

additional prac

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NYU
EATS

The Weekly EATS



To wish you well on finals, NYU Eats will be hosting a **complimentary breakfast for meal plan holders TONIGHT, May 10th** from **10 pm - 12 am** at NYU Eats at **Downstein and Palladium**. Kick start your study session with us!

NYU Eats at Downstein Menu

Scrambled Eggs
Vegetable Egg Whites

1

```
#include <bits/stdc++.h>
```

```
using namespace std;
```

```
void print(vector<int> a)
```

```
{
```

```
    for(int i = 0; i < a.size(); ++i)
```

```
        printf("%d ", a[i]);
```

```
    printf("\n");
```

```
}
```

```
int cmp_dec(int i, int j)
```

```
{/* sort in decreasing order */
```

```
    return i > j;
```

```
}
```

```
int max_ball(vector<int> diameter, vector<int> length)
```

```
{
```

```
    sort(diameter.begin(), diameter.end(), cmp_dec); print(diameter);
```

```
    sort(length.begin(), length.end(), cmp_dec); print(length);
```

```
    int i = 0, j = 0, count = 0;
```

```
    while(i < diameter.size())
```

```
    {
```

```
        if(diameter[i] <= length[j])
```

```
        {
```

```
            ++count; ++i; ++j;
```

```
        }
```

```
    else
```

```
        ++i;
```

```
}
```

```

        return count;
    }

int main()
{
    vector<int> diameter = {4, 5, 5, 9, 1, 10, 2, 7};
    vector<int> length = {9, 9, 10, 5, 3, 1, 1, 1, 2};
    printf("max_ball = %d\n", max_ball(diameter, length));
}

```

2

```

graham(set q of points):
vector<point> p
p[0] = point with min y-coordinate
p[1], p[2], ..., p[m] are remaining points
sorted by polar angle in counterclockwise order around p[0]
stack<point> s
s.push(p[0])
s.push(p[1])
s.push(p[2])
for i = 3 to m
    while angle formed by points
        s.next_to_top(), s.top(), p[i] /* next_to_top() returns the point one
entry below the top of stack s */
        make a nonleft (counterclockwise) turn
            s.pop()
            s.push(p[i])
return s

```

operations in graham:

1. sort $n - 1$ points
2. add each point once to stack
3. remove points at most once from stack

max if

1. input points are given in reversely sorted order
2. n operations cannot be increased or decreased
3. # removals becomes maximized when almost all points are removed, need at least 3 points to remain in stack since at least 3 points on boundary of convex hull

worst case for graham:

set of n points whose convex hull consists of 3 points

3

/ split all edges of weight 2 into two edges of weight 1 each */*

`new_graph(G)`

`G' = (V', E')`

`for vertex u in V(G)`

`u' = u`

`V'.add(u')`

`for edge e(u, v) in E(G)`

`if weight(e) == 2`

`V'.add(vertex uv)`

`E'.add(e'(u', uv))`

`E'.add(e'(uv, v'))`

`else`

```

        E'.add(e'(u', v'))
    return G'

```

```

G' = new_graph(G)
bfs(G', src vertex s')

```

4

```

/* prim tc =  $O(E \lg V)$  since
 $|V| \cdot \text{extract\_min}() = V \lg V$ 
 $|E| \cdot \text{decrease\_key}() = E \lg V$  */
instead of min priority queue q using heap
consider set of edges whose weight is 1, 10, then 25
in extract_min(), extract element from 1-set, 10-set, 25-set take  $O(1)$  each
time
 $O(|V| + |E|)$ 

```

```

mst_prim(G, source vertex s)
    s.key = 0
    min heap queue Q := G.V
    vector<edge> T

    while Q is not empty
        r = Q.extract_min()
        T.push_back(edge (v, v.parent))
        for each vertex u in adj[v] /*  $O(|E|)$  */
            if u is in Q & u.key > weight(u, v)
                u.key = weight(u, v)
                u.parent = v

```

return T