

$$\text{LCS}(X[1..N], Y[1..M]) \triangleq \text{LCS}(N, M)$$

$$\text{LCS}(X[1..i], Y[1..j]) \triangleq \text{LCS}(i, j)$$

$$i \leq N$$

$$j \leq M$$

$$|\text{LCS}(N, M)| \geq |\text{LCS}(i, j)|$$



$$\text{LCS}(N, M) \neq$$

$$\text{LCS}(N-1, M) \checkmark$$

$$\text{LCS}(N-1, M-1) \times$$

$$\text{LCS}(N, M-1) \checkmark$$

$A[i, j]$ as the set of common subsequence...

between $X[1..i]$ and $Y[1..j]$ ^{requires}

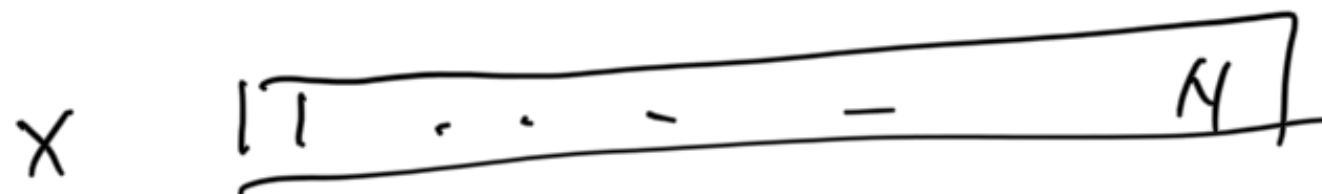
$$A[i_1, j_1] \subseteq A[i_2, j_2] \text{ if } i_1 \leq i_2, j_1 \leq j_2$$

$$A[N, M]$$

$$A[N-1, M] \subseteq A[N, M]$$

$$A[N, M-1] \subseteq A[N, M]$$

$$\underline{A[N, M]} \neq \underline{A[N-1, M]} \cup \underline{A[N, M-1]}$$



$$CS \in A[N, M], \quad \underline{CS \notin A[N-1, M]}, \quad CS \notin A[N, M-1]$$

CS is a ~~subseq~~ subseq of $Y[1..M]$

not a ~~subseq~~ subseq.

of $X[1..N-1]$

$X[1..N]$

$$CS(k) = X[N]$$

$$CS(k) = Y[M]$$

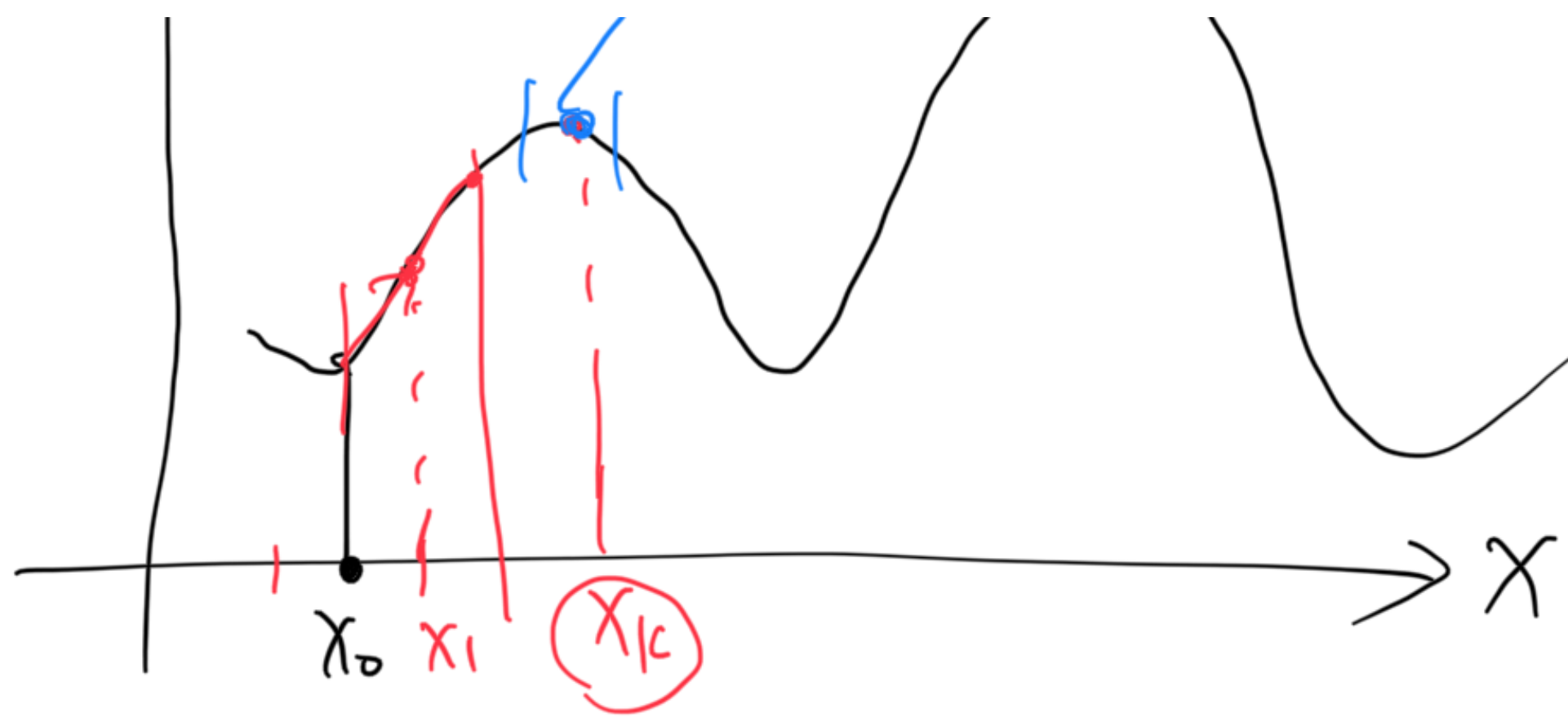
Structure

$$A[N, M] =$$

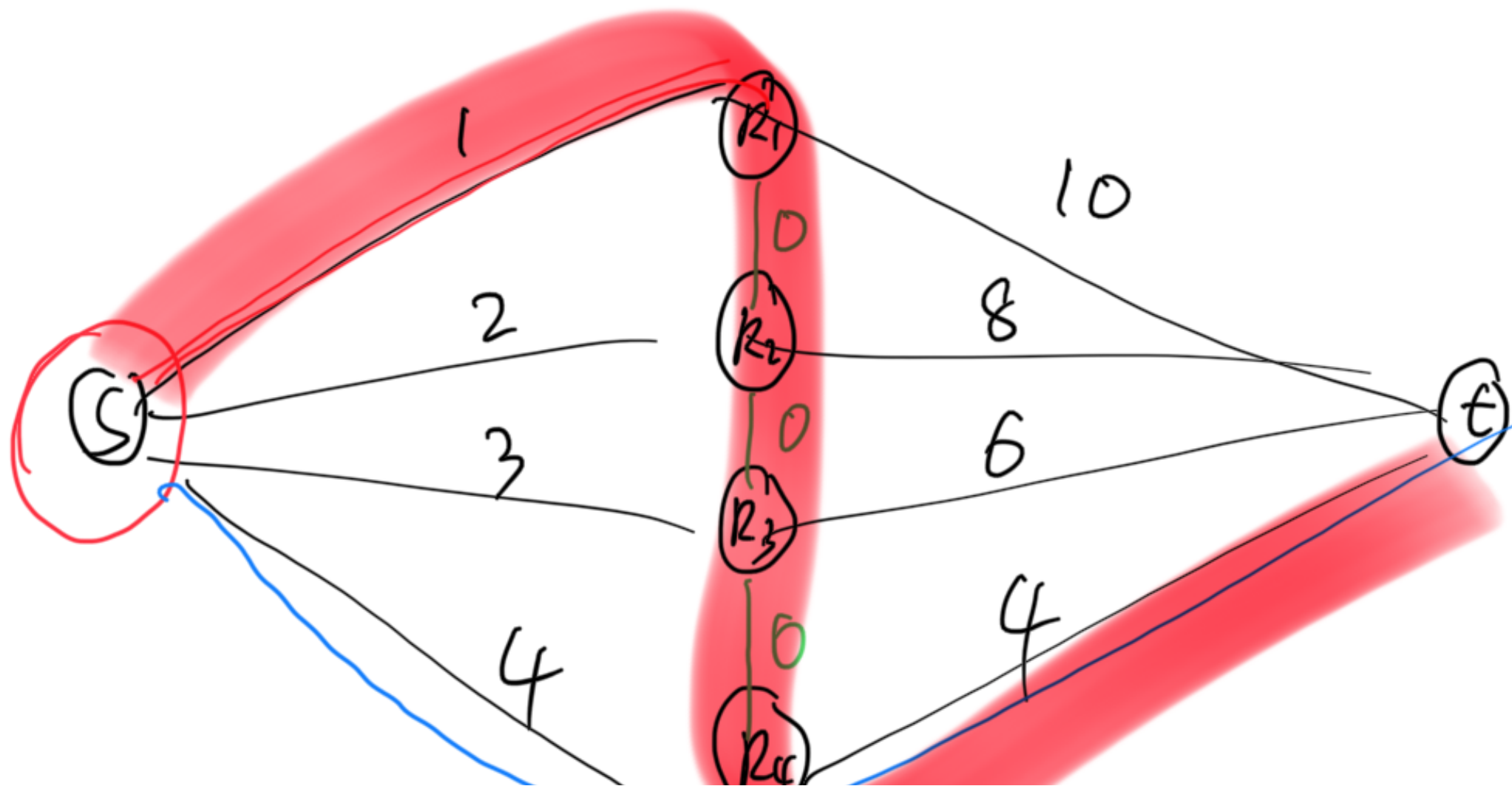
$$\begin{cases} A[N-1, M] \cup A[N, M-1] & \text{if } X[N] \neq Y[M] \\ A[N-1, M] \cup A[N, M-1] & \text{if } X[N] = Y[M] \\ \cup [A[N-1, M-1] + X[N]] \end{cases}$$

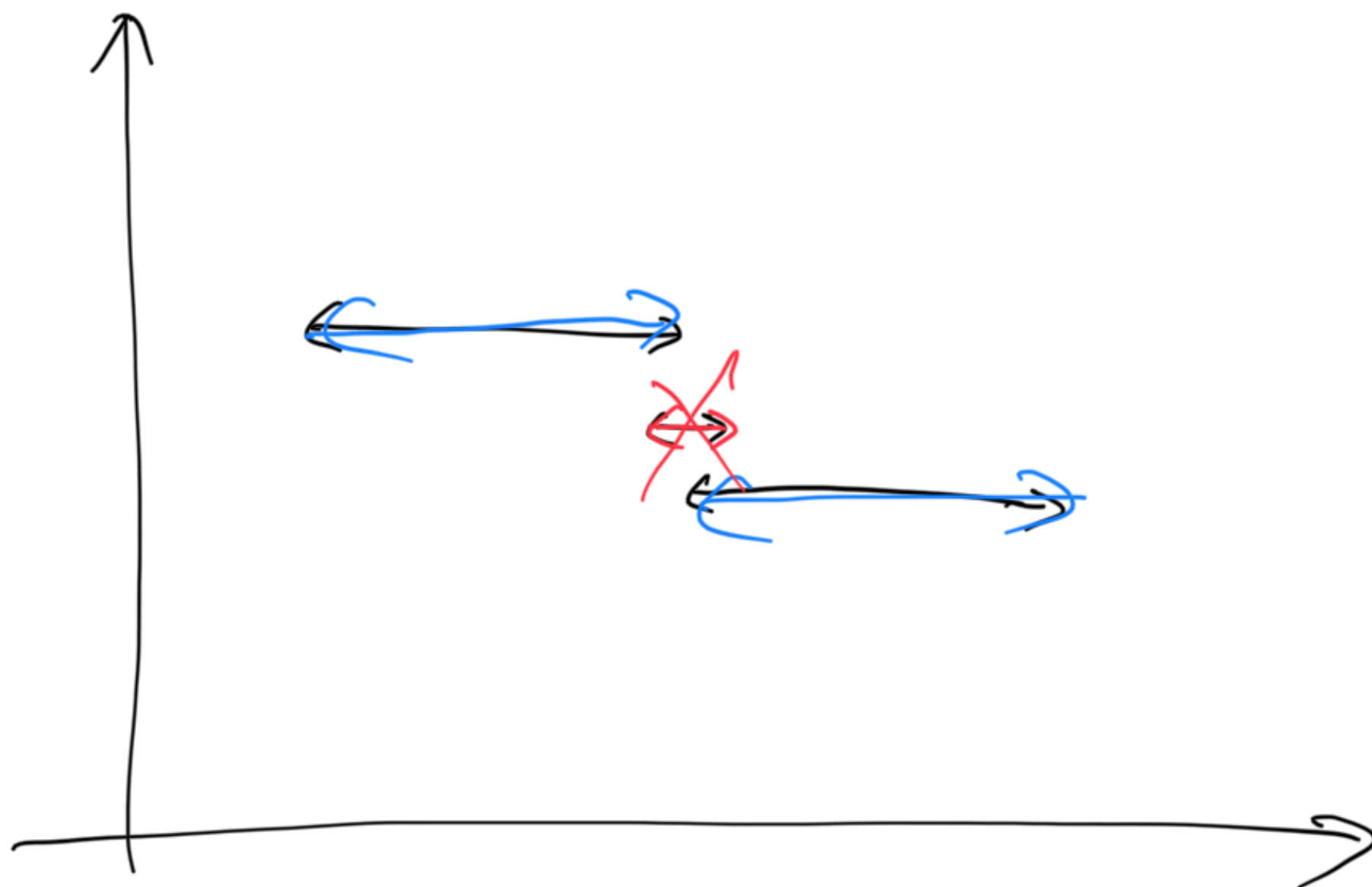
$\uparrow f(x)$

hill climbing
local maximum?

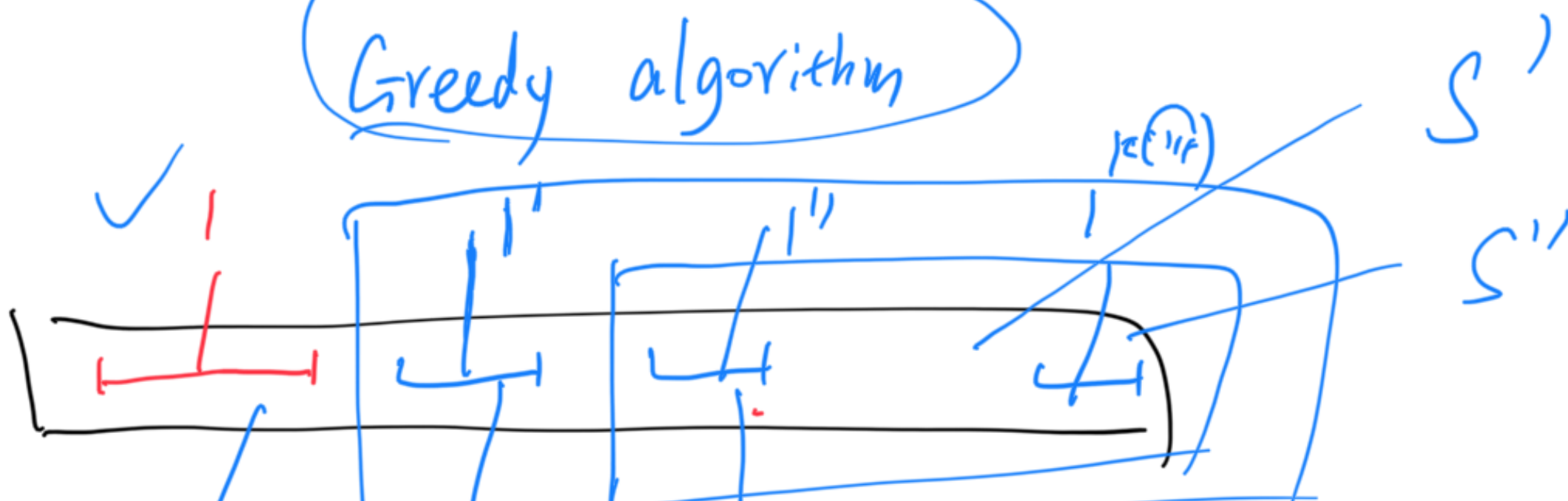


Local optimum vs. global optimum





Greedy algorithm



(A)



$w, 1 \dots n$ — S

$w - w_j, 1 \dots j-1, j+1, \dots, n$ — $S' = S - \{j\}$

~~X~~ S' is not optimal

S'' is optimal

$S'' + \{j\}$

will be a better packing than S

$\text{weight}(S'') \leq w - w_j$

$$\text{weight}(S'' + \{j\}) \leq w$$

$$\begin{aligned} \text{value}(S'' + \{j\}) &= \underline{\text{value}(S'')} + \text{value}(j) \\ &> \text{value}(S') + V_j \\ &= \cancel{\text{value}(S)} \end{aligned}$$