

Dokl. Akad. Nauk SSSR  
Tom 179 (1968), No. 2

Soviet Math. Dokl.  
Vol. 9 (1968), No. 2

# NORMAL SEQUENCES AND FINITE AUTOMATA

UDC 519.95:519.211:511

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A series of works [1-6] contains studies of infinite binary sequences, which are "random" in one sense or another. The concept of a random sequence considered in [1-3] is that of a collective, introduced by von Mises [1] and made more precise by Wald [2]. In papers [5, 6] normal sequences are studied. The present note clarifies the connection between these concepts in terms of finite automata.

1. Collective. A predicate defined on the set of all finite binary sequences (including the empty sequence  $\Lambda$ ) will be called a strategy. A strategy  $f$  defines the following procedure of singling out the (possibly finite) subsequence  $y_1 y_2 y_3 \dots$  from the infinite binary sequence  $x_1 x_2 x_3 \dots$ . If  $z_1 z_2 z_3 \dots$  is the sequence of values of the strategy  $f$  corresponding to the sequence  $x_1 x_2 x_3 \dots$ :  $z_1 = f(\Lambda)$ ,  $z_2 = f(x_1)$ ,  $z_3 = f(x_1 x_2)$ ,  $\dots$  and if  $i_1 < i_2 < i_3 < \dots$  are the indices for which  $z_{i_k} = 1$ , then  $y_1 = x_{i_1}$ ,  $y_2 = x_{i_2}$ ,  $y_3 = x_{i_3}$ ,  $\dots$ . A sequence  $x_1 x_2 x_3 \dots$  is called a collective relative to a set of strategies  $F$  (containing the strategy identically equal to one), if, for all infinite subsequences singled out from  $x_1 x_2 x_3 \dots$  by the strategies of the set  $F$ , there exists a limit (the same for all these subsequences) of the relative frequency of the ones  $p = \lim_{n \rightarrow \infty} (1/n) \sum_{i=1}^n x_i$  ( $0 < p < 1$ ).

2. Normal sequence. An infinite binary sequence  $x_1 x_2 x_3 \dots$  is called normal if there exists a number  $p$  ( $0 < p < 1$ ) such that for any positive integer  $n$  and any word  $a_1 a_2 \dots a_n$  of length  $n$  in the alphabet  $\{0, 1\}$  the following condition is satisfied: in the sequence of words  $x_1 x_2 \dots x_n$ ,  $x_{n+1} x_{n+2} \dots x_{2n}$ ,  $x_{2n+1} x_{2n+2} \dots x_{3n}$ ,  $\dots$  the relative frequency of appearance of the word  $a_1 a_2 \dots a_n$  has a limit equal to  $p^k (1-p)^{n-k}$ , where  $k = \sum_{i=1}^n a_i$  is the number of ones in the word  $a_1 a_2 \dots a_n$ .

3. A normal sequence as a collective. In [7] it is established that a sequence is normal if and only if it is a collective relative to the set of strategies possessing the following property: the strategy singles out or does not single out a letter according to what the  $l$  preceding letters are, with  $l$  fixed for every given strategy.

It is obvious that for every such strategy  $f$  there exists a finite automaton  $\langle \{0, 1\}, Q, Q^*, \phi, q \rangle$  with input alphabet  $\{0, 1\}$ , a set of states  $Q$ , a subset of singled out states  $Q^* \subseteq Q$ , a transition function  $\phi: Q \times \{0, 1\} \rightarrow Q^*$  and an initial state  $q \in Q$ , which computes the strategy  $f$  in the following manner:  $f(x_1 x_2 \dots x_n) = 1 \iff \phi(\phi(\dots \phi(\phi(q, x_1), x_2) \dots), x_n) \in Q^*$ . Hence it follows that a collective relative to the set of strategies computed by finite automata is a normal sequence. The question naturally arises whether the converse assertion is valid. A positive answer to this question is given by the following

**Theorem.** *A normal sequence is a collective relative to the set of strategies computed by finite automata.*

The proof of the theorem utilizes the ergodic property of a finite automaton, whose input is a sequence of independent random variables and the following property of "nondecrease of measure", possessed by a transformation defined by strategies on words. We call the  $p$ -weight of the word  $a_1 \dots a_n$  the number  $p^l (1-p)^{n-l}$ , where  $l$  is the number of ones in the word  $a_1 \dots a_n$  and  $0 < p < 1$ ;

the  $p$ -weight of a subset of words of fixed length is the sum of the  $p$ -weights of all the words of this set. Then for any strategy  $f$  and any word  $a_1 \dots a_k$  the set of words of length  $n \geq k$ , from which  $f$  singles out words beginning with the word  $a_1 \dots a_k$ , has a  $p$ -weight not larger than the  $p$ -weight of the word  $a_1 \dots a_k$ .

In conclusion I would like to express my gratitude to Professor B. A. Trahtenbrot, who suggested the problem and made a number of valuable remarks.

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Received 14/MAY/67

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