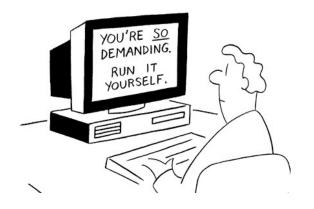


Operating Systems Memory Management I

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Programmer's Dream

Memory

- Private
- Infinitely large
- Infinitely fast
- Non-volatile
- Inexpensive

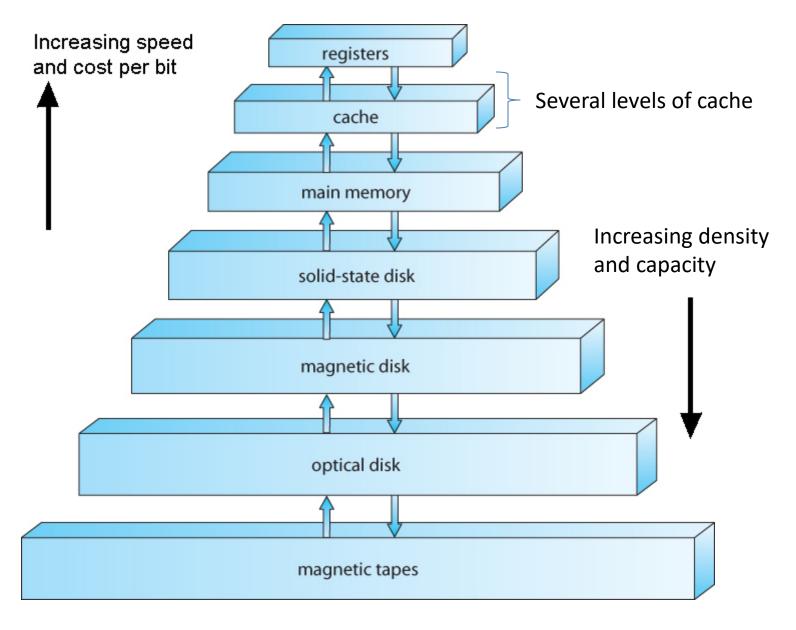
Programmer's Dream

Memory

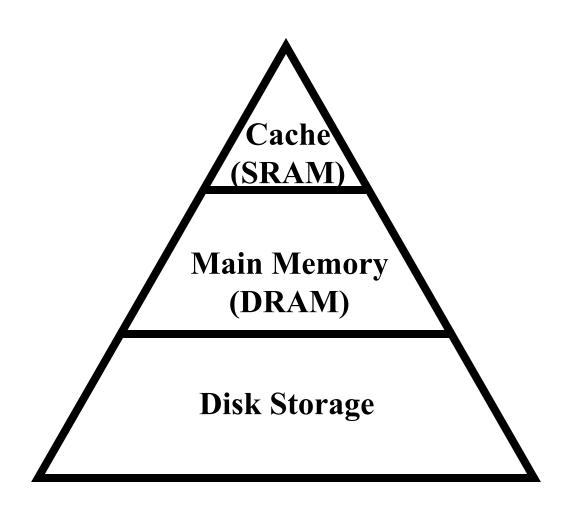
- Private
- Infinitely large
- Infinitely fast
- Non-volatile
- Inexpensive

Programs are getting bigger faster than memories.

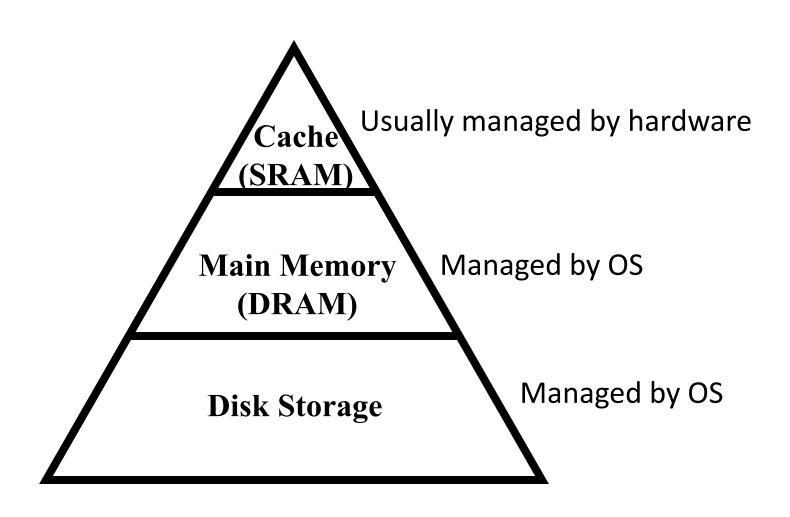
Storage/Memory Hierarchy



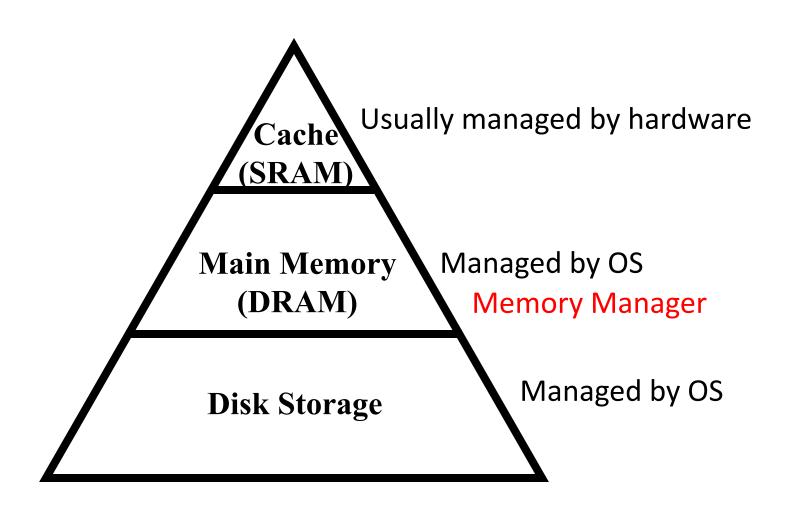
Memory Hierarchy



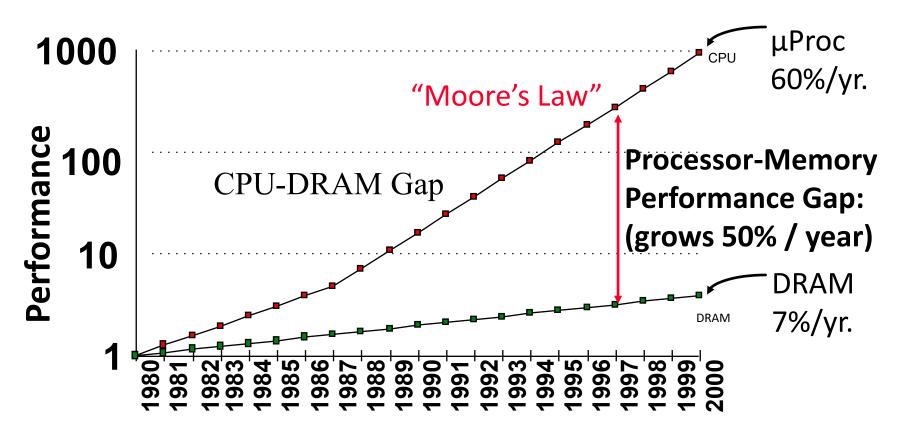
Memory Hierarchy



Memory Hierarchy

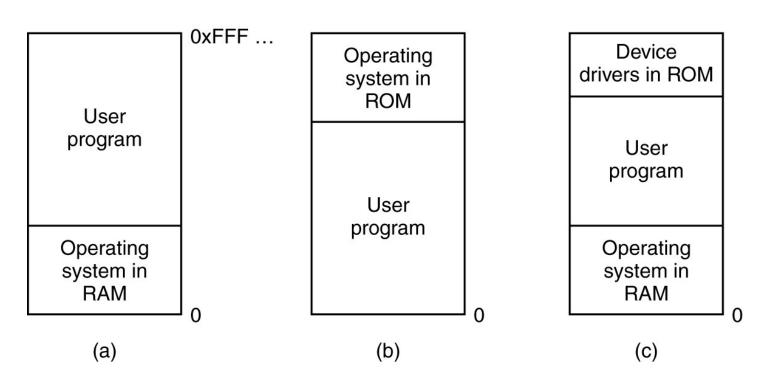


Question: Who Cares About the Memory Hierarchy?

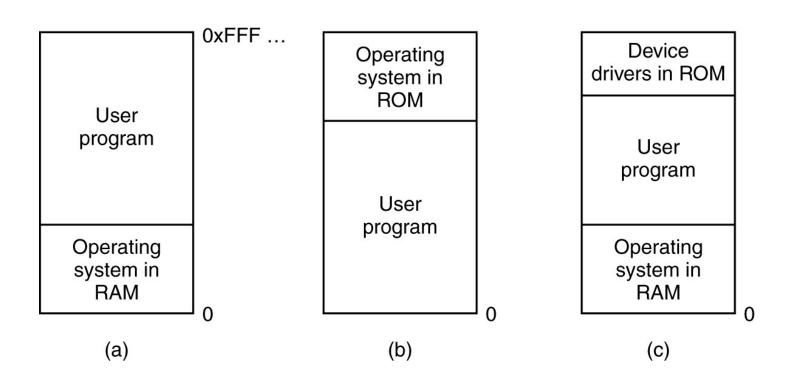


- The hardware and OS memory manager make you see the memory as a single contiguous entity
- How do they do that?
 - Abstraction

Is abstraction necessary?

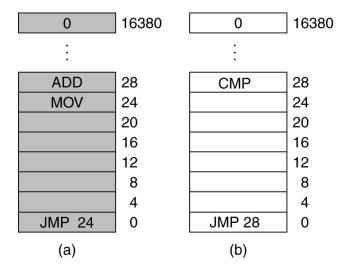


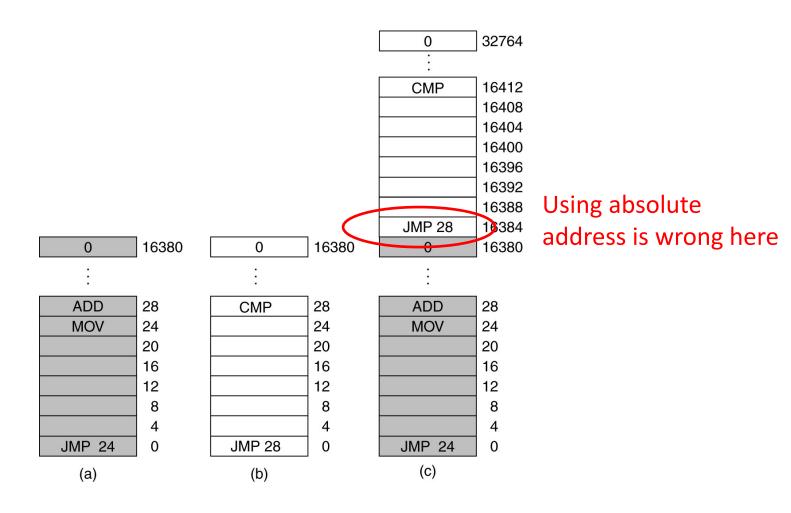
Even with no abstraction, we can have several setups!

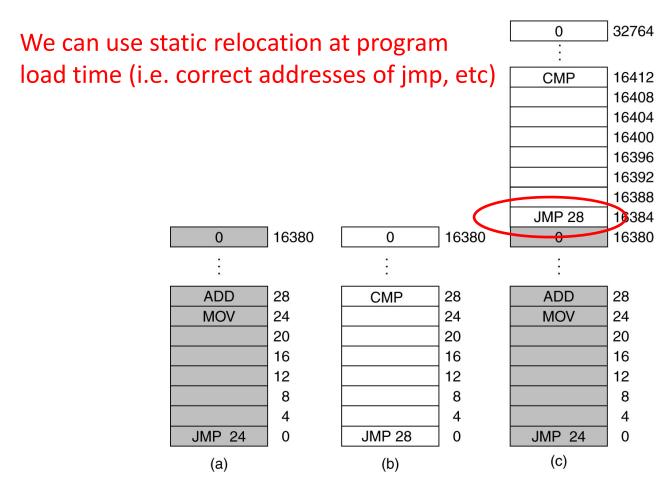


Only one process at a time can be running

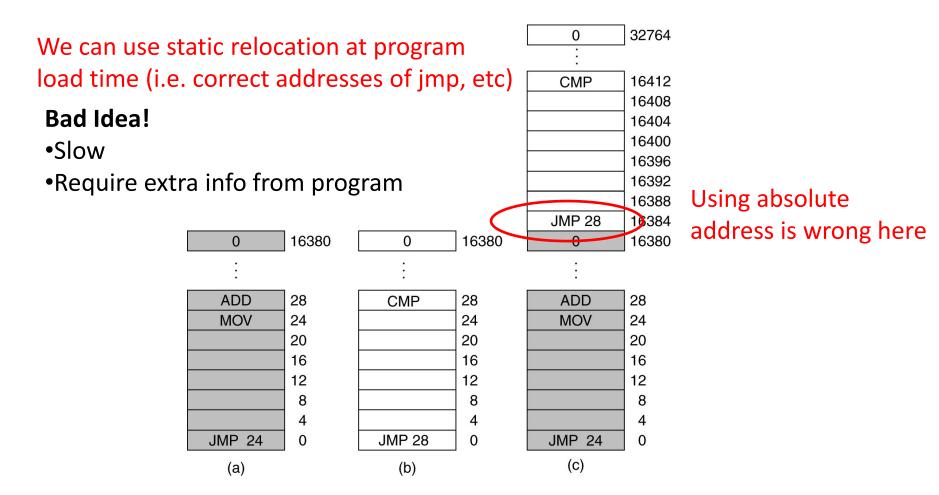
- · What if we want to run multiple programs?
 - OS saves entire memory on disk
 - OS brings next program
 - OS runs next program
- We can use swapping to run multiple programs concurrently
 - Memory divided into blocks
 - Each process assigned to a block.
 - Example: IBM 360







Using absolute address is wrong here



Bottom line: Memory abstraction is needed!

- To allow several programs to co-exist in memory we need
 - Protection
 - Relocation
 - Sharing
 - Logical organization
 - Physical organization
- A new abstraction for memory: Address Space
- Address space = set of addresses that a process can use to address memory

Protection

- Processes need to acquire permission to reference memory locations for reading or writing purposes.
- Location of a program in main memory is unpredictable.
- Memory references generated by a process must be checked at run time.

Relocation

- Programmers typically do not know in advance which other programs will be resident in main memory at the time of execution of their program.
- Active processes need to be able to be swapped in and out of main memory in order to maximize processor utilization.
- Specifying that a process must be placed in the same memory region when it is swapped back in would be limiting
 - may need to relocate the process to a different area of memory

Sharing

- It is advantageous to allow each process access to the same copy of the program/library/... rather than have their own separate copy.
- Memory management must allow controlled access to shared areas of memory without compromising protection.

Logical Organization

- We see memory as linear one-dimensional address space.
- A program = code + data + ... = modules
- Those modules must be organized in that logical address space

Physical Organization

- Memory is really a hierarchy
 - Several levels of caches
 - Main memory
 - Disk
- Managing the different modules of different programs in such a way as:
 - To give illusion of the logical organization
 - To make the best use of the above hierarchy

All of this must be done while ensuring:

- Transparency: the running programs must not know that all of this is happening.
- Efficiency: both in terms of time (speed) and space (not wasting a lot of memory)
- Protection: as we saw, protecting processes from each other

Address Space: Base and Limit

- Map each process address space onto a different part of physical memory
- Two registers: Base and Limit
 - Base: start address of a program in physical memory
 - Limit (sometimes called bound): length of the program
- For every memory access
 - Base is added to the address
 - Result compared to Limit
 - Who is doing this? A piece of hardware inside the processor called the memory management unit (MMU).
- Only OS can modify Base and Limit

Address Space: Base and Limit

Main drawbacks:

Need to add and compare for each memory address

What if memory space is not enough for all programs?

Address Space: Base and Limit

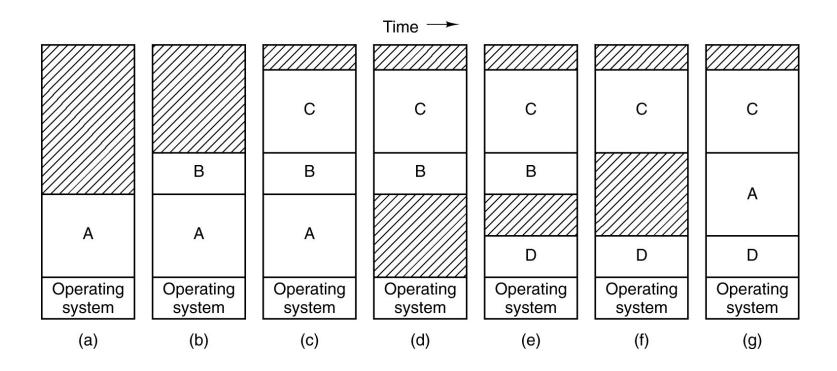
Main drawbacks:

Need to add and compare for each memory address

What if memory space is not enough for all programs?

Then we may need to swap some programs out of the memory.

Swapping



Swapping

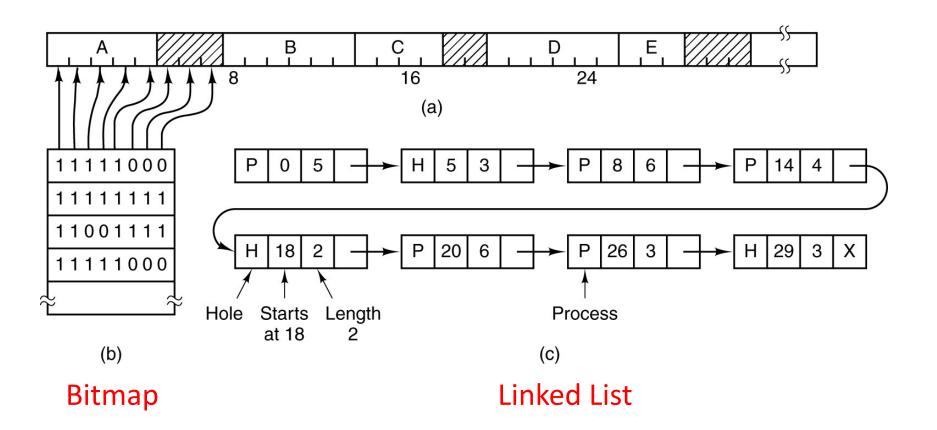
- Programs move in and out of memory.
- Holes are created.
- Holes can be combined -> memory compaction
- · What if a process needs more memory?
 - If a hole is adjacent to the process, it is allocated to it
 - Process must be moved to a bigger hole
 - Process suspended till enough memory is there

Managing Free Memory

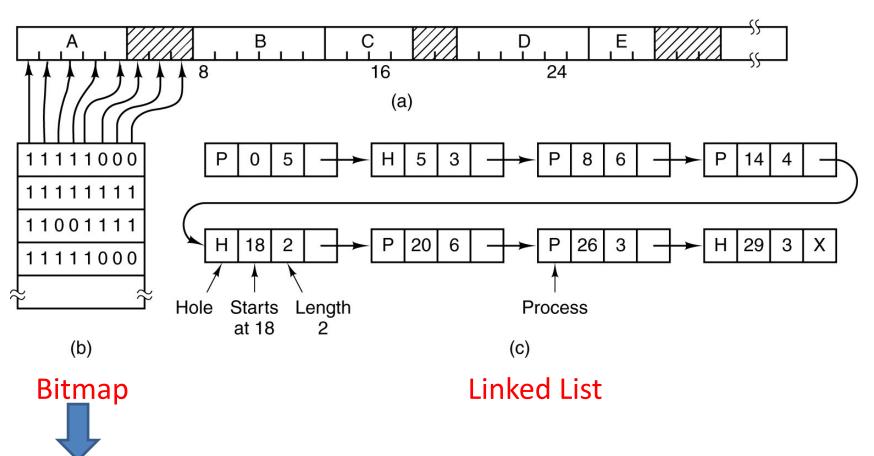
• Bitmap:

- Memory is divided into allocation units of equal size.
- Each unit has a corresponding bit in the bitmap.
- 0 = unit is free 1 = unit is occupied (or vice versa, doesn't matter).
- Linked List: of allocated and free memory segments.
 - Segments are of different sizes.

Managing Free Memory



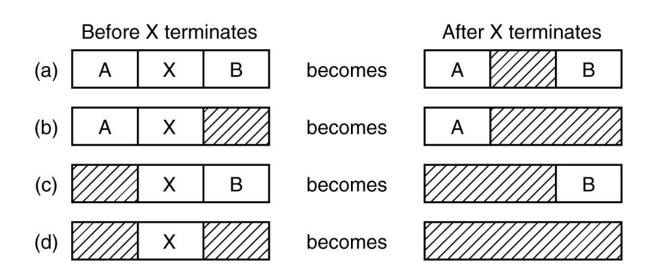
Managing Free Memory



Slow: To find k-consecutive 0s for a new process

Managing Free Memory: Linked List

- Linked list of allocated and free memory segments
- More convenient be double-linked list



Managing Free Memory: Linked List

- How to allocate?
 - First fit
 - Best fit
 - Next fit
 - Worst fit

— ...

Memory Management Techniques

- Memory management brings processes into main memory for execution by the processor
 - involves virtual memory
 - -based on:
 - segmentation (variable size parts) or
 - paging (fixed size parts)

Conclusions

- Process is CPU abstraction
- · Address space is memory abstraction
 - OS memory manager and the hardware helps providing this abstraction
- Two main tasks needed from OS regarding memory management:
 - managing free space
 - making best use of the memory hierarchy