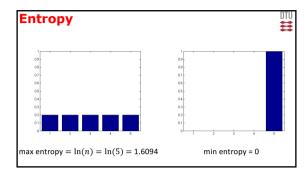
Entropy and related information theoretical measures: Examples on application in data analysis

### Allan Aasbjerg Nielsen

**Technical University of Denmark DTU Compute - Department of Applied Mathematics and Computer Science** alan@dtu.dk

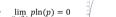
IACG Christmas WS 17 Dec 2013



# Entropy

- Discrete stochastic variable X with probability density function (pdf)  $p(X = x_i)$ i = 1, ..., n; n is number of possible outcomes (or bins)
- Measure of *information* (or *surprise*)  $h(X = x_i)$ :  $x_i$  very probable, i.e.,  $p(X = x_i)$  is high then  $h(X = x_i)$  should be low, and v.v.
- Realizations  $x_i$  and  $y_i$  of independent X and Y, i.e., the two-dimensional pdf  $p(X = x_i, Y = y_i)$  equals the product of the one-dimensional marginal pdfs  $p(X = x_i)p(Y = y_i)$ , we would like the joint information content to equal the sum of the marginal information contents, i.e.,  $h(X = x_i, Y = y_i) = h(X = x_i) + h(Y = y_i)$
- Feasible with  $h(X = x_i) = -\ln(p(X = x_i))$
- The expectation H(X) of the information content is termed the (Shannon) entropy:  $H(X) = -\sum_{i=1}^{n} p(X = x_i) \ln(p(X = x_i)) = -\mathbb{E}\{\ln(p)\}$

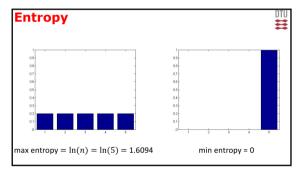
# Entropy



- Unit: nat, nit or nepit (natural log); bit (log base 2); ban or dit (log base 10)
- For continuous stochastic variable  $H(X) = -\int p(x) \ln(p(x)) dx$
- Differential entropy; since p(x) can be >1, H(X) can be negative (or infinite) Gaussian pdf has max entropy for continuous distributions with finite variance.  $H(X) = \{1 + \ln(2\pi\sigma^2)\}/2; H(X) = \ln(|(2\pi e)\Sigma|)/2$
- Discrete X: be careful if binning is applied
- Empirical entropy  $\widehat{H}(X) = -\sum_{i=1}^{N} \ln(p(X = x_i)) / N$ ; N is number of obs
- Rényi entropy:  $H_{\alpha}(X) = \frac{1}{1-\alpha} \ln(\sum_{i=1}^n p(X=x_i)^{\alpha}), \, \alpha > 0$  and  $\alpha \neq 1$
- $\alpha \rightarrow 1$ : Shannon entropy
- $\alpha=2$ : Collision or **Rényi entropy**:  $H_2(X)=-\ln(\sum_{i=1}^n p(X=x_i)^2)=-\ln(\mathbb{E}\{p\})$

## Entropy

- Shannon entropy:  $-\mathbf{E}\{\ln(p)\}$
- *Rényi entropy:*  $-ln(E\{p\})$  $E\{p\}$  is energy
- Entropy is the average amount of information (or surprise) obtained from obs
- Entropy is a measure of order
- Physical system: heat source/sink, ability to do work, low
- Source/sink same temperature: no work done, max entropy Histogram of temperature of source/sink
- A bit like my (low entropy) office: a lot of potential for work



## Entropy

- Claude Elwood Shannon (1916-2001), A Mathematical Theory of Communication. Bell System Technical Journal, 1948: application areas communication, compression - "father of information theory"
- Alfréd Rényi (1921-1970), contributions in combinatory, graph theory, number theory but mostly in probability theory
- John von Neumann (1903-1957), major contributions in mathematics, physics, economics, computer science, and statistics (source http://en.wikipedia.org/wiki/John\_von\_Neumann): use the term entropy not only because of the link to physics, but also because "nobody knows what entropy really is, so in any discussion you

will always have an advantage", (Christopher Michael Bishop, 2006)

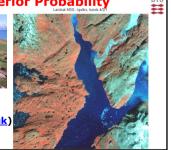
#### Entropy

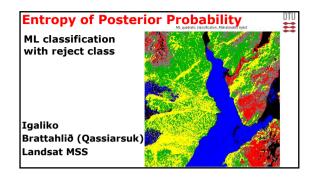
- Early use in maximum entropy (image) reconstruction (B. Roy Frieden, 1972)
- · Entropy as regularizer in inverse problems to
- · suppress low-intensity ripple
- · sharpen point sources
- · in e.g. astronomical data
- Equality of distributions of saleries/income
- Gini index first order series expansion of ln(x) in Shannon entropy around x = 1,  $\ln(x) \cong x - 1$ : Gini  $= \sum_{i=1}^{n} p(X = x_i)(1 - p(X = x_i)) = \mathbb{E}\{1 - p\}$
- None of this described further here, on to simpler things

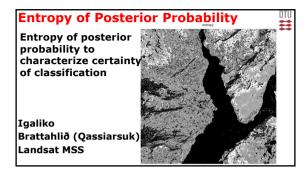
# Intropy of Posterior Probability

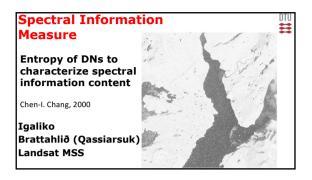


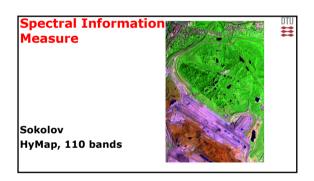


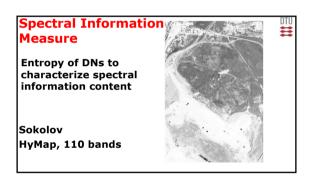


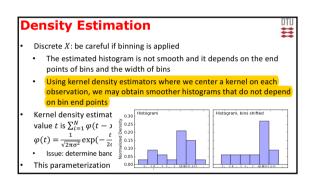


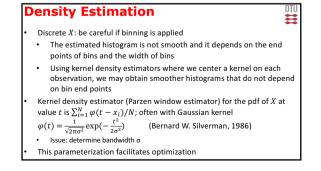


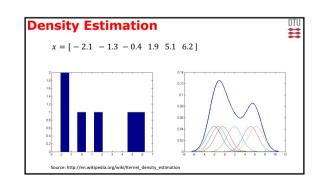


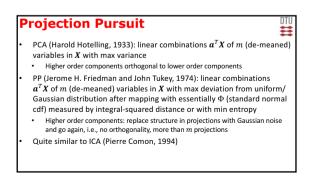


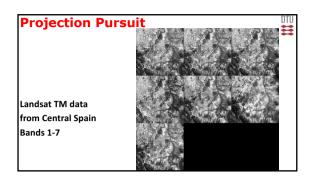


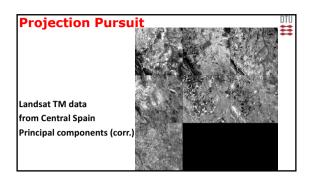


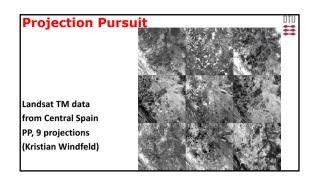






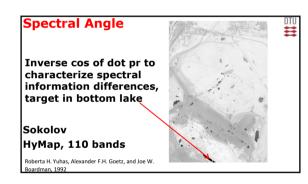




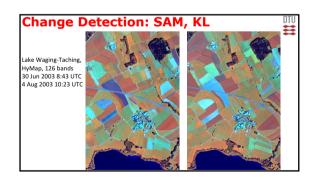


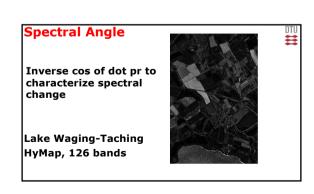
# Relative Entropy/KL divergence

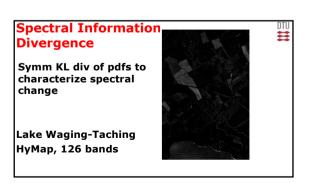
- Relative entropy aka Kullback-Leibler divergence between two pdfs  $p(X=x_i)$  and  $q(X=x_i)$  defined on the same set of outcomes (or bins): the expectation of the logarithmic difference between p and q  $D_{KL}(p,q) = \sum_{i=1}^n p(X=x_i) \ln \frac{p(X=x_i)}{q(X=x_i)}$
- Measure of proximity, satisfies Gibbs' inequality  $D_{KL}(p,q) \ge 0$ , = 0 for  $p(X = x_i) = q(X = x_i)$  only; not symmetric in p and q (not a metric)
- Symmetrized measure:  $D_{KL}(p,q) + D_{KL}(q,p)$
- Rényi divergence  $D_{\alpha}(p,q) = \frac{1}{\alpha-1} \ln \left( \sum_{i=1}^{n} \left( \frac{p(X=x_{i})^{\alpha}}{q(X=x_{i})^{\alpha-1}} \right) \right) = \frac{1}{\alpha-1} \ln (\sum_{i=1}^{n} p(X=x_{i})^{\alpha} q(X=x_{i})^{1-\alpha})$







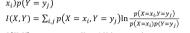




#### **Mutual Information**

Kullback-Leibler divergence  $D_{KL}(p,q) = \sum_{i=1}^{n} p(X=x_i) \ln \frac{p(X=x_i)}{q(X=x_i)}$ 

• **Mutual information**: measure of extent to which X and Y are independent, KL divergence between  $p(X=x_i,Y=y_j)$  and  $p(X=x_i)p(Y=y_j)$ 



• I(X,Y) symmetric in X and Y (a metric)

- I(X,Y) = H(X) + H(Y) H(X,Y)
- A normed version as supplement to correlation



DTU

#### **Mutual Information**

CCA: linear combinations  $a^TX$  and  $b^TY$  of m and k (de-meaned) variables in X and Y with max correlation (Harold Hotelling, 1936)

Higher order components orthogonal to lower order components

Replace correlation in CCA with mutual information (Xiangrong Yin, 2004; Karasuyama & Sugiyama, 2012; Vestergaard & Nielsen: CIA, 2015)

- Higher order components: min MI with lower order components; no
- orthogonality; more than  $\max(m, k)$  components
  CIA handles large datasets
- Examples
- CIA paper

