
        (let*
                          ((temp (car args))
                           (it temp))
          (if temp
             (cadr args)
             (caddr args))))
(let* ((temp (calculate))
       (it temp))
  (if temp
    (print it)
    (error "does not compute")))
```

Surround it with parentheses

```
(aif (calculate)
 (print it)
  (error "does not compute"))
(defmacro aif args
  (list 'let* (list (list 'temp (car args))
                    (list 'it 'temp))
    (list 'if 'temp
             (cadr args)
             (caddr args))))
(let* ((temp (calculate))
       (it temp))
  (if temp
    (print it)
    (error "does not compute")))
```

Quasiquote

```
(aif (calculate)
  (print it)
  (error "does not compute"))
(defmacro aif args
        (let*
                          ((temp ....)
                           (it temp))
          (if
               temp
             . . . . . . . . . . . .
             ....)))
(let* ((temp (calculate))
       (it temp))
  (if temp
    (print it)
    (error "does not compute")))
```

Quasiquote

```
(aif (calculate)
  (print it)
  (error "does not compute"))
(defmacro aif args
       '(let*
                          ((temp ....)
                           (it temp))
          (if
               temp
             . . . . . . . . . . . .
             .....)))
(let* ((temp (calculate))
       (it temp))
  (if temp
    (print it)
    (error "does not compute")))
```

Unquote

```
(aif (calculate)
  (print it)
  (error "does not compute"))
(defmacro aif args
       '(let*
                          ((temp ,(car args))
                           (it temp))
          (if temp
            ,(cadr args)
            ,(caddr args))))
(let* ((temp (calculate))
       (it temp))
  (if temp
    (print it)
    (error "does not compute")))
```

Unquote

```
(aif (calculate)
 (print it)
 (error "does not compute"))
(defmacro aif args
 '(let* ((temp ,(car args))
         (it temp))
    (if temp
        ,(cadr args)
        ,(caddr args))))
(let* ((temp (calculate))
       (it temp))
  (if temp
    (print it)
    (error "does not compute")))
```

Macro by example (MBE)

```
(aif (calculate)
 (print it)
  (error "does not compute"))
(defmacro+ aif
  (aif cond then else)
  (let* ((temp cond)
         (it temp))
    (if temp
        then
        else)))
(let* ((temp (calculate))
       (it temp))
  (if temp
    (print it)
    (error "does not compute")))
```

Interlude

- ▶ Macros are regular functions that happen to work with syntax objects
- Quasiquotes = static templates + dynamic holes

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Anaphoric if

- ► So far macros are simple: define a function, recognize pieces of text and replace them with a template
- ▶ This is so immediately useful, that we could wrap up right now

But actually

The aif macro has two bugs

```
(aif (calculate)
 (print it)
 (error "does not compute"))
(defmacro+ aif
  (aif cond then else)
 (let* ((temp cond)
         (it temp))
    (if temp then else)))
(let* ((temp (calculate))
       (it temp))
 (if temp
    (print it)
    (error "does not compute")))
```

Bug #1: Violation of hygiene

```
(let ((temp 451°F))
  (aif (calculate)
    (print it)
    (print temp)))
(defmacro+ aif
  (aif cond then else)
  (let* ((temp cond)
         (it temp))
    (if temp then else)))
(let ((temp 451°F))
  (let* ((temp (calculate))
         (it temp))
    (if temp
      (print it)
      (print temp))))
```

Bug #2: Violation of referential transparency

```
(let ((if hijacked))
  (aif (calculate)
    (print it)
    (error "does not compute")))
(defmacro+ aif
  (aif cond then else)
  (let* ((temp cond)
         (it temp))
    (if temp then else)))
(let ((if hijacked))
  (let* ((temp (calculate))
         (it temp))
    (if temp
      (print it)
      (error "does not compute"))))
```

Old school

Old school

And please don't rename core forms

Interlude

- Cross-pollination of scopes can lead to inadvertent variable capture
- ▶ Violation of hygiene = def site harms call site
- ▶ Violation of referential transparency = call site harms def site

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Template Haskell

```
aif :: Q Exp -> Q Exp -> Q Exp aif cond then' else' =
  [| let temp = $cond
                it = temp
    in if temp /= O then $then' else $else' |]
```

Template Haskell

Template Haskell

```
$(aif [| calculate |]
  [| putStrLn (show it) |]
  [| error "does not compute" |])
aif :: Q Exp -> Q Exp -> Q Exp -> Q Exp
aif cond then' else' =
  [| let temp = $cond
         it = temp
     in if temp /= 0 then $then' else $else' |]
let temp_a1mx = calculate
    it_a1my = temp_a1mx
in if (temp_a1mx /= 0)
   then putStrLn (show it)
   else error "does not compute"
```

Not in scope: 'it'

The Q monad

```
aif cond then' else' =
  [| let temp = $cond
          it = temp
  in if temp /= 0 then $then' else $else' |]
```

```
The Q monad
aif cond then' else' =
  [| let temp = $cond
         it = temp
     in if temp /= 0 then $then' else $else' |]
aif :: Q Exp -> Q Exp -> Q Exp -> Q Exp
aif cond' then', else' =
    do { temp <- newName "temp"
        : it <- newName "it"
        : cond <- cond'
        : then' <- then',
        : else' <- else''
        ; let notEq = mkNameG_v "ghc-prim" "GHC.Classes" "/="
         in return
           (LetE [ValD (VarP temp) (NormalB cond) [],
                 ValD (VarP it) (NormalB (VarE temp)) []]
                 (CondE (... (VarE notEq) ...) then' else'))
```

Perils of hygiene

```
$(aif [| calculate |]
  [| putStrLn (show $(dyn "it")) |]
  [| error "does not compute" |])

aif :: Q Exp -> Q Exp -> Q Exp -> Q Exp
aif cond then' else' =
  [| let temp = $cond
         it = temp
  in if temp /= O then $then' else $else' |]
```

Summary: Template Haskell

- Quasiquotes in Template Haskell are automatically hygienic and referentially transparent
- ► That's because are translated into Q monad, which takes care of gensymming and fully qualified names
- ▶ When hygiene fails, we are forced to drop to low level

Nemerle

```
macro aif(cond, then, else_) {
    <[
        def temp = $cond;
        def it = temp;
        if (temp != 0) $then else $else_
     ]>
}
```

Nemerle

```
aif(calculate,
  WriteLine(it),
  throw Exception("does not compute"))
macro aif(cond, then, else_) {
  <[
    def temp = $cond;
    def it = temp;
    if (temp != 0) $then else $else_
  ]>
```

Nemerle

error: unbound name 'it'

```
aif(calculate,
  WriteLine(it),
  throw Exception("does not compute"))
macro aif(cond, then, else_) {
  ۲۲
    def temp = $cond;
    def it = temp;
    if (temp != 0) $then else $else_
  1>
def calculate = 42;
def temp_1087 = calculate;
def it_1088 = temp_1087;
if (temp_1087 != 0) WriteLine(it) else throw Exception("...")
```

Coloring algorithm

```
def calculate = 42;
aif(calculate,
  WriteLine(it).
  throw Exception("does not compute"))
macro aif(cond, then, else_) {
  <۲
    def temp = $cond;
    def it = temp;
    if (temp != 0) $then else $else_
 ]>
def calculate = 42:
def temp = calculate;
def it = temp;
if (temp != 0) WriteLine(it) else throw Exception("...")
```

Coloring algorithm: normal code gets a vanilla color

```
def calculate = 42;
                                     // vanilla color
aif(calculate,
  WriteLine(it),
  throw Exception("does not compute"))
macro aif(cond, then, else_) {
  <۲
    def temp = $cond;
    def it = temp;
    if (temp != 0) $then else $else_
 ]>
def calculate = 42:
def temp = calculate;
def it = temp;
if (temp != 0) WriteLine(it) else throw Exception("...")
```

Coloring algorithm: each expansion gets unique colors

```
def calculate = 42;
                                    // vanilla color
aif(calculate,
  WriteLine(it),
  throw Exception("does not compute"))
macro aif(cond, then, else_) { // expansion color
  <۲
    def temp = $cond;
    def it = temp;
    if (temp != 0) $then else $else_
 ]>
def calculate = 42:
def temp = calculate;
def it = temp;
if (temp != 0) WriteLine(it) else throw Exception("...")
```

Coloring algorithm: at the end of the day

```
def calculate = 42;
                                    // vanilla color
aif(calculate,
  WriteLine(it),
  throw Exception("does not compute"))
macro aif(cond, then, else_) { // expansion color
  <۲
    def temp = $cond;
    def it = temp;
    if (temp != 0) $then else $else_
 ]>
def calculate = 42:
                                    // bind using colors
def temp = calculate;
def it = temp;
if (temp != 0) WriteLine(it) else throw Exception("...")
```

Coloring algorithm: inherit use site

```
def calculate = 42;
                                    // vanilla color
aif(calculate,
  WriteLine(it),
  throw Exception("does not compute"))
macro aif(cond, then, else_) { // expansion color
  <۲
    def temp = $cond;
    def $("it": usesite) = temp;
    if (temp != 0) $then else $else_
 ]>
def calculate = 42:
                                    // bind using colors
def temp = calculate;
def it = temp;
if (temp != 0) WriteLine(it) else throw Exception("...")
```

Coloring algorithm: polychromatic

```
def calculate = 42;
                                    // vanilla color
aif(calculate,
  WriteLine(it),
  throw Exception("does not compute"))
macro aif(cond, then, else_) { // expansion color
  <۲
    def temp = $cond;
    def $("it": dyn) = temp;
    if (temp != 0) $then else $else_
 ]>
def calculate = 42:
                                    // bind using colors
def temp = calculate;
def it = temp;
if (temp != 0) WriteLine(it) else throw Exception("...")
```

Summary: Nemerle

- ► Nemerle takes care of hygiene with a coloring algorithm of impressive simplicity and power
- ▶ No complex translation algorithms are necessary
- ► As another bonus programmer can fine-tune colors with MacroColors
- Referential transparency works as well

A Lisp, descendent from Scheme
25 years of hygienic macros, a bunch of macro systems
Language features written using macros (classes, modules, etc)
Q: "How to turn macros into proper abstractions?"

```
(aif (calculate)
  (print it)
  (error "does not compute"))
(define-syntax (aif stx)
  (syntax-case stx ()
    ((aif cond then else)
       #'(let ((temp cond)
               (it temp)))
           (if temp then else)))))
```

```
(aif (calculate)
 (print it)
 (error "does not compute"))
(define-syntax (aif stx)
  (syntax-case stx ()
    ((aif cond then else)
       #'(let ((temp cond)
               (it temp)))
           (if temp then else)))))
(let* ((temp (calculate))
       (it temp))
  (if temp
    (print it)
    (error "does not compute")))
```

```
(aif (calculate)
 (print it)
  (error "does not compute"))
(define-syntax (aif stx)
  (syntax-case stx ()
    ((aif cond then else)
     (with-syntax ((it (datum->syntax #'aif 'it)))
       #'(let ((temp cond)
               (it temp)))
           (if temp then else))))))
(let* ((temp (calculate))
       (it temp))
  (if temp
    (print it)
    (error "does not compute")))
```

```
(define-syntax (aunless stx)
  (syntax-case stx ()
      ((aunless cond then else)
      #'(aif (not cond) then else))))
```

```
(aunless (not (calculate))
  (print it)
  (error "does not compute"))

(define-syntax (aunless stx)
  (syntax-case stx ()
        ((aunless cond then else)
        #'(aif (not cond) then else))))
```

```
(aunless (not (calculate))
  (print it)
  (error "does not compute"))
(define-syntax (aunless stx)
  (syntax-case stx ()
    ((aunless cond then else)
    #'(aif (not cond) then else))))
(let* ((temp (not (not (calculate))))
       (it temp))
  (if temp
    (print it)
    (error "does not compute")))
```

```
(aunless (not (calculate))
  (print it)
  (error "does not compute"))
(define-syntax (aunless stx)
  (syntax-case stx ()
    ((aunless cond then else)
     #'(aif (not cond) then else))))
(let* ((temp (not (not (calculate))))
       (it temp))
 (if temp
    (print it)
    (error "does not compute")))
```

```
(aunless (not (calculate))
  (print it)
  (error "does not compute"))
(define-syntax (aunless stx)
  (syntax-case stx ()
    ((aunless cond then else)
    #'(aif (not cond) then else))))
(let* ((temp (not (not (calculate))))
       (it temp))
 (if temp
    (print it)
    (error "does not compute")))
```

Solution: dynamic variables

Solution: dynamic variables

Summary: overall

- ► There are algorithms that take care of hygiene and referential transparency
- ► These algorithms can work in automatic mode, but are flexible enough to give the programmer full control
- ► Like a silver bullet
- ▶ Nevertheless sometimes even better solutions come from integration with language features

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scalamacros.org

- ► Since this semester Scala has macros
- ► Even better: macros are an official part of the language in the next production release 2.10.0
- Now it's time to put the pens down and think about the future

 ${\tt def \ serialize[T](x: \ T): \ Pickle}$

```
trait Serializer[T] {
  def write(pickle: Pickle, x: T): Unit
}
def serialize[T](x: T)(s: Serializer[T]): Pickle
```

```
trait Serializer[T] {
  def write(pickle: Pickle, x: T): Unit
}
def serialize[T](x: T)(implicit s: Serializer[T]): Pickle
```

```
trait Serializer[T] {
  def write(pickle: Pickle, x: T): Unit
}

def serialize[T](x: T)(implicit s: Serializer[T]): Pickle

implicit object ByteSerializer extends Serializer[Byte] {
  def write(pickle: Pickle, x: Byte) = pickle.writeByte(x)
}
```

```
trait Serializer[T] {
  def write(pickle: Pickle, x: T): Unit
}
def serialize[T](x: T)(implicit s: Serializer[T]): Pickle
implicit def generator: Serializer[T] = macro impl[T]
```

Research proposal

- ► Macros + functions = programmable inlining, specialization, fusion
- ► Macros + annotations = code contracts
- ► Macros + path-dependent types = controlled effects
- ► Macros + implicits = theorem prover
- **.**..