区间问题

笔记本: 0_leetcode

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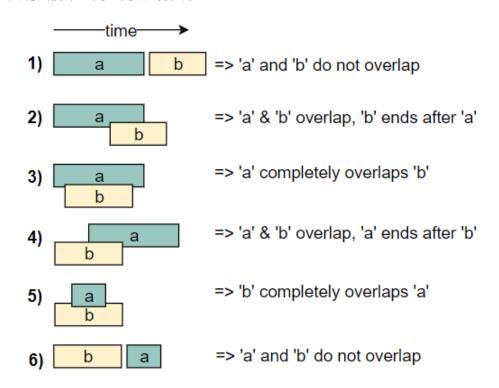
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区间问题

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0. 介绍

首先要明白俩个区间之间的6种关系。



1. 实战

1.1 合并区间

给定一系列区间,将相交的区间合并成一个大的区间。区间之间没有顺序。

e.g.1

```
Intervals: [[1,4], [2,5], [7,9]]
Output: [[1,5], [7,9]]
Explanation: Since the first two intervals [1,4] and [2,5] overlap, we merged them into one [1,5].
```

```
Intervals: [[6,7], [2,4], [5,9]]
Output: [[2,4], [5,9]]
```

Explanation: Since the intervals [6,7] and [5,9] overlap, we merged them into one [5,9].

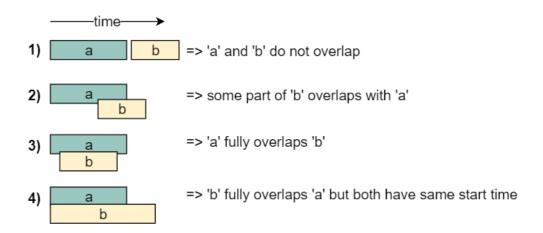
e.g.3

```
Intervals: [[1,4], [2,6], [3,5]]
Output: [[1,6]]
Explanation: Since all the given intervals overlap, we merged them into one.
```

思路

- 1. 区间是无序的,第一想到的是排序。很明显,应该直接使用区间start作为比较值。
- 2. 判断相交。因为排序了,所以保证a.start <= b.start。

那么只会出现如下四种情况。



判断相交就是b.start <= a.end。合并后的区间就是c

c.start = a.start

c.end = max(a.end, b.end)

3. 很明显,类似双指针,我们应该持有一个指向当前区间的interval,和一个指向下一个区间的指针。

如果相交,当前interval被merged。当区间不相交时,更新当前区间,并将之前的区间放到结果集合中。

```
using namespace std;

#include <algorithm>
#include <iostream>
#include <vector>

class Interval {
```

```
int start = 0;
    int end = 0;
    Interval( int start, int end ) {
        this->start = start;
        this->end = end;
class MergeIntervals {
    static vector<Interval> merge( vector<Interval> &intervals ) {
        vector<Interval> mergedIntervals;
        sort(intervals.begin(), intervals.end(), [](Interval &il, Interval
&i2 ) {
            if (i1. start == i2. start)
                return i1. end > i2. end;
            return i1. start <= i2. start;
        } );
        Interval curr interval = intervals[0];
        Interval *next interval;
        for ( int i = 1; i < intervals.size(); i++) {
            next interval = &intervals[i];
            if ( next interval->start <= curr interval.end ) {</pre>
                curr interval.end = max( curr interval.end, next interval-
>end);
                mergedIntervals.push back( curr interval );
                curr interval = *next interval;
        mergedIntervals.push back( curr interval );
        return mergedIntervals;
int main( int argc, char *argv[] ) {
    vector\langle Interval \rangle input = {{1, 3}, {2, 5}, {7, 9}};
    for ( auto interval : MergeIntervals::merge( input ) ) {
        cout << "[" << interval.start << "," << interval.end << "] ";</pre>
```

```
cout << endl;
input = {{6, 7}, {2, 4}, {5, 9}};
cout << "Merged intervals: ";
for ( auto interval : MergeIntervals::merge( input ) ) {
      cout << "[" << interval. start << "," << interval. end << "] ";
}
cout << endl;
input = {{1, 4}, {2, 6}, {3, 5}};
cout << "Merged intervals: ";
for ( auto interval : MergeIntervals::merge( input ) ) {
      cout << "[" << interval. start << "," << interval. end << "] ";
}
cout << endl;
system( "pause" );
}</pre>
```

1.2 插入区间

给定一系列不相交的区间,且使用start排序。插入一个新的区间,求合并后的区间。

e.g.1

```
Input: Intervals=[[1,3], [5,7], [8,12]], New Interval=[4,6]
Output: [[1,3], [4,7], [8,12]]
Explanation: After insertion, since [4,6] overlaps with [5,7], we merged them into one [4,7].
```

```
Input: Intervals=[[1,3], [5,7], [8,12]], New Interval=[4,10]
Output: [[1,3], [4,12]]
Explanation: After insertion, since [4,10] overlaps with [5,7] & [8,12], we merged them into [4,12]. Example 3:
```

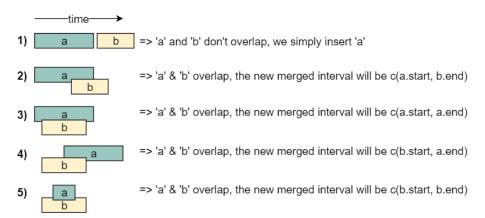
```
Input: Intervals=[[2,3],[5,7]], New Interval=[1,4]
Output: [[1,4], [5,7]]
Explanation: After insertion, since [1,4] overlaps with [2,3], we merged them into one [1,4].
```

思路

区间已经排好序,我们只需要将新的区间插入到即可。自己可以构造几种情况来模拟,就知道如何插入了。

1. 如何跳过区间。**很明显,新区间出现在intervals[i]右边,且不相交。就可以跳过。 intervals[i].end** < **newInterval.start**。明白这个条件是这道题的关键。

直接画图。一目了然。



a为newInterval

```
if(intervals[i].end < newInterval.start)
skip;
else{
    可能发生相交assert(newInterval.start <= intervals.end)
    if(newInterval.end < intervals[i].end)
    不相交
else
    相交
    curr_interval.start = max(...)
    curr_interval.end = max(...)
}
```

```
#include <algorithm>
class Interval {
    int start = 0;
    int end = 0;
    Interval() = default;
    Interval( int start, int end ) {
        this->start = start;
        this->end = end;
class InsertInterval {
    static vector (Interval) insert (const vector (Interval) & intervals,
Interval newInterval ) {
        vector<Interval> mergedIntervals;
        Interval curr interval;
        int record = 0;
        for ( int i = 0; i < intervals.size(); i++) {
            if ( intervals[i].end < newInterval.start ) {</pre>
                mergedIntervals.push back( intervals[i] );
                // 五种情况,实际上就是相交不相交俩种情况
                if ( newInterval.end < intervals[i].start ) {</pre>
                    mergedIntervals.push_back( newInterval );
                    for ( int k = i; k < intervals. size(); k++)
                        mergedIntervals.push back( intervals[k] );
                    return mergedIntervals;
                    curr interval.start = min( intervals[i].start,
newInterval.start);
                    curr_interval.end = max( intervals[i].end,
newInterval.end );
                    record = i;
```

```
const Interval *next interval;
        for ( int i = record + 1; i < intervals.size(); i++ ) {
            next interval = &intervals[i];
            if ( next interval->start >= curr interval.start &&
next interval->start <= curr interval.end ) {</pre>
                curr interval.end = max( curr interval.end, next interval-
>end );
                mergedIntervals.push back( curr interval );
                curr interval = *next interval;
        mergedIntervals.push back( curr interval );
        return mergedIntervals;
int main( int argc, char *argv[] ) {
    vector\langle Interval \rangle input = \{\{1, 3\}, \{5, 7\}, \{8, 12\}\}\};
    cout << "Intervals after inserting the new interval: ";</pre>
    for (auto interval : InsertInterval::insert(input, {4, 6})) {
        cout << "[" << interval.start << "," << interval.end << "] ";</pre>
    cout << "Intervals after inserting the new interval: ";</pre>
    for (auto interval::InsertInterval::insert(input, {4, 10})) {
        cout << "[" << interval.start << "," << interval.end << "] ";</pre>
    input = \{\{2, 3\}, \{5, 7\}\};
    cout << "Intervals after inserting the new interval: ";</pre>
    for (auto interval::InsertInterval::insert(input, {1, 4})) {
        cout << "[" << interval.start << "," << interval.end << "] ";</pre>
    input = \{\{2, 3\}, \{6, 7\}\};
```

```
cout << "Intervals after inserting the new interval: ";
for ( auto interval : InsertInterval::insert( input, {4, 5} ) ) {
    cout << "[" << interval.start << "," << interval.end << "] ";
}
cout << endl;
system("pause");
}</pre>
```

1.3 合并俩个区间链表 (经历过的面试题)

合并俩个有序的区间数组, 找交集

e.g.1

```
Input: arr1=[[1, 3], [5, 6], [7, 9]], arr2=[[2, 3], [5, 7]]
Output: [2, 3], [5, 6], [7, 7]
Explanation: The output list contains the common intervals between the two lists.
```

e.g.2

```
Input: arr1=[[1, 3], [5, 7], [9, 12]], arr2=[[5, 10]]
Output: [5, 7], [9, 10]
Explanation: The output list contains the common intervals between the two lists.
```

思路

这里有点像二路归并。 因为已经按照start排好序。

所以思路如下

- 1. 比较俩个数组的第一个区间。是否相交。
- 2. 如果相交。合并成一个大的区间。
- 3. 然后将相交的的区间插入即可。

```
using namespace std;

#include <iostream>
```

```
#include <algorithm>
class Interval {
    int start = 0;
    int end = 0;
    Interval ( int start = 0, int end = 0 ) {
        this->start = start;
        this->end = end;
   bool operator( const Interval &rhs )const {
        if ( start == rhs. start )
            return end < rhs. end;
        return start < rhs. start;</pre>
class IntervalsIntersection {
    static vector(Interval) merge (const vector(Interval) & arrl, const
vector<Interval> &arr2 ) {
        vector<Interval> result;
        for ( ; x < arr1. size() && y < arr2. size(); ) {</pre>
            if (arr1[x].start >= arr2[y].start && arr1[x].start <=
arr2[y].end ||
                    arr2[x].start >= arr1[y].start && arr2[x].start <=
arr1[y].end ) {
                Interval cur interval;
                cur interval.start = max( arr1[x].start, arr2[y].start );
                cur interval.end = min( arr1[x].end, arr2[y].end );
                result.push back( cur interval );
            if (arr1[x] < arr2[y]) {
        return result;
```

```
}
};
int main( int argc, char *argv[] ) {
    vector<Interval> input1 = {{1, 3}, {5, 6}, {7, 9}};
    vector<Interval> input2 = {{2, 3}, {5, 7}};
    vector<Interval> result = IntervalsIntersection::merge( input1, input2 );
    cout << "Intervals Intersection: ";
    for ( auto interval : result ) {
        cout << "[" << interval. start << "," << interval. end << "] ";
    }
    cout << endl;

input1 = {{1, 3}, {5, 7}, {9, 12}};
    input2 = {{5, 10}};
    result = IntervalsIntersection::merge( input1, input2 );
    cout << "Intervals Intersection: ";
    for ( auto interval : result ) {
        cout << "[" << iinterval. start << "," << iinterval. end << "] ";
    }
    system( "pause" );
}</pre>
```

1.4 合并俩个区间链表 (经历过的面试题)

合并俩个有序的区间数组, 找并集

e.g.1

```
Input: arr1=[[1, 3], [5, 6], [7, 9]], arr2=[[2, 3], [5, 7]]
Output: [1, 3], [5, 9]
Explanation: The output list contains the common intervals between the two lists.
```

```
Input: arr1=[[1, 3], [5, 7], [9, 12]], arr2=[[5, 10]]
Output: [1, 3], [5, 12]
Explanation: The output list contains the common intervals between the two lists.
```

思路

这里有点像二路归并。 因为已经按照start排好序。 所以思路如下

- 1. 比较俩个数组的第一个区间。是否相交。
- 2. 如果相交。合并成一个大的区间。

这里的思路是二路归并。因此需要重载<,来定义区间比较。即获得较小的区间。 最后需要处理如果余下一路的问题。

```
#include <algorithm>
class Interval {
    int start = 0;
    int end = 0;
    Interval( int start = 0, int end = 0 ) {
        this->start = start;
        this->end = end;
   bool operator ( const Interval &rhs ) const {
        if ( start == rhs. start )
            return end < rhs. end;
       return start < rhs. start;</pre>
class IntervalsIntersection {
    static vector (Interval) merge (const vector (Interval) & arrl, const
vector<Interval> &arr2 ) {
        vector<Interval> result;
        Interval curr_interval( 0, 0 );
        if (arr1[0] < arr2[0]) {
```

```
curr interval = arr1[0];
            curr interval = arr2[0];
        for (; x < arr1.size() && y < arr2.size(); ) {
            Interval to compare;
            if ( arr1[x] < arr2[y] ) {</pre>
                to compare = arr1[x];
                to_compare = arr2[y];
            if ( to compare.start <= curr interval.end ) {</pre>
                curr interval.end = std::max( to compare.end,
curr interval.end );
                result.push back( curr interval );
                curr_interval = to_compare;
        for (; x != arr1. size(); x++) {
            if ( arr1[x].start <= curr interval.end ) {</pre>
                curr interval.end = std::max( arr1[x].end, curr interval.end
        for (; y != arr2. size(); y++) {
            if ( arr2[y].start <= curr interval.end ) {</pre>
                curr interval.end = std::max(arr2[y].end, curr interval.end
        result.push back( curr interval );
        for (; x != arr1. size(); x++) {
            result.push back(arr1[x]);
        for (; y != arr2. size(); y++) {
            result.push_back( arr2[y] );
```

```
return result;
}
};
```

1-5 冲突的会议

给定一系列区间,代表会议时间。判断一个人是否能参加所有的会议。

e.g.1

```
Appointments: [[1,4], [2,5], [7,9]]
Output: false
Explanation: Since [1,4] and [2,5] overlap, a person cannot attend both of these appointments.
```

e.g.2

```
Appointments: [[6,7], [2,4], [8,12]]
Output: true
Explanation: None of the appointments overlap, therefore a person can attend all of them.
```

e.g.3

```
Appointments: [[4,5], [2,3], [3,6]]
Output: false
Explanation: Since [4,5] and [3,6] overlap, a person cannot attend both of these appointments.
```

思路

俩个思路

- 1. 按照start排序,然后遍历查看是否存在冲突。时间复杂度是O(logN)
- 2. 如果将所有时间点排在一个时间线上,可以发现,同一时间,只能经历一个start。

这里实现思路2。但是思路2存在局限。就是需要给定最大的结束时间。

```
#include <algorithm>
class Interval {
    int start;
    int end;
    Interval( int start, int end ) {
        this->start = start;
        this->end = end;
class ConflictingAppointments {
    static bool canAttendAllAppointments( vector<Interval> &intervals ) {
       vector<int> arr(1000, 0);
        int count = 0;
        for ( int i = 0; i < intervals.size(); i++) {
            int start = intervals[i].start;
            int end = intervals[i].end;
            if (!arr[start])
               arr[start] = 1;
            if (!arr[end])
               arr[end] = -1;
        for ( int i = 0; i < arr. size(); i++ ) {
            count += arr[i];
            if (count > 1)
```

```
int main( int argc, char *argv[] ) {
    vector<Interval> intervals = {{1, 4}, {2, 5}, {7, 9}};
    bool result = ConflictingAppointments::canAttendAllAppointments(
intervals);
    cout << "Can attend all appointments: " << result << endl;

intervals = {{6, 7}, {2, 4}, {8, 12}};
    result = ConflictingAppointments::canAttendAllAppointments( intervals );
    cout << "Can attend all appointments: " << result << endl;

intervals = {{4, 5}, {2, 3}, {3, 6}};
    result = ConflictingAppointments::canAttendAllAppointments( intervals );
    cout << "Can attend all appointments::canAttendAllAppointments( intervals );
    cout << "Can attend all appointments: " << result << endl;

system("pause");
}</pre>
```

1-6 会议安排

这算是比较经典的问题了。给定一些列区间。代表会议时间。问需要的最小会议室数 目。

e.g.1

```
Meetings: [[1,4], [2,5], [7,9]]
Output: 2
Explanation: Since [1,4] and [2,5] overlap, we need two rooms to hold these two meetings. [7,9] can occur in any of the two rooms later
```

e.g.2

```
Meetings: [[6,7], [2,4], [8,12]]
Output: 1
Explanation: None of the meetings overlap, therefore we only need one room to hold all meetings.
```

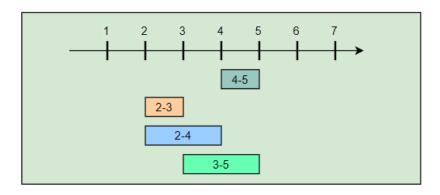
```
Meetings: [[1,4], [2,3], [3,6]]
Output:2
```

Explanation: Since [1,4] overlaps with the other two meetings [2,3] and [3,6], we need two rooms to hold all the meetings.

e.g.4

```
Meetings: [[4,5], [2,3], [2,4], [3,5]]
Output: 2
Explanation: We will need one room for [2,3] and [3,5], and another room for [2,4] and [4,5]. Here is a visual representation of Example 4:
```

思路



从这幅图,我们可以直接将这个建模成为在同一时间,重叠的最多区间数目。

假设我们从时间1开始遍历。

遇到2。cnt = 1。即遇到时间的起点。

遇到2。cnt = 2

遇到3。 cnt = 1。约到结束点了

遇到3。cnt = 2。从这里我们可以判断结束点的优先级高于起始点。(这里将会影响排序的优先级)

依次类推。发现我们的最大cnt = 2。

```
using namespace std;

#include <algorithm>
#include <iostream>
#include <queue>
#include <vector>

class Meeting {
   public:
    int start = 0;
```

```
int end = 0;
    Meeting( int start, int end ) {
        this->start = start;
        this->end = end;
struct Timetype {
    int time;
    int type;
class MinimumMeetingRooms {
    static int findMinimumMeetingRooms( vector<Meeting> &meetings ) {
        vector<Timetype> time vec;
        time vec.reserve(meetings.size() * 2);
        for ( int i = 0; i < meetings.size(); i^{++} ) {
            time vec.push back( { meetings[i].start, 0 } );
            time vec.push back( { meetings[i].end, 1 } );
        std::sort( time vec.begin(), time vec.end(), []( const Timetype &x1,
const Timetype& x2 ) {
            if (x1. time = x2. time)
                return x1. type > x2. type;
            return x1. time < x2. time;
        } );
        int max meeting rooms = 0;
        for ( int i = 0; i < time vec.size(); i++ ) {</pre>
            if ( time vec[i].type == 0 )
            max meeting rooms = std::max( max meeting rooms, cnt );
        return max meeting rooms;
int main( int argc, char *argv[] ) {
    vector<Meeting> input;
```

```
int result;
input = {{1, 4}, {2, 5}, {7, 9}};
result = MinimumMeetingRooms::findMinimumMeetingRooms( input );
cout << "Minimum meeting rooms required: " << result << endl;

input = {{6, 7}, {2, 4}, {8, 12}};
result = MinimumMeetingRooms::findMinimumMeetingRooms( input );
cout << "Minimum meeting rooms required: " << result << endl;

input = {{1, 4}, {2, 3}, {3, 6}};
result = MinimumMeetingRooms::findMinimumMeetingRooms( input );
cout << "Minimum meeting rooms required: " << result << endl;

input = {{4, 5}, {2, 3}, {2, 4}, {3, 5}};
result = MinimumMeetingRooms::findMinimumMeetingRooms( input );
cout << "Minimum meeting rooms required: " << result << endl;
system( "pause" );
}</pre>
```

思路2

对于[2 3] [2 4] [3 5] [4 5] 来说。对于这种题,我们一般喜欢构造一个很长的例子,来 寻找规律。

- 1. 将[2, 3] 安排到第一个会议室
- 2. [2,4]和[2,3]冲突,所以需要安排到一个新的会议室。怎么判断冲突呢,发现start:2 < active room::end:3
- 3. 安排[3, 5]的时候,我们发现我们需要查看已经安排会议室的是否有空闲出来。从这里开始,思路就出来了。怎么判断空闲呢? start:3 >= active_room::min_end:3。很明显,这里应该while循环,将所有空闲的全部腾出去。

```
active_list.push_back(meeting[0]);
for(int i = 1; i < size; i++)
{
    if(active_list.is_expire(meeting[0].start))
        ; // 查看是否有会议过期
    if(当前会议是否和active_list中的冲突)
        push(meeting[i]);//这里只要start <活动会议中的最小的end就有冲突
    update(min_room);
}
```

```
#include <algorithm>
class Meeting {
    int start = 0;
    int end = 0;
    Meeting( int start, int end ) {
        this->start = start;
        this->end = end;
class MinimumMeetingRooms {
    struct endCompare {
        bool operator()( const Meeting &x, const Meeting &y ) {
            return x. end > y. end;
    static int findMinimumMeetingRooms( vector<Meeting> &meetings ) {
        if ( meetings.empty() ) {
        sort( meetings.begin(), meetings.end(),
        [](const Meeting &x, const Meeting &y) {
            return x. start < y. start;
        } );
        int minRooms = 0;
        priority_queue<Meeting, vector<Meeting>, endCompare> minHeap;
        for ( auto meeting : meetings ) {
            while (!minHeap.empty() && meeting.start >= minHeap.top().end)
                minHeap.pop();
```

```
minHeap.push( meeting );
            minRooms = max( minRooms, ( int )minHeap.size() );
        return minRooms;
int main( int argc, char *argv[] ) {
    vector\langle Meeting \rangle input = {{4, 5}, {2, 3}, {2, 4}, {3, 5}};
    int result = MinimumMeetingRooms::findMinimumMeetingRooms( input );
    cout << "Minimum meeting rooms required: " << result << endl;</pre>
    input = \{\{1, 4\}, \{2, 5\}, \{7, 9\}\};
    result = MinimumMeetingRooms::findMinimumMeetingRooms(input);
    cout << "Minimum meeting rooms required: " << result << endl;</pre>
    input = \{\{6, 7\}, \{2, 4\}, \{8, 12\}\};
    result = MinimumMeetingRooms::findMinimumMeetingRooms(input);
    cout << "Minimum meeting rooms required: " << result << endl;</pre>
    input = \{\{1, 4\}, \{2, 3\}, \{3, 6\}\};
    result = MinimumMeetingRooms::findMinimumMeetingRooms(input);
    cout << "Minimum meeting rooms required: " << result << endl;</pre>
    input = \{\{4, 5\}, \{2, 3\}, \{2, 4\}, \{3, 5\}\}\};
    result = MinimumMeetingRooms::findMinimumMeetingRooms(input);
    cout << "Minimum meeting rooms required: " << result << endl;</pre>
```

1-7 max CPU load

给定一系列job。有3个field——start, end, CPU load。寻找最大的CPU, load。多个任务可以同时工作。

```
Jobs: [[1,4,3], [2,5,4], [7,9,6]]
Output: 7
Explanation: Since [1,4,3] and [2,5,4] overlap, their maximum CPU load (3+4=7) will be when both the jobs are running at the same time i.e., during the time interval (2,4).
```

e.g.2

```
Jobs: [[6,7,10], [2,4,11], [8,12,15]]
Output: 15
Explanation: None of the jobs overlap, therefore we will take the maximum load of any job which is 15.
```

e.g.3

```
Jobs: [[1,4,2], [2,4,1], [3,6,5]]
Output: 8
Explanation: Maximum CPU load will be 8 as all jobs overlap during the time interval [3,4].
```

思路

和上道题完全一模一样,只不过上道题的代价是1,这里新增加了代价。

```
using namespace std;

#include <algorithm>
#include <iostream>
#include <queue>
#include <vector>

class Job {
   public:
     int start = 0;
     int end = 0;
     int cpuLoad = 0;

Job( int start, int end, int cpuLoad ) {
```

```
this->start = start;
        this->end = end;
        this->cpuLoad = cpuLoad;
struct JobCompare {
    bool operator()(const Job &x1, const Job &x2) {
        return x1. start < x2. start;
class MaximumCPULoad {
    static int findMaxCPULoad( vector<Job> &jobs ) {
        int maxCPULoad = 0;
        sort( jobs.begin(), jobs.end(), []( const Job &x1, const Job &x2 ) {
            return x1. start < x2. start;
        } );
        priority queue (Job, vector (Job), JobCompare > min heap;
        int min heap load = 0;
        min heap.push(jobs[0]);
        min heap load += jobs[0].cpuLoad;
        for ( int i = 1; i < jobs. size(); i++ ) {
            while (!min heap.empty() && jobs[i].start > min heap.top().end)
                min heap load -= min heap. top().cpuLoad;
                min heap.pop();
            min heap.push(jobs[i]);
            min heap load += jobs[i].cpuLoad;
            maxCPULoad = std::max( maxCPULoad, min heap load );
        return maxCPULoad;
int main( int argc, char *argv[] ) {
    vector\langle Job \rangle input = {{1, 4, 3}, {7, 9, 6}, {2, 5, 4}};
    cout << "Maximum CPU load at any time: " <<</pre>
MaximumCPULoad::findMaxCPULoad( input ) << endl;</pre>
    input = \{\{6, 7, 10\}, \{8, 12, 15\}, \{2, 4, 11\}\};
```

```
cout << "Maximum CPU load at any time: " <<
MaximumCPULoad::findMaxCPULoad( input ) << endl;

input = {{1, 4, 2}, {3, 6, 5}, {2, 4, 1}};
    cout << "Maximum CPU load at any time: " <<
MaximumCPULoad::findMaxCPULoad( input ) << endl;
    system( "pause" );
}</pre>
```

1-8 员工的空闲时间

给K个员工的**工作区间**。找到所有员工都**空闲的时间段**。每个员工的工作时间都是经过start排序后的。

e.g.1

```
Input: Employee Working Hours=[[[1,3], [5,6]], [[2,3], [6,8]]]
Output: [3,5]
Explanation: Both the employess are free between [3,5].
```

e.g.2

```
Input: Employee Working Hours=[[[1,3], [9,12]], [[2,4]], [[6,8]]]
Output: [4,6], [8,9]
Explanation: All employess are free between [4,6] and [8,9].
```

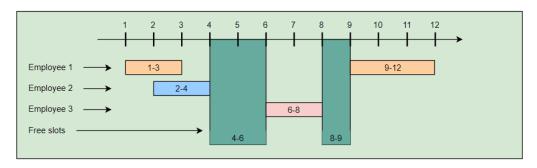
e.g.3

```
Input: Employee Working Hours=[[[1,3]], [[2,4]], [[3,5], [7,9]]]
Output: [5,7]
Explanation: All employess are free between [5,7].
```

思路

注意, 题目给的是工作时间。去找空闲时间。

我们将每个员工的区间画在时间点上。



不难看出,实际上我们就是在找没有区间覆盖的地方,和之前的做法刚好相反。

- 1. 直接将所有区间放到一个数组中,然后去排序。再按之前的思路去做。
- 2. 但实际上,这里实际上是K个排好序的区间。一般K个有序数组直接会让我们思考到K路归并——最小堆。

提供一个经典自作聪明解法。这个解法虽然能解出来。但是相比于上面俩个思路的时间复杂度要高。

如果你数据结构学的好的话,一眼看出来这里的时间复杂度是O(N * logN * logN) 因为每次插入一次数据,都要进行一次调整。共调整N次。每次调整i次,i = 0-N-1。显然是logN * logN的复杂度。

```
class Interval {
    int start = 0;
    int end = 0;
    Interval (int start, int end, int employee = 0, int n work interval = 0)
        this->start = start;
        this->end = end;
        this->employee = employee;
        this->n_work_interval = n_work_interval;
   bool is overlap (const Interval &rhs) {
        return rhs. start <= this->end;
    int employee;
    int n_work_interval;
```

```
struct IntervalCompare {
    bool operator()(const Interval &x1, const Interval &x2) {
        return x1. start > x2. start;
class EmployeeFreeTime {
    static vector<Interval> findEmployeeFreeTime( const
vector<vector<Interval>> &schedule ) {
        vector<Interval> result;
        priority queue<Interval, vector<Interval>, IntervalCompare> min heap;
        for ( int i = 0; i < schedule.size(); i++) {
            for ( int k = 0; k < schedule[i].size(); k++ ) {
                min_heap.push( Interval( schedule[i][k].start, schedule[i]
[k].end, i, k);
        Interval curr interval = min heap. top();
        min heap. pop();
        while (!min heap.empty()) {
            Interval next interval = min heap. top();
            min heap.pop();
            if ( curr interval.is overlap( next interval ) ) {
                if( next interval.end > curr interval.end )
                    curr interval = next interval;
                Interval res( curr interval.end, next interval.start );
                result.push back(res);
                curr interval = next interval;
        return result;
int main( int argc, char *argv[] ) {
    vector\langle \text{Interval} \rangle \rangle input = {{{1, 3}, {5, 6}}, {{2, 3}, {6, 8}}};
    vector<Interval> result = EmployeeFreeTime::findEmployeeFreeTime( input
);
    cout << "Free intervals: ";</pre>
```

```
for ( auto interval : result ) {
    cout << "[" << interval.start << ", " << interval.end << "] ";
}
cout << endl;

input = {{{1, 3}, {9, 12}}, {{2, 4}}, {{6, 8}}};

result = EmployeeFreeTime::findEmployeeFreeTime( input );
cout << "Free intervals: ";
for ( auto interval : result ) {
    cout << "[" << interval.start << ", " << interval.end << "] ";
}
cout << endl;

input = {{{1, 3}}, {{2, 4}}, {{3, 5}, {7, 9}}};

result = EmployeeFreeTime::findEmployeeFreeTime( input );
cout << "Free intervals: ";
for ( auto interval : result ) {
    cout << "[" << iinterval.start << ", " << iinterval.end << "] ";
}
system( "pause" );
}</pre>
```

真正的K路归并堆只做K大小。其时间复杂度只有O(N * logK) 对比使用排序的O(N * logN)。非常值得使用了。

```
using namespace std;

#include <iostream>
#include <queue>
#include <vector>

class Interval {
  public:
    int start = 0;
    int end = 0;

    Interval( int start, int end, int employee = 0, int n_work_interval = 0 )
    {
        this->start = start;
        this->employee = employee;
        this->n_work_interval = n_work_interval;
    }
}
```

```
bool is overlap (const Interval &rhs) {
        return rhs. start <= this->end;
    int employee;
    int n work interval;
struct IntervalCompare {
   bool operator()(const Interval &x1, const Interval &x2) {
        return x1. start > x2. start;
class EmployeeFreeTime {
    static vector<Interval> findEmployeeFreeTime( const
vector<vector<Interval>> &schedule ) {
        vector<Interval> result;
       priority queue<Interval, vector<Interval>, IntervalCompare> min heap;
        for ( int i = 0; i < schedule.size(); i++ ) {</pre>
            min_heap.push( Interval( schedule[i][0].start, schedule[i]
[0]. end, i, 0);
        Interval curr interval = min heap. top();
        min heap. pop();
        if (curr interval.n work interval + 1 <
schedule[curr interval.employee].size() )
            min heap. push (
                Interval(
                    schedule[curr interval.employee]
[curr_interval.n_work_interval + 1].start,
                    schedule[curr interval.employee]
[curr interval.n work interval + 1].end,
                    curr interval. employee,
                    curr interval.n work interval + 1
            );
        while (!min heap.empty()) {
            // 选取下一个最小
            Interval next interval = min heap. top();
            min_heap.pop();
```

```
if ( curr interval.is overlap( next interval ) ) {
                if( next interval.end > curr interval.end )
                     curr interval = next interval;
                Interval res( curr interval.end, next interval.start );
                result.push back(res);
                curr interval = next interval;
            if (curr interval.n work interval + 1 <
schedule[curr interval.employee].size() )
                min heap. push (
                     Interval(
                         schedule[curr interval.employee]
[curr interval.n work interval + 1]. start,
                         schedule[curr interval.employee]
[curr interval.n work interval + 1].end,
                         curr interval. employee,
                         curr interval.n work interval + 1
                );
        return result;
int main( int argc, char *argv[] ) {
    vector\langle \text{Interval} \rangle \rangle input = { {{1, 3}, {5, 6}}, {{2, 3}, {6, 8}} };
    vector<Interval> result = EmployeeFreeTime::findEmployeeFreeTime( input
);
    for ( auto interval : result ) {
        cout << "[" << interval.start << ", " << interval.end << "] ";</pre>
    input = \{\{1, 3\}, \{9, 12\}\}, \{\{2, 4\}\}, \{\{6, 8\}\}\};
    result = EmployeeFreeTime::findEmployeeFreeTime( input );
    for ( auto interval : result ) {
        cout << "[" << interval.start << ", " << interval.end << "] ";</pre>
```

```
cout << endl;

input = { {{1, 3}}, {{2, 4}}, {{3, 5}, {7, 9}} };

result = EmployeeFreeTime::findEmployeeFreeTime( input );

cout << "Free intervals: ";

for ( auto interval : result ) {
    cout << "[" << interval.start << ", " << interval.end << "] ";
}

cout << endl;
input = { {{1, 3}}, {{2, 4}}, {{3, 5}, {7, 9}}, {{6, 7}, {11, 12}} };

result = EmployeeFreeTime::findEmployeeFreeTime( input );

cout << "Free intervals: ";

for ( auto interval : result ) {
    cout << "[" << interval.start << ", " << interval.end << "] ";
}

cout << endl;
system( "pause" );
}
</pre>
```

2. 总结

- 1. 明白区间之间的关系。6种小情况。重点掌握相交的4种情况,以及当按照start排序时候,相交的情况判断。只需要比较arr[i+1].start <= arr[i].end就存在交集。
- 2. 学会使用自定义的sort排序,在很多时候,我们需要定义这种排序,经常以start作为排序。有时候还会附加其他东西。当按start排序时,情况一般就会比较明了。
- 3. 学会自定义区间的比较。
- 4. 学会使用优先级队列。注意优先级队列的插入,和删除都是logN的操作复杂度。
- 5. 在多数组问题上,区间排好序后,可以等价为K路归并问题。