# Metaprogramming with Macros

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### Macros

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#### **Macros**

Macros realize the notion of textual abstraction.

#### Textual abstraction:

- Recognize pieces of text that match a specification
- ► Replace them according to a procedure

## Why macros?

- Deeply embedded DSLs (database access, testing)
- Optimization (programmable inlining, fusion)
- Analysis (integrated proof-checker)
- Effects (effect containment and propagation)
- **.**..

#### Outline

#### The prelude of macros

The tale of bindings

The trilogy of tongues

The vision of the days to come

# if in Lisp

```
(if (calculate)
  (print "success")
  (error "does not compute"))
```

### Anaphoric if

```
(aif (calculate)
  (print it)
  (error "does not compute"))
```

```
(let* ((temp (calculate))
          (it temp))
  (if temp
          (print it)
          (error "does not compute")))
```

#### The aif macro

```
(aif (calculate)
  (print it)
  (error "does not compute"))
(defmacro aif args
```

```
(let* ((temp (calculate))
          (it temp))
  (if temp
          (print it)
          (error "does not compute")))
```

### Low-level implementation

```
(aif (calculate)
 (print it)
  (error "does not compute"))
(defmacro aif args
  (list 'let* (list (list 'temp (car args))
                     (list 'it 'temp))
     (list 'if 'temp
       (cadr args)
       (caddr args))))
(let* ((temp (calculate))
       (it temp))
  (if temp
    (print it)
    (error "does not compute")))
```

## Quasiquoting: static template

```
(aif (calculate)
 (print it)
  (error "does not compute"))
(defmacro aif args
 '(let* ((temp .....)
          (it temp))
     (if temp
       . . . . . . . . . . . .
       .....)
(let* ((temp (calculate))
       (it temp))
  (if temp
    (print it)
    (error "does not compute")))
```

## Quasiquoting: dynamic holes

```
(aif (calculate)
 (print it)
  (error "does not compute"))
(defmacro aif args
  '(let* ((temp ,(car args))
          (it temp))
     (if temp
       ,(cadr args)
       (caddr args))
(let* ((temp (calculate))
       (it temp))
  (if temp
    (print it)
    (error "does not compute")))
```

# Macro by example (MBE)

```
(aif (calculate)
 (print it)
  (error "does not compute"))
(defmacro+ aif
  (aif cond then else)
  (let* ((temp cond)
         (it temp))
    (if temp
        then
        else)))
(let* ((temp (calculate))
       (it temp))
  (if temp
    (print it)
    (error "does not compute")))
```

#### Interlude

- ▶ Macros are regular functions that happen to work with syntax objects
- ► Quasiquotes = static templates + dynamic holes

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### Anaphoric if

- ► So far macros are simple: define a function, recognize pieces of text and replace them with a template
- ▶ This is so immediately useful, that we could wrap up right now

# But actually

The aif macro has two bugs

## What's wrong?

## What's wrong?

### What's wrong?

```
(aif (calculate)
 (print it)
 (error "does not compute"))
(defmacro+ aif
  (aif cond then else)
 (let* ((temp cond)
         (it temp))
    (if temp then else)))
(let* ((temp (calculate))
       (it temp))
 (if temp
    (print it)
    (error "does not compute")))
```

## Bug #1: Violation of hygiene

```
(let ((temp 451°F))
  (aif (calculate)
    (print it)
    (print temp)))
(defmacro+ aif
  (aif cond then else)
  (let* ((temp cond)
         (it temp))
    (if temp then else)))
(let ((temp 451°F))
  (let* ((temp (calculate))
         (it temp))
    (if temp
      (print it)
      (print temp))))
```

## Bug #2: Violation of referential transparency

```
(let ((if hijacked))
  (aif (calculate)
    (print it)
    (error "does not compute")))
(defmacro+ aif
  (aif cond then else)
  (let* ((temp cond)
         (it temp))
    (if temp then else)))
(let ((if hijacked))
  (let* ((temp (calculate))
         (it temp))
    (if temp
      (print it)
      (error "does not compute"))))
```

#### Interlude

- ► Cross-pollination of scopes can lead to inadvertent variable capture
- ▶ Violation of hygiene = def site harms call site
- ▶ Violation of referential transparency = call site harms def site

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### Template Haskell

```
$(aif [| calculate |]
  [| putStrLn (show it) |]
  [| error "does not compute" |])

aif :: Q Exp -> Q Exp -> Q Exp -> Q Exp
aif cond then' else' =
  [| let temp = $cond
         it = temp
  in if temp /= O then $then' else $else' |]
```

## Template Haskell

Not in scope: 'it'

```
$(aif [| calculate |]
  [| putStrLn (show it) |]
  [| error "does not compute" |])
aif :: Q Exp -> Q Exp -> Q Exp -> Q Exp
aif cond then' else' =
  [| let temp = $cond
         it = temp
     in if temp /= 0 then $then' else $else' |]
let temp_a1mx = calculate
    it_a1my = temp_a1mx
in if (temp_a1mx /= 0)
   then putStrLn (show it)
   else error "does not compute"
```

### The Q monad

```
aif cond then' else' =
  [| let temp = $cond
          it = temp
    in if temp /= 0 then $then' else $else' |]
```

```
The Q monad
aif cond then' else' =
  [| let temp = $cond
         it = temp
     in if temp /= 0 then $then' else $else' |]
aif :: Q Exp -> Q Exp -> Q Exp
aif cond' then', else' =
    do { temp <- newName "temp"
       : it <- newName "it"
       : cond <- cond'
       : then' <- then',
       : else' <- else''
       ; let notEq = mkNameG_v "ghc-prim" "GHC.Classes" "/="
         in return
          (LetE [ValD (VarP temp) (NormalB cond) [],
                 ValD (VarP it) (NormalB (VarE temp)) []]
                (CondE (... (VarE notEq) ...) then' else'))
```

## Breaking hygiene

```
$(aif [| calculate |]
  [| putStrLn (show $(dyn "it")) |]
  [| error "does not compute" |])
aif :: Q Exp -> Q Exp -> Q Exp -> Q Exp
aif cond then' else' =
  [ let temp = $cond
         it = temp
     in if temp /= 0 then $then' else $else' []
let temp_a1mx = calculate
    it_a1my = temp_a1mx
in if (temp_a1mx /= 0)
   then putStrLn (show it_a1my)
   else error "does not compute"
```

## Summary: Template Haskell

- Quasiquotes in Template Haskell are automatically hygienic and referentially transparent
- ► That's because they are lifted into the Q monad, which takes care of names (fresh names for locals, fully qualified names for globals)
- ▶ Sometimes we need to break hygiene

#### Nemerle

```
aif(calculate,
  WriteLine(it),
  throw Exception("does not compute"))
macro aif(cond, then, else_) {
  <[
    def temp = $cond;
    def it = temp;
    if (temp != 0) $then else $else_
  ]>
```

#### Nemerle

```
aif(calculate,
  WriteLine(it),
  throw Exception("does not compute"))
macro aif(cond, then, else_) {
  ۲۲
    def temp = $cond;
    def it = temp;
    if (temp != 0) $then else $else_
  1>
def calculate = 42;
def temp_1087 = calculate;
def it_1088 = temp_1087;
if (temp_1087 != 0) WriteLine(it) else throw Exception("...")
```

## Coloring algorithm

```
def calculate = 42;
aif(calculate,
  WriteLine(it),
  throw Exception("does not compute"))
macro aif(cond, then, else_) {
 <[
    def temp = $cond;
    def it = temp;
    if (temp != 0) $then else $else_
 ]>
def calculate = 42:
def temp = calculate;
def it = temp;
if (temp != 0) WriteLine(it) else throw Exception("...")
```

# Coloring algorithm: normal code gets a vanilla color

```
def calculate = 42;
                                     // vanilla color
aif(calculate,
  WriteLine(it),
  throw Exception("does not compute"))
macro aif(cond, then, else_) {
  <۲
    def temp = $cond;
    def it = temp;
    if (temp != 0) $then else $else_
 ]>
def calculate = 42:
def temp = calculate;
def it = temp;
if (temp != 0) WriteLine(it) else throw Exception("...")
```

## Coloring algorithm: each expansion gets unique colors

```
def calculate = 42;
                                    // vanilla color
aif(calculate,
  WriteLine(it),
  throw Exception("does not compute"))
macro aif(cond, then, else_) { // expansion color
  <۲
    def temp = $cond;
    def it = temp;
    if (temp != 0) $then else $else_
 ]>
def calculate = 42:
def temp = calculate;
def it = temp;
if (temp != 0) WriteLine(it) else throw Exception("...")
```

# Coloring algorithm: at the end of the day

```
def calculate = 42;
                                    // vanilla color
aif(calculate,
  WriteLine(it),
  throw Exception("does not compute"))
macro aif(cond, then, else_) { // expansion color
  <۲
    def temp = $cond;
    def it = temp;
    if (temp != 0) $then else $else_
 ]>
def calculate = 42:
                                    // bind using colors
def temp = calculate;
def it = temp;
if (temp != 0) WriteLine(it) else throw Exception("...")
```

## Coloring algorithm: inherit use site

```
def calculate = 42;
                                    // vanilla color
aif(calculate,
  WriteLine(it),
  throw Exception("does not compute"))
macro aif(cond, then, else_) { // expansion color
  <۲
    def temp = $cond;
    def $("it": usesite) = temp;
    if (temp != 0) $then else $else_
 ]>
def calculate = 42:
                                    // bind using colors
def temp = calculate;
def it = temp;
if (temp != 0) WriteLine(it) else throw Exception("...")
```

# Summary: Nemerle

- ► Nemerle takes care of hygiene with a coloring algorithm of impressive simplicity and power
- No complex translation algorithms are necessary
- ► As another bonus programmer can fine-tune colors with MacroColors
- Referential transparency works as well

A Lisp, descendent from Scheme
25 years of hygienic macros, a bunch of macro systems
Language features written using macros (classes, modules, etc)
Q: "How to turn macros into proper abstractions?"

```
(aif (calculate)
  (print it)
  (error "does not compute"))
(define-syntax (aif stx)
  (syntax-case stx ()
    ((aif cond then else)
       #'(let ((temp cond)
               (it temp)))
           (if temp then else)))))
```

```
(aif (calculate)
 (print it)
 (error "does not compute"))
(define-syntax (aif stx)
  (syntax-case stx ()
    ((aif cond then else)
       #'(let ((temp cond)
               (it temp)))
           (if temp then else)))))
(let* ((temp (calculate))
       (it temp))
  (if temp
    (print it)
    (error "does not compute")))
```

```
(aif (calculate)
 (print it)
 (error "does not compute"))
(define-syntax (aif stx)
  (syntax-case stx ()
    ((aif cond then else)
     (with-syntax ((it (datum->syntax #'aif 'it)))
       #'(let ((temp cond)
               (it temp)))
           (if temp then else))))))
(let* ((temp (calculate))
       (it temp))
  (if temp
    (print it)
    (error "does not compute")))
```

### Doesn't scale

```
(aunless (not (calculate))
  (print it)
  (error "does not compute"))
(define-syntax (aunless stx)
  (syntax-case stx ()
    ((aunless cond then else)
    #'(aif (not cond) then else))))
(let* ((temp (not (not (calculate))))
       (it temp))
  (if temp
    (print it)
    (error "does not compute")))
```

### Doesn't scale

```
(aunless (not (calculate))
  (print it)
  (error "does not compute"))
(define-syntax (aunless stx)
  (syntax-case stx ()
    ((aunless cond then else)
     #'(aif (not cond) then else))))
(let* ((temp (not (not (calculate))))
       (it temp))
 (if temp
    (print it)
    (error "does not compute")))
```

#### Doesn't scale

```
(aunless (not (calculate))
  (print it)
  (error "does not compute"))
(define-syntax (aunless stx)
  (syntax-case stx ()
    ((aunless cond then else)
    #'(aif (not cond) then else))))
(let* ((temp (not (not (calculate))))
       (it temp))
 (if temp
    (print it)
    (error "does not compute")))
```

# Solution: dynamic variables

# Summary: overall

- ► There are algorithms that take care of hygiene and referential transparency
- ► These algorithms can work in automatic mode, but are flexible enough to give the programmer full control
- ► Like a silver bullet
- ▶ Nevertheless sometimes even better solutions come from integration with language features

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### scalamacros.org

- ► Since this semester Scala has macros
- ► Even better: macros are an official part of the language in the next production release 2.10.0
- Now it's time to put the pens down and think about the future

def serialize[T](x: T): Pickle

```
trait Serializer[T] {
  def write(pickle: Pickle, x: T): Unit
}
def serialize[T](x: T)(s: Serializer[T]): Pickle
```

```
trait Serializer[T] {
  def write(pickle: Pickle, x: T): Unit
}

def serialize[T](x: T)(implicit s: Serializer[T]): Pickle

implicit object ByteSerializer extends Serializer[Byte] {
  def write(pickle: Pickle, x: Byte) = pickle.writeByte(x)
}
```

```
trait Serializer[T] {
  def write(pickle: Pickle, x: T): Unit
}
def serialize[T](x: T)(implicit s: Serializer[T]): Pickle
implicit def generator: Serializer[T] = macro impl[T]
```

## Research proposal

### Marry macros and high-level language features:

- ► Macros + functions = programmable inlining, specialization, fusion
- ► Macros + annotations = code contracts
- ► Macros + path-dependent types = controlled effects
- ► Macros + implicits = static verification
- •