

# EL9343 Homework 3

Due: Feb.13th 5:00 p.m.

1. True or False questions:

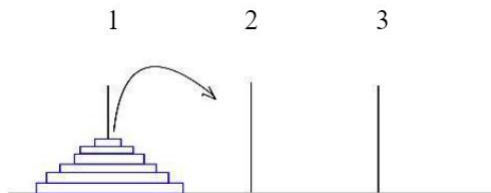
- (a) One can find the maximum sub-array of an array with  $n$  elements within  $O(n \log n)$  time.
- (b) In the maximum sub-array problem, combining solutions to the sub-problems is more complex than dividing the problem into sub-problems.
- (c) Bubble sort is stable.
- (d) When input size is very large, a divide-and-conquer algorithm is always faster than an iterative algorithm that solves the same problem.
- (e) It takes  $O(n)$  time to check if an array of length  $n$  is sorted or not.
- (f) Insertion sort is NOT in-place.
- (g) The running time of merge-sort in worst-case is  $O(n \log n)$ .

2. Consider sorting  $n$  numbers stored in array  $A$ , indexed from 1 to  $n$ . First find the smallest element of  $A$  and exchanging it with the element in  $A[1]$ . Then find the second smallest element of  $A$ , and exchange it with  $A[2]$ . Continue in this manner for the first  $n - 1$  elements of  $A$ .

- (a) Write pseudo-code for this algorithm, which is known as *selection sort*.
- (b) What loop invariant does this algorithm maintain?
- (c) Give the best-case and worst-case running times of selection sort in  $\Theta$ -notation.

3. The Tower of Hanoi consists of 3 rods and a number of disks of various diameters, which can slide onto any rod. The puzzle begins with **n** disks stacked on a **start** rod in order of decreasing size, the smallest at the top, thus approximating a conical shape. The objective of the puzzle is to move the entire stack to the **end** rod, obeying the following rules:

- i Only one disk may be moved at a time;
- ii Each move consists of taking the top disk from one of the rods and placing it on top of another rod or on an empty rod;
- iii No disk may be placed on top of a disk that is smaller than it.



Please design a `MOVE(n, start, end)` function to illustrate the minimum number of steps of moving  $n$  disks from start rod to the end rod.

You **MUST** use the following functions and format, otherwise you will not get full points of part (a):

```
def PRINT(origin , destination):  
    print("Move the top disk from rod", origin , "to rod", destination)  
  
def MOVE(n, start , end): # TODO: you need to design this function  
    pass
```

For example, the output of `MOVE(2, 1, 3)` should be:

```
Move the top disk from rod 1 to rod 2  
Move the top disk from rod 1 to rod 3  
Move the top disk from rod 2 to rod 3
```

- (a) Fill in the function `MOVE(n, start, end)` shown above. You can use Python, C/C++ or pseudo-code

form, as you want.

- (b) Suppose the minimum number of moves of  $\text{MOVE}(n, 1, 3)$  is  $f(n)$ . What's the relationship between  $f(n)$  and  $f(n-1)$ ? Please write the recurrence and then show  $f(n) = 2^n - 1$ .

4. Consider a variation of the Towers of Hanoi problem mentioned in Q3, in which there are 4 rods, instead of 3. All the other rules are the same.

Design a strategy to move the  $n$  disks from the first rod to the last rod with at most  $2^{O(\sqrt{n})}$  moves, that is, the number of moves should be bounded by  $2^{c\sqrt{n}}$  for some constant  $c$ . You are free to use the conclusion in Q3: there is a strategy to solve the 3-rod version problem in  $f(n) = 2^n - 1$  moves.

- (a) Suppose in the 4-rod version, to move  $n$  disks to another rod requires  $T(n)$  moves, and we are to show that  $T(n) \leq 2^{c\sqrt{n}}$  for some constant  $c$ . Now consider we have  $x$  disks. First, we move  $x - a$  disks to the third rod, then move the remaining  $a$  disks to the last rod, and finally move the  $x - a$  disks from the third rod to the last rod. Following this way, how can you express  $T(x)$ ? You are supposed to use  $T(\cdot)$ ,  $x$  and  $a$  in the expression.
- (b) Now we can generate a sequence of recurrences by letting  $a = \sqrt{n}$  (suppose  $n$  is some  $k^2$  to avoid corner cases), and  $x = n, n-a, n-2a, \dots$ . Please solve for the  $T(n)$  and show that  $c = 2$  is sufficient for  $T(n) \leq 2^{c\sqrt{n}}$ .