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//-----createFile-----
int createFile(char* fileName, int permission){
  struct Block* directoryBlock = getCurrentDirectory();
  if(FileAlreadyExist(fileName)) return -1;
  struct iNode newINode = {permission, size = 0};
  if current cluster is full on data block || full on INode
    find nearest cluster with enough space (inode + one block);
    int iNodeIndex = allocateINodeInThatCluster(newINode);
    if allocate fails return -1;
  else
    int iNodeIndex = allocateINodeInCurrentCluster(newINode);
    if allocate fails return -1;
  // Do preallocation, all files will automatically have one block (4k) at creation
  allocateFirstDataBlock(iNodeIndex);
  if(isFull(directoryBlock))
    Append a newBlock after current directory blocks in its parent directory;
    newBlock.append({filename, iNodeIndex});
  else
    directoryBlock.append({filename, iNodeIndex});
  updateINodeInCache(); // Write changed directory to the disk
          // And load Inode table of this cluster in cache
  return 0;
}
//-----openFile-----
int openFile(char* fileName){
  if (file is already opened)
    return getFileDescriptor(fileName);
  if fileDescriptorTable is full return -1;
  int newFD = findEmptyEntry(fileDescriptorTable);
  int iNodeIndex = findINodeIndex(fileName);
  if iNode not in cache:
    cluster = find the cluster containing the INode in Disk;
    loadINodeFromDisk(cluster);
  if not a single cluster contain this INode
    // file does not exist yet. Try create first
    if (createFile(fileName) failed) return -1;
  fileDescriptorTable[newFD].iNodeIndex = iNodeIndex;
  fileDescriptorTable[newFD].readPtr = FILE START;
  if(peekFilePermission(iNodeIndex) == READ ONLY)
    fileDescriptorTable[newFD].writePtr = NULL;
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else
     fileDescriptorTable[newFD].writePtr = FILE START;
  return newFD;
}
//----seek-----seek------
int seek(int FDindex, int offset, char readOrWrite){
  // the seek updates the selected pointer by number specified in offset
  // if update is successful, this method return the pointer value prior to the update
  // (Analogy to sbrk())
  int pointerToChange;
  if(readOrWrite == "r")
     fileDescriptorTable[FDindex].readPtr = fileDescriptorTable[FDindex].readPtr + offset
     return pointerToChange;
  else if (readOrWrite == "w")
     fileDescriptorTable[FDindex].writePtr = fileDescriptorTable[FDindex].writePtr + offset
     return pointerToChange;
  else
     return -1;
}
//----read-----
int read(int fdID, char* buf, int length){
  // check if file descriptor ID is valid
  if(fdID < 0 || fdID > Max_size_of_fdt){
     return -1;
  }
  int INodeID = fileDescriptorTable[fdID].INodeIndex
  int readPointer = fileDescriptorTable[fdID].readPtr
  // unlikely to happen since the prior open() will bring INode to cache
  // But to cover extreme condition (like the file is opened but its iNode got replaced in cache)
  // we added the condition below. Same thing in write
  if (iNodeldx exist in cache)
    curlNode = get INode in cache;
  else
    curlNode = find corresponding cluster and loadINodeFromDisk(cluster);
  // check if read outside of file
  if (readPointer+length > file size){
     length = file size - readPointer; //if outside, change length to read to end of file
  }
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int read_end = readPointer+length; //end position of read
  char* buf tracker = buf
                           //track position in buf
  int byteAlreadyRead = 0;
  while(readPointer != read_end){
    int block_Position_In_INode = readPointer/BLOCK_SIZE;
    int bytesToRead = min(BLOCK SIZE - readPointer % BLOCK SIZE,
length-byteAlreadyRead);
    if(isSingleIndirect(block_Position_In_INode)){
       Target_Block_toRead = single_indirect_pointer[block_Position_In_INode-12];
    }else if(isDoubleIndirect(block Position In INode)){
       int first_level_indirect_location = (i-12-1024)/1024
       int second_level_indirect_location = (i-12-1024)%1024
       Target Block toRead =
double_indirect_pointer[first_level_indirect_location][second_level_indirect_location];
    }else{
       Target_Block_toRead = direct_pointer[block_Position_In_INode];
    memcpy(Target_Block_toRead + readPointer % BLOCK_SIZE, buf_tracker,
bytesToRead);
    readPointer += bytesToRead; //update pointers
    buf tracker += bytesToRead;
    byteAlreadyRead += bytesToRead;
  }
  updateReadpointer(); //update read pinter in I-node
  return length;
}
//-----writeFile-----
int writeFile(int fd, const char* buf, int len){
  // Note every file "borns with one block (4k)"
  // this is to reduce the time of updating/syncing inode
  bool inodeHasChanged = false;
  char bufferForOneBlock [BLOCK_SIZE];
  memset(buffer, 0 ,BLOCK_SIZE); //intialize buffer
  int byteAlreadyWritten = 0;
  iNodeldx = fileDescriptorTable[fd].iNodeIndex;
  if iNodeldx is null or invalid
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return -1; // there is error in opening the file
  writePtr = fileDescriptorTable[fd].writePtr;
  if (iNodeldx exist in cache)
     INode = get INode in cache;
  else
     find corresponding cluster holding INode and loadINodeFromDisk(cluster);
  int original_File_Size_Before_Write = INode.filesize;
  while(byteAlreadyWritten < len){
     int block_Position_In_INode = ptr/BLOCK_SIZE;
     int block idx;
     if(is In_12_DirectBlocks(block_Position_In_INode)){
       // One of the 12 direct node
       block idx = directBlocksInINode[block_Position_In_INode];
       if(getBlock(block idx) is null){
          block_idx = allocate a new data block and link to inode else break;
          inodeHasChanged = true;
       }
     }else if (is In Single Indirect Region(block Position In INode)){
       // writeptr is in single indir region
       if (single indirect index block has not existed yet){
          allocate a new index block and link to INode else break;
          indexBlock_idx = get index of the index block;
          inodeHasChanged = true;
       loadIndexBoxIntoBuffer(indexBlock idx, bufferForOneBlock);
       indexBlock_SingleIndir = (int[1024])bufferForOneBlock; // now this block buffer holds
1024 entry for block
       block idx = indexBlock_SingleIndir[block_Position_In_INode -12];
       if block idx points to null
          // Need to allocate the target block
          block idx = allocate a new data block else break;
          indexBlock_SingleIndir[block_Position_In_INode -12] = block_idx;
     }
     else{
       // based on block_Position_In_INode, the writePtr is targeting at the double indirect
pointer
       // 1. make sure single block has exist, if not, create a single indir index block first
       if(single indirect index block has not existed yet){
          allocate a new index block and link to INode else break;
          inodeHasChanged = true;
       }
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// after making sure there are no sparse space in the single indir region in INode
       // we start touching the double indir pointer
       if(double indirect index block has not existed yet){
          // if there is no double indir block yet, allocate one
          indexBlock idx = allocate a index block fo double indir in available cluster else break;
          inodeHasChanged = true;
       loadIndexBoxIntoBuffer(indexBlock idx, bufferForOneBlock);
       indexBlock DoubleIndir = (int[1024])bufferForOneBlock;
       // every entry of double indir index block should be a index block (as tutorial, B-->B, not
A/C)
       // the corresponding double indir index is calculated through (block Position In INode -
12 - 1024)/1024
       // (block_Position_In_INode - 12 - 1024) is current block index. One entry on double indir
holds 1024 blocks.
       if (indexBlock DoubleIndir[(block Position In INode - 12 - 1024)/1024] is not an index
block) break;
       // targeted index block is what the current entry of double Indir block pointing to
       targeted index block = indexBlock DoubleIndir[(block Position In INode - 12 -
1024)/1024];
       // check if the target single indir block has not existed yet
       else if( targeted_index_block is null){
          targeted_index_block = allocate an index block else break;
          indexBlock_DoubleIndir[(block_Position_In_INode - 12 - 1024)/1024] =
targeted_index_block;
       // The two lines below access one level down the index blocks and get the data block
       indexBlock SingleIndir = getIndexBlock(targeted index block)
       block idx = indexBlock SingleIndir[(block Position In INode - 12 - 1024)%1024];
       if(block idx points to null){
          block idx = allocate a data block for write else break;
          targeted_index_block[(block_Position_In_INode - 12 - 1024)%1024] = block_idx;
       }
     bufferForOneBlock = getWriteTargetBlock(block idx);
     // determine how many bytes to write in this iteration based on the leftover space in block
     int bytesToWrite = = min(BLOCK_SIZE - writePtr % BLOCK_SIZE , len-
byteAlreadyWritten);
     // Keep writing from the input buf to the disk data block buffer
     memcpy(bufferForOneBlock + writePtr % BLOCK_SIZE, buf + byteAlreadyWritten,
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bytesToWrite );
    byteAlreadyWritten += bytesToWrite;
    writeptr += bytesToWrite;
  }
  if(inodeHasChanged == true) updateINodeInCache();
  if updateClusterToDisk(original File Size Before Write, writeptr) == -1
    return -1; // Update to disk failed, the write is not good
  return byteAlreadyWritten;
}
//-----close-----
int CloseFile(fd){
  if(IsFileOpen(fd)){
    openFileTable[fd] = OpenFD_unused;
    return 0:
  }
  return -1;
}
//------Helper Method--------
// To improve performance, our implementation follows below principles syncing cache to disk:
// 1. we are using clustered INode system, thus everytime we always load complete INode table
from one cluster to make use of locality
// 2. We update INode in cache only when there is a change. The changed INode is synced
back to disk when it is evicted from cache
// 3. To improve performance, we keep all updating of index blocks at the end of writeFile()
using helper method updateClusterToDisk()
// please check updateClusterToDisk() for detail.
int updateClusterToDisk(original File Size Before Write, writeptr){
  // Our design choose clustered INode system
  // Principle 1 is one file must have all its data blocks in the same cluster
  // Principle 2 is that all files are dense (sparse space will be filled with data blocks of all zero,
  // also empty entries on index blocks will be filled)
  // This method first checks how many extra blocks needed to fill sparse region due to write
  // then based on the extra space, it determines whether to allocate all those blocks in current
cluster
  // or to move the entire file with all data block and INode to a new cluster (since very unlikely,
but current cluster is nearly full)
  dense_Block_Position_In_INode = original_File_Size_Before_Write/BLOCK_SIZE;
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end Block Position In INode = writeptr/BLOCK SIZE;
  numberOfAdditionBlockNeeded = end_Block_Position_In_INode -
dense Block Position In INode;
  if (numberOfAdditionBlockNeeded * BLOCK SIZE > CLUSTER SIZE - currentCluster.size){
     // current cluster is full and can not be write directly back to disk
     newCluster = find clostest cluster with enough size and bring to cache else return -1;
     copy current file I-node and data blocks there;
     fillSparseSpaceAt(newCluster, writeptr, numberOfAdditionBlockNeeded, fileInode);
  }else{
     fillSparseSpaceAt(currentCluster, writeptr, numberOfAdditionBlockNeeded,fileInode);
  }
  syncDataBlocksInCacheToDisk();
  return 0;
}
void fillSparseSpaceAt(selectedCluster, writeptr, numberOfAdditionBlockNeeded){
  newIntermediateBlocks = allocate numberOfAdditionBlockNeeded of data blocks in
selectedCluster:
  if(isIn 12Direct Block(writeptr)){
     initialize newIntermediateBlocks and link to fileInode;
  }else if(is_In_Single_Indirect_Region(writeptr)){
     initialize newIntermediateBlocks and link to fileInode and single indir index block;
  }else{
     // in second indir region
     for every single indir index block linked by the double indie index block:
       initialize newIntermediateBlocks and link to fileInode and single indir index block;
  }
}
void loadINodeFromDisk(newlyEnteringCluster){
  // according to our memory structure, the cache has enough space to hold
  // over 600 cluster's complete INode tables
  // EveryTime a new cluster is called, its inodes will be brought to cache
  // The cache replace cluster's INode table using LRU principle
  if cache is full on cluster INode table entries:
     replacedOne = find the cluster which has not been called for longest time;
     Write INode content in replaceOne back to disk;
     Evict replaceOne;
     Load newlyEnteringCluster.INodeTable in cache;
  else:
     // Cache is not full, need not to replace
     Load newlyEnteringCluster.INodeTable in cache;
```