

- Predicate Testing

Predicate Testing

- Introduction
- Basic Concepts
- Predicate Coverage
- Summary

- **Predicates** are expressions that can be evaluated to **a boolean value**, i.e., true or false.
- Many decision points can be encoded as a predicate, i.e., which action should be taken under what condition?
- **Predicate-based testing** is about ensuring those predicates are implemented correctly.

- **Program-based**: Predicates can be identified from the branching points of the source code
 - e.g.: `if ((a > b) || c) { ... } else { ... }`
- **Specification-based**: Predicates can be identified from both formal and informal requirements as well as behavioral models such as FSM
 - “if the printer is ON and has paper then send the document for printing”
- **Predicate testing** is required by US FAA for safety critical avionics software in commercial aircraft

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- A **predicate** is an expression that evaluates to a Boolean value
- Predicates may contain:
 - **Boolean** variables
 - **Non-boolean** variables that are compared with the relational operators $\{>, <, =, \geq, \leq, \neq\}$
 - **Boolean** function calls
- The internal structure is created by **logical operators**:
 $\neg, \wedge, \vee, \rightarrow, \oplus, \leftrightarrow$

Logical operators

- \neg – the *negation* operator
- \wedge – the *and* operator
- \vee – the *or* operator
- \rightarrow – the *implication* operator
- \oplus – the *exclusive or* operator
- \leftrightarrow – the *equivalence* operator

- A **clause** is a predicate that **does not** contain any of the logical operators
- Example: $(a = b) \vee C \wedge p(x)$ has three clauses:
 - a **relational expression** $(a = b)$,
 - a **boolean variable** C ,
 - a **boolean function** call $p(x)$

Predicate Faults

- An incorrect **Boolean** *operator* is used
- An incorrect **Boolean** *variable* is used
- **Missing** or **extra** Boolean **variables**
- An incorrect relational operator is used
- **Parentheses** are used **incorrectly**

Example

- Assume that $(a < b) \vee (c > d) \wedge e$ is a correct Boolean expression:
 - $(a < b) \wedge (c > d) \wedge e$
 - $(a < b) \vee (c > d) \wedge f$
 - $(a < b) \vee (c > d)$
 - $(a = b) \vee (c > d) \wedge e$
 - $(a = b) \vee (c \leq d) \wedge e$
 - $(a < b \vee c > d) \wedge e$

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- Program-Based Predicate Testing
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- P is the set of predicates
- p is a single predicate in P
- C is the set of clauses in P
- C_p is the set of clauses in predicate p
- c is a single clause in C

Predicate Coverage (PC)

- The first (and simplest) two criteria require that each **predicate** and each **clause** be evaluated to both **true** and **false** and each clause be evaluated to both **true** and **false**
- For each predicate **p**, TR contains two requirements: **p** evaluates to **true**, and **p** evaluates to **false**.
- Example: $p = ((a > b) \vee C) \wedge p(x)$

| | a | b | C | p(x) |
|---|---|---|-------|-------|
| 1 | 5 | 4 | true | true |
| 2 | 5 | 6 | false | false |

Predicate Coverage Example

$$p = ((a < b) \vee D) \wedge (m \geq n * o)$$

Predicate = true

$a = 5, b = 10, D = \text{true}, m = 1, n = 1, o = 1$
 $= (5 < 10) \vee \text{true} \wedge (1 \geq 1 * 1)$
 $= \text{true} \vee \text{true} \wedge \text{TRUE}$
 $= \text{true}$

Predicate = false

$a = 10, b = 5, D = \text{false}, m = 1, n = 1, o = 1$
 $= (10 < 5) \vee \text{false} \wedge (1 \geq 1 * 1)$
 $= \text{false} \vee \text{false} \wedge \text{TRUE}$
 $= \text{false}$

Clause Coverage (CC)

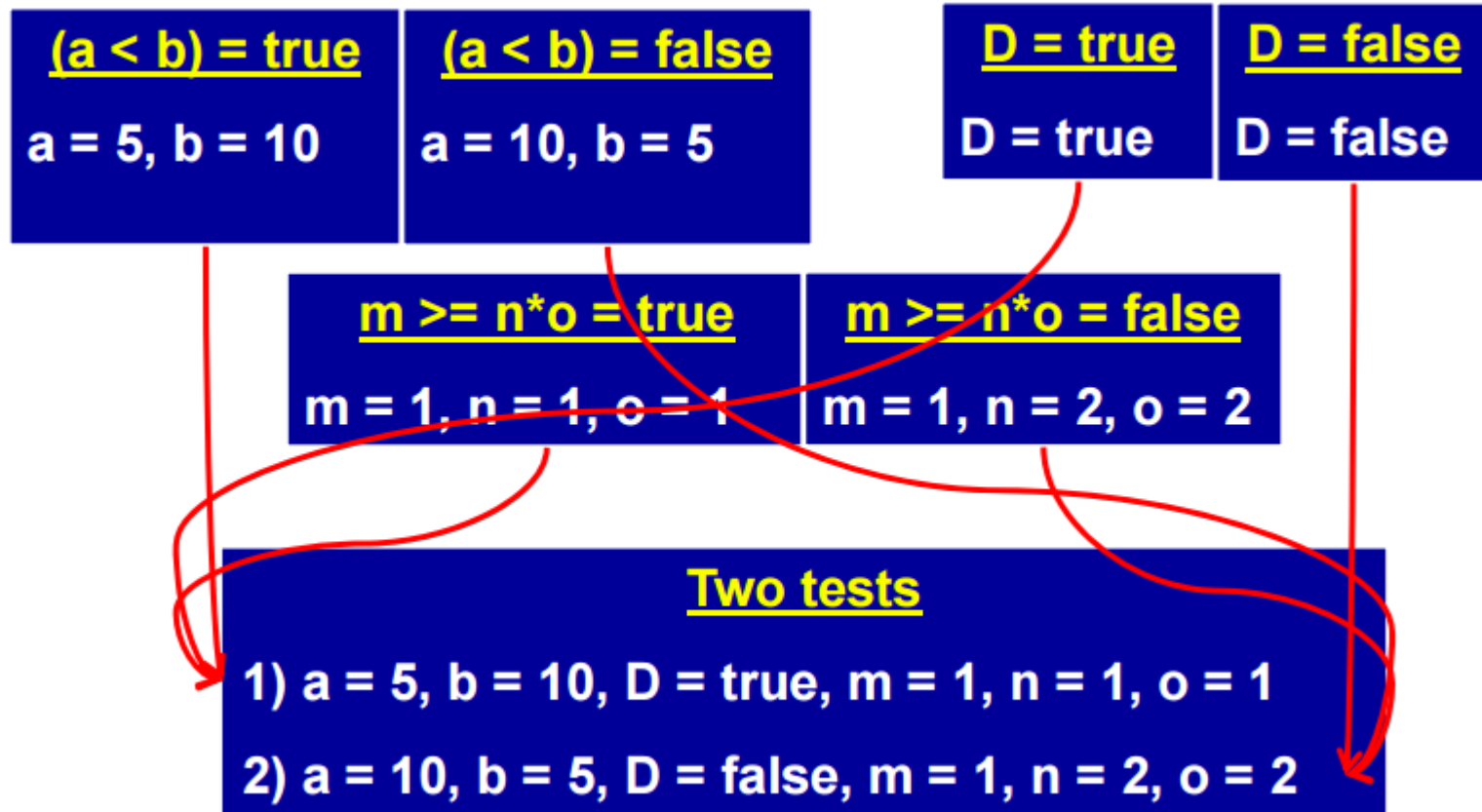
- For each clause c , TR contains two requirements: c evaluates to **true**, and c evaluates to **false**.
- Example: $((a > b) \vee C) \wedge p(x)$

| | a | b | C | p(x) |
|---|---|---|-------|-------|
| 1 | 5 | 4 | true | true |
| 2 | 5 | 6 | false | false |

“condition coverage” in literature

Clause Coverage Example

$$P = ((a < b) \vee D) \wedge (m \geq n * o)$$



Predicate vs. Clause Coverage

- Does **predicate** coverage subsume **clause** coverage? Does **clause** coverage subsume **predicate** coverage?
- Example: $p = a \vee b$

| | a | b | $a \vee b$ |
|---|---|---|------------|
| 1 | T | T | T |
| 2 | T | F | T |
| 3 | F | T | T |
| 4 | F | F | F |

Naturally, we want to test both the predicate and individual clauses.

- CC does not always ensure PC
- This is, we can satisfy CC without causing the predicate to be both true or false
- This is definitely not what we want!!
 - We need to come up with other approaches

Combinatorial Coverage (CoC)

- For each predicate p , TR has test requirements for the clauses in p to evaluate to each possible combination of **truth values**
- Example: $(a \vee b) \wedge c$

CoC requires every possible combination

| | a | b | c | $(a \vee b) \wedge c$ |
|---|---|---|---|-----------------------|
| 1 | T | T | T | T |
| 2 | T | T | F | F |
| 3 | T | F | T | T |
| 4 | T | F | F | F |
| 5 | F | T | T | T |
| 6 | F | T | F | F |
| 7 | F | F | T | F |
| 8 | F | F | F | F |

Combinatorial Coverage (CoC) – Exercise -1

Write all the clauses and the CoC of the given predicate:

$$P = ((a > b) \vee C) \wedge p(x)$$

Combinatorial Coverage (CoC) – Exercise -1

$$P = ((a > b) \vee C) \wedge p(x)$$

$$P = (X \vee Y) \wedge Z$$

| | X | Y | Z | Predicate |
|---|---|---|---|-----------|
| 1 | F | F | F | F |
| 2 | F | F | T | F |
| 3 | F | T | F | F |
| 4 | F | T | T | T |
| 5 | T | F | F | F |
| 6 | T | F | T | T |
| 7 | T | T | F | F |
| 8 | T | T | T | T |

Z is more important clause in this predicate than the others

What is the problem with combinatorial coverage ?

Combinatorial coverage is very expensive if we have multiple clauses in the predicate.

- 2^n possibilities, which n is number of independent clauses.

- **Major clause**
 - The clause which is being focused upon;
- **Minor clause**
 - All other clauses in the predicate (everything else).
- **Determination:**
 - A clause c_i in predicate p , called the major clause, determines p if and only if the values of the remaining minor clauses c_i are such that changing c_i changes the value of p (c_i Controls the behavior)
- *In the previous example, if we chose Z as **Major clause**, when it has value of “False”, it doesn’t matter what the other clauses are, but when it is “True”, it does matter what other clauses are*

Determining Predicates

$$\underline{P = A \vee B}$$

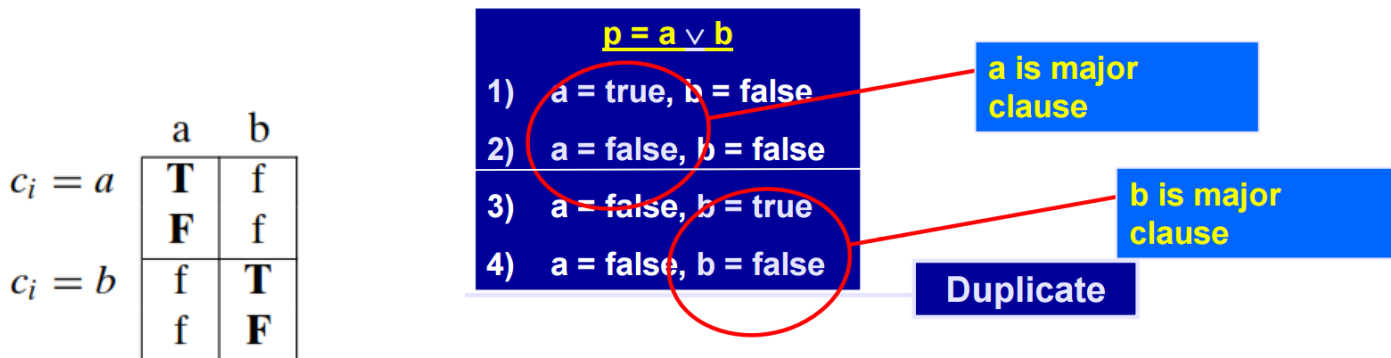
if $B = \text{true}$, p is always true.
so if $B = \text{false}$, A determines p .
if $A = \text{false}$, B determines p .

$$\underline{P = A \wedge B}$$

if $B = \text{false}$, p is always false.
so if $B = \text{true}$, A determines p .
if $A = \text{true}$, B determines p .

Active Clause Coverage (ACC)

- For each predicate p and each major clause c of p , choose minor clauses so that c determines p . TR has two requirements for each c : c evaluates to **true** and c evaluates to **false**.



Two of these requirements are identical, so we end up with **three distinct** test requirements for active clause coverage for the predicate $a \vee b$, namely, $\{(a = \text{true}, b = \text{false}), (a = \text{false}, b = \text{true}), (a = \text{false}, b = \text{false})\}$

Active Clause Coverage (ACC) – Example (1)

```
public static void printHonorRollStatus(double cumulativeGPA,  
    double termGPA, int creditsCompleted, boolean fullTimeStatus) {  
    // Determine if the student is on the deans list.  
    if ((creditsCompleted > 30) && (cumulativeGPA > 3.20)  
        && (fullTimeStatus == true) && (termGPA > 2.0)) {  
        System.out.println("You are on the dean's list.");  
    } else if ((creditsCompleted > 30) && (cumulativeGPA > 3.70)  
        && (fullTimeStatus == true) && (termGPA > 2.0)) {  
        System.out.println("You are on the high honors dean's list.");  
    } else if ((creditsCompleted > 30) && (cumulativeGPA > 2.0)  
        && (fullTimeStatus == true) && (termGPA > 3.2)) {  
        System.out.println("You are on the honor list.");  
    } else {  
        // ...  
    }  
}
```

Active Clause Coverage (ACC) – Example (2)

| | CC | CGPA | FT | TGPA | Predicate |
|---|----|------|-------|------|-----------|
| | 29 | 3.3 | True | 2.4 | False |
| → | 31 | 3.3 | True | 2.4 | True |
| | 31 | 3.0 | True | 2.4 | False |
| → | 31 | 3.3 | True | 2.4 | True |
| | 31 | 3.3 | False | 2.4 | False |
| → | 31 | 3.3 | True | 2.4 | True |
| | 31 | 3.3 | True | 1.9 | False |
| → | 31 | 3.3 | True | 2.4 | True |

- The Green cells indicate active clauses,
- The Orange color cells indicate minor clauses
- We can test this with only **5 test cases**

Resolving the Ambiguity

$p = a \vee (b \wedge c)$

Major clause : a

a = true, b = false, c = true

a = false, b = false, c = false

Is this allowed ?

The diagram shows a blue box containing the logical expression $p = a \vee (b \wedge c)$ and two sets of variable assignments. The first set is 'a = true, b = false, c = true' and the second is 'a = false, b = false, c = false'. The expression 'c = false' in the second set is circled in blue. A callout box with a blue border and white background, containing the text 'Is this allowed ?', has a line pointing to the circled 'c = false'.

- This question caused confusion among testers for years
- Considering this carefully leads to three separate criteria :
 - Minor clauses do not need to be the same (GACC)
 - Minor clauses do need to be the same (RACC)
 - Minor clauses force the predicate to become both true and false (CACC)

General Active Clause Coverage (GACC)

- The same as ACC, and it **does not require** the minor clauses have the same values when the major clause evaluates to **true** and **false**.
- Does **GACC** subsume **predicate** coverage?

| | a | b | c | $a \wedge (b \vee c)$ |
|---|---|---|---|-----------------------|
| 1 | T | T | T | T |
| 2 | T | T | F | T |
| 3 | T | F | T | T |
| 4 | T | F | F | F |
| 5 | F | T | T | F |
| 6 | F | T | F | F |
| 7 | F | F | T | F |
| 8 | F | F | F | F |

Correlated Active Clause Coverage (CACC)

- The same as ACC, but it requires the **entire predicate** to be **true** for one value of the major clause and **false** for the other.

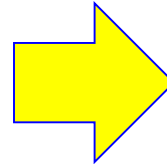
CACC (2)

- Example: $p = a \wedge (b \vee c)$

For a to determine the value of p , the expression $b \vee c$ must be true

This can be achieved in three ways: b true and c false, b false and c true, and both b and c true

| | a | b | c | $a \wedge (b \vee c)$ |
|---|---|---|---|-----------------------|
| 1 | T | T | T | T |
| 2 | T | T | F | T |
| 3 | T | F | T | T |
| 4 | T | F | F | F |
| 5 | F | T | T | F |
| 6 | F | T | F | F |
| 7 | F | F | T | F |
| 8 | F | F | F | F |



| | a | b | c | $a \wedge (b \vee c)$ |
|---|---|---|---|-----------------------|
| 1 | T | T | T | T |
| 2 | T | T | F | T |
| 3 | T | F | T | T |
| 5 | F | T | T | F |
| 6 | F | T | F | F |
| 7 | F | F | T | F |

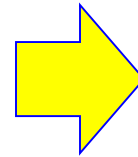
Rows 4 and 8 are missing because a is not active in the two rows.

(1,5) (1,6) (1,7)
 (2,5)(2,6)(2,7)
 (3,5)(3,6)(3,7)

Restricted Active Clause Coverage (RACC)

- The same as ACC, but it requires the minor clauses have the **same values** when the major clause evaluates to **true** and **false**.
- Example: $p = a \wedge (b \vee c)$

| | a | b | c | $a \wedge (b \vee c)$ |
|---|---|---|---|-----------------------|
| 1 | T | T | T | T |
| 2 | T | T | F | T |
| 3 | T | F | T | T |
| 4 | T | F | F | F |
| 5 | F | T | T | F |
| 6 | F | T | F | F |
| 7 | F | F | T | F |
| 8 | F | F | F | F |



| | a | b | c | $a \wedge (b \vee c)$ |
|---|---|---|---|-----------------------|
| 1 | T | T | T | T |
| 5 | F | T | T | F |
| 2 | T | T | F | T |
| 6 | F | T | F | F |
| 3 | T | F | T | T |
| 7 | F | F | T | F |

RACC can only be satisfied by one of the three pairs
(1,5) (2,6) (3,7)

CACC and RACC

| | a | b | c | $a \wedge (b \vee c)$ |
|---|---|---|---|-----------------------|
| 1 | T | T | T | T |
| 2 | T | T | F | T |
| 3 | T | F | T | T |
| 4 | T | F | F | F |
| 5 | F | T | T | F |
| 6 | F | T | F | F |
| 7 | F | F | T | F |
| 8 | F | F | F | F |

major clause

$P_a : b = \text{true} \text{ or } c = \text{true}$

CACC can be satisfied by choosing any of rows 1, 2, 3 AND any of rows 5, 6, 7 – a total of nine pairs

| | a | b | c | $a \wedge (b \vee c)$ |
|---|---|---|---|-----------------------|
| 1 | T | T | T | T |
| 2 | T | T | F | T |
| 3 | T | F | T | T |
| 4 | T | F | F | F |
| 5 | F | T | T | F |
| 6 | F | T | F | F |
| 7 | F | F | T | F |
| 8 | F | F | F | F |

RACC can only be satisfied by row pairs (1, 5), (2, 6), or (3, 7)
Only three pairs

- **Active clause coverage** criteria ensure that “major” clauses do affect the predicates
- **Inactive clause coverage** takes the opposite approach - major clauses do not affect the predicates
- For each predicate **p** and each major clause **c** of **p**, choose minor clauses so that **c** does not determine **p**.
TR has four requirements for each clause **c**:
 - 1) **c** evaluates to **true** with **p** = **true**
 - 2) **c** evaluates to **false** with **p** = **true**
 - 3) **c** evaluates to **true** with **p** = **false**
 - 4) **c** evaluates to **false** with **p** = **false**.

- **GICC** does not require the values of the minor clauses to be the same when the major clause evaluates to **true** and **false**.
- **RICC** requires the values the minor clauses to be the same when the major clause evaluates to **true** and **false**.

Making Clauses Determine a Predicate

- Let p be a predicate and c a clause. Let $p_{c=\text{true}}$ (or $p_{c=\text{false}}$) be the predicate obtained by replacing every occurrence of c with true (or false)
- The following expression describes the exact conditions under which the value of c determines the value of p :

$$p_c = p_{c=\text{true}} \oplus p_{c=\text{false}}$$

Example-1

$$\underline{p = a \vee b}$$

$$\begin{aligned} p_a &= p_{a=\text{true}} \oplus p_{a=\text{false}} \\ &= (\text{true} \vee b) \text{ XOR } (\text{false} \vee b) \\ &= \text{true XOR } b \\ &= \neg b \end{aligned}$$

$$\underline{p = a \wedge b}$$

$$\begin{aligned} p_a &= p_{a=\text{true}} \oplus p_{a=\text{false}} \\ &= (\text{true} \wedge b) \oplus (\text{false} \wedge b) \\ &= b \oplus \text{false} \\ &= b \end{aligned}$$

$$\underline{p = a \vee (b \wedge c)}$$

$$\begin{aligned} p_a &= p_{a=\text{true}} \oplus p_{a=\text{false}} \\ &= (\text{true} \vee (b \wedge c)) \oplus (\text{false} \vee (b \wedge c)) \\ &= \text{true} \oplus (b \wedge c) \\ &= \neg (b \wedge c) \\ &= \neg b \vee \neg c \end{aligned}$$

“NOT b v NOT c” means either b or c can be false

Example-2

Compute (and simplify) the conditions under which each of the clauses determines predicate p .

$$p = a \wedge (\neg b \vee c)$$

$$p_a = \neg b \vee c$$

$$p_b = a \wedge \neg c$$

$$p_c = a \wedge b$$

Note that the last step in the simplification may not be immediately obvious. If it is not, try constructing the truth table. For instance, for $(a \wedge c) \oplus a$. The computation for p_c is equivalent and yields the solution $a \wedge \neg c$

Example-3

Compute (and simplify) the conditions under which each of the clauses determines predicate p .

Consider $p = (a \wedge b) \vee (a \wedge \neg b)$

$p_a = \dots$

$p_b = \dots$

Example-3(1)

$$p = (a \wedge b) \vee (a \wedge \neg b)$$

$$\begin{aligned} p_a &= p_{a=\text{true}} \oplus p_{a=\text{false}} \\ &= ((\text{true} \wedge b) \vee (\text{true} \wedge \neg b)) \oplus ((\text{false} \wedge b) \vee (\text{false} \wedge \neg b)) \\ &= (b \vee \neg b) \oplus \text{false} \\ &= \text{true} \oplus \text{false} \\ &= \text{true} \end{aligned}$$

$$p = (a \wedge b) \vee (a \wedge \neg b)$$

$$\begin{aligned} p_b &= p_{b=\text{true}} \oplus p_{b=\text{false}} \\ &= ((a \wedge \text{true}) \vee (a \wedge \neg \text{true})) \oplus ((a \wedge \text{false}) \vee (a \wedge \neg \text{false})) \\ &= (a \vee \text{false}) \oplus (\text{false} \vee a) \\ &= a \oplus a \\ &= \text{false} \end{aligned}$$

- ***a* always determines the value of this predicate**
- ***b* never determines the value – *b* is irrelevant !**

Example-4

$$p = a \wedge (\neg b \vee c)$$

| | a | b | c | p | p_a | p_b | p_c |
|---|---|---|---|---|-------|-------|-------|
| 1 | T | T | T | T | T | F | T |
| 2 | T | T | F | F | F | T | T |
| 3 | T | F | T | T | T | F | F |
| 4 | T | F | F | T | T | T | F |
| 5 | F | T | T | F | T | F | F |
| 6 | F | T | F | F | F | F | F |
| 7 | F | F | T | F | T | F | F |
| 8 | F | F | F | F | T | F | F |

- All pairs of rows satisfying GACC
 - a: {1,3,4} x {5,7,8}, b: {(2,4)}, c: {(1,2)}
- All pairs of rows satisfying CACC
 - Same as GACC
- All pairs of rows satisfying RACC
 - a: {(1,5),(3,7),(4,8)}
 - Same as CACC pairs for b, c
- GICC
 - a: {(2,6)} for p=F, no feasible pair for p=T
 - b: {5,6}x{7,8} for p=F, {(1,3)} for p=T
 - c: {5,7}x{6,8} for p=F, {(3,4)} for p=T
- Conditions under which each of the clauses determines p
 - $p_a: (\neg b \vee c)$
 - $p_b: a \wedge \neg c$
 - $p_c: a \wedge b$
- RICC
 - a: same as GICC
 - b: {(5,7),(6,8)} for p=F, {(1,3)} for p=T
 - c: {(5,6),(7,8)} for p=F, {(3,4)} for p=T

Boolean Algebra Laws

- **Commutativity Laws**

$$a \vee b = b \vee a$$

$$a \wedge b = b \wedge a$$

$$a \oplus b = b \oplus a$$

- **Associativity Laws**

$$(a \vee b) \vee c = a \vee (b \vee c)$$

$$(a \wedge b) \wedge c = a \wedge (b \wedge c)$$

$$(a \oplus b) \oplus c = a \oplus (b \oplus c)$$

- **Distributive Laws**

$$a \wedge (b \vee c) = (a \wedge b) \vee (a \wedge c)$$

$$a \vee (b \wedge c) = (a \vee b) \wedge (a \vee c)$$

- **DeMorgan's Laws**

$$\neg(a \vee b) = \neg a \wedge \neg b$$

$$\neg(a \wedge b) = \neg a \vee \neg b$$

- **Negation Laws**

$$\neg(\neg a) = a$$

$$\neg a \vee a = \text{true}$$

$$\neg a \wedge a = \text{false}$$

$$a \vee \neg a \wedge b = a \vee b$$

- **AND Identity Laws**

$$\text{false} \wedge a = \text{false}$$

$$\text{true} \wedge a = a$$

$$a \wedge a = a$$

$$a \wedge \neg a = \text{false}$$

- **OR Identity Laws**

$$\text{false} \vee a = a$$

$$\text{true} \vee a = \text{true}$$

$$a \vee a = a$$

$$a \vee \neg a = \text{true}$$

- **XOR Identity Laws**

$$\text{false} \oplus a = a$$

$$\text{true} \oplus a = \neg a$$

$$a \oplus a = \text{false}$$

$$a \oplus \neg a = \text{true}$$

- **XOR Equivalence Laws**

$$a \oplus b = (a \wedge \neg b) \vee (\neg a \wedge b)$$

$$a \oplus b = (a \vee b) \wedge (\neg a \vee \neg b)$$

$$a \oplus b = (a \vee b) \wedge \neg(a \wedge b)$$

Infeasible Test Requirements

- Consider the predicate: $(a > b \wedge b > c) \vee c > a$
 $(a > b) = \text{true}, (b > c) = \text{true}, (c > a) = \text{true}$ is infeasible
- As with graph-based criteria, infeasible test requirements have to be recognized and ignored
- Recognizing infeasible test requirements is hard, and in general, undecidable
- Software testing is inexact – engineering, not science

Finding Satisfying Values

How to choose values that satisfy a given coverage goal?

Example (1)

- Consider $p = (a \vee b) \wedge c$:

| | |
|---|--------------------|
| a | $x < y$ |
| b | done |
| c | List.contains(str) |

How to choose values to satisfy predicate coverage?

Example (2)

| | a | b | c | p |
|---|---|---|---|---|
| 1 | t | t | t | t |
| 2 | t | t | f | f |
| 3 | t | f | t | t |
| 4 | t | f | f | f |
| 5 | f | t | t | t |
| 6 | f | t | f | f |
| 7 | f | f | t | f |
| 8 | f | f | f | f |

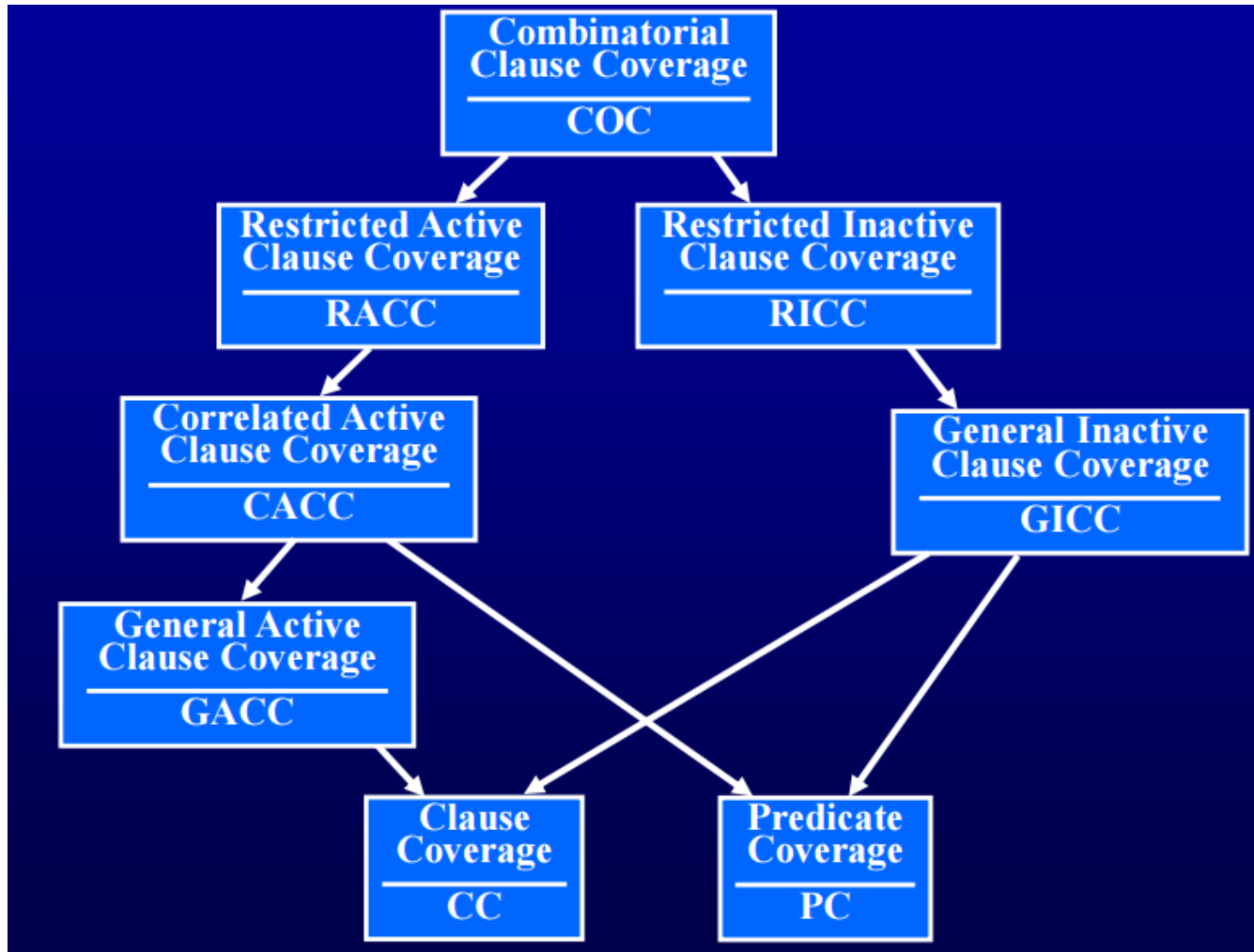
$$\{1, 3, 5\} \times \{2, 4, 6, 7, 8\}$$

Suppose we choose {1, 2}.

| a | b | c |
|-----------------|-------------|--|
| $x = 3 \ y = 5$ | done = true | List = ["Rat", "cat", "dog"] str = "cat" |
| $x = 0, y = 7$ | done = true | List = ["Red", "White"] str = "Blue" |

Logic Coverage Criteria

Subsumption



Recap

- **Predicate testing** is about ensuring that each decision point is implemented correctly.
- If we flip the value of an **active** clause, we will change the value of the entire predicate.
- Different active clause criteria are defined to clarify the requirements on the values of the minor clauses.