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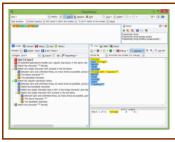
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## Lookahead and Lookbehind Zero-Length Assertions

Lookahead and lookbehind, collectively called "lookaround", are zero-length assertions just like the <u>start a line</u>, and <u>start and end of word</u> anchors explained earlier in this tutorial. The difference is that lookaroun matches characters, but then gives up the match, returning only the result: match or no match. That is are called "assertions". They do not consume characters in the string, but only assert whether a match is or not. Lookaround allows you to create regular expressions that are impossible to create without ther would get very longwinded without them.

# **Positive and Negative Lookahead**

Negative lookahead is indispensable if you want to match something not followed by something elsexplaining character classes, this tutorial explained why you cannot use a negated character class to match followed by a  $\overline{\underline{u}}$ . Negative lookahead provides the solution:  $\overline{\underline{q(?!u)}}$ . The negative lookahead construct is to parentheses, with the opening parenthesis followed by a question mark and an exclamation point. I lookahead, we have the trivial regex  $\overline{u}$ .

Positive lookahead works just the same.  $\overline{q(?=u)}$  matches a q that is followed by a u, without making the the match. The positive lookahead construct is a pair of parentheses, with the opening parenthesis follo question mark and an equals sign.

You can use any regular expression inside the lookahead (but not lookbehind, as explained below). regular expression can be used inside the lookahead. If it contains <u>capturing groups</u> then those groups w as normal and backreferences to them will work normally, even outside the lookahead. (The only except which treats all groups inside lookahead as non-capturing.) The lookahead itself is not a capturing group included in the count towards numbering the backreferences. If you want to store the match of the <u>regestate lookahead</u>, you have to put capturing parentheses around the regex inside the lookahead, like this: <u>(?=(</u>) The other way around will not work, because the lookahead will already have discarded the regex mat time the capturing group is to store its match.

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## **Regex Engine Internals**

First, let's see how the engine applies  $\overline{q(?!u)}$  to the string  $\overline{1raq}$ . The first token in the regex is the <u>literal</u> already know, this causes the engine to traverse the string until the  $\overline{q}$  in the string is matched. The posit string is now the void after the string. The next token is the lookahead. The engine takes note that it is lookahead construct now, and begins matching the regex inside the lookahead. So the next token is  $\overline{u}$ . not match the void after the string. The engine notes that the regex inside the lookahead failed. Bec lookahead is negative, this means that the lookahead has successfully matched at the current positio point, the entire regex has matched, and  $\overline{q}$  is returned as the match.

Let's try applying the same regex to  $\overline{\underline{quit}}$ .  $\overline{\underline{q}}$  matches  $\overline{\underline{q}}$ . The next token is the  $\overline{\underline{u}}$  inside the lookahead. character is the  $\overline{\underline{u}}$ . These match. The engine advances to the next character:  $\overline{\underline{1}}$ . However, it is done with inside the lookahead. The engine notes success, and discards the regex match. This causes the engine back in the string to  $\overline{\underline{u}}$ .

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Because the lookahead is negative, the successful match inside it causes the lookahead to fail. Since the other permutations of this regex, the engine has to start again at the beginning. Since  $\frac{1}{9}$  cannot match  $\frac{1}{9}$  else, the engine reports failure.

Let's take one more look inside, to make sure you understand the implications of the lookahead. Let  $\underline{q(?=u)}$   $\underline{1}$  to  $\underline{qu}$   $\underline{i}$   $\underline{t}$ . The lookahead is now positive and is followed by another token. Again,  $\underline{q}$  matches matches  $\underline{u}$ . Again, the match from the lookahead must be discarded, so the engine steps back from  $\underline{i}$  in to  $\underline{u}$ . The lookahead was successful, so the engine continues with  $\underline{i}$ . But  $\underline{i}$  cannot match  $\underline{u}$ . So this matc fails. All remaining attempts fail as well, because there are no more q's in the string.

## **Positive and Negative Lookbehind**

Lookbehind has the same effect, but works backwards. It tells the regex engine to temporarily step bac the string, to check if the text inside the lookbehind can be matched there. (?<!a)b matches a "b" to preceded by an "a", using negative lookbehind. It doesn't match  $\overline{cab}$ , but matches the  $\overline{b}$  (and only the  $\overline{b}$ )  $\overline{debt}$ . (?<=a)b (positive lookbehind) matches the  $\overline{b}$  (and only the  $\overline{b}$ ) in  $\overline{cab}$ , but does not match  $\overline{bed}$  or  $\overline{de}$ 

The construct for positive lookbehind is <a href="(?<=text)">(?<=text)</a>: a pair of parentheses, with the opening parenthesis by a question mark, "less than" symbol, and an equals sign. Negative lookbehind is written as <a href="(?<!text)">(?<!text)</a>, exclamation point instead of an equals sign.

# **More Regex Engine Internals**

Let's apply (?<=a)b to thingamabob. The engine starts with the lookbehind and the first character in the this case, the lookbehind tells the engine to step back one character, and see if a can be matched the engine cannot step back one character because there are no characters before the a. So the lookbehind the engine starts again at the next character, the a (Note that a negative lookbehind would have succeed Again, the engine temporarily steps back one character to check if an "a" can be found there. It finds a positive lookbehind fails again.

The lookbehind continues to fail until the regex reaches the  $\underline{m}$  in the string. The engine again steps character, and notices that the  $\underline{a}$  can be matched there. The positive lookbehind matches. Because i length, the current position in the string remains at the  $\underline{m}$ . The next token is  $\underline{b}$ , which cannot match here. character is the second  $\underline{a}$  in the string. The engine steps back, and finds out that the  $\underline{m}$  does not match  $\underline{a}$ .

The next character is the first  $\frac{1}{b}$  in the string. The engine steps back and finds out that  $\frac{1}{a}$  satisfies the look matches  $\frac{1}{b}$ , and the entire regex has been matched successfully. It matches one character: the first  $\frac{1}{b}$  in the

## **Important Notes About Lookbehind**

The bad news is that most regex flavors do not allow you to use just any regex inside a lookbehind, because cannot apply a regular expression backwards. The regular expression engine needs to be able to figure many characters to step back before checking the lookbehind. When evaluating the lookbehind, the regedetermines the length of the regex inside the lookbehind, steps back that many characters in the subject and then applies the regex inside the lookbehind from left to right just as it would with a normal regex.

Many regex flavors, including those used by <u>Perl</u> and <u>Python</u>, only allow fixed-length strings. You can use <u>text</u>, <u>character escapes</u>, <u>Unicode escapes</u> other than <u>X</u>, and <u>character classes</u>. You cannot use <u>qual backreferences</u>. You can use <u>alternation</u>, but only if all alternatives have the same length. These flavors lookbehind by first stepping back through the subject string for as many characters as the lookbehind not then attempting the regex inside the lookbehind from left to right.

PCRE is not fully Perl-compatible when it comes to lookbehind. While Perl requires alternatives inside look have the same length, PCRE allows alternatives of variable length. PHP, Delphi, R, and Ruby also allow t alternative still has to be fixed-length. Each alternative is treated as a separate fixed-length lookbehind.

<u>Java</u> takes things a step further by allowing finite repetition. You still cannot use the <u>star</u> or <u>plus</u>, but you the <u>question mark</u> and the <u>curly braces</u> with the <u>max</u> parameter specified. <u>Java</u> determines the minin maximum possible lengths of the lookbehind. The lookbehind in the regex  $(?<!ab\{2,4\}c\{3,5\}d)$  te possible lengths. It can be between 7 to 11 characters long. When Java (version 6 or later) tries to r lookbehind, it first steps back the minimum number of characters (7 in this example) in the string evaluates the regex inside the lookbehind as usual, from left to right. If it fails, Java steps back one more and tries again. If the lookbehind continues to fail, Java continues to step back until the lookbehind either or it has stepped back the maximum number of characters (11 in this example). This repeated stept through the subject string kills performance when the number of possible lengths of the lookbehind gro this in mind. Don't choose an arbitrarily large maximum number of repetitions to work around the lack

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quantifiers inside lookbehind. Java 4 and 5 have bugs that cause lookbehind with alternation or variable q to fail when it should succeed in some situations. These bugs were fixed in Java 6.

The only regex engines that allow you to use a full regular expression inside lookbehind, including infinite and backreferences, are the <u>JGsoft engine</u> and the <u>NET framework RegEx classes</u>. These regex engil apply the regex inside the lookbehind backwards, going through the regex inside the lookbehind and this subject string from right to left. They only need to evaluate the lookbehind once, regardless of how many possible lengths it has.

Finally, flavors like JavaScript and Tcl do not support lookbehind at all, even though they do support looka

### **Lookaround Is Atomic**

The fact that lookaround is zero-length automatically makes it <u>atomic</u>. As soon as the lookaround co satisfied, the regex engine forgets about everything inside the lookaround. It will not backtrack in lookaround to try different permutations.

The only situation in which this makes any difference is when you use <u>capturing groups</u> inside the loc Since the regex engine does not backtrack into the lookaround, it will not try different permutations of the groups.

For this reason, the regex  $(?=(\d+))\w+\1$  never matches  $123\times12$ . First the lookaround captures 12:  $\w+\d+$  then matches the whole string and backtracks until it matches only 1. Finally,  $\w+\d+$  fails since  $\w+\d+$  matched at any position. Now, the regex engine has nothing to backtrack to, and the overall regex backtracking steps created by  $\w+\d+$  have been discarded. It never gets to the point where the lookahead only 12.

Obviously, the regex engine does try further positions in the string. If we change the subject string,  $\frac{(?=(\d+))\w+\1}{\d+1}$  does match  $\frac{56x56}{\d+1}$  in  $\frac{456x56}{\d+1}$ .

If you don't use capturing groups inside lookaround, then all this doesn't matter. Either the lookaround can be satisfied or it cannot be. In how many ways it can be satisfied is irrelevant.

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