CSE351: Section 5

Procedures, Stacks, Structs and Unions

October 27, 2011

Recursive Factorial

```
int rfact(int x)
{
  int rval;
  if (x <= 1)
    return 1;
  rval = rfact(x-1);
  return rval * x;
}</pre>
```

Recursive Factorial

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```
rfact:
   pushl %ebp
   movl %esp,%ebp
   pushl %ebx
   movl 8(%ebp),%ebx
   cmpl $1,%ebx
    jle .L78
   leal -1(%ebx),%eax
   pushl %eax
   call rfact
    imull %ebx,%eax
    jmp .L79
    .align 4
.L78:
   movl $1,%eax
.L79:
   mov1 -4(%ebp),%ebx
   movl %ebp, %esp
   popl %ebp
    ret
```

Recursive Factorial

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int rfact(int x)
{
  int rval;
  if (x <= 1)
    return 1;
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}</pre>
```

Registers

%ebx used, but saved at beginning & restored at end

%eax used without first saving

- expect caller to save
- pushed onto stack as parameter for next call . L79:
- used for return value

Convention dictates behavior

```
rfact:
   pushl %ebp
   movl %esp,%ebp
   pushl %ebx
   mov1 8(%ebp), %ebx
   cmpl $1,%ebx
    jle .L78
    leal -1(%ebx), %eax
   pushl %eax
   call rfact
    imull %ebx, %eax
    jmp .L79
    .align 4
.L78:
   movl $1, %eax
   movl -4(%ebp),%ebx
   movl %ebp, %esp
   popl %ebp
    ret
```

Passing by Reference with Pointers

Recursive Procedure

```
void s_helper
  (int x, int *accum)
{
  if (x <= 1)
    return;
  else {
    int z = *accum * x;
    *accum = z;
    s_helper (x-1,accum);
  }
}</pre>
```

Top-Level Call

```
int sfact(int x)
{
  int val = 1;
  s_helper(x, &val);
  return val;
}
```

Passing a pointer to a function allows the function to modify the contents of the memory being pointed to

Creating & Initializing Pointer

```
int sfact(int x)
{
  int val = 1;
  s_helper(x, &val);
  return val;
}
```

- Variable val must be stored on stack
- Need to pass a pointer to s_helper
- Compute pointer as -4 (%ebp)
- Push on stack as second argument

Initial part of sfact

```
_sfact:
   pushl %ebp
   movl %esp,%ebp
   subl $16,%esp
   movl 8(%ebp),%edx
   movl $1,-4(%ebp)
```

```
8 x

4 Rtn adr

0

-4

-8

-12

-16
```

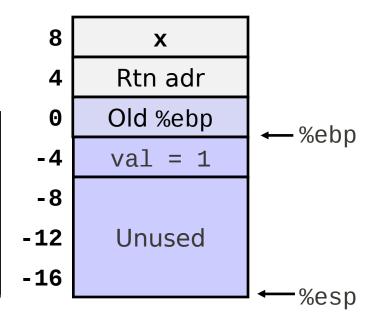
Creating & Initializing Pointer

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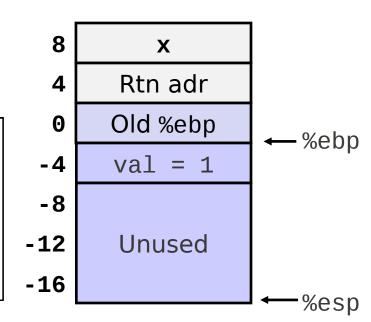


Creating & Initializing Pointer

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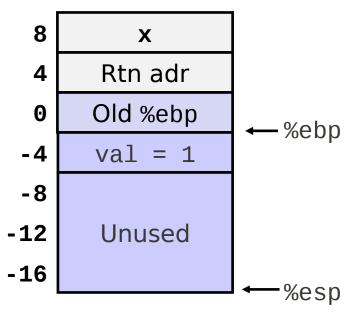
Initial part of sfact



```
int sfact(int x)
{
  int val = 1;
  s_helper(x, &val);
  return val;
}
```

Calling s helper from sfact

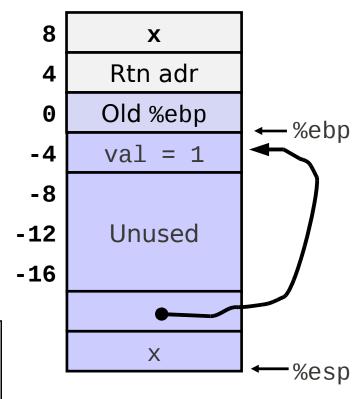
```
leal -4(%ebp),%eax
pushl %eax
pushl %edx
call s_helper
movl -4(%ebp),%eax
• • •
```



```
int sfact(int x)
{
  int val = 1;
  s_helper(x, &val);
  return val;
}
```

Calling s_helper **from** sfact

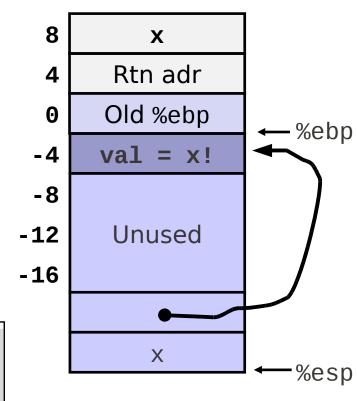
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call s helper
movl -4(%ebp), %eax
• • •
```



```
int sfact(int x)
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  s_helper(x, &val);
  return val;
}
```

Calling s_helper **from** sfact

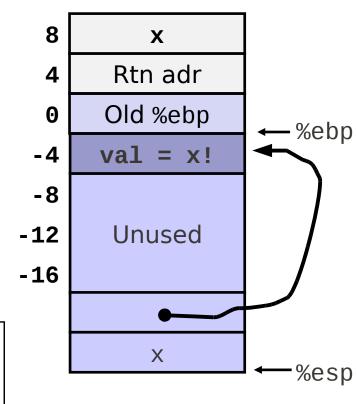
```
leal -4(%ebp),%eax
pushl %eax
pushl %edx
call s helper
movl -4(%ebp),%eax
• • •
```



```
int sfact(int x)
{
  int val = 1;
  s_helper(x, &val);
  return val;
}
```

Calling s_helper **from** sfact

```
leal -4(%ebp),%eax # Compute &val
pushl %eax # Push on stack
pushl %edx # Push x
call s_helper # call
movl -4(%ebp),%eax # Return val
• • • # Finish
```



IA 32 Procedure Summary

- Stack makes recursion work
- Private storage for each *instance* of procedure call
 - Instantiations don't clobber each other
 - Addressing of locals + arguments can be relative to stack positions
- Managed by stack discipline
 - Procedures return in inverse order of calls
- IA32 procedures
 - Combination of Instructions + Conventions
 - call / ret instructions
 - Register usage conventions
 - caller / callee save
 - **%ebp** and **%esp**
 - Stack frame organization conventions

Caller **Frame** %ebp -Local **Variables**

Arguments Return Addr Old %ebp Saved Registers

Argument Build

%esp

x86-64 Registers: Conventions

%rax	Return value
%rbx	Callee saved
%rcx	Argument #4
%rdx	Argument #3
%rsi	Argument #2
%rdi	Argument #1
%rsp	Stack pointer
%rbp	Callee saved

%r8	Argument #5
%r9	Argument #6
%r10	Caller saved
%r11	Caller Saved
%r12	Callee saved
%r13	Callee saved
%r14	Callee saved
%r15	Callee saved

Some Differences between x8664 and IA32

- More general purpose registers
- First six function arguments passed via registers
 - Why?
- Sometimes we don't need a frame pointer
 - Why?
 - How are local variables accessed?
- Misc. differences in instructions, too

Using Nested Arrays

Strengths

- C compiler handles doubly subscripted arrays
- Generates very efficient code
- Avoids multiply in index computation

Limitation

 Only works for fixed array size

```
#define N 16
typedef int fix_matrix[N][N];
```

```
/* Compute element i,k of
   fixed matrix product */
int fix_prod_ele
(fix_matrix a, fix_matrix b,
 int i, int k)
  int j;
  int result = 0;
  for (j = 0; j < N; j++)
    result += a[i][j]*b[j][k];
  return result;
```

```
b
          a
                                 j-th column
                    X
i-th row
```

Dynamic Nested Arrays

Strength

 Can create matrix of any size

Programming

Must do index computation explicitly

<u>Performance</u>

- Accessing single element costly
- Must do multiplication

```
int * new_var_matrix(int n)
{
   return (int *)
    calloc(sizeof(int), n*n);
}
```

```
int var_ele
  (int *a, int i, int j, int n)
{
   return a[i*n+j];
}
```

Dynamic Nested Arrays

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Programming

Must do index computation explicitly

```
int * new_var_matrix(int n)
{
   return (int *)
     calloc(sizeof(int), n*n);
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```

```
int var_ele
  (int *a, int i, int j, int n)
{
   return a[i*n+j];
}
```

Performance

```
movl 12(%ebp),%eax # i
movl 8(%ebp),%edx # a
imull 20(%ebp),%eax # n*i
addl 16(%ebp),%eax # n*i+j
movl (%edx,%eax,4),%eax # Mem[a+4*(i*n+j)]
```

Arrays of Structures

Each element in the array must be properly aligned.

```
struct S2 {
  double v;
  int i[2];
  char c;
} a[10];
```

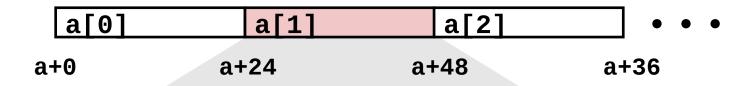


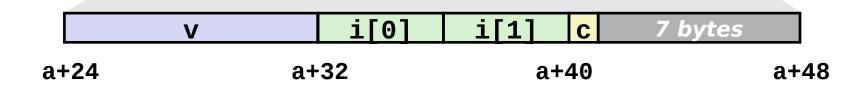
Arrays of Structures

Each element in the array must be properly aligned.

True data length is 8 + 2*4 + 1, but actually uses 8 + 2*4 + 8

```
struct S2 {
  double v;
  int i[2];
  char c;
} a[10];
```





Accessing Array Elements

Struct S3 is 8 bytes, but requires 12 bytes for padding

To access the ith element in a, we compute the offset as 12*i

```
struct S3 {
   short i;
   float v;
   short j;
} a[10];
```

Accessing Array Elements

To get to member j:

- Compute array offset 12i
- Compute offset 8 with structure
- Assembler gives offset a+8

```
struct S3 {
    short i;
    float v;
    short j;
} a[10];
```

```
short get_j(int idx)
{
    return a[idx].j;
// return (a + idx)->j;
}
```

```
# %eax = idx
leal (%eax,%eax,2),%eax # 3*idx
movswl a+8(,%eax,4),%eax
```

Unions

```
struct rec {
   int i;
   int a[3];
   int *p;
};
```

```
union U1 {
   int i;
   int a[3];
   int *p;
} *up;
```

Concept

- Allow same regions of memory to be referenced as different types
- Aliases for the same memory location

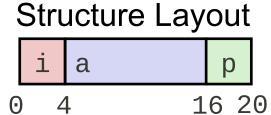
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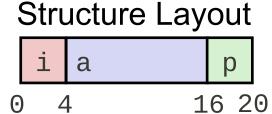
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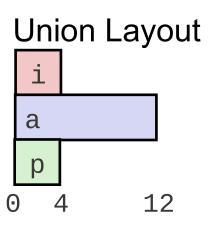
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  int i;
  int a[3];
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   int i;
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} *up;
```

Concept

- Allow same regions of memory to be referenced as different types
- Aliases for the same memory location





Union Allocation

- Size determined by the largest element
- Can only use one field at a time

```
union U1 {
  char c;
  int i[2];
  double v;
} *up;
```

```
c
i[0] i[1]
v
up+0 up+4 up+8
```

```
struct S1 {
  char c;
  int i[2];
  double v;
} *sp;
```



Using Unions to Access Bit Patterns

```
typedef union {
  float f;
  unsigned u;
} bit_float_t;
```

```
u
f
0 4
```

```
float bit2float(unsigned u) {
  bit_float_t arg;
  arg.u = u;
  return arg.f;
}
```

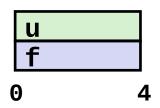
```
unsigned float2bit(float f) {
  bit_float_t arg;
  arg.f = f;
  return arg.u;
}
```

Same as (float)u?

Same as (unsigned)f?

Using Unions to Access Bit Patterns

```
typedef union {
  float f;
  unsigned u;
} bit_float_t;
```



```
float bit2float(unsigned u) {
  bit_float_t arg;
  arg.u = u;
  return arg.f;
}
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unsigned float2bit(float f) {
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Same as (float)u?

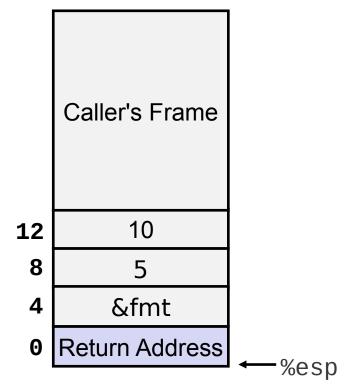
Same as (unsigned)f?

No! Casts actually trigger a bit conversion

- How many arguments does printf take?
 - As many as we want!
- What does the function signature look like?
 - int printf(const char *fmt, ...)
 - The "..." tells compiler to expect a variable number of arguments
- How do we pass an arbitrary number of arguments?
 - Just push 'em all on the stack like before
- How does printf know how many arguments it received?
 - The format string tells it what to expect and in what order

• Example: printf("%d %d\n", 5, 10)

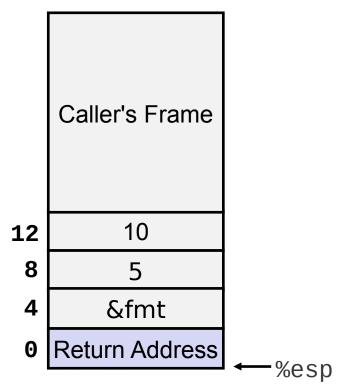
Stack at time of call



Output: "5 10"

• Example: printf("%d %d %d %d\n", 5, 10)

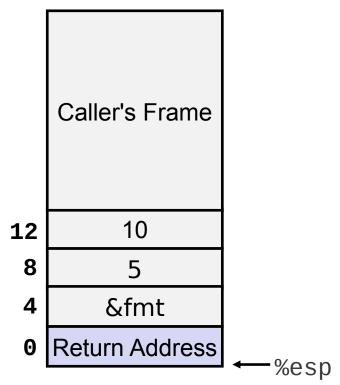
Stack at time of call



Output?

• Example: printf("%d %d %d %d\n", 5, 10)

Stack at time of call



Output: "5 10 ??????? ??????"