

阿里云基础设施网络 智能运维研究

翟恩南（恩南）

网络研究团队



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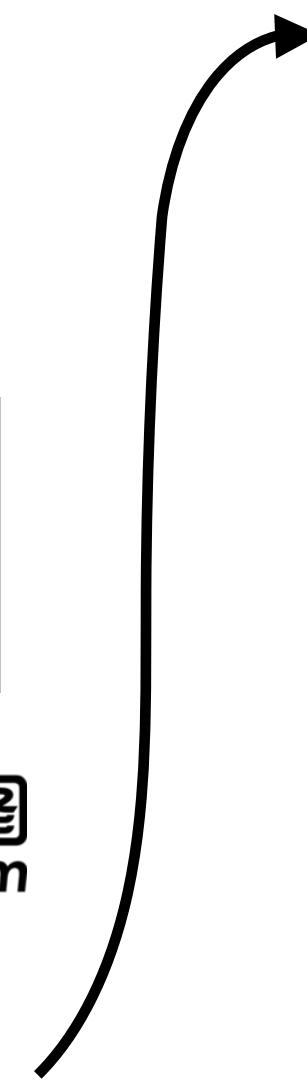
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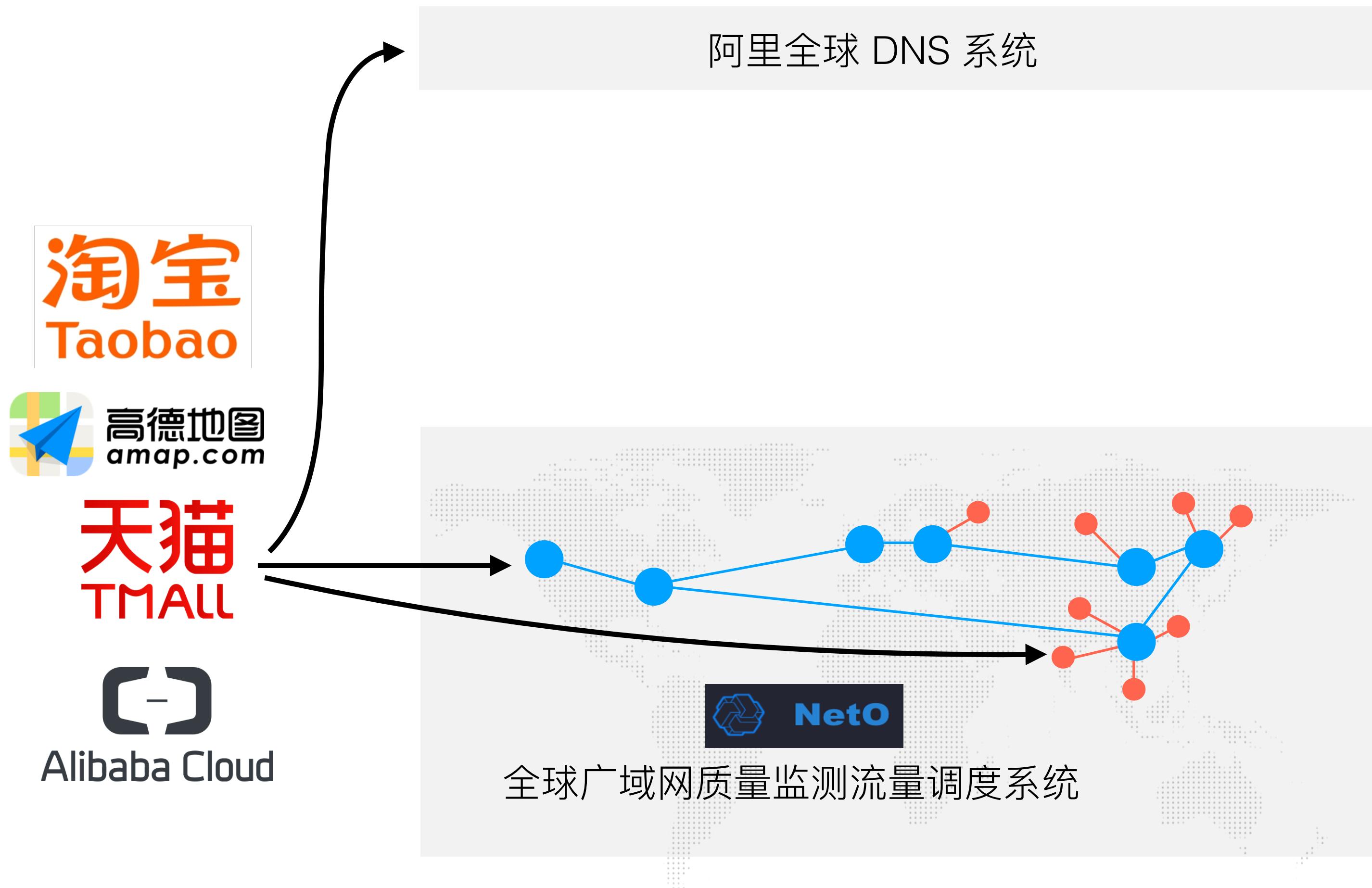
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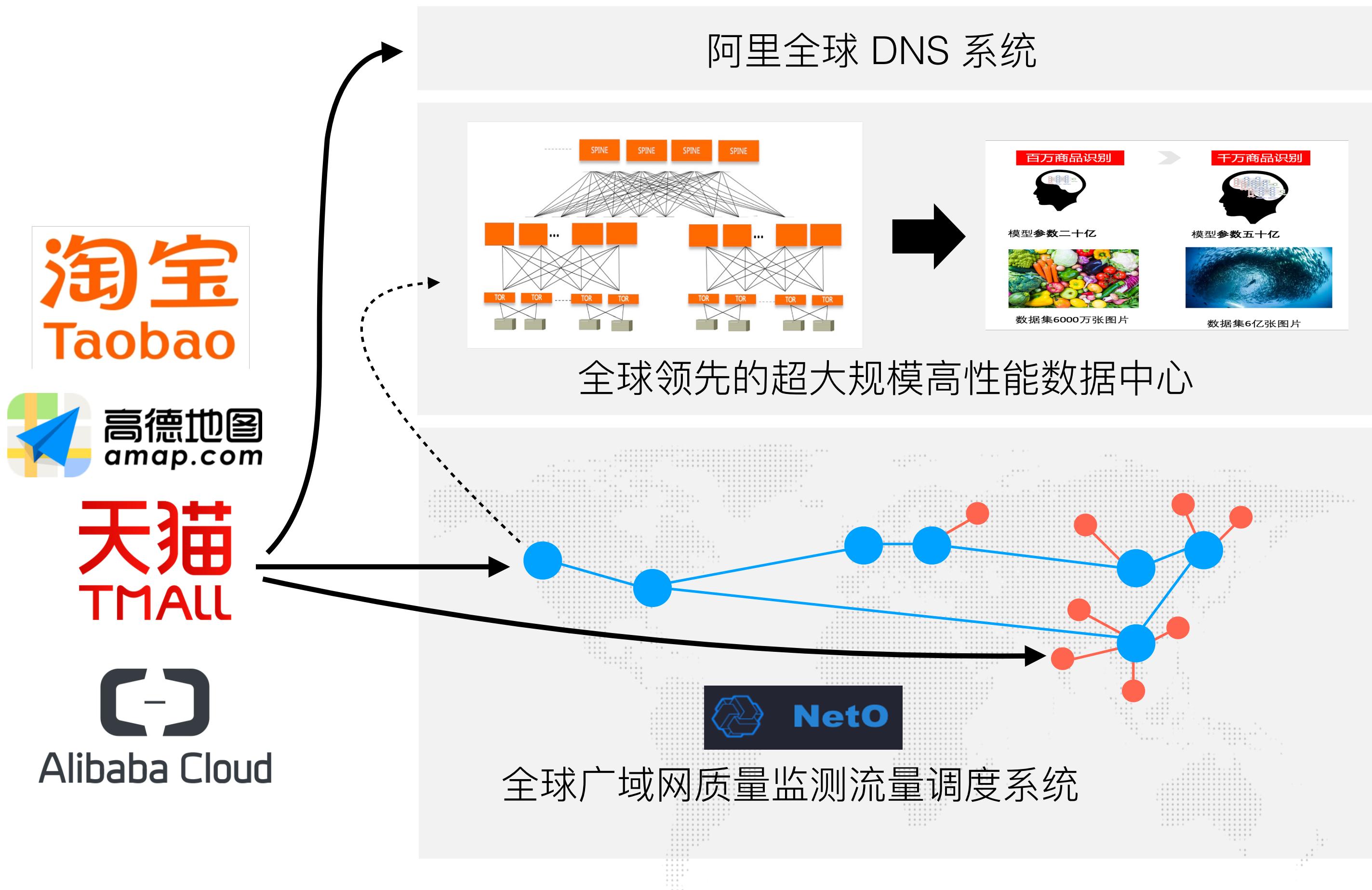
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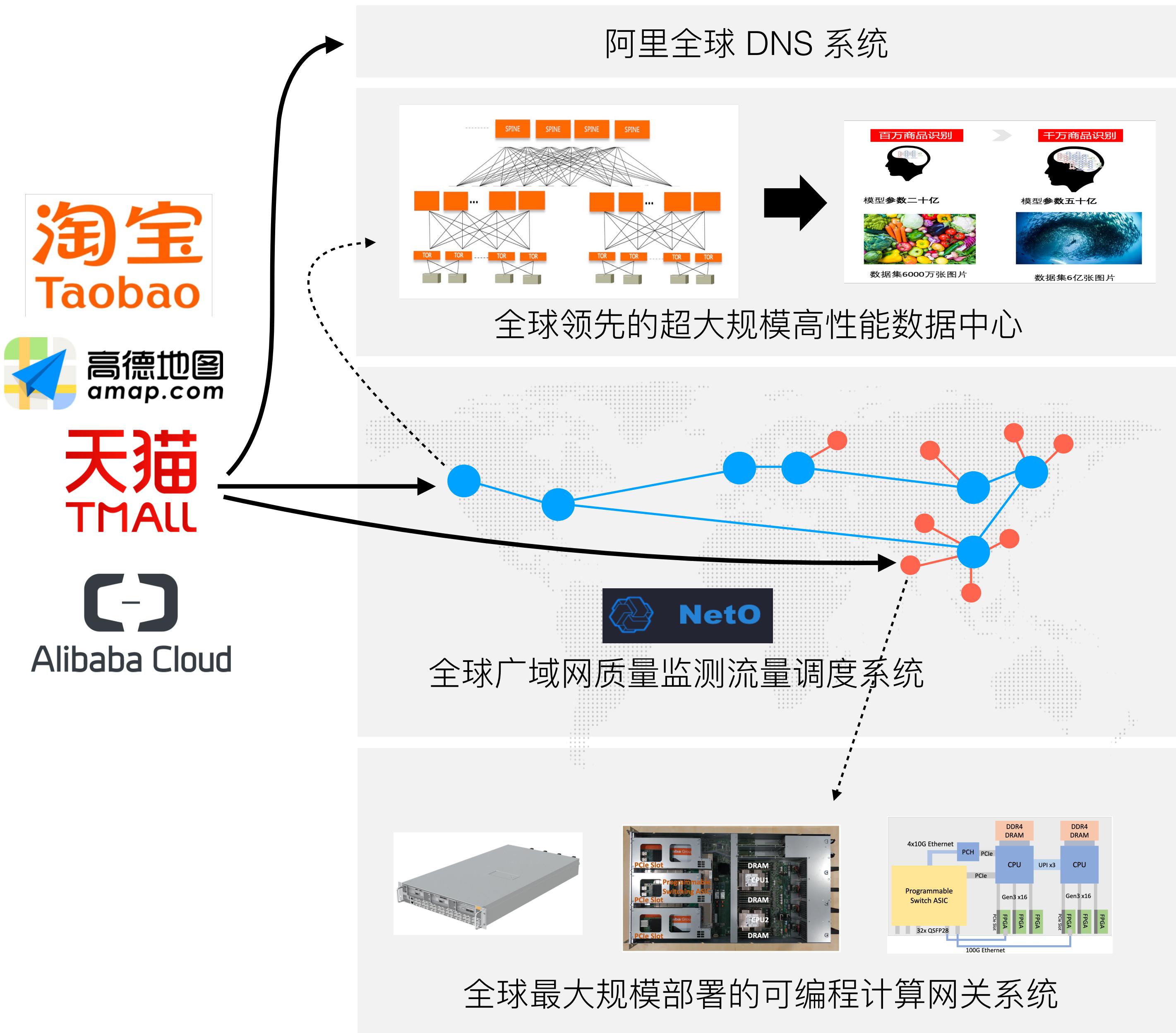
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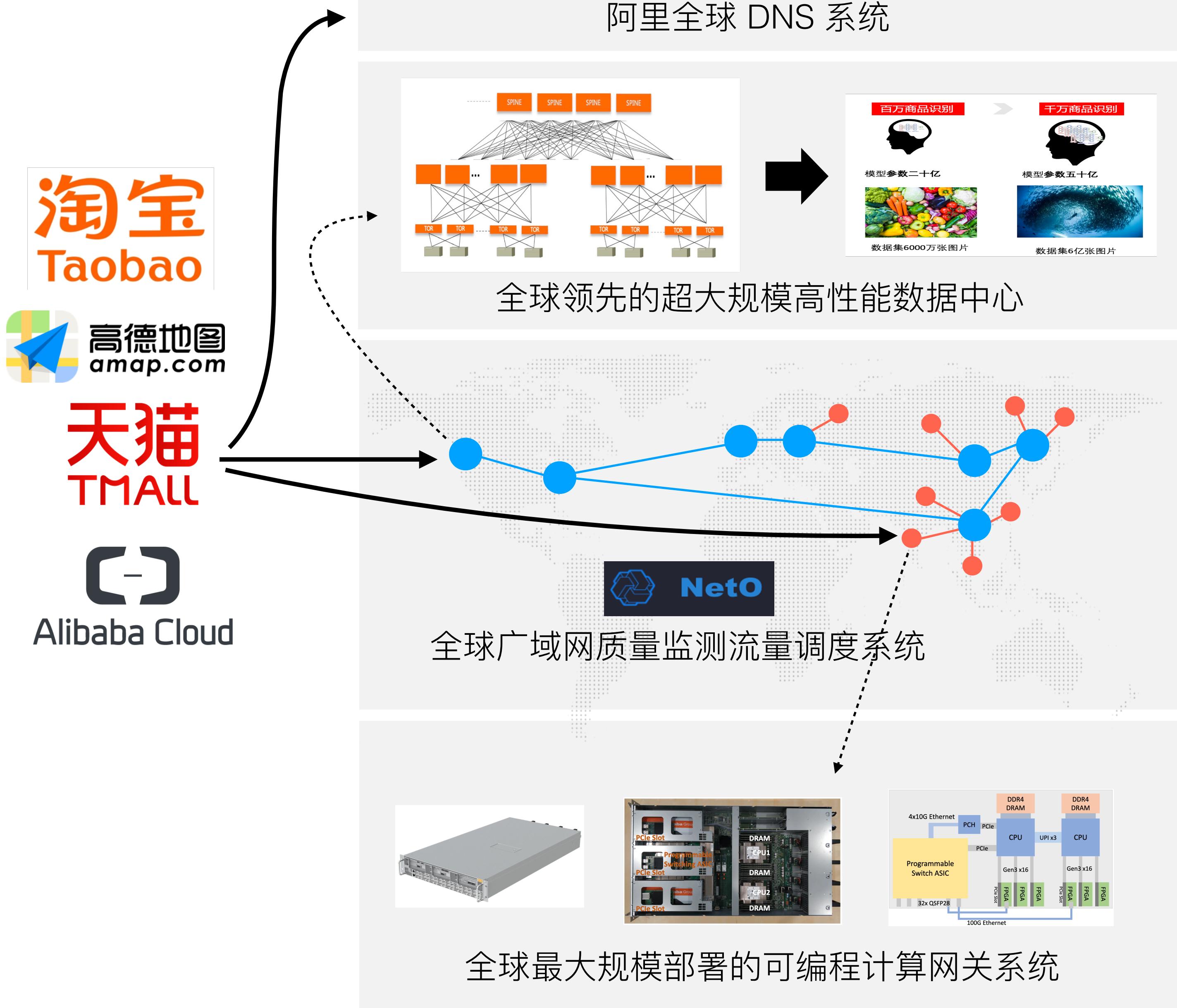
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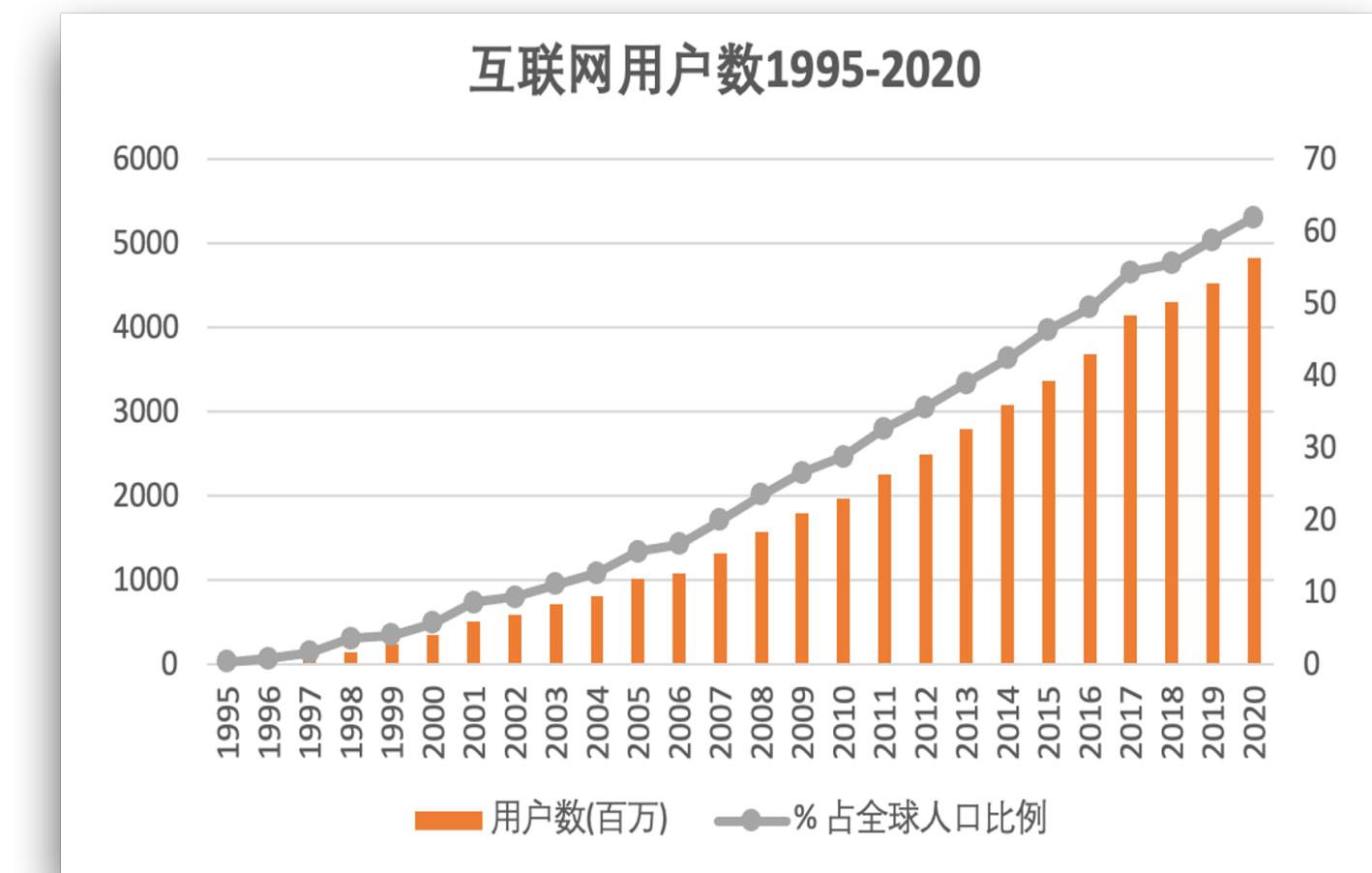
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阿里云全球基础设施网络系统



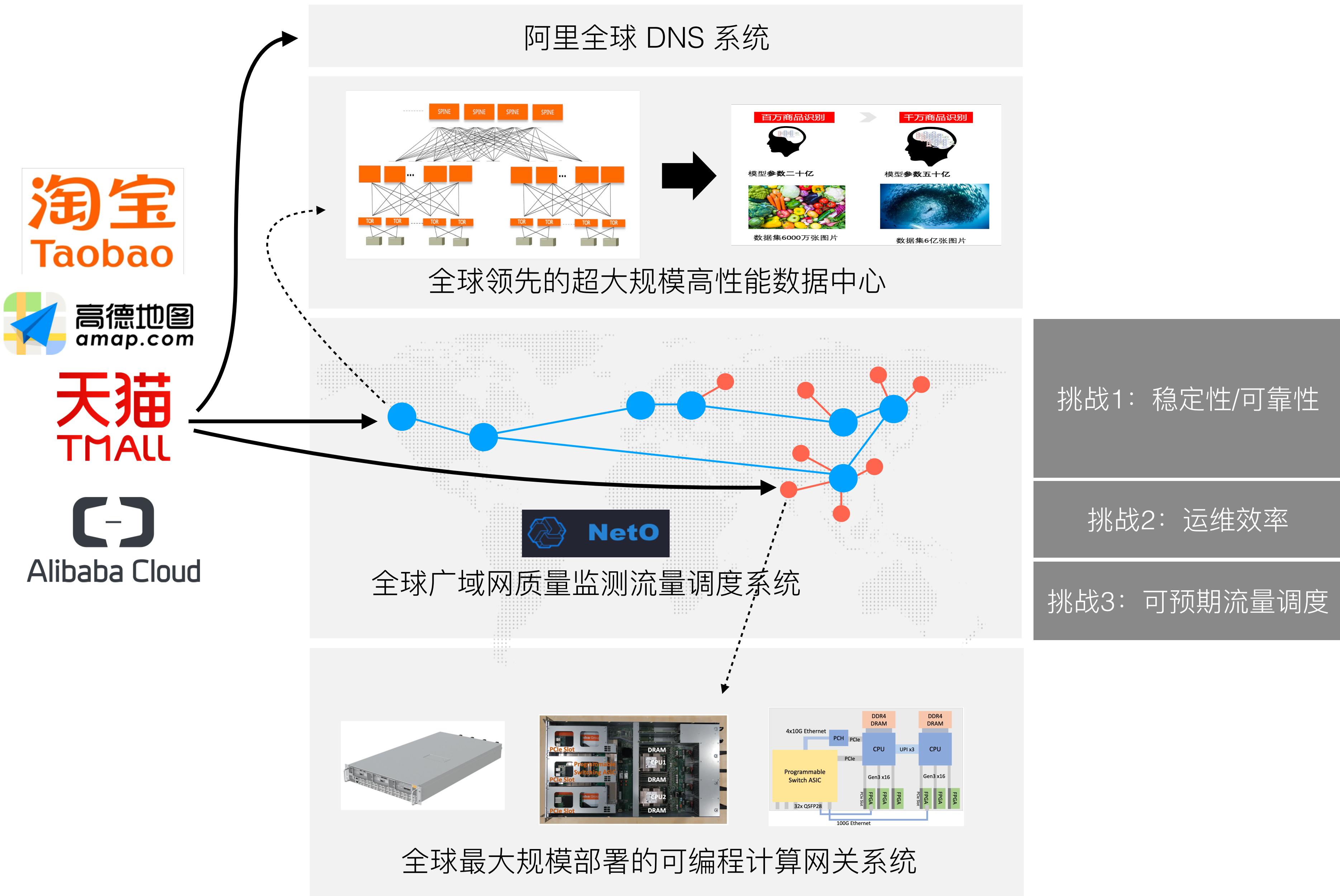
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在网络规模持续增长的情况下，如何保证稳定性和高性能

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阿里云全球基础设施网络系统

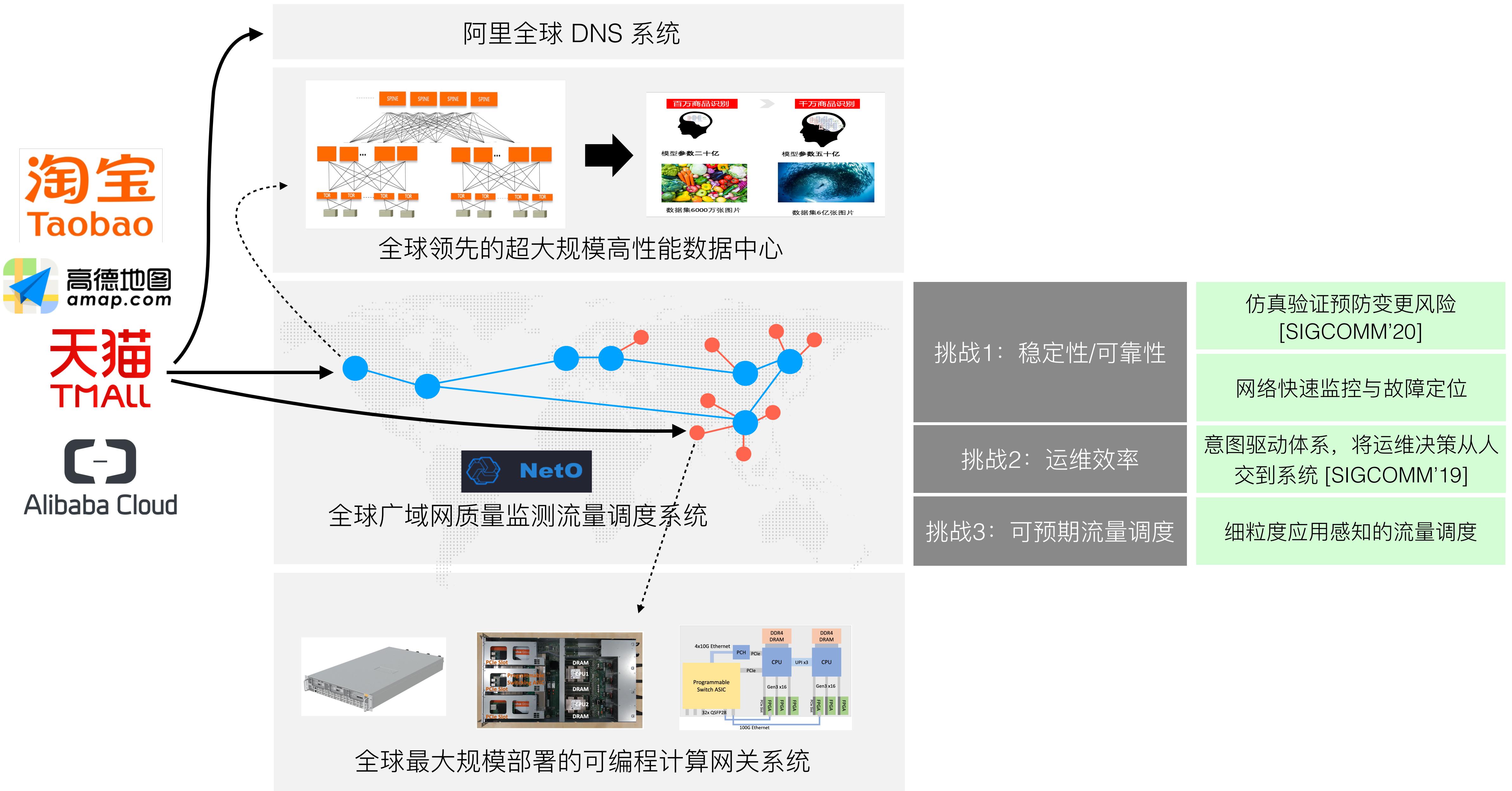


挑战1：稳定性/可靠性

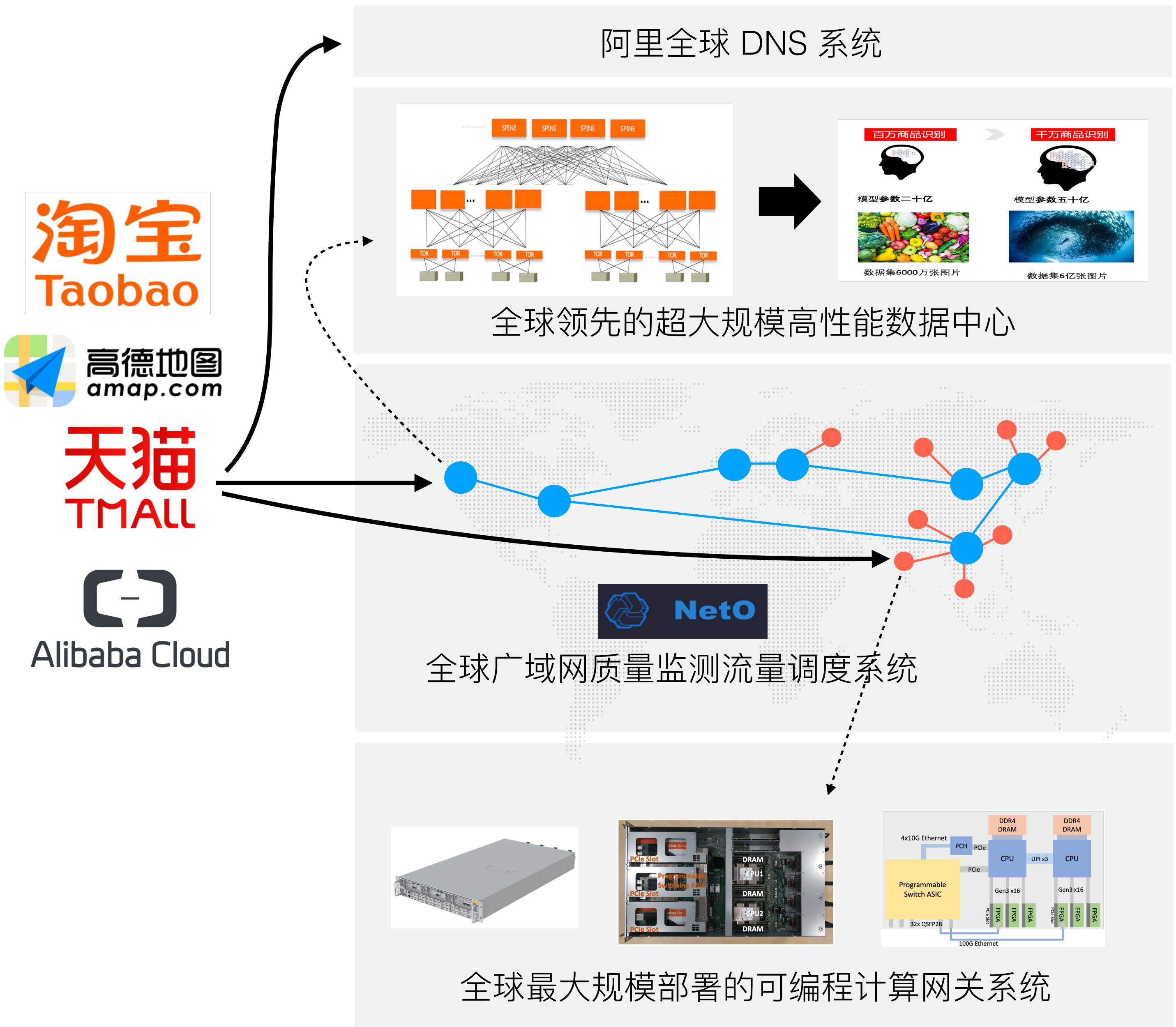
挑战2：运维效率

挑战3：可预期流量调度

阿里云全球基础设施网络系统



阿里云全球基础设施网络系统



挑战1：弹性可扩展

挑战2：稳定性/可靠性

挑战3：高性能算力

挑战1：稳定性/可靠性

挑战2：运维效率

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仿真验证预防变更风险

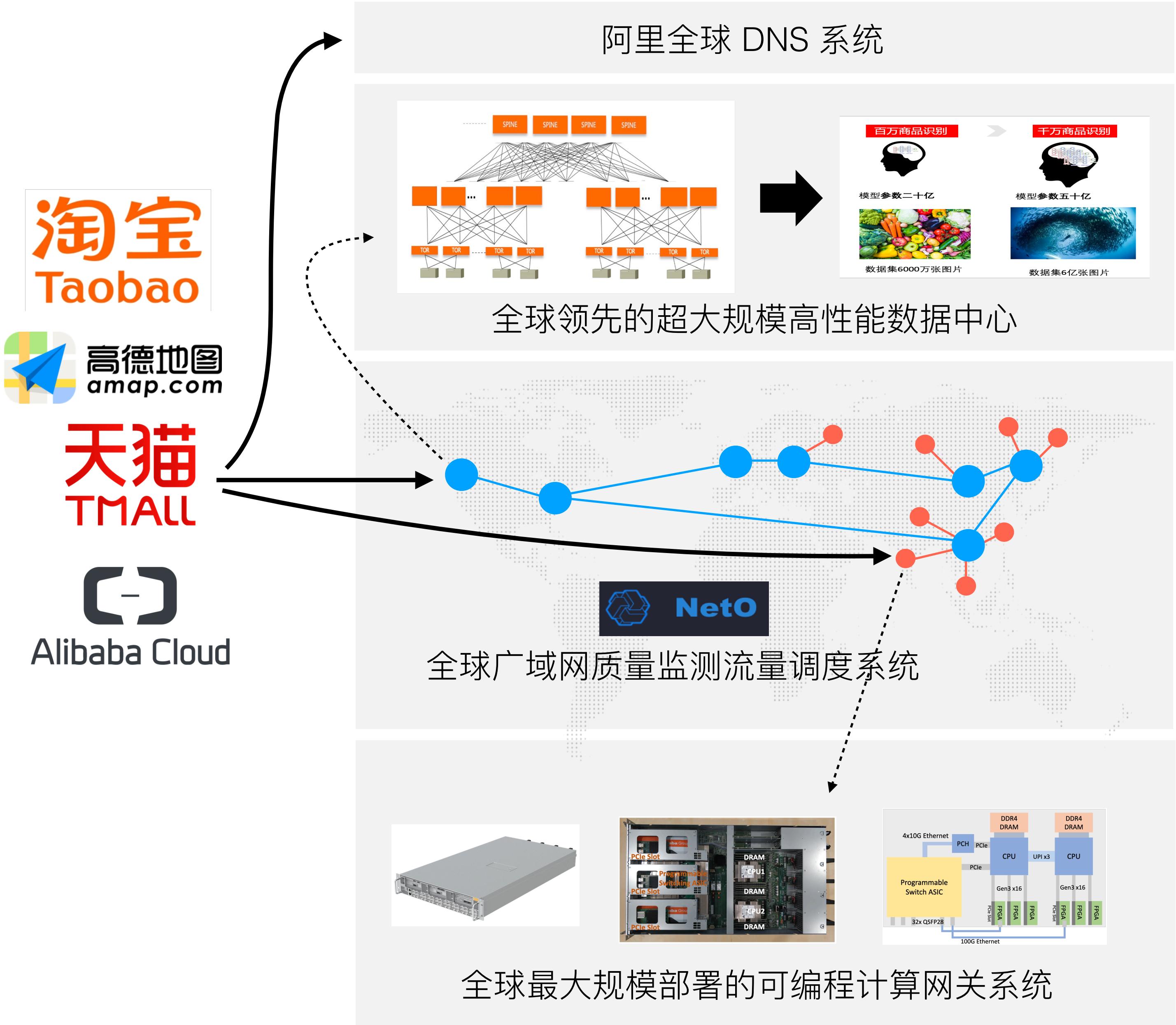
[SIGCOMM'20]

网络快速监控与故障定位

意图驱动体系，将运维决策从人
交到系统 [SIGCOMM'19]

细粒度应用感知的流量调度

阿里云全球基础设施网络系统



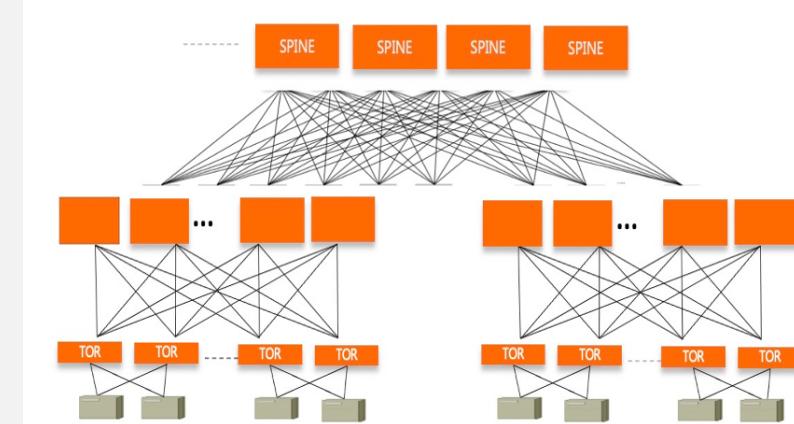
淘宝
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高德地图
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天猫
TMALL

Alibaba Cloud

阿里全球 DNS 系统



全球领先的超大规模高性能数据中心

挑战1：弹性可扩展

自研交换机，协议简化，体系化融合

挑战2：稳定性/可靠性

实时遥测精准定位故障 [SIGCOMM'20]

挑战3：高性能算力

精确反馈拥塞控制、端网协同低延时 [SIGCOMM'19, 22]

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仿真验证预防变更风险 [SIGCOMM'20]

挑战2：运维效率

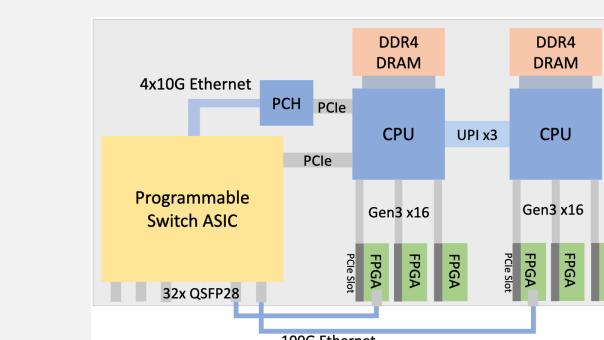
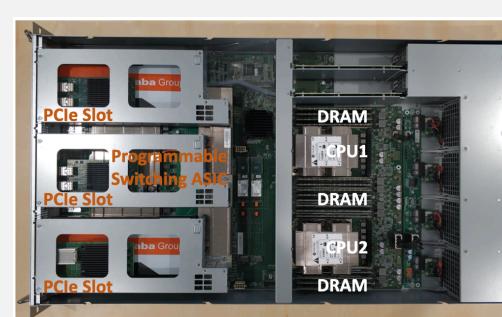
意图驱动体系，将运维决策从人交到系统 [SIGCOMM'19]

挑战3：可预期流量调度

细粒度应用感知的流量调度

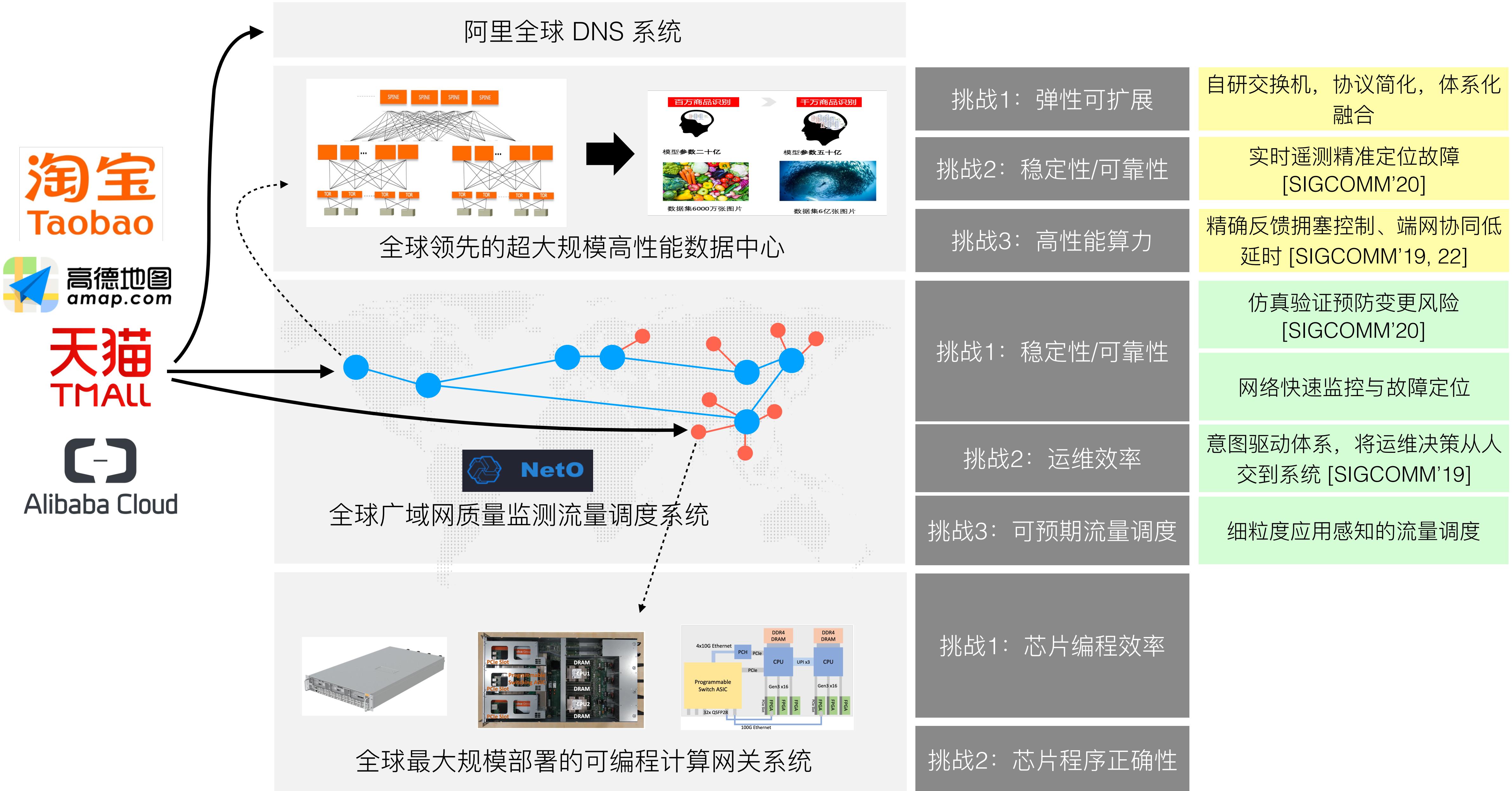
NetO

全球广域网质量监测流量调度系统

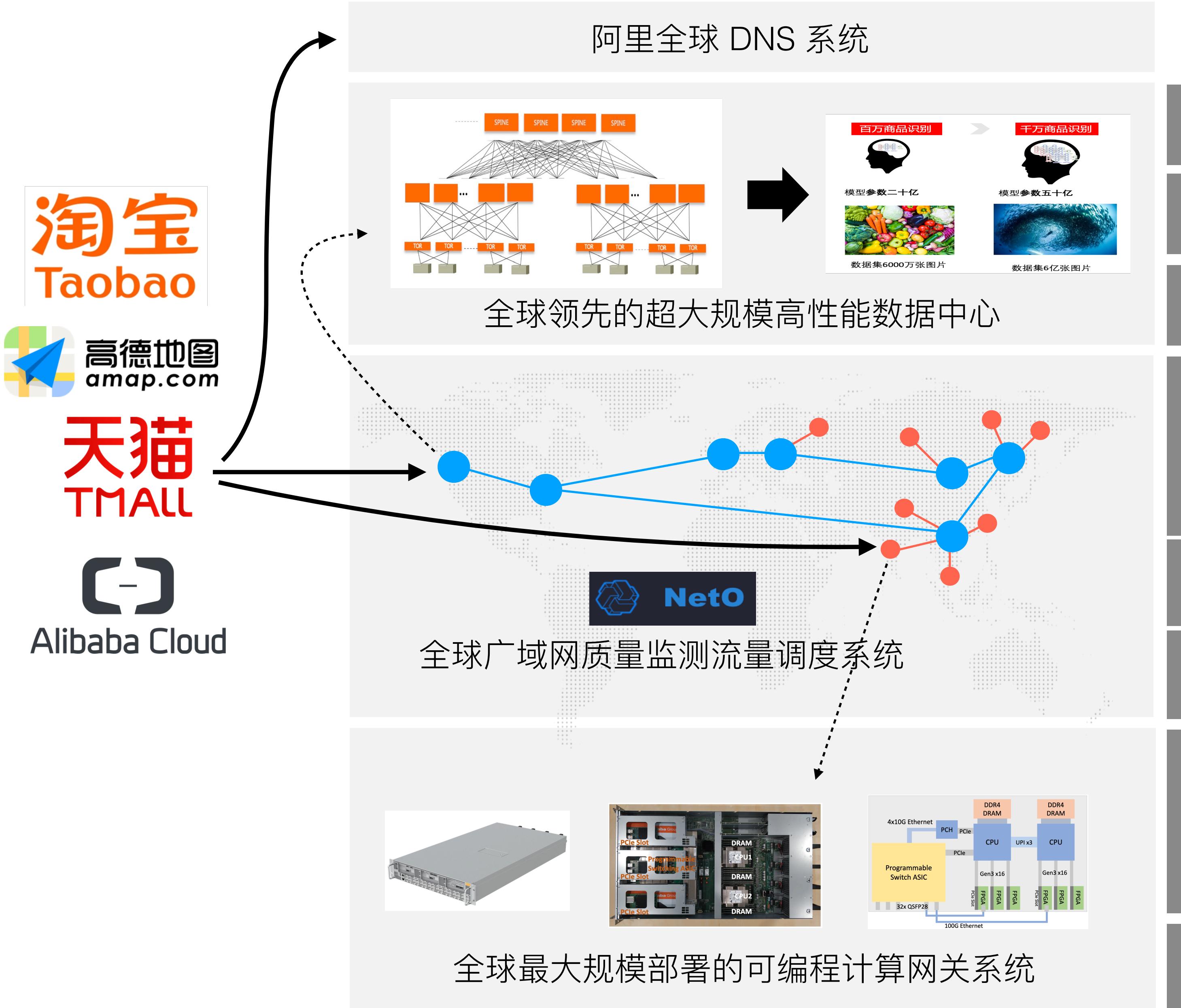


全球最大规模部署的可编程计算网关系统

阿里云全球基础设施网络系统



阿里云全球基础设施网络系统



挑战1：弹性可扩展

自研交换机，协议简化，体系化融合

挑战2：稳定性/可靠性

实时遥测精准定位故障 [SIGCOMM'20]

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挑战3：可预期流量调度

细粒度应用感知的流量调度

挑战1：芯片编程效率

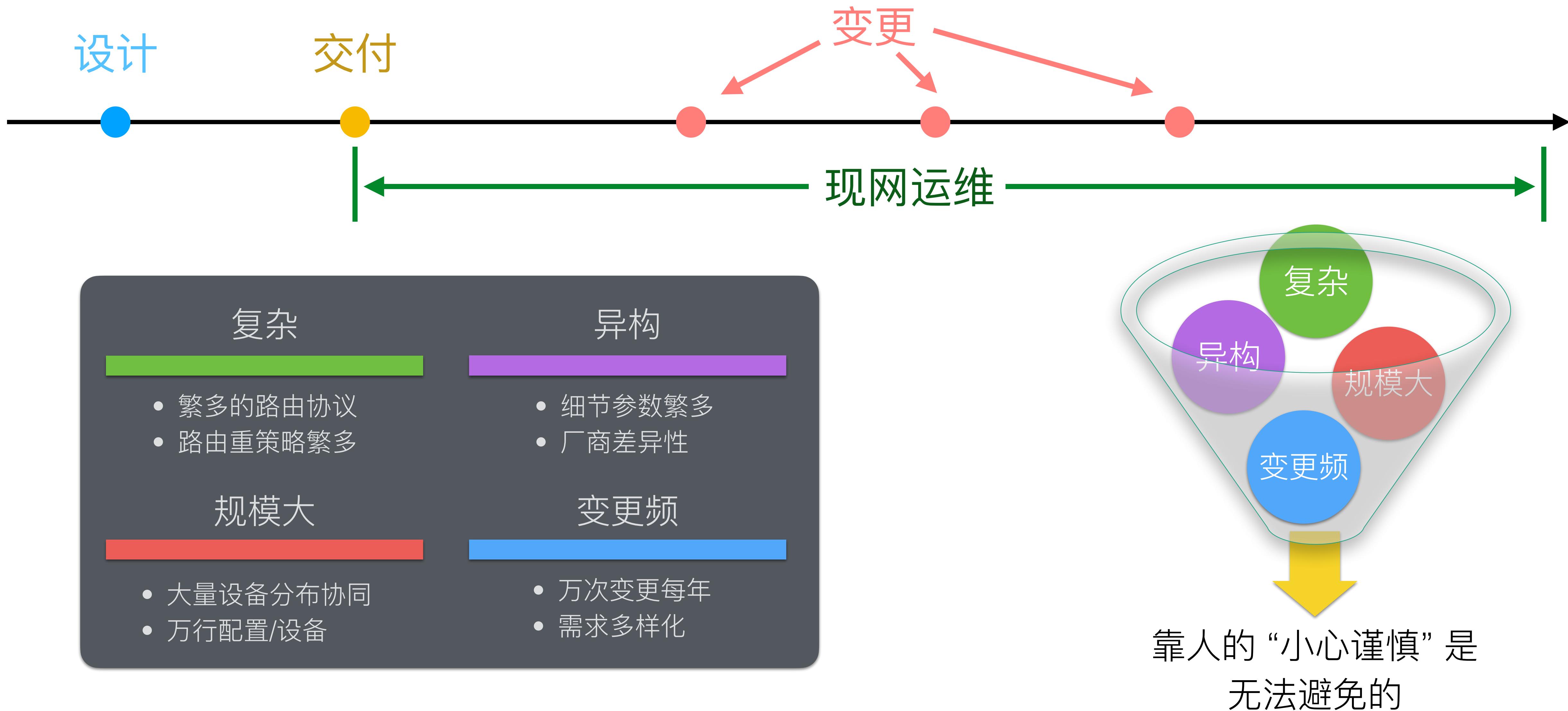
Lyra 异构芯片的程序编译 [SIGCOMM'20]

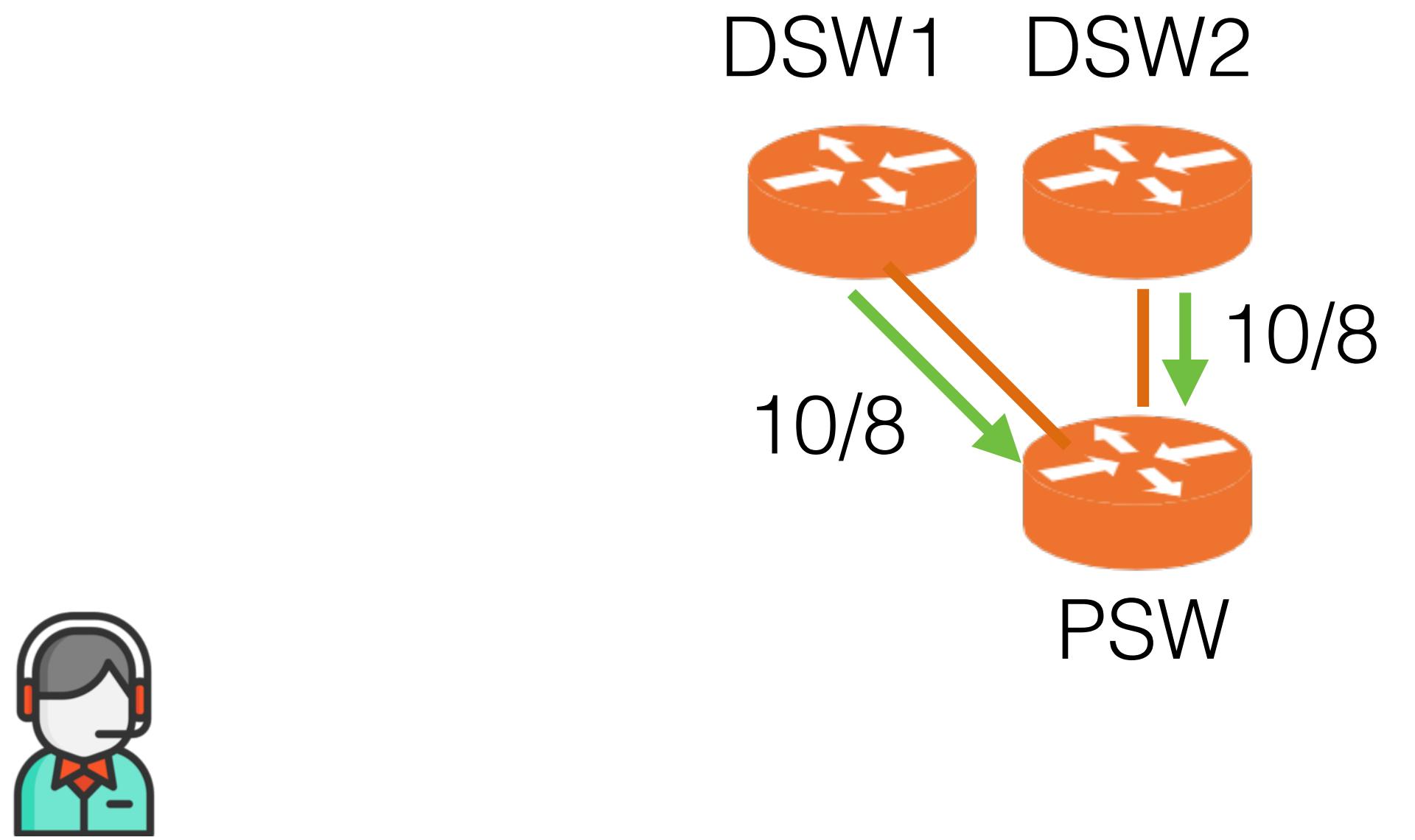
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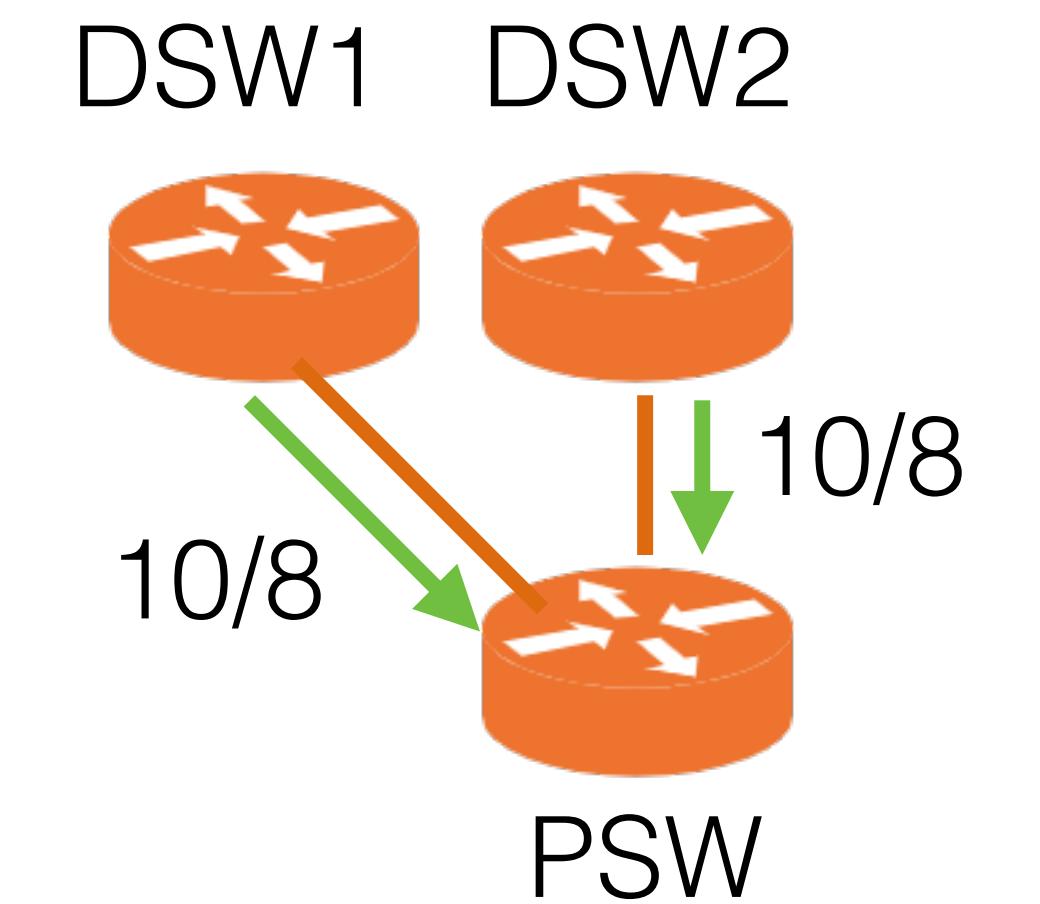
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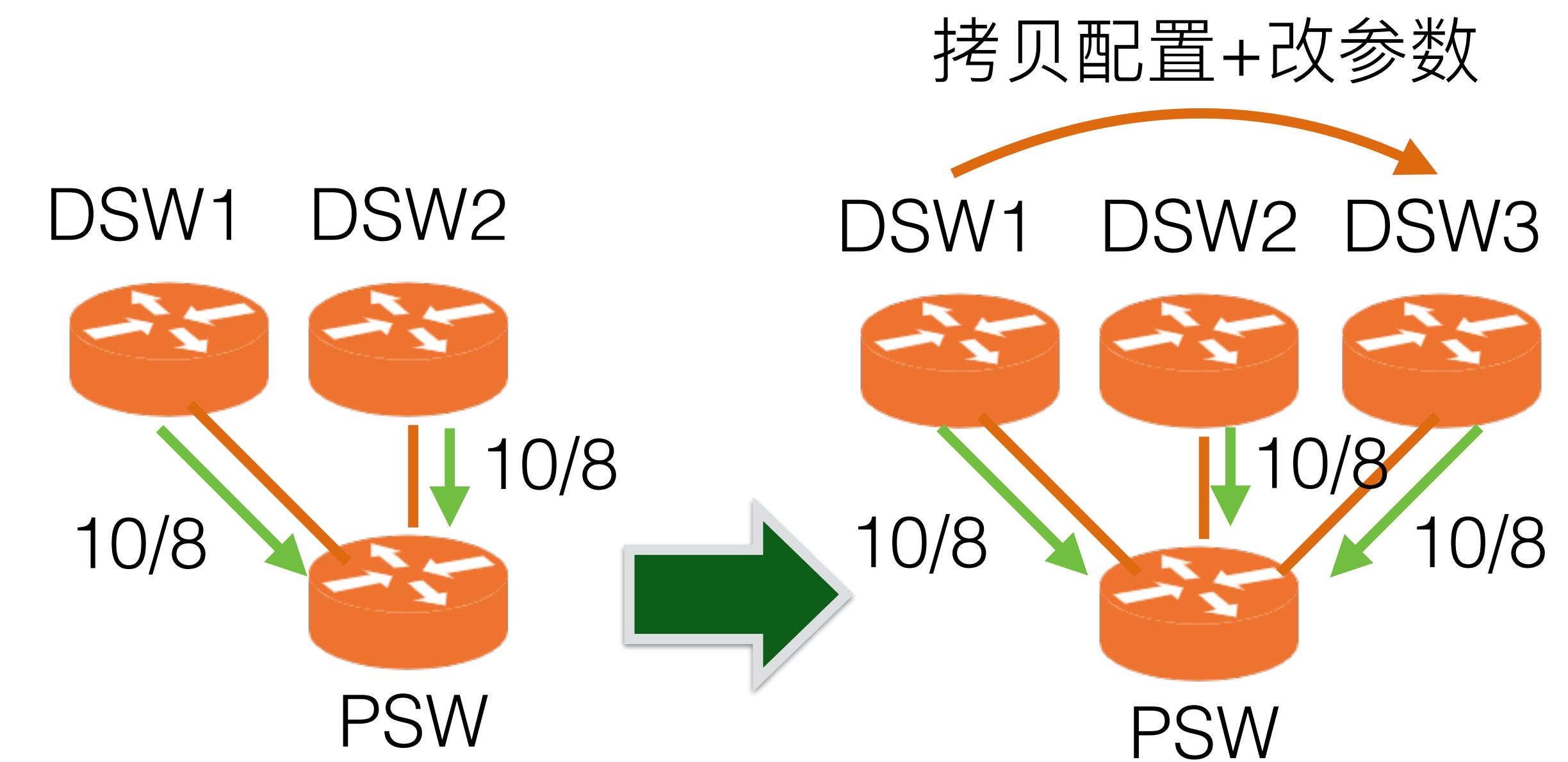
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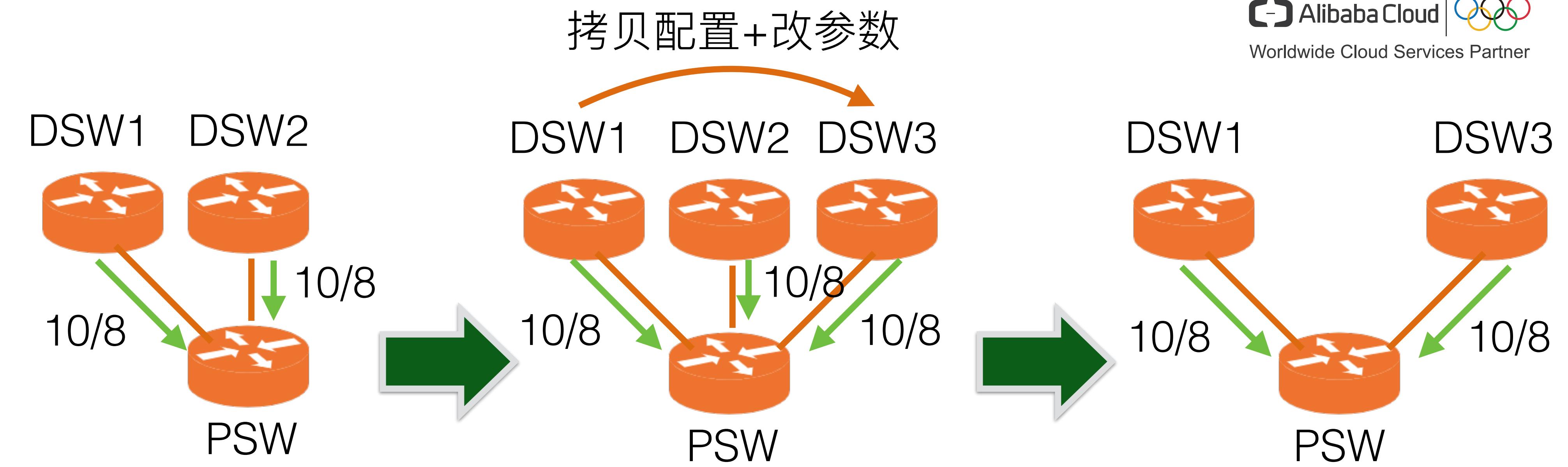
网络规划和运维的复杂







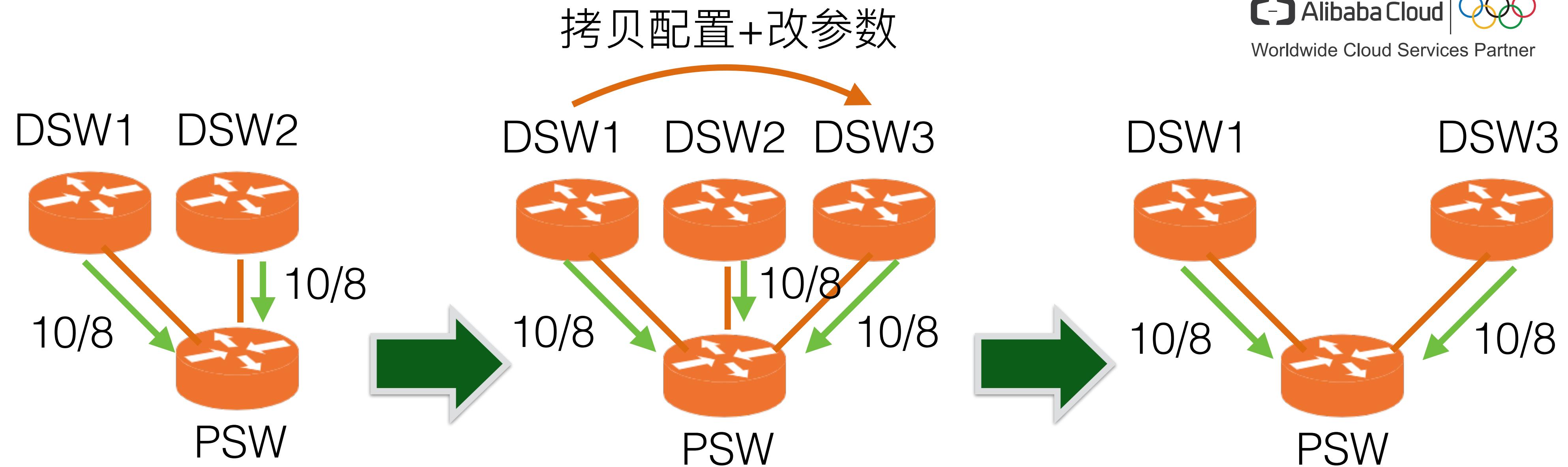




意图：
DSW3 替换 DSW2



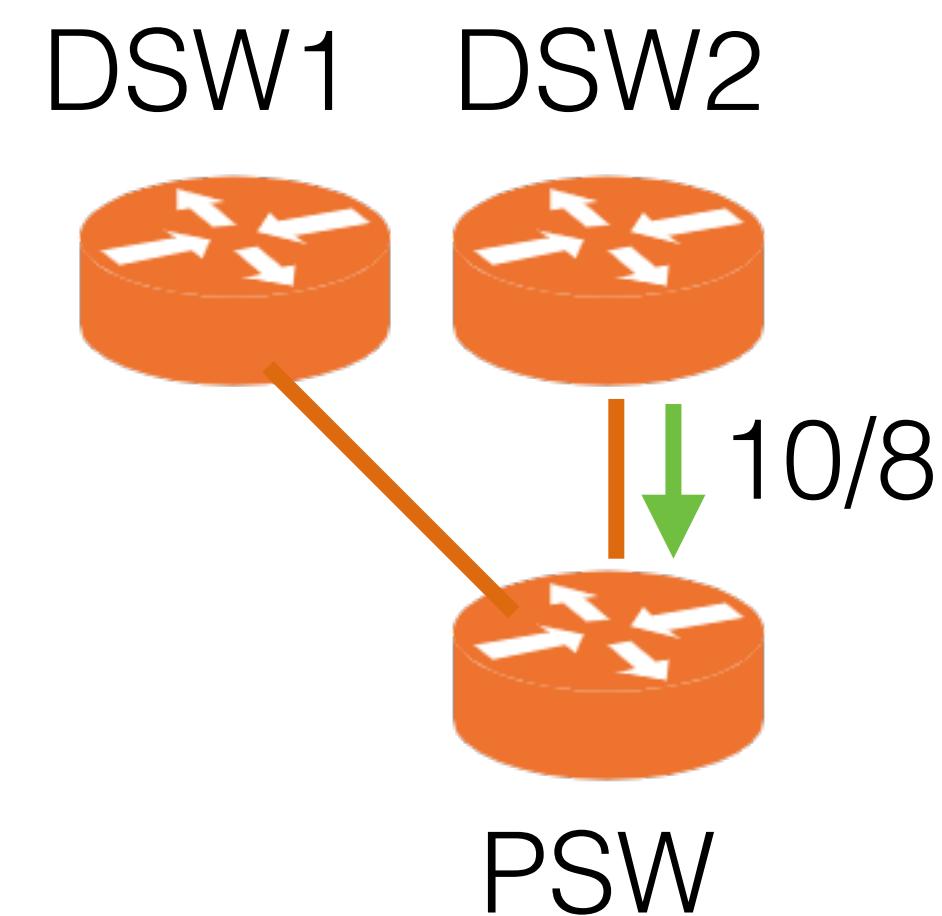
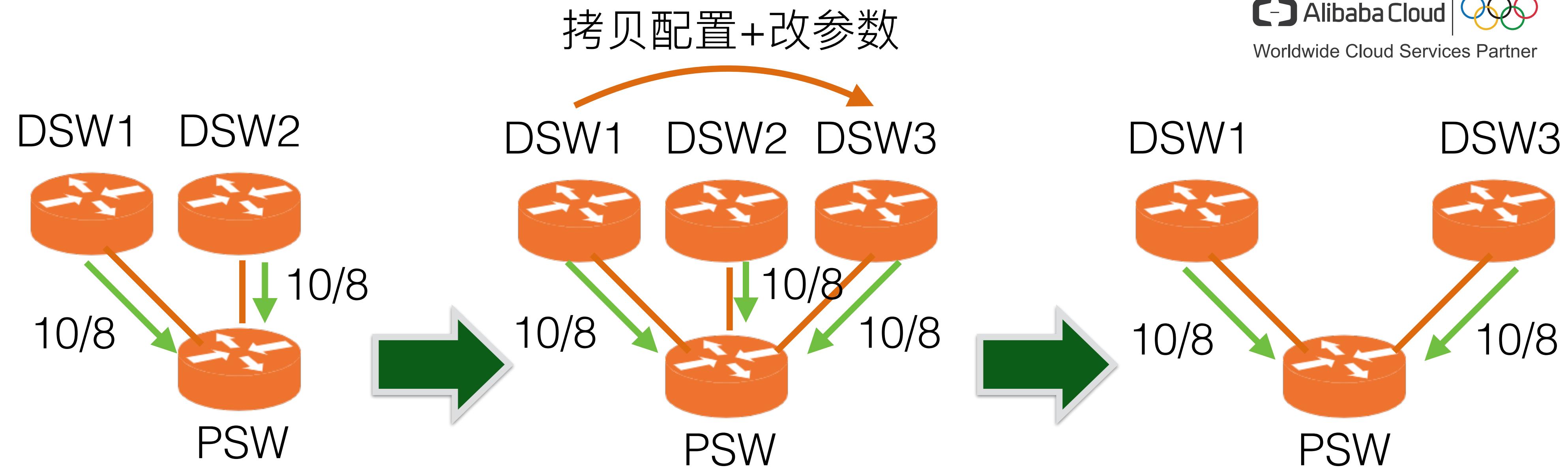
我们认为
实际情况



意图：
DSW3 替换 DSW2



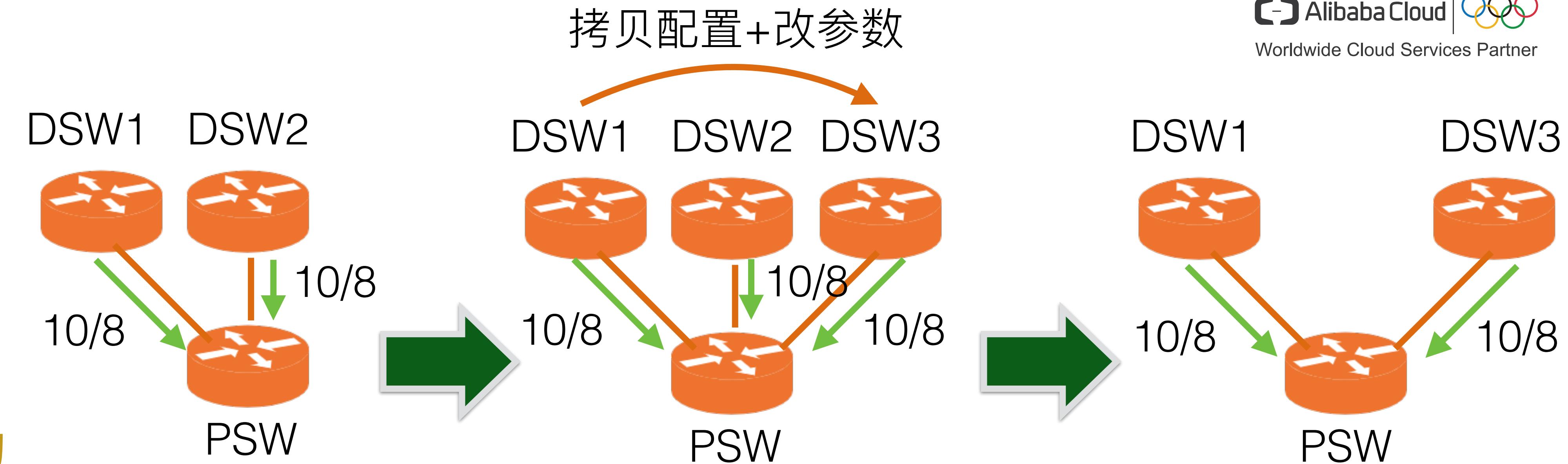
我们认为
实际情况



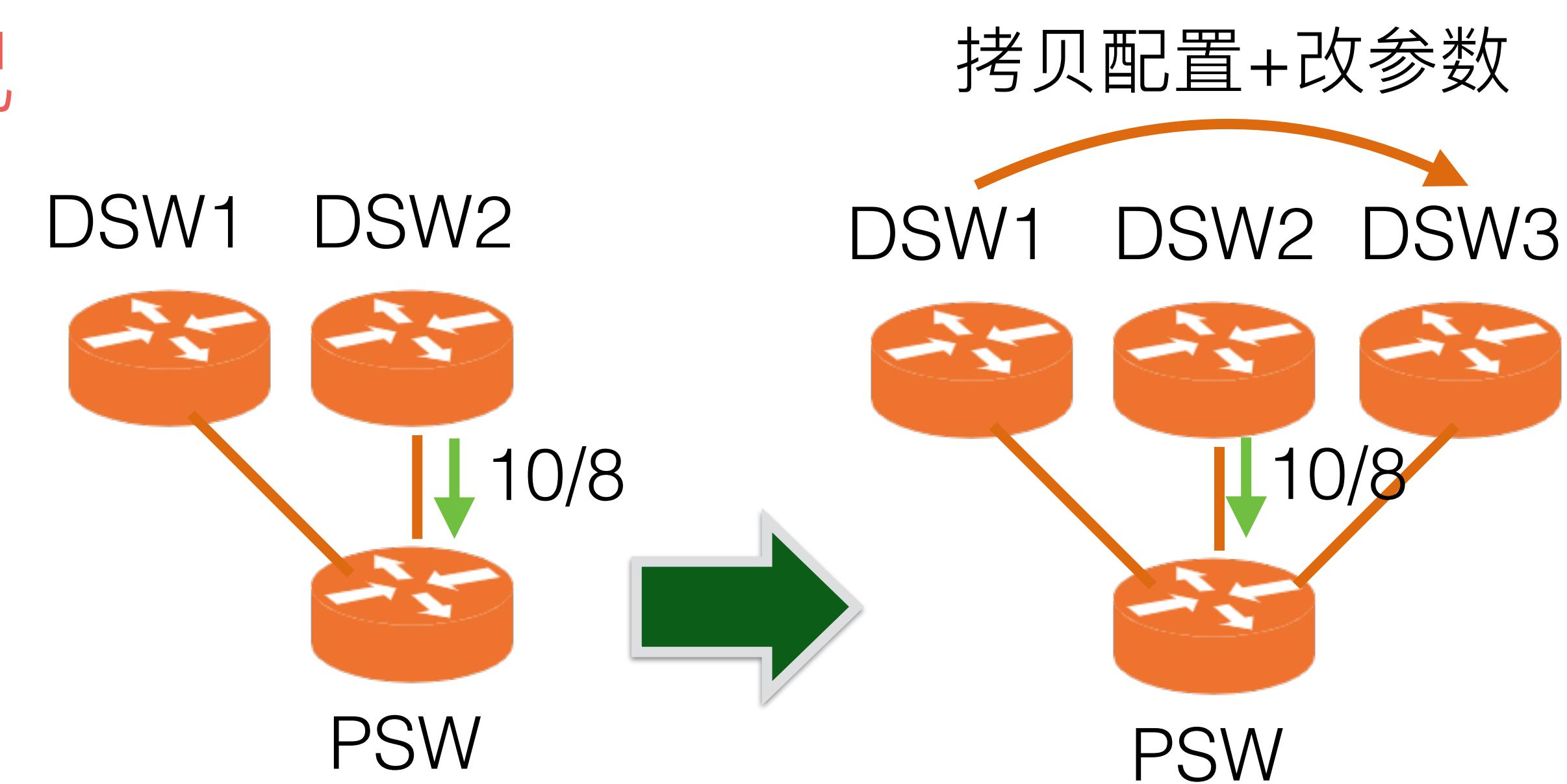
意图：
DSW3 替换 DSW2



我们认为
实际情况



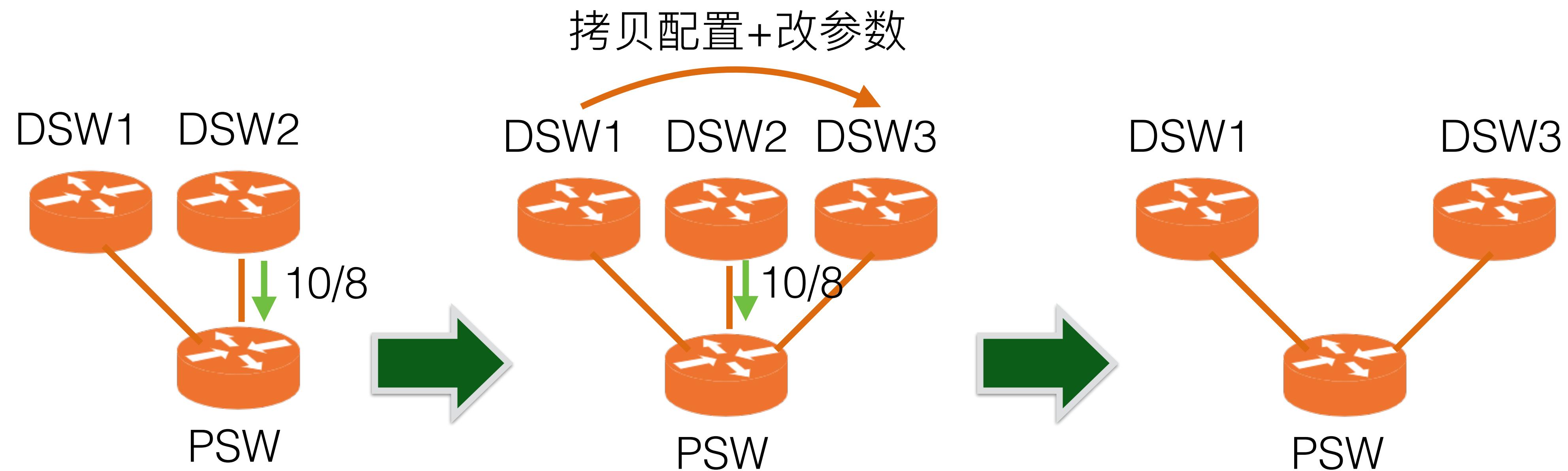
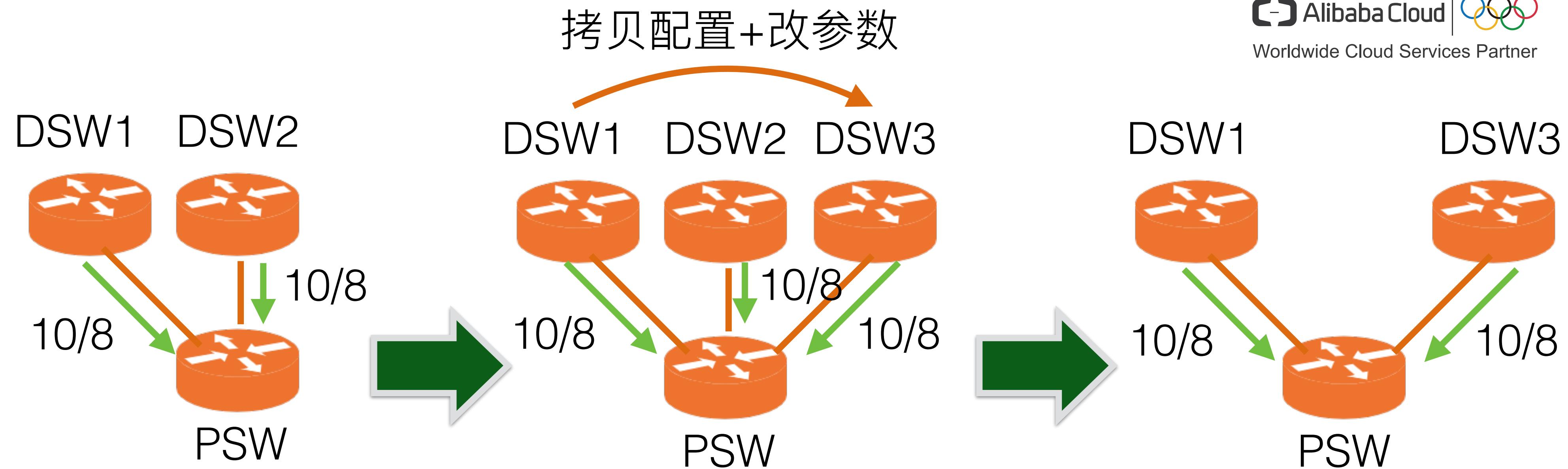
拷贝配置+改参数



意图：
DSW3 替换 DSW2



我们认为
实际情况



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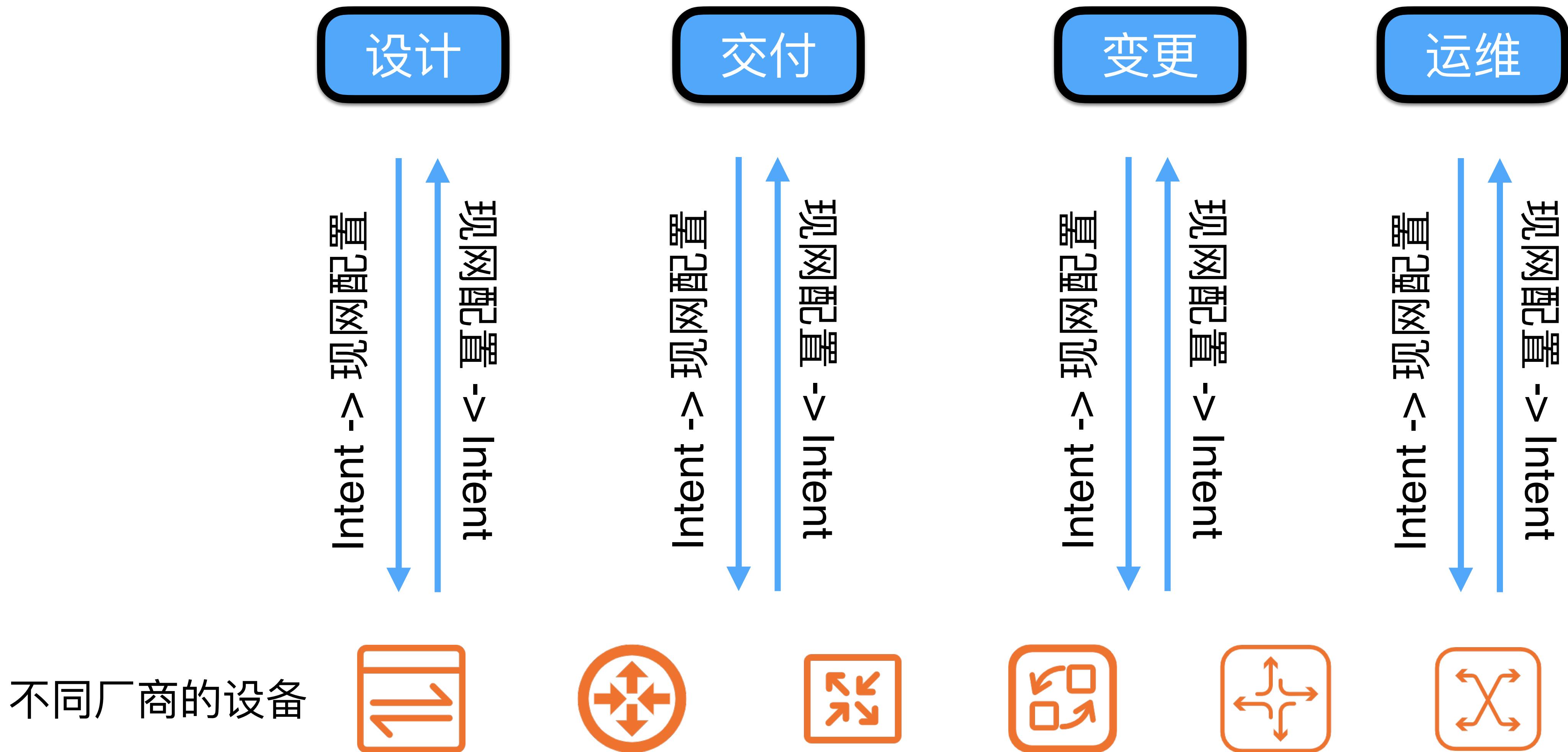
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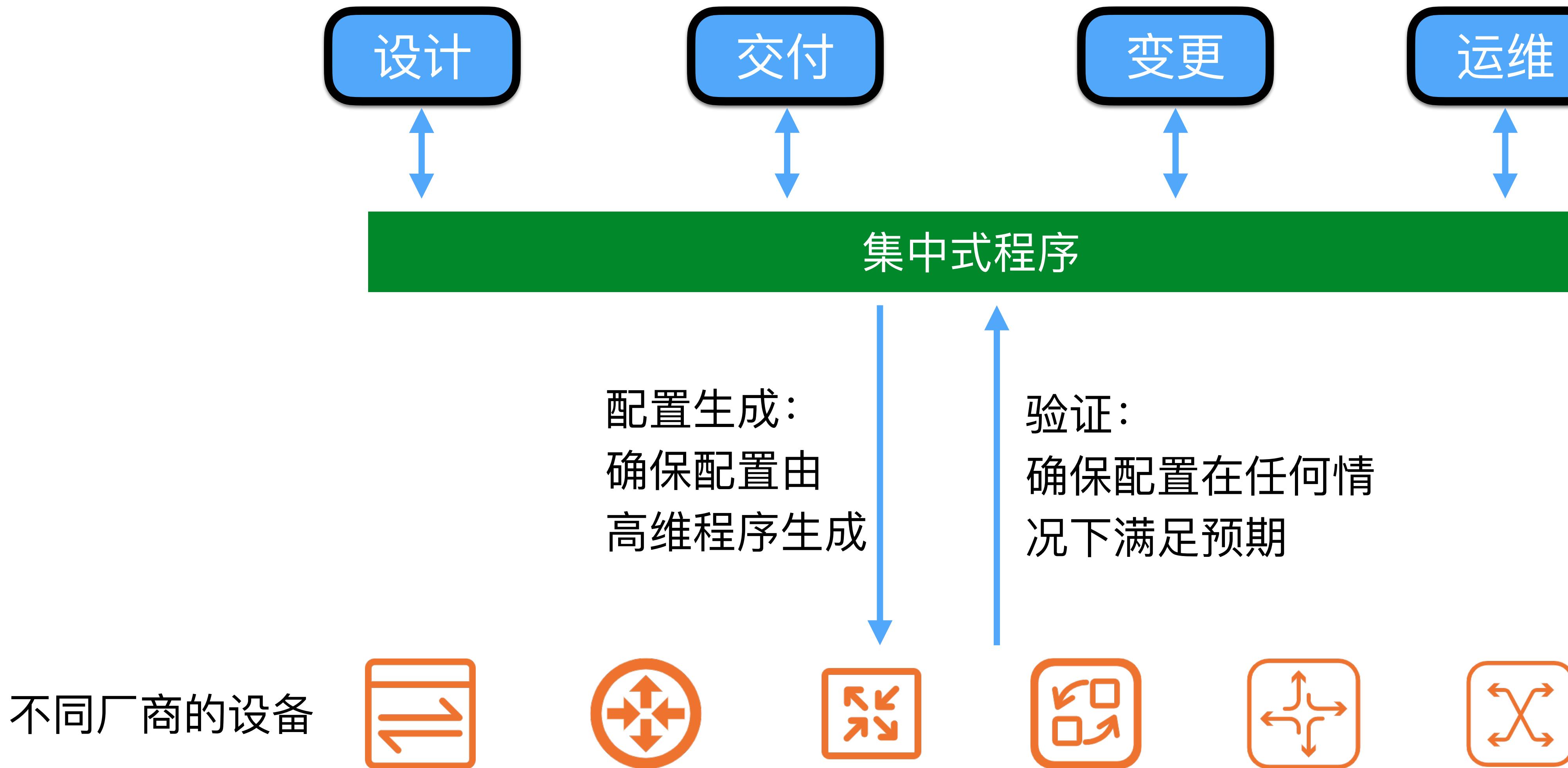
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传统基于配置的思想



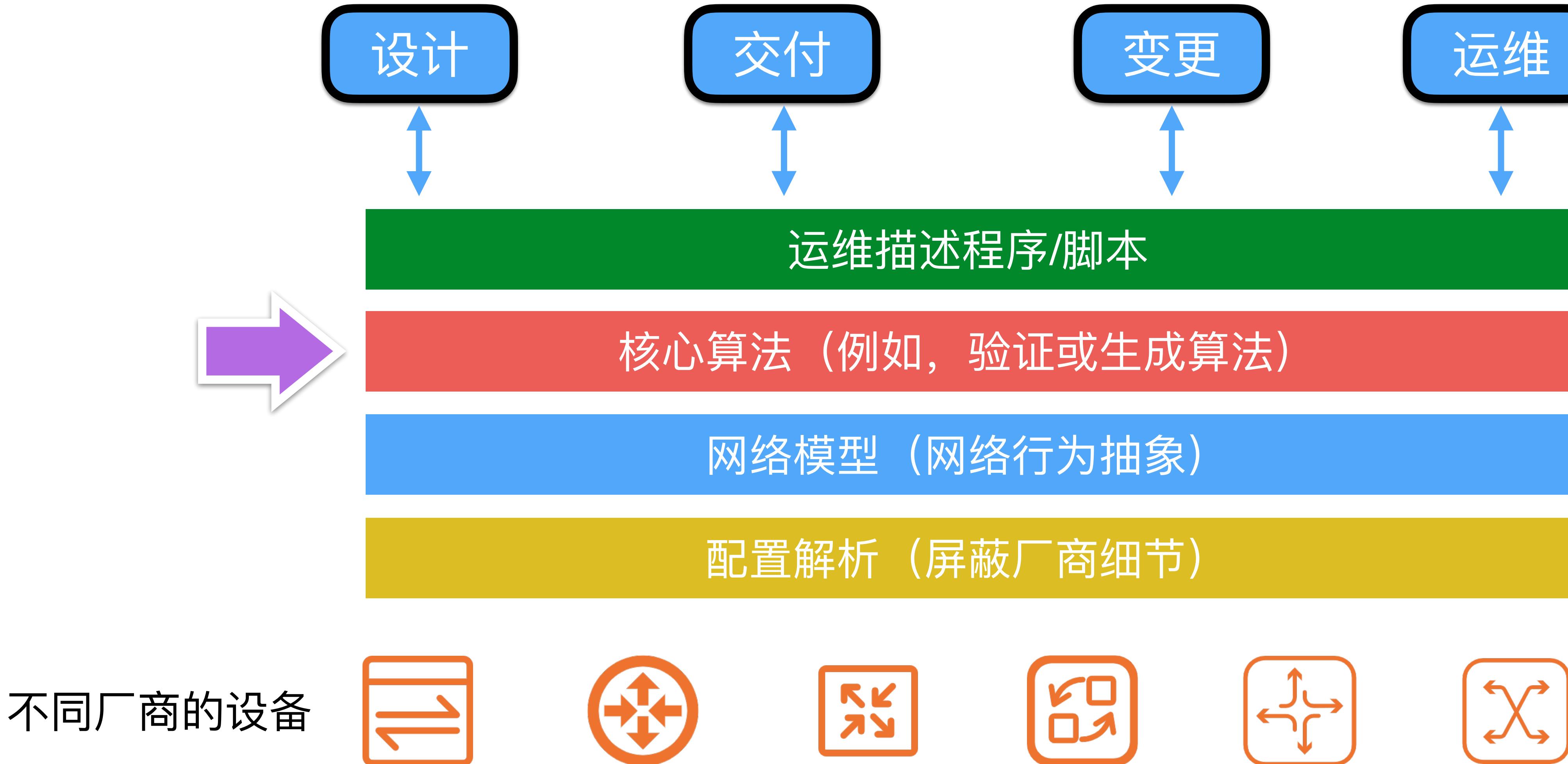
IBN 思想：高级抽象



IBN 思想



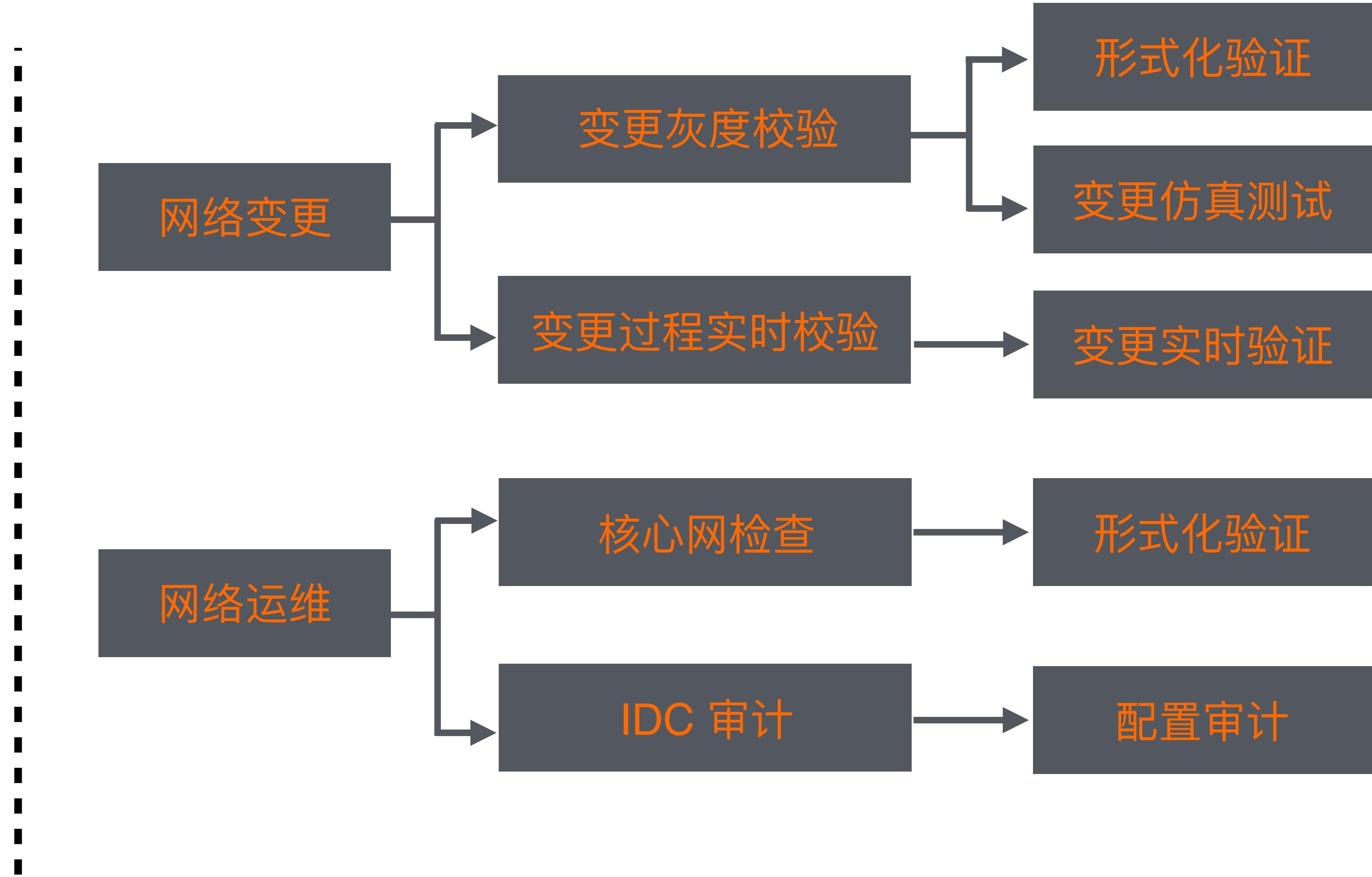
集中式网络规划体系



核心技术之一：验证



挑战1：有多准?
挑战2：有多快?



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Network configuration verification in Alibaba



Monkey King: Fire & Golden Eye (The ability to see the truth)

Accuracy, Scalability, Coverage - A Practical Configuration Verifier on a Global WAN
Fangdan Ye (Tsinghua University); Da Yu (Brown University); Ennan Zhai, Hongqiang Harry Liu (Alibaba Group); Bingchuan Tian (Nanjing University); Qiaobo Ye, Chunsheng Wang, Xin Wu, Tianchen Guo, Cheng Jin, Duncheng She, Qing Ma, Biao Cheng, Hui Xu, Ming Zhang (Alibaba Group); Zhiliang Wang (Tsinghua University); Rodrigo Fonseca (Brown University)

3:30 - 3:45 pm EDT
2:30 - 2:45 am EDT

Abstract: This paper presents Hoyan-- the first reported large scale deployment of configuration verification in a global-scale wide area network (WAN). Hoyan has been running in production for more than two years and is currently used for all critical configuration auditing and updates on the WAN. We highlight our innovative designs and real-life experience to make Hoyan accurate and scalable in practice. For accuracy under the inconsistencies of devices' vendor-specific behaviors (VSBs), Hoyan continuously discovers the flaws in device behavior models, thus aiding the operators in fixing the models. For scalability to verify our global WAN, Hoyan introduces a "global-simulation & local formal-modeling" strategy to model uncertainties in small scales and perform aggressive pruning of possibilities during the protocol simulations. Hoyan achieves near-100% verification accuracy after it detected and fixed 0.01 VSBs on our WAN. Hoyan has prevented many potential service failures resulting from misconfiguration and reduced the failure rate of updates of our WAN by more than half in 2019.

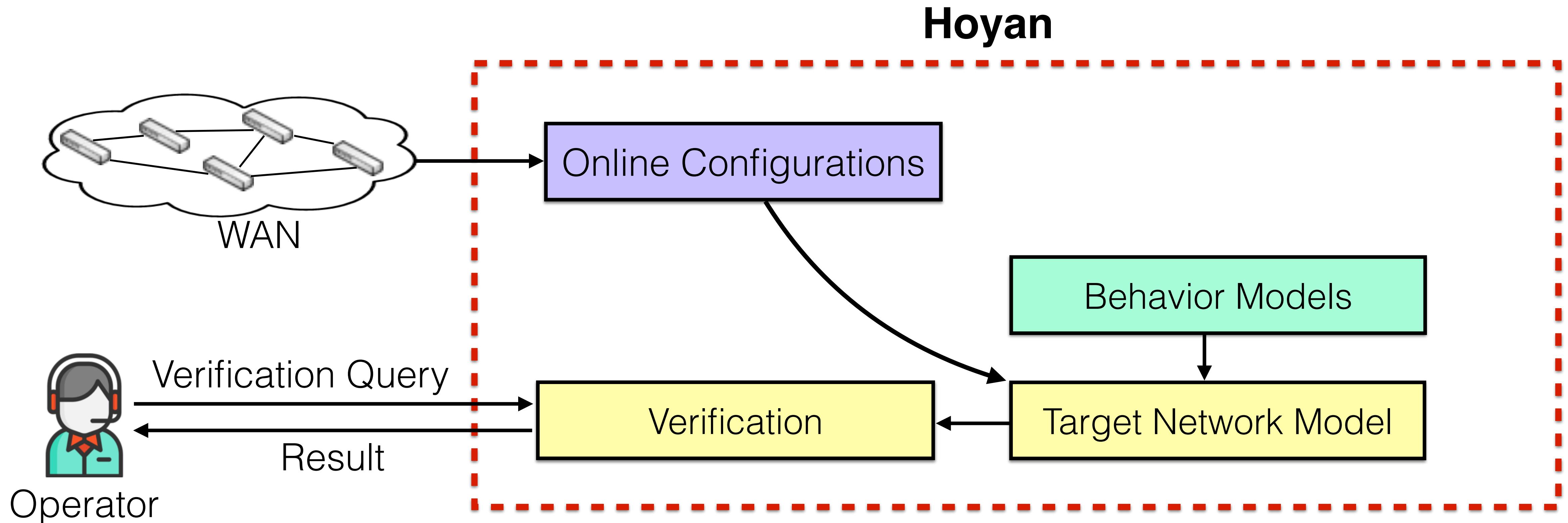
Safely and Automatically Updating In-Network ACL Configurations with Intent Language
Bingchuan Tian (Nanjing University), Xinyi Zhang (University of California Santa Barbara), Ennan Zhai, Hongqiang Harry Liu, Qiaobo Ye, Chunsheng Wang, Xin Wu, Zhiming Ji, Yihong Sang, Ming Zhang (Alibaba Group), Da Yu (Brown University), Chen Tian (Nanjing University), Haitao Zheng, Ben Y. Zhao (University of Chicago)

Abstract: In-network Access Control List (ACL) is an important technique in ensuring network-wide connectivity and security. As cloud-scale WANs today constantly evolve in size and complexity, in-network ACL rules are becoming increasingly more complex. This presents a great challenge to the updating process of ACL configurations: network operators are frequently required to update "tangled" ACL rules across thousands of devices to meet diverse business requirements, and even a single ACL misconfiguration may lead to network disruptions. Such increasing challenges call for an automated system to improve the efficiency and correctness of ACL updates. This paper presents Jinjing, a system that aids Alibaba's network operators in automatically and correctly updating ACL configurations in Alibaba's global WAN. Jinjing allows the operators to express in a declarative language, named LAI, their update intent (e.g., ACL migration and traffic control). Then, Jinjing automatically synthesizes ACL update plans that satisfy their intent. At the heart of Jinjing, we develop a set of novel verification and synthesis techniques to rigorously guarantee the correctness of update plans. In Alibaba, our operators have used Jinjing to efficiently update their ACLs and have thus prevented significant service downtime.

Hoyan (火眼)
Routing configuration verification ←
[ACM SIGCOMM'20]



Jinjing (金睛)
In-Network ACL verification →
[ACM SIGCOMM'19]

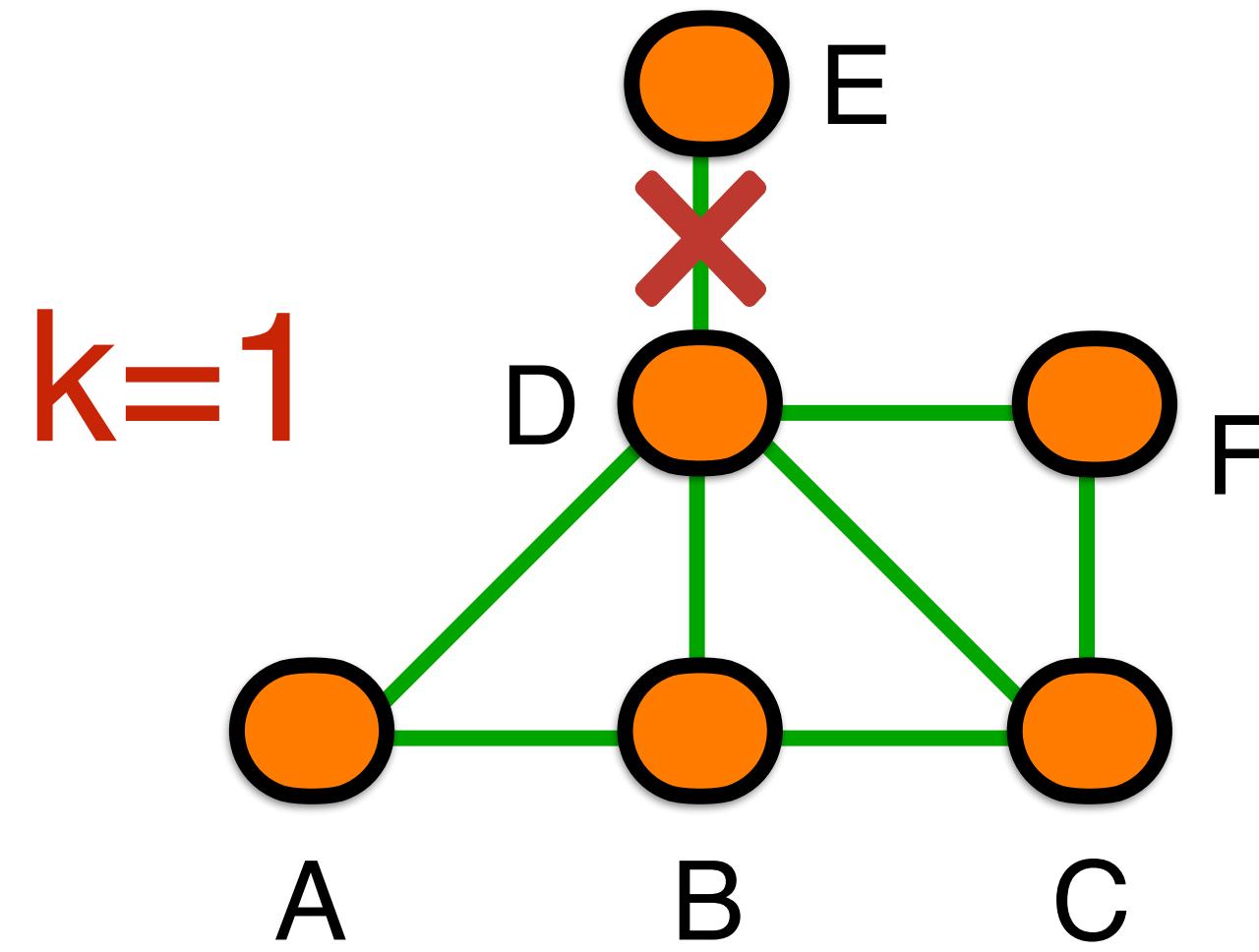


Technical challenges

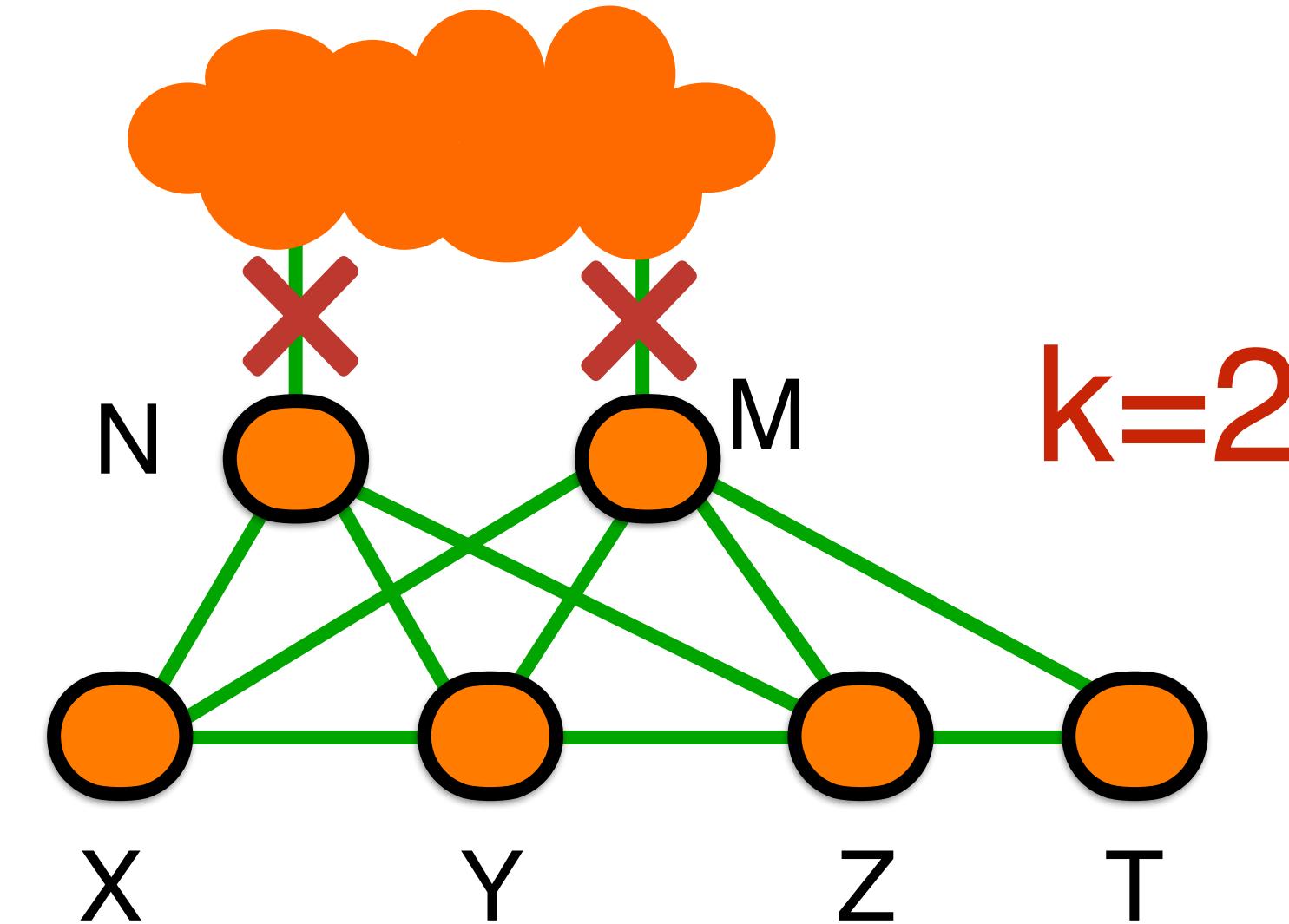
Challenge 1: Scalability with failure coverage in our global WAN

Technical challenges

Challenge 1: Scalability with failure coverage in our global WAN



$k=1$

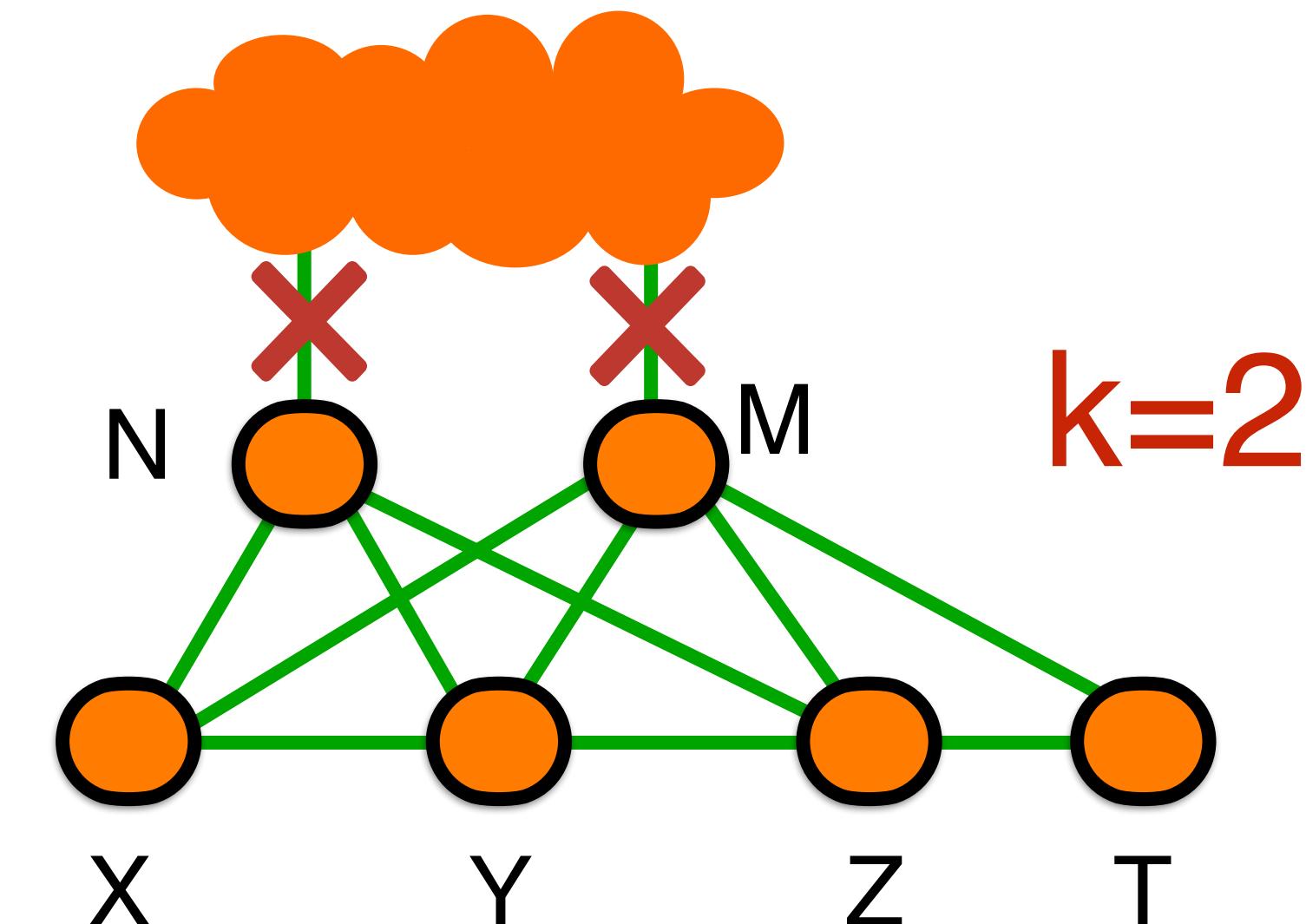
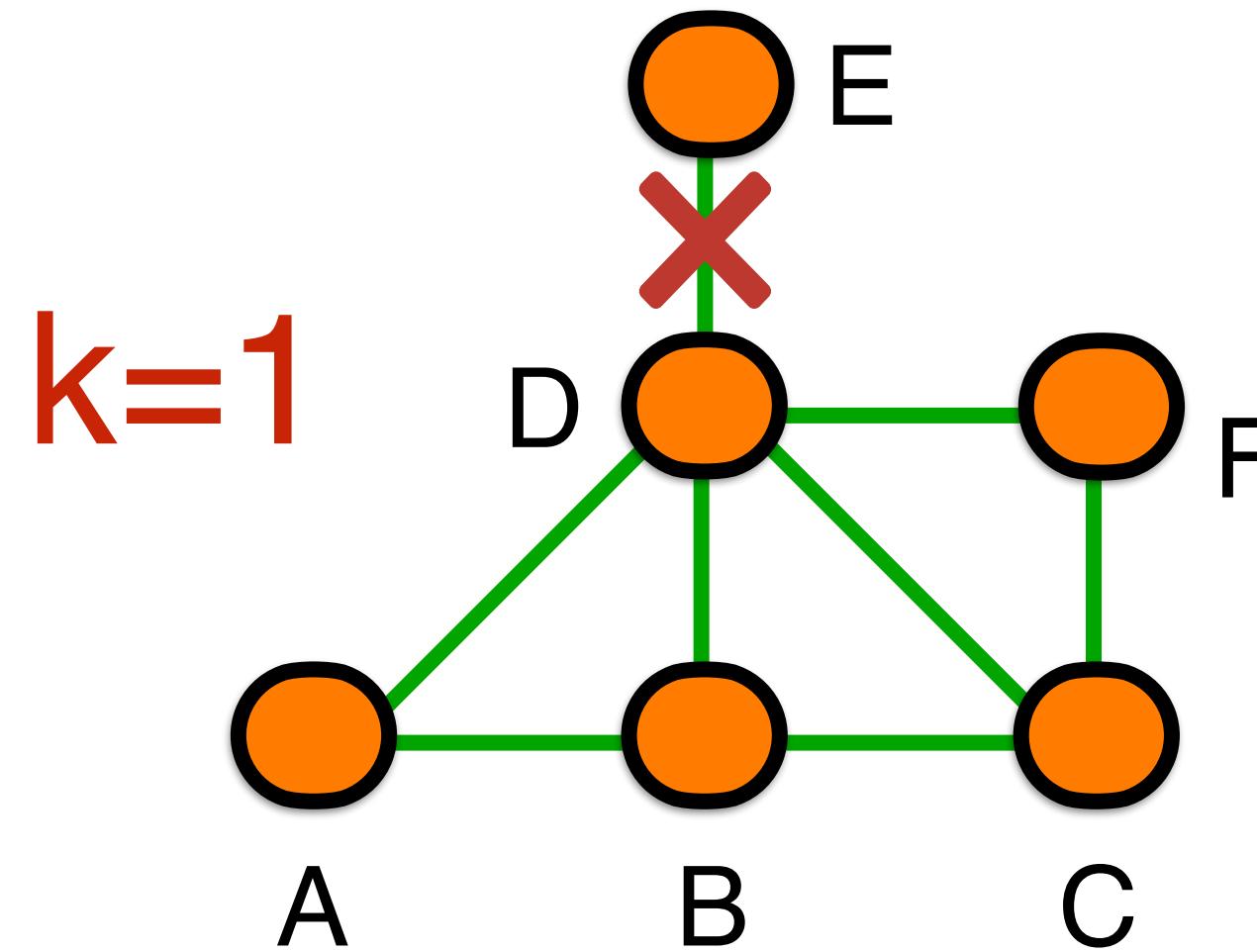


$k=2$

Technical challenges

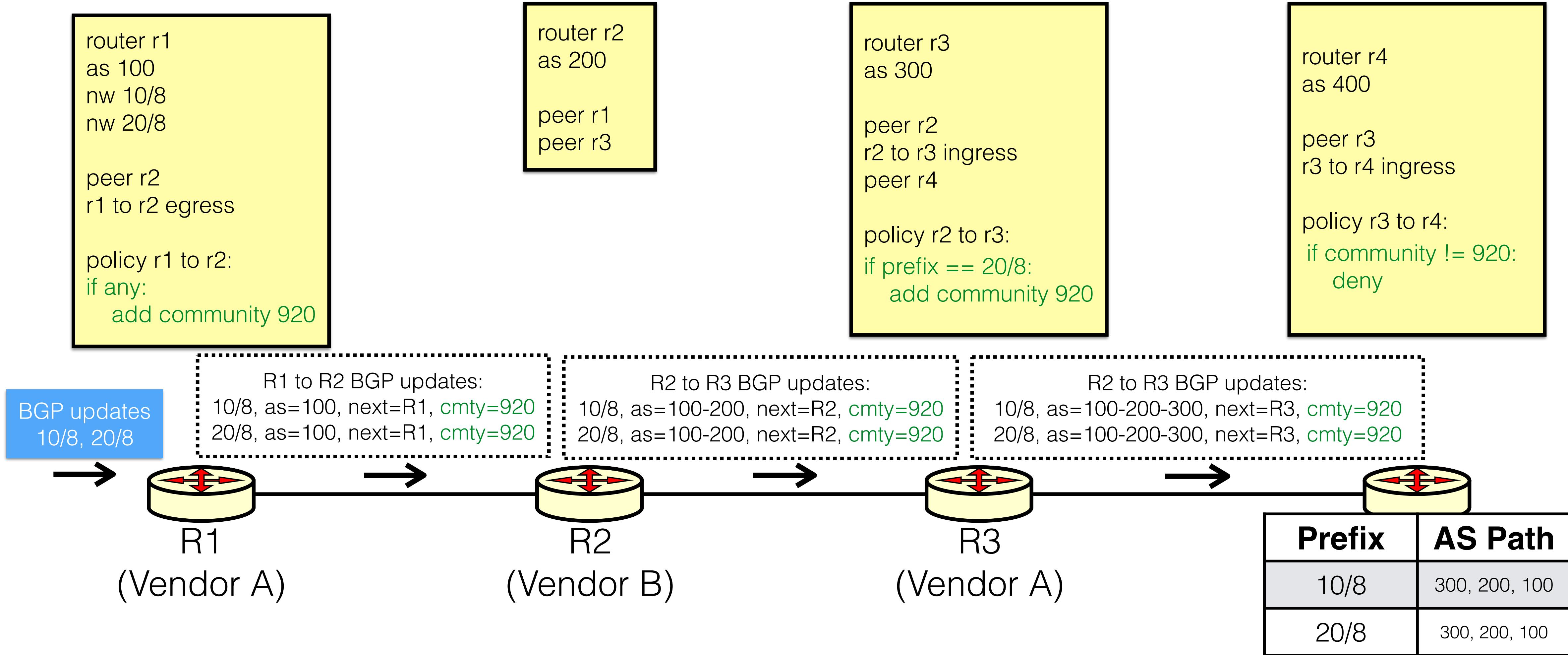
Challenge 1: Scalability with failure coverage in our global WAN

- Simulation-based verification efforts, e.g., Batfish, need to enumerate C_n^k cases for checking under arbitrary k failures out of n links
- Formal model approaches, e.g., Minesweeper, are not scalable since a large WAN generates a huge SMT formula
- Asymmetry of our WAN invalidates symmetry-based simplification approaches



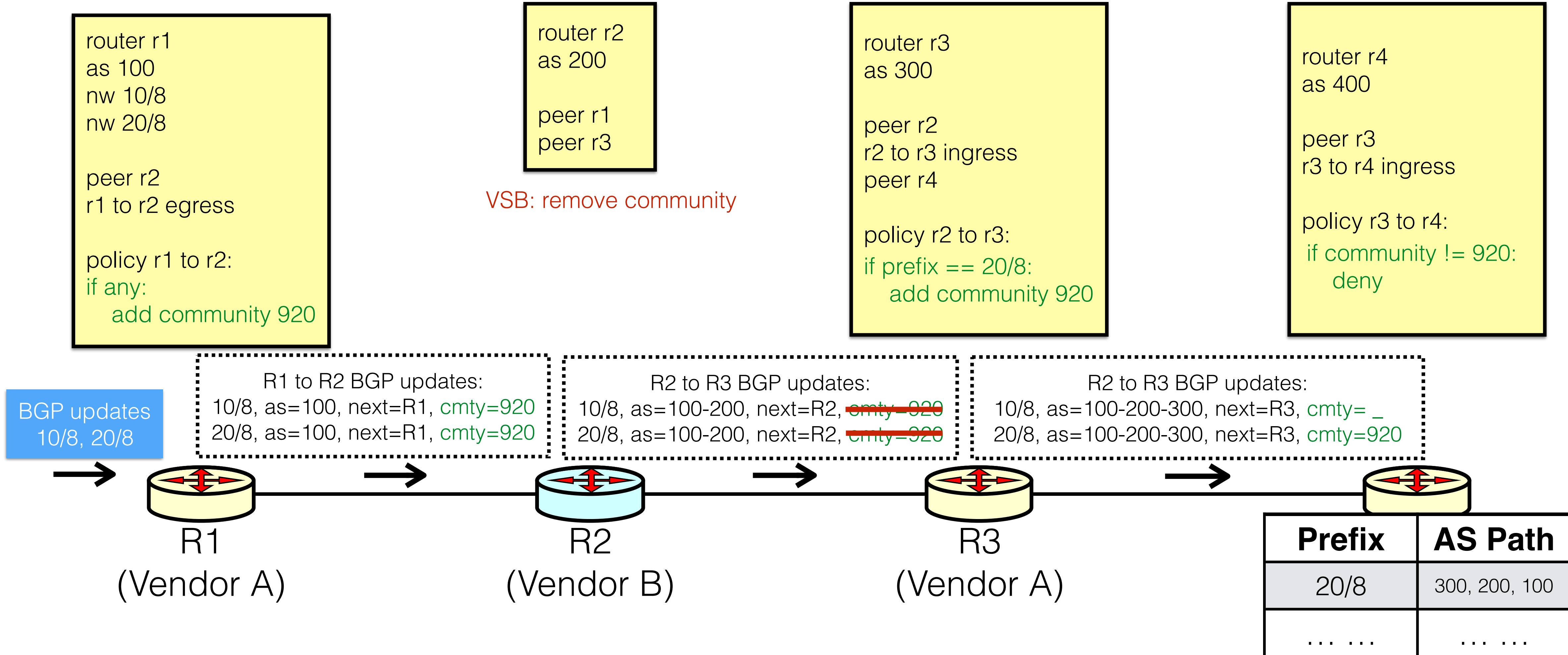
Technical challenges

Challenge 2: Faithfulness of device behavior model



Technical challenges

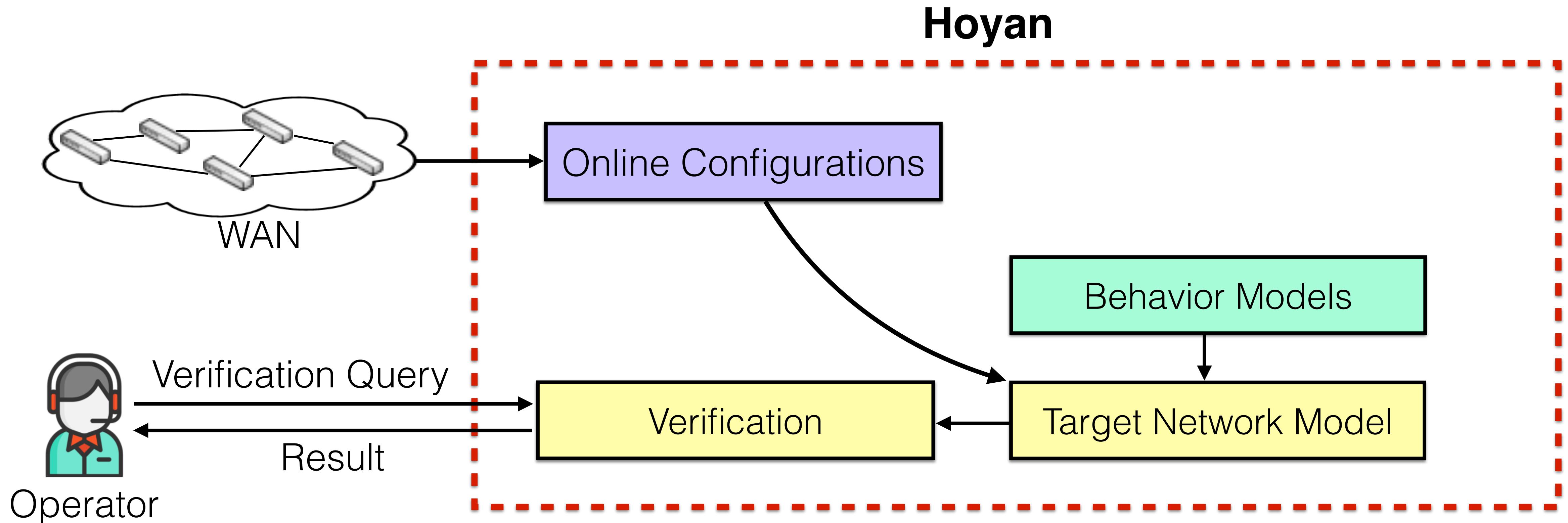
Challenge 2: Faithfulness of device behavior model

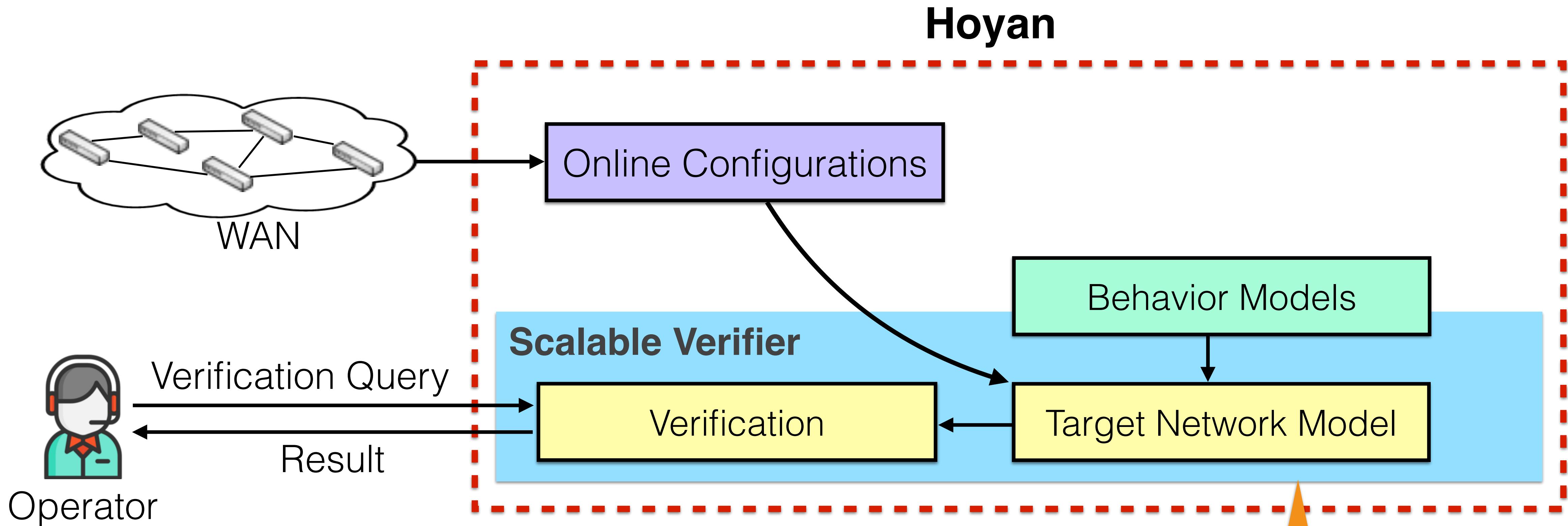


Technical challenges

Challenge 2: Faithfulness of device behavior model

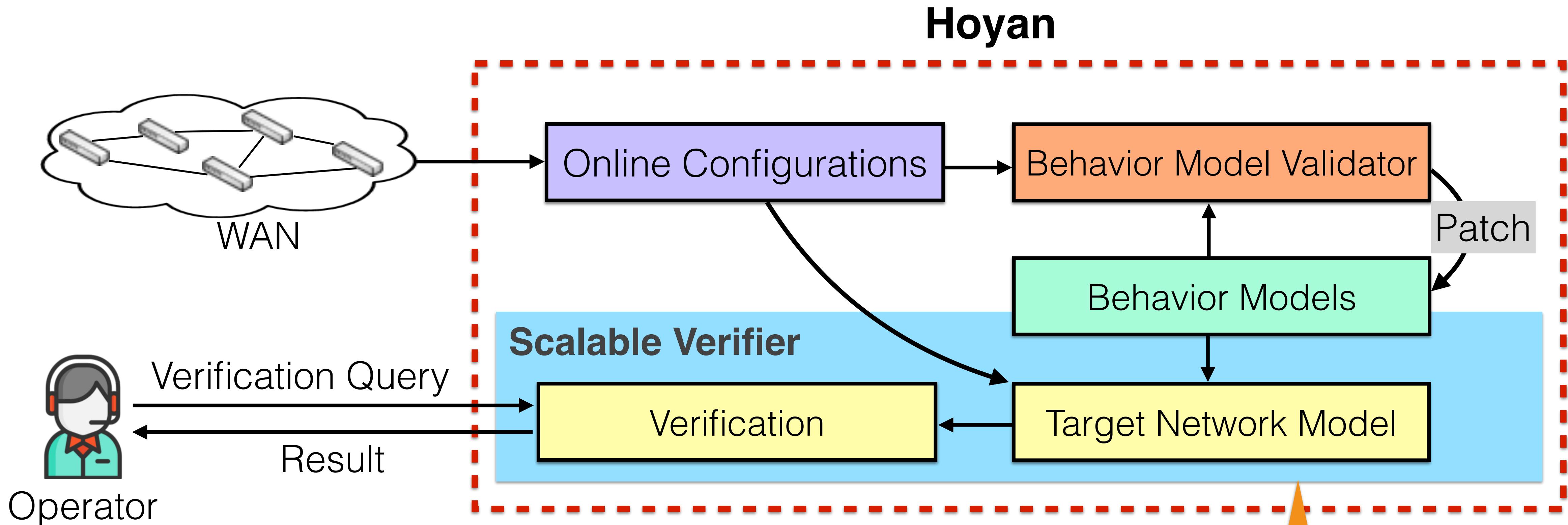
- Vendors implemented the same protocols in various ways due to the ambiguity of RFC
- Vendor-specific behaviors (VSBs) significantly affect the accuracy of network modeling
- No existing efforts can detect or tune the VSBs in network model





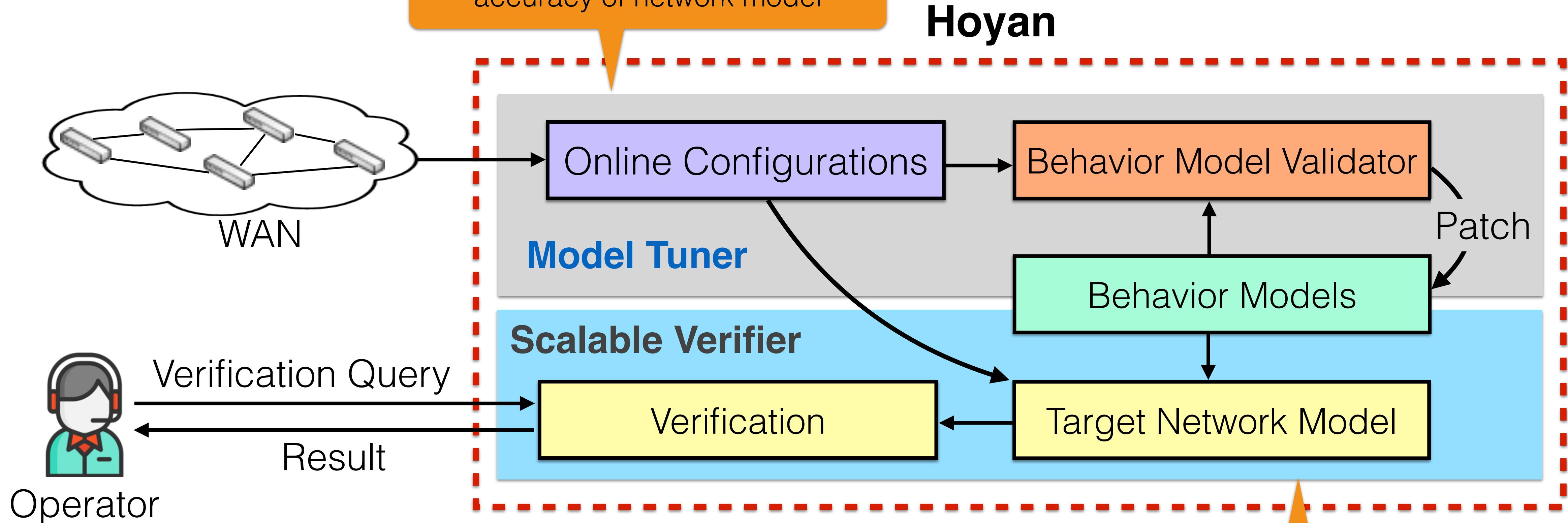
Innovation 1: Scalable Verifier

A set of scalable algorithms to reason about properties of interest



Innovation 1: Scalable Verifier

A set of scalable algorithms to reason about properties of interest

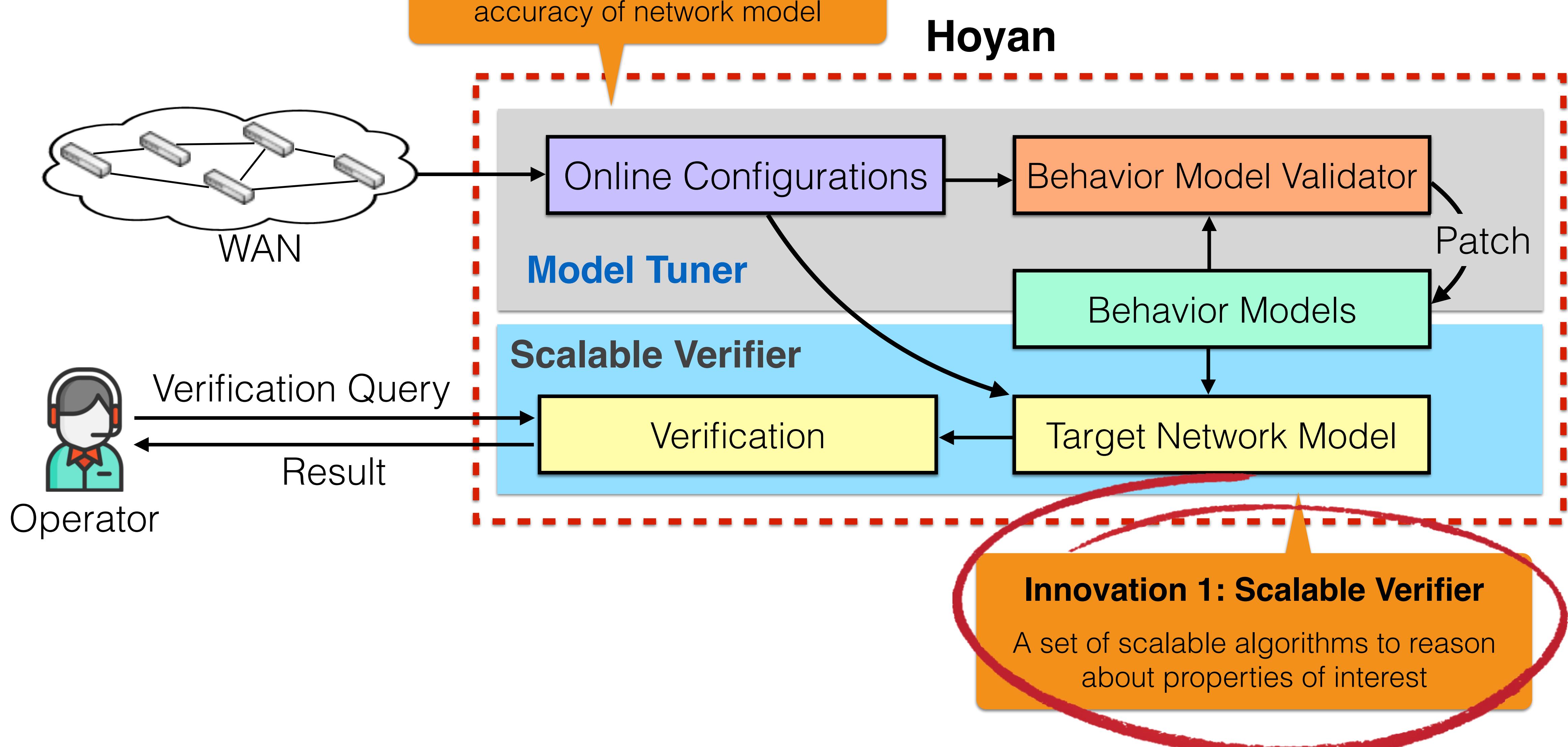


Innovation 1: Scalable Verifier

A set of scalable algorithms to reason about properties of interest

Innovation 2: Model Tuner

A systematic way to tune the accuracy of network model



How to address the scalability challenge?

Key insight:

- Simulation-based approaches are scalable but fail to cover all possible cases
- Formal model approaches can encode all possibilities and conditions as a formula but are not scalable because solving such a huge formula is expensive

How to address the scalability challenge?

Key insight:

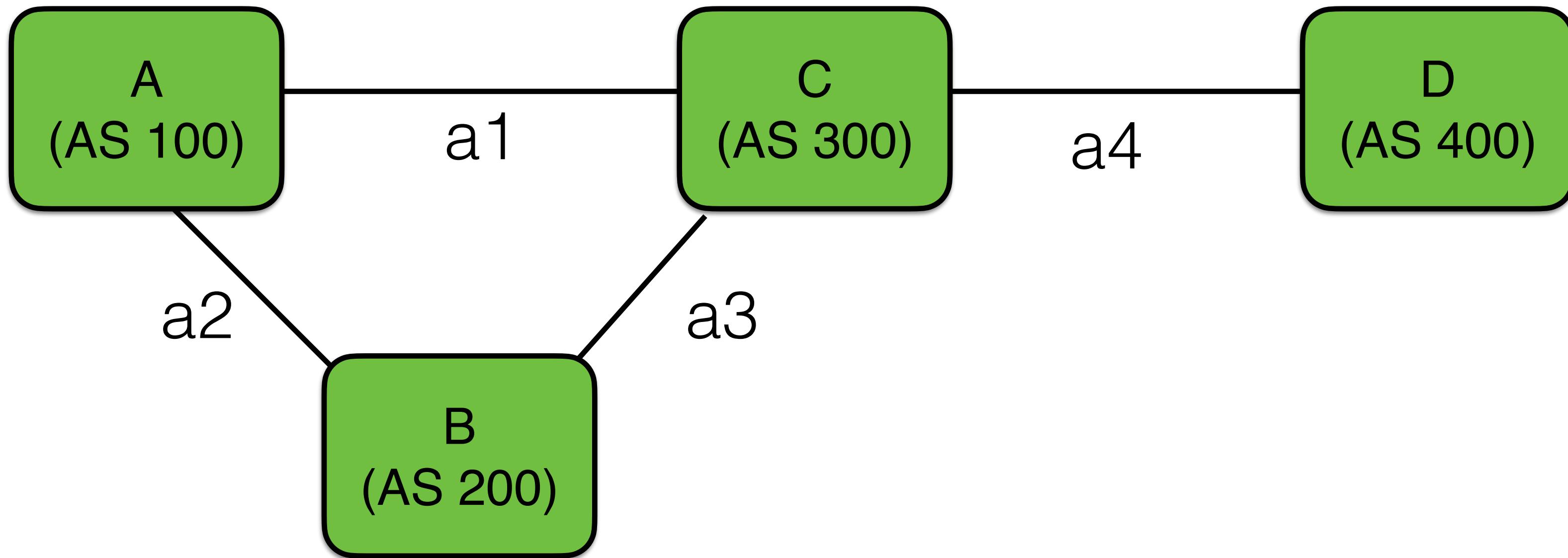
- Simulation-based approaches are scalable but fail to cover all possible cases
- Formal model approaches can encode all possibilities and conditions as a formula but are not scalable because solving such a huge formula is expensive

Solution: Global simulation & locally formal modeling approach

- We run overall simulation for global scalability
- Each router node locally encodes its all possible routing conditions that are transmitted along with the simulation process

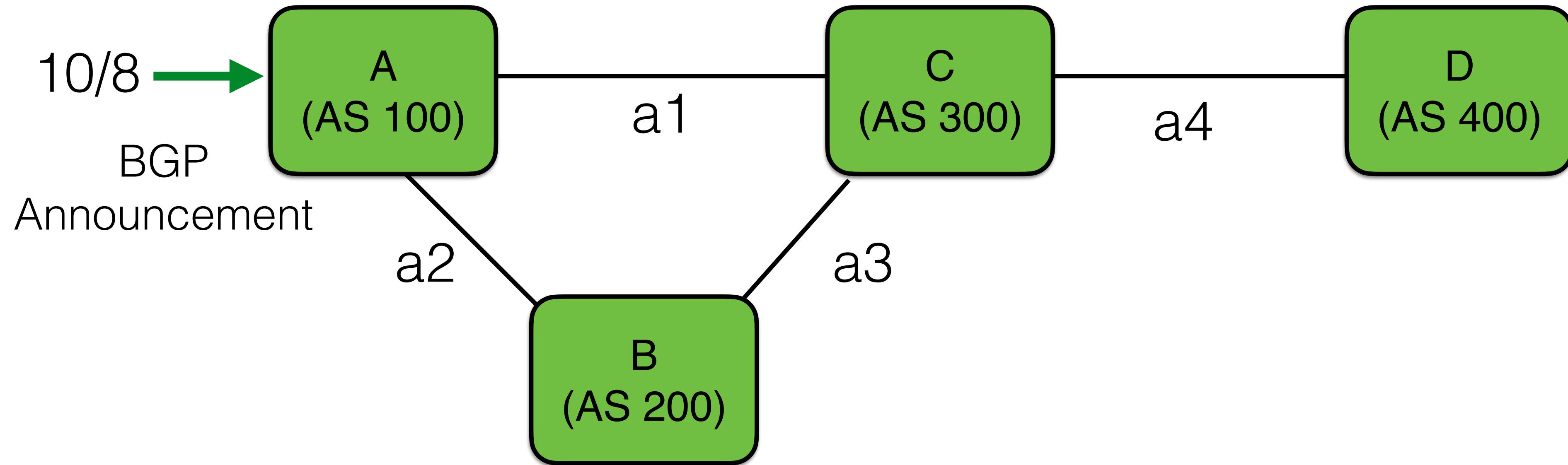
Global simulation & local formal-modeling

Checking k-failure tolerance between A and D

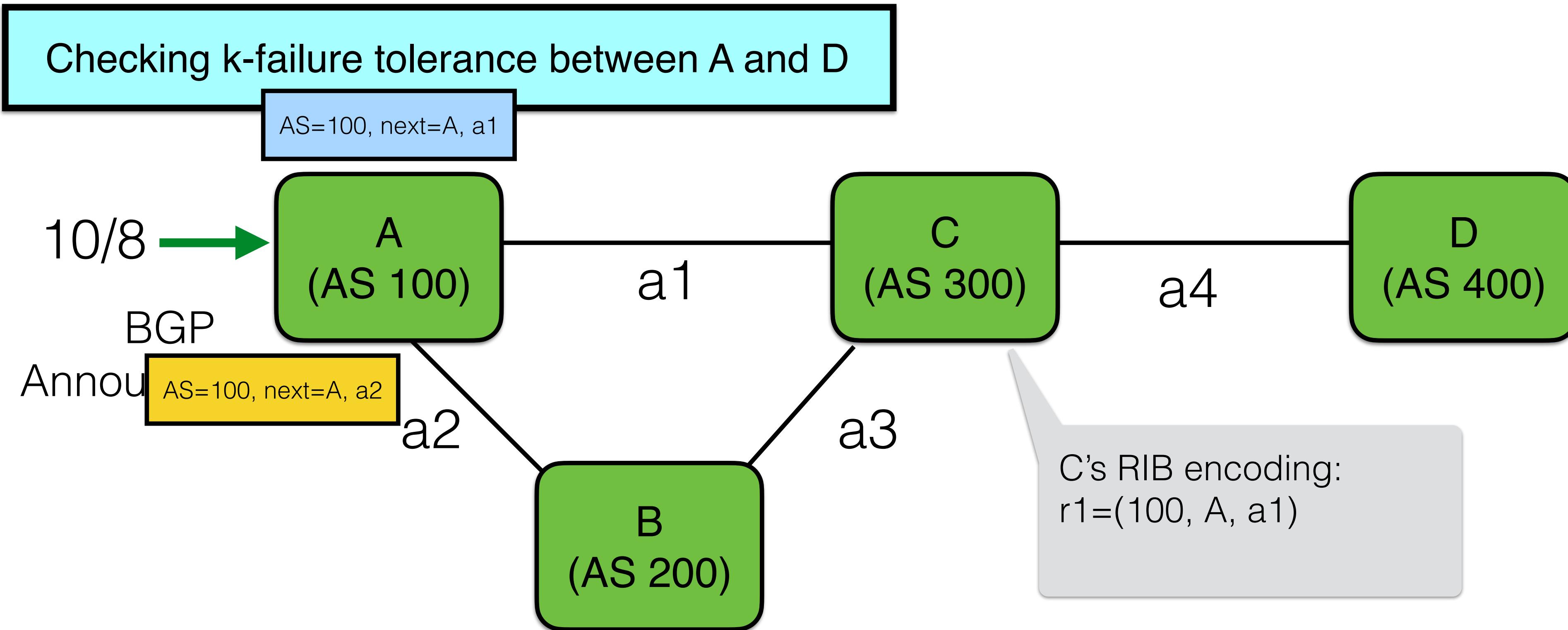


Global simulation & local formal-modeling

Checking k-failure tolerance between A and D

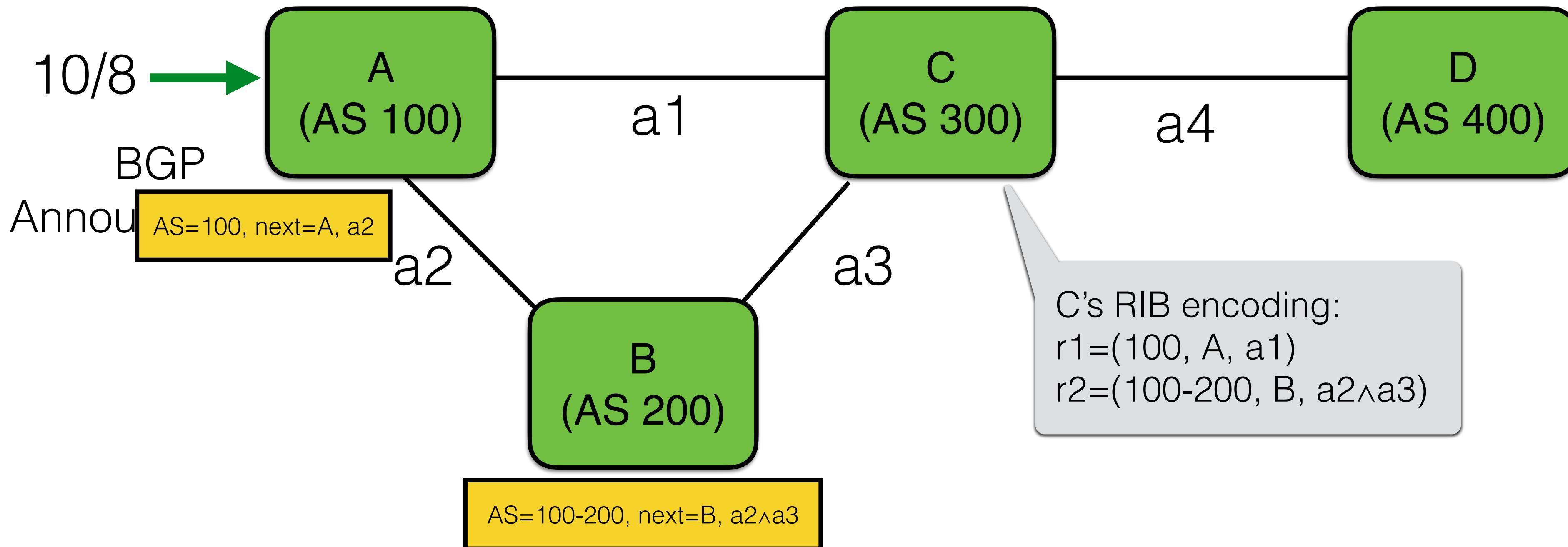


Global simulation & local formal-modeling

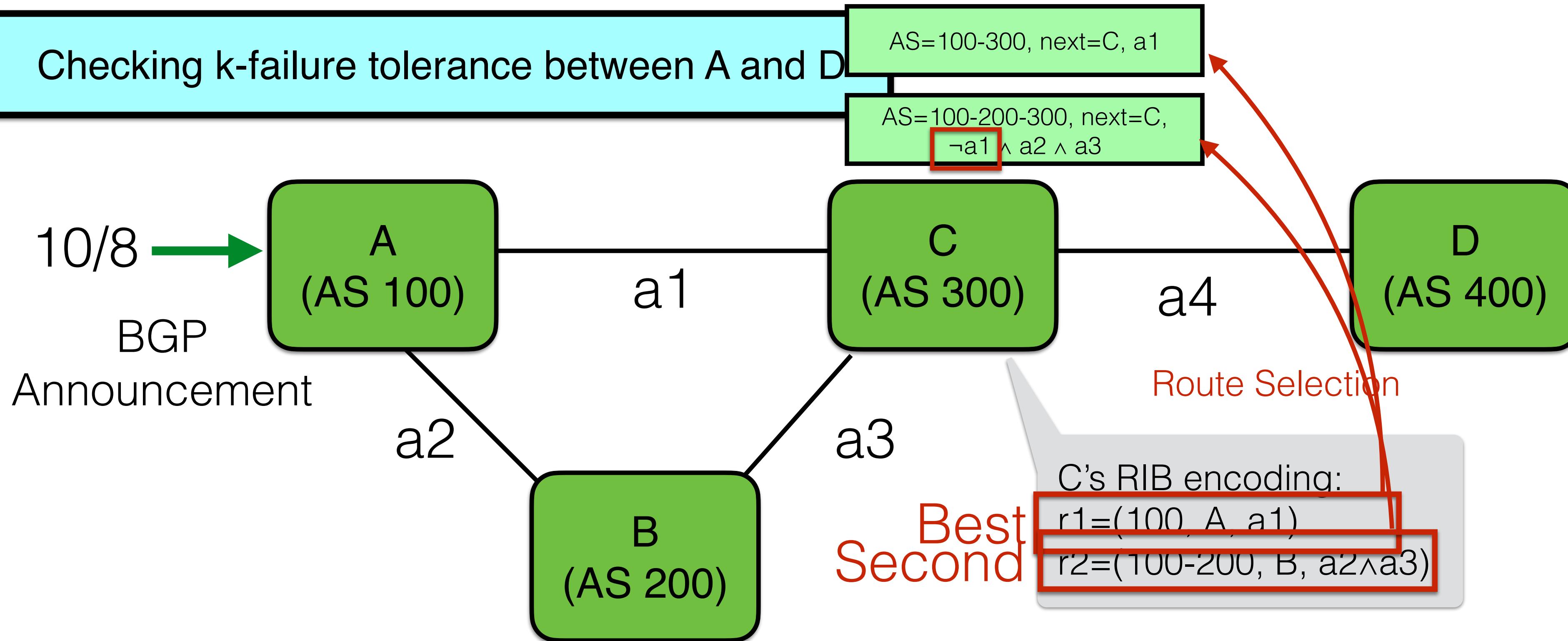


Global simulation & local formal-modeling

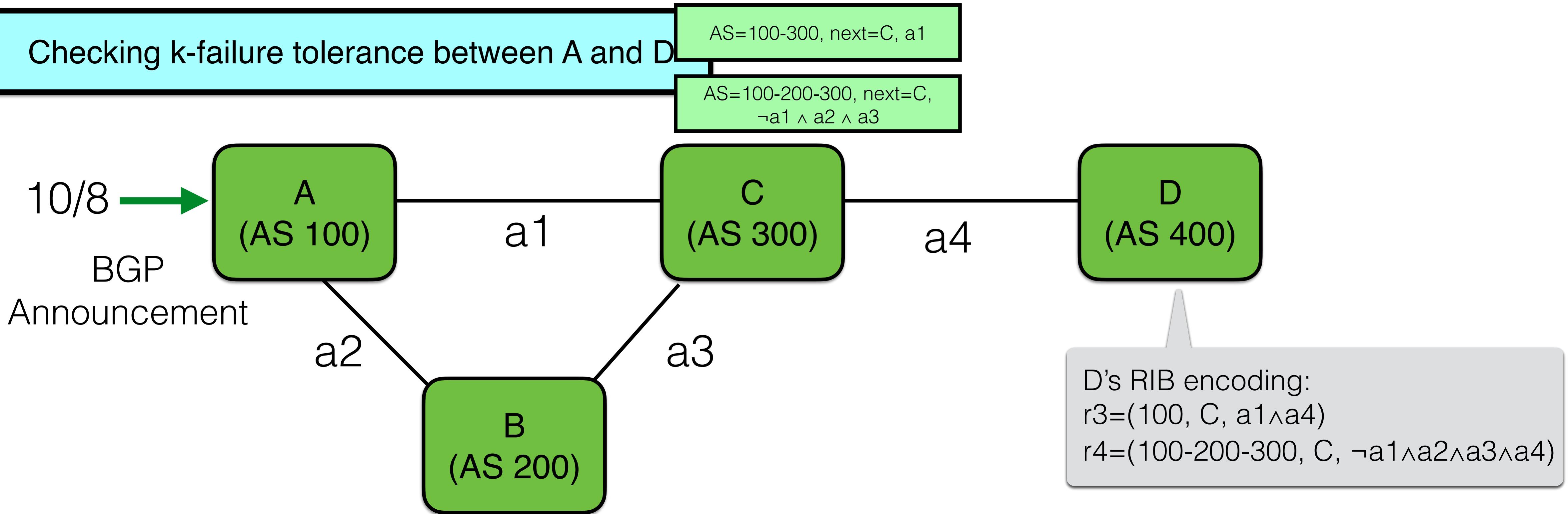
Checking k-failure tolerance between A and D



Global simulation & local formal-modeling

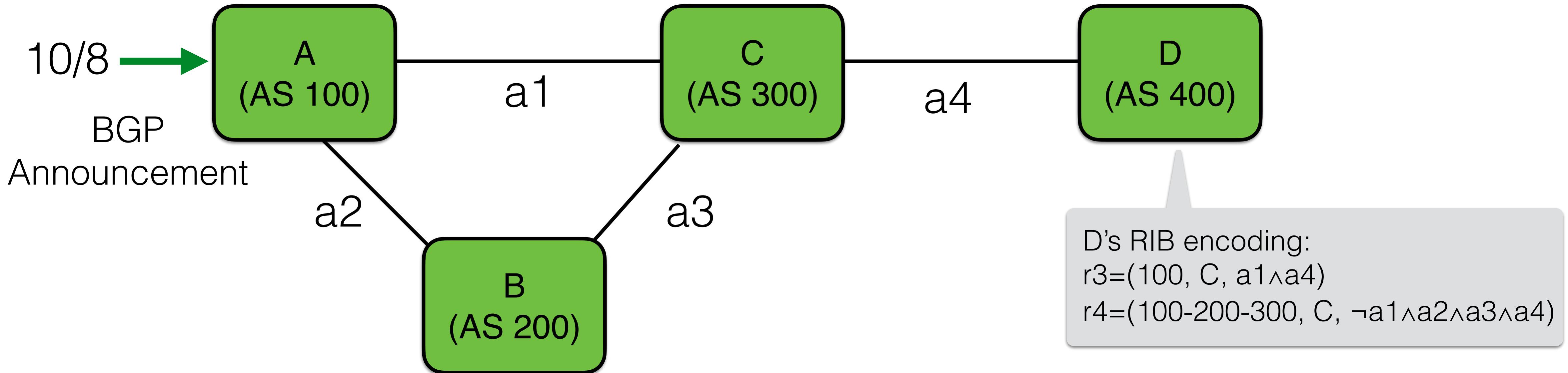


Global simulation & local formal-modeling



Global simulation & local formal-modeling

Checking k-failure tolerance between A and D

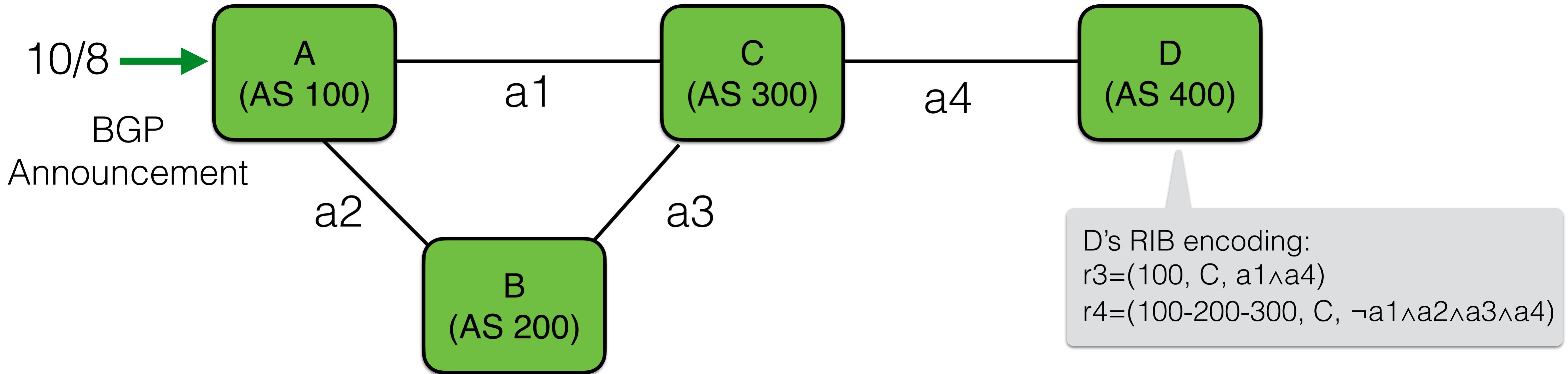


The local formula of “route of A reach D” is $(a1 \wedge a4) \vee (\neg a1 \wedge a2 \wedge a3 \wedge a4)$

← $\neg a4$

Global simulation & local formal-modeling

Checking k-failure tolerance between A and D



The good scalability comes from:

Dropping more than k cases

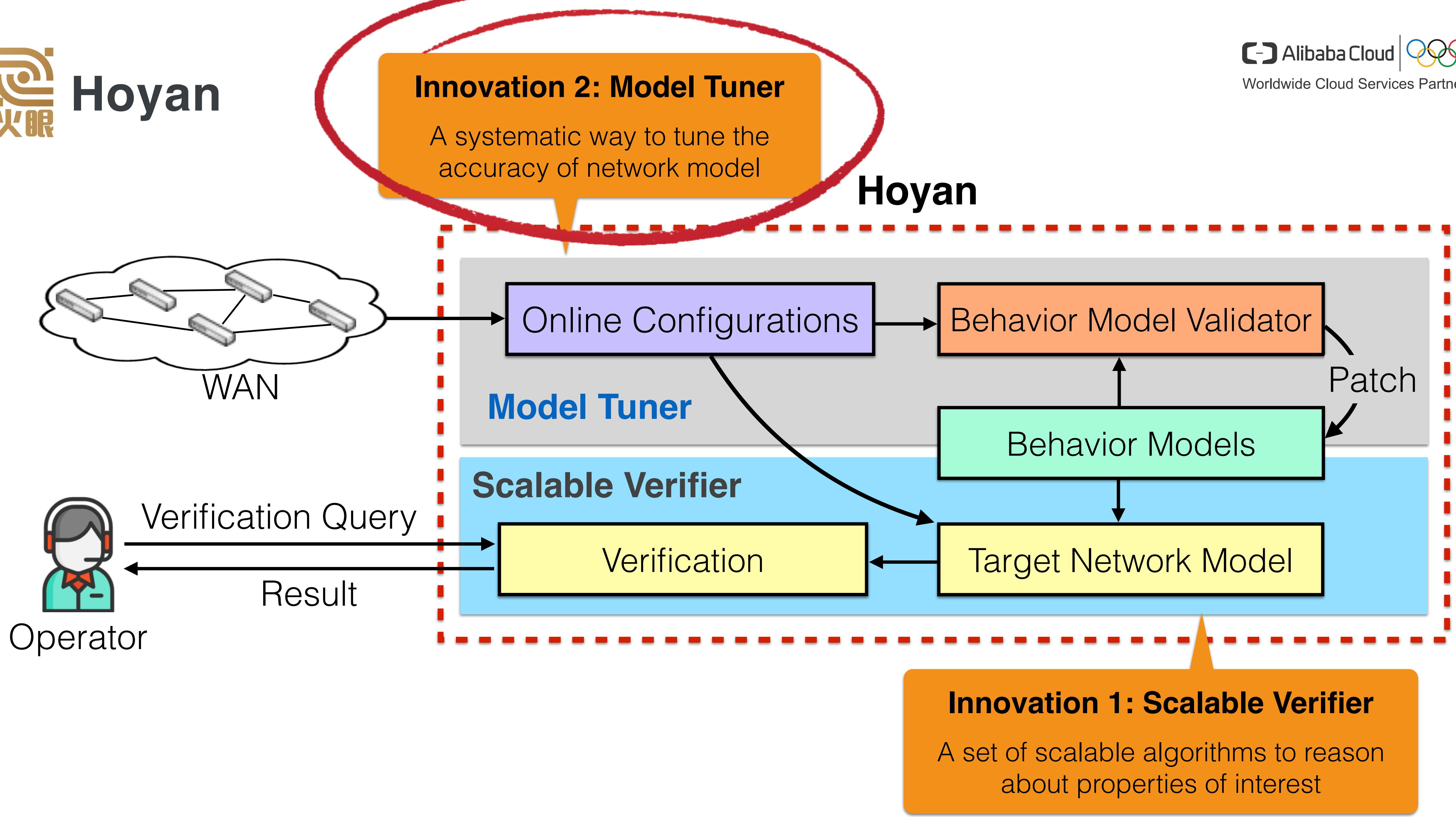
Dropping impossible cases

Simplifying logical formulas

Comparing Hoyan, Batfish and Minesweeper

A sub-network: 80 routers with thousands of lines of configurations

Network Properties		Hoyan	Minesweeper	Batfish
Reachability	k=0	14 seconds	84,043 seconds	683 seconds
	k=1	22 seconds	> 24 hours	> 24 hours
	k=2	43 seconds	> 24 hours	> 24 hours
	k=3	1176 seconds	> 24 hours	> 24 hours
Role Equivalence		4 seconds	> 24 hours	-



How to address the model faithfulness challenge?

Key insight:

- Using production network RIBs as a reference to debug simulated model

How to address the model faithfulness challenge?

Key insight:

- Using production network RIBs as a reference to debug simulated model

Strawman solution:

- Localizing vendor-specific behaviors by comparing real and simulated RIBs

Real-world routing tables

prefix	as path
10/8	i
20/8	i

prefix	as path
10/8	100
20/8	100

prefix	as path
10/8	200
	100
20/8	200
	100

prefix	as path
20/8	300
	200
	100

Diff

Simulated routing tables

Normal RIB

prefix	as path
10/8	i
20/8	i

prefix	as path
10/8	100
20/8	100

prefix	as path
10/8	200
	100
20/8	200
	100

prefix	as path
10/8	300
	200
	100
20/8	300
	200
	100

R1

R2

R3

R4

How to address the model faithfulness challenge?

Key insight:

- Using production network RIBs as a reference to debug simulated model

Strawman solution:

- Localizing vendor-specific behaviors by comparing real and simulated RIBs

Real-world routing tables

prefix	as path
10/8	i
20/8	i

prefix	as path
10/8	100
20/8	100

prefix	as path
10/8	200
	100
20/8	200
	100

prefix	as path
20/8	300
	200
	100

Simulated routing tables

Normal RIB

prefix	as path
10/8	i
20/8	i

prefix	as path
10/8	100
20/8	100

prefix	as path
10/8	200
	100
20/8	200
	100

prefix	as path
10/8	300
	200
10/8	200
	100
20/8	300
	200
	100

Same

Same

Same

Diff

R1

R2

R3

R4

How to address the model faithfulness challenge?

Key insight:

- Using production network RIBs as a reference to debug simulated model

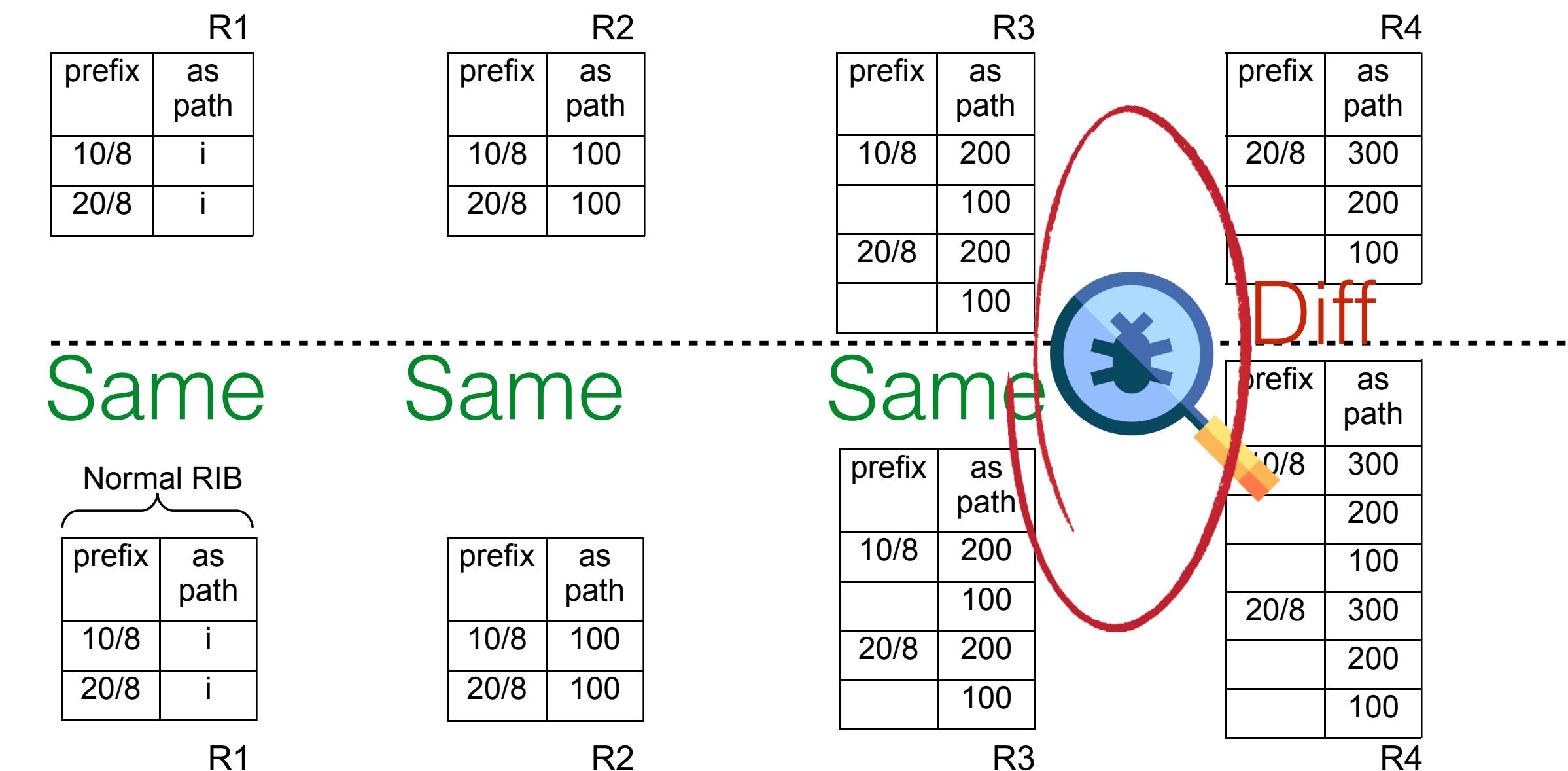
Strawman solution:

- Localizing vendor-specific behaviors by comparing real and simulated RIBs

Real-world routing tables

However, using normal RIBs cannot localize the correct VSBs!

Simulated routing tables



How to address the model faithfulness challenge?

Key insight:

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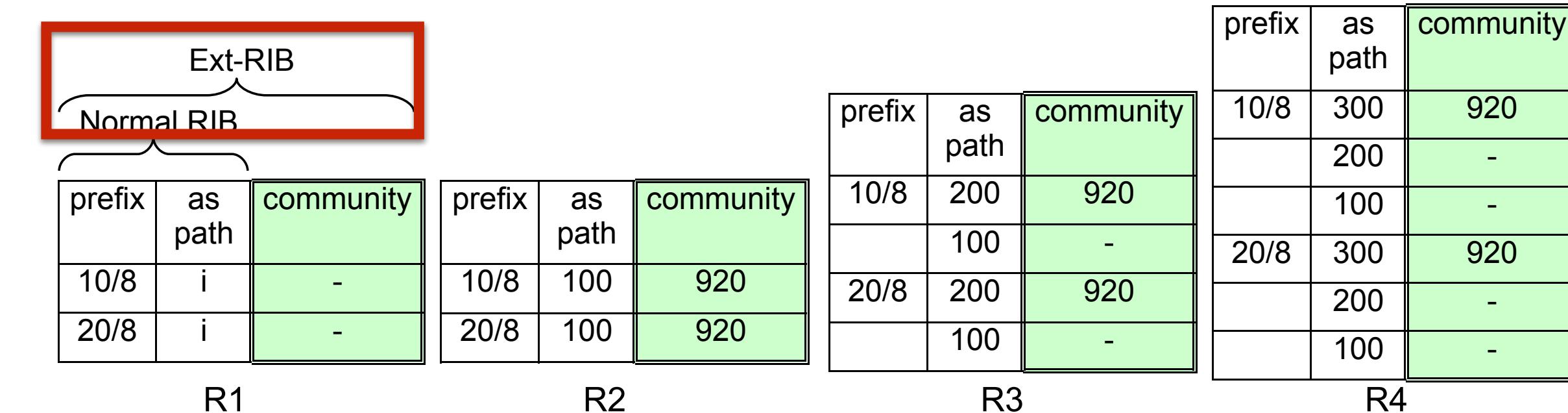
Our solution:

- Hoyan combines all the attributes of a route relevant for routing into an extended RIB
 - The first mismatch between real and simulated ext-RIBs is vendor-specific behavior

Real-world routing tables

R1			R2			R3			R4		
prefix	as path	community									
10/8	i	-	10/8	100	920	10/8	200	-	20/8	300	920
20/8	i	-	20/8	100	920		100	-		200	-
						20/8	200	920		100	-
							100	-			

Simulated routing tables



How to address the model faithfulness challenge?

Key insight:

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Our solution:

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Real-world routing tables

R1			R2			R3			R4		
prefix	as path	community									
10/8	i	-	10/8	100	920	10/8	200	-	20/8	300	920
20/8	i	-	20/8	100	920		100	-		200	-

Same

Same

Diff

Simulated routing tables

Ext-RIB

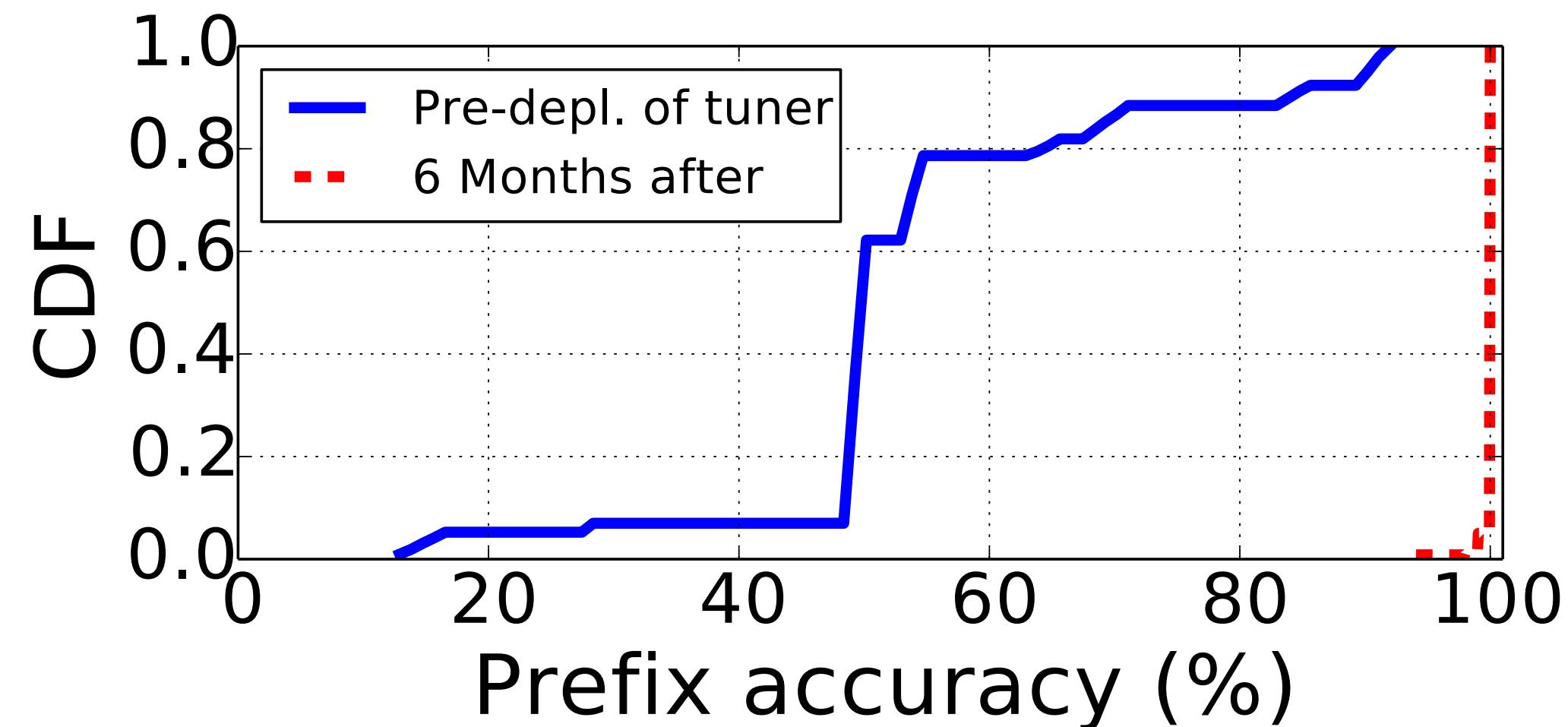
Normal RIB

Ext-RIB			Normal RIB			Ext-RIB			Normal RIB		
prefix	as path	community	prefix	as path	community	prefix	as path	community	prefix	as path	community
10/8	i	-	10/8	100	920	10/8	200	920	10/8	300	920
20/8	i	-	20/8	100	920		100	-		200	-

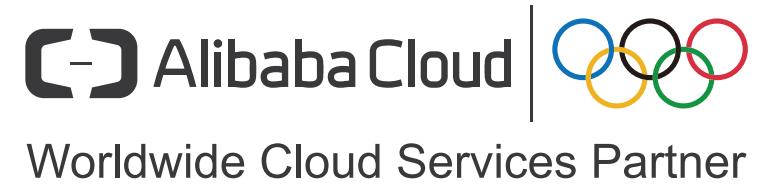
R1 R2 R3 R4

Device behavior model tuner

- Eight VSBs detected
 - Default ACL, route policy, removing private AS, etc.;
 - VSB localization within 10 lines of configuration



Hoyan: Deployment experience

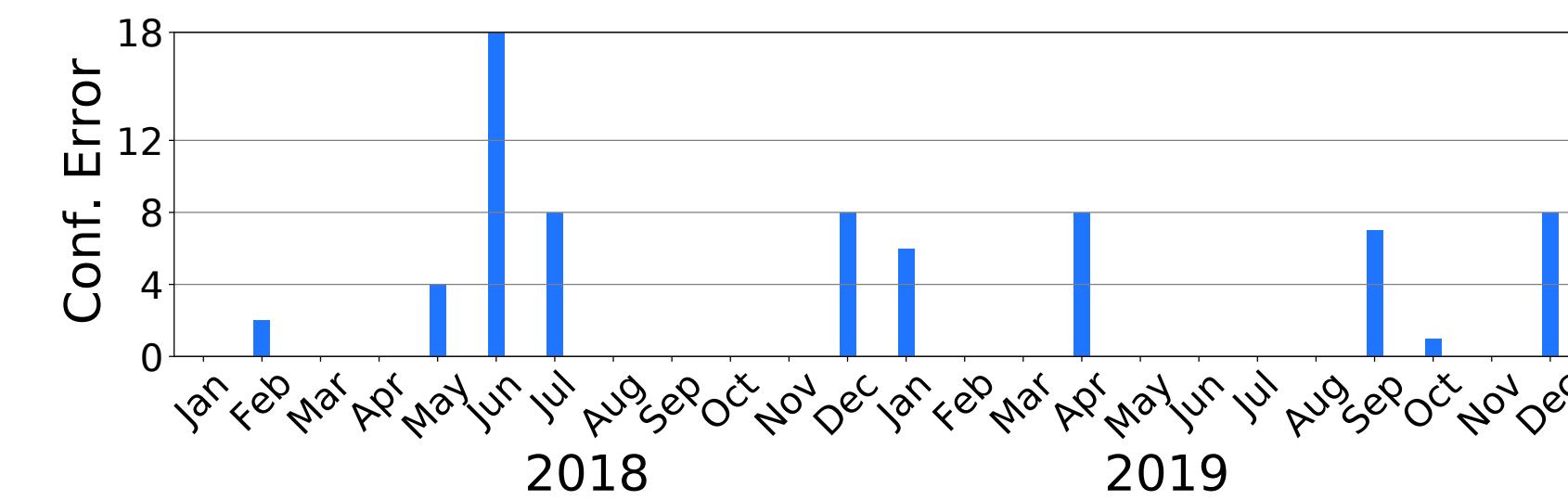


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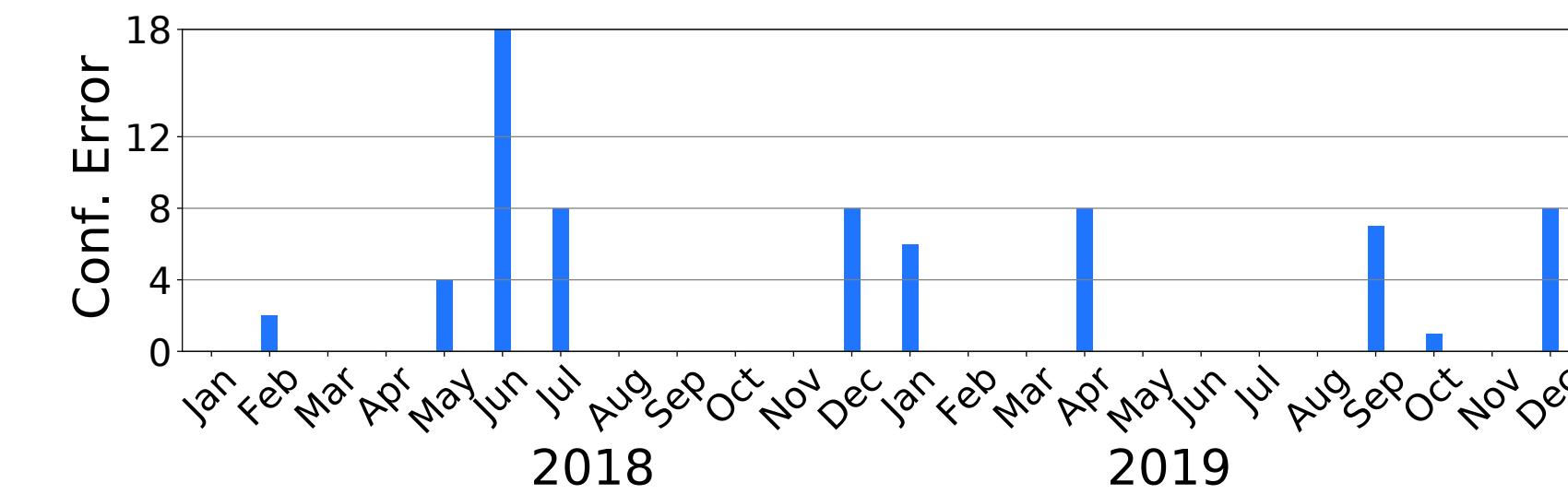
- Configuration auditing verification
(Jan 2018 - Dec 2019)



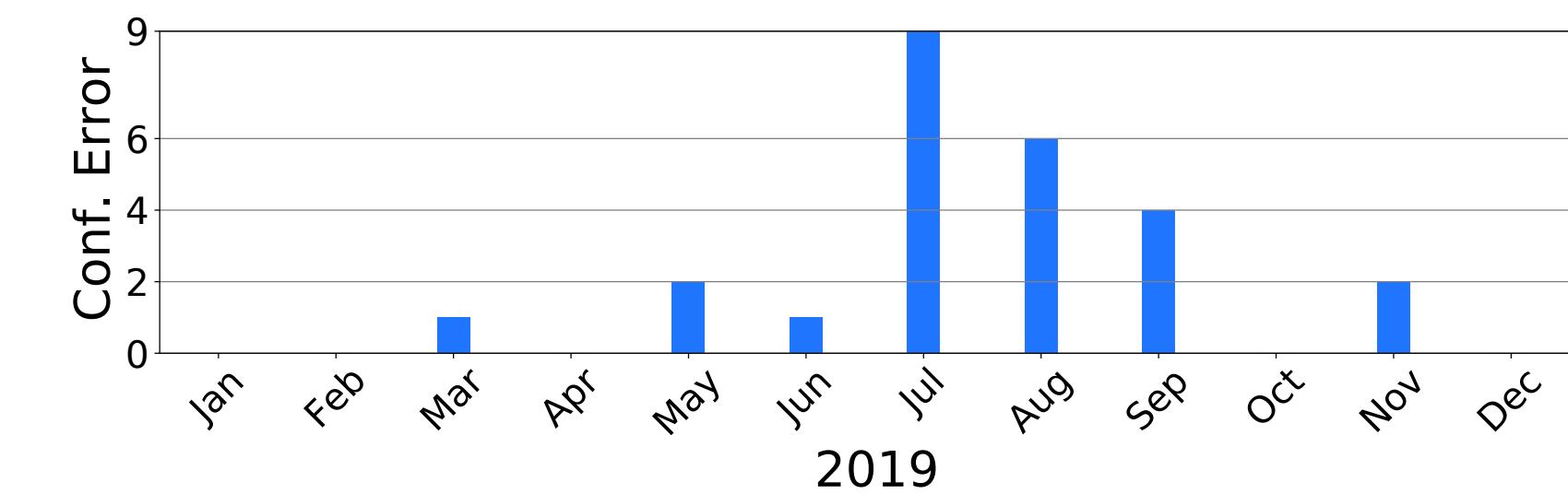
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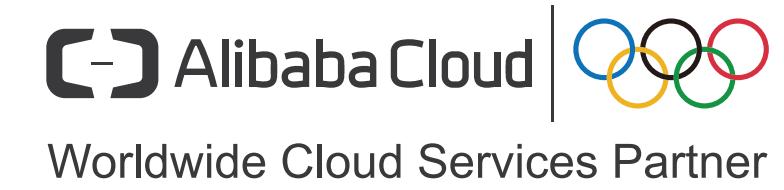
- Configuration auditing verification
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- Configuration update verification
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Hoyan's verification capability summary



What protocols are covered by Hoyan?

- BGP (eBGP and iBGP), IGP (IS-IS and OSPF), static routes, route redistribution and ACL

Hoyan's verification capability summary

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- BGP (eBGP and iBGP), IGP (IS-IS and OSPF), static routes, route redistribution and ACL

What properties can Hoyan verify?

- Reachability (route and packet reachability)
- K-failure tolerance
- Non-deterministic routing update race
- Multipath consistency (equality)

目录

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03 具体例子：网络验证

04 更多的方向和未来的思考

Our lessons and open questions

- Verification on WAN is more needed than DCNs
 - WAN is more complex than DCNs in routing policies and dependencies
 - WAN is not symmetric and configuration is not such “modular”
 - Tricky errors, e.g., non-deterministic issues

Our lessons and open questions

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 - WAN is more complex than DCNs in routing policies and dependencies
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- Expressing intent and intent extraction
 - Multiple security domains and implicit business needs
 - Does example-based specification help? e.g., example spec verification or synthesis

Our lessons and open questions

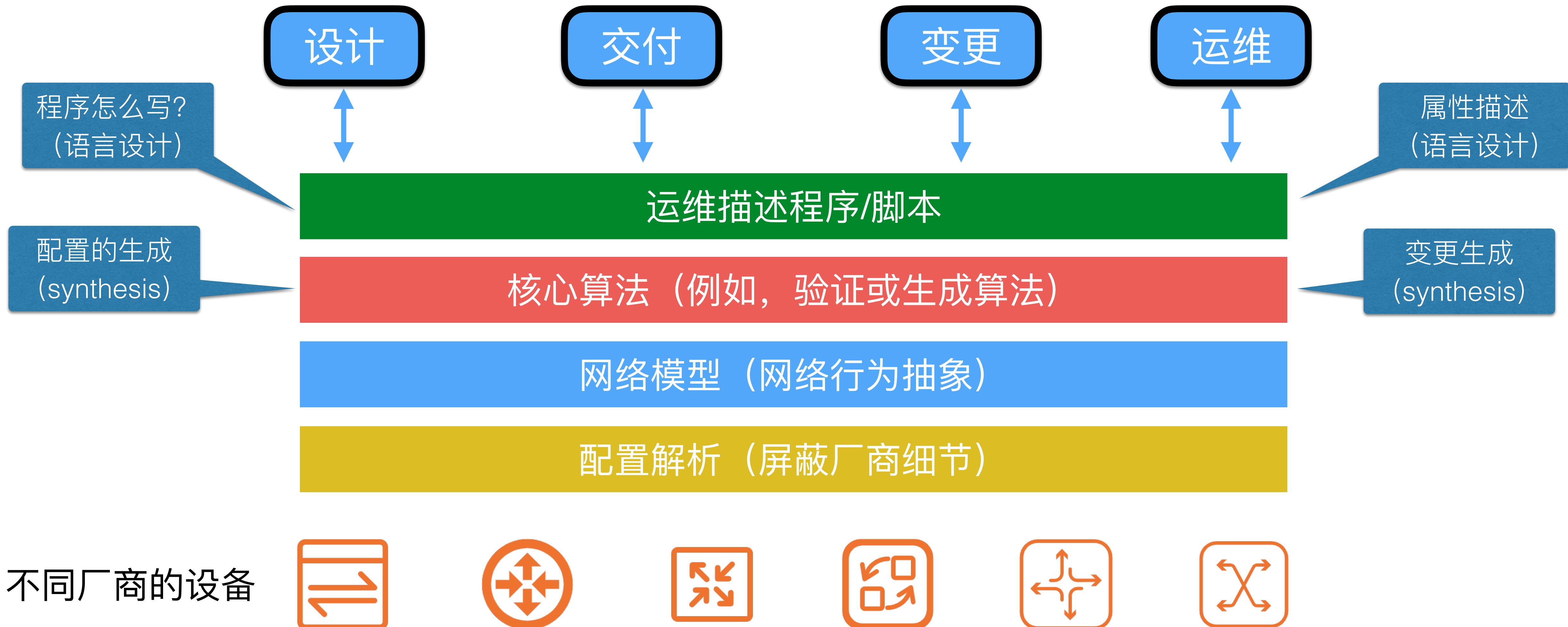
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Our lessons and open questions

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 - How to express the fixing purpose?

- Vendor diversity
 - Vendor diversity can prevent cascading failures
 - How to quantify diversity? How to detect my network diversity is low?
 - Can we transform a “non-diverse” network to a diverse network?

集中式网络规划体系



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THANKS