

# C477: Mathematical Introduction to Optimisation

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# Outline

## • Topics

- ▶ Global vs local optimisation
- ▶ Neighbourhoods & Other properties of sets
  - ★ Boundary vs interior points; Closed vs open sets; Compact sets
- ▶ Weierstrass Theorem or *How to assert the existence of an optimum*
- ▶ Min vs Argmin
- ▶ Minimisation vs maximisation
- ▶ Defining types of optimisation problems
- ▶ Formulating a nonlinear optimisation problem

## • Example

- ▶ Enclosing points

## • Reading

- ▶ Chap. 6.1 (Introduction), 19.1 (Introduction), & 4.4 (Neighbourhoods) in *An Introduction to Optimization*, Chong & Zak, Third Edition.

## • Acknowledgements

- ▶ Parts of these slides were originally developed by Benoit Chachuat and Panos Parpas.  $\text{\LaTeX}$  design and proof reading by Miten Mistry. Mistakes by Ruth Misener.

# Refresh: Definition: Mathematical Optimisation

Optimisation models (a.k.a. **mathematical programs**) represent problem choices as **decision variables** and seek values that minimise (or maximise) **objective functions** of the decision variables subject to **constraints** on variable values expressing the limits on possible decision choices

$$\begin{array}{ll} \min_{\mathbf{x}} f(\mathbf{x}) & \longleftarrow \text{Objective function} \\ \text{s.t. } \mathbf{h}(\mathbf{x}) = \mathbf{0} & \longleftarrow \text{Equality constraints} \\ \mathbf{g}(\mathbf{x}) \leq \mathbf{0} & \longleftarrow \text{Inequality constraints} \\ \mathbf{x} \in [\mathbf{x}^L, \mathbf{x}^U] = X \subset \mathbb{R}^n & \longleftarrow \text{Variable bounds} \\ \mathbf{h} : \mathbb{R}^n \rightarrow \mathbb{R}^m, \mathbf{g} : \mathbb{R}^n \rightarrow \mathbb{R}^p & \end{array}$$

- **Analytic expressions** of the objective and constraint functions **may or may not be available**
- If we want to find the absolute best set of admissible decisions:  
**Global optimisation**

# Defining (Global) Optimality

## Feasible Set

The feasible set  $S$  (or feasible region) of an optimisation model is the collection of choices for decision variables satisfying **all** of the model

constraints:  $S \triangleq \{\mathbf{x} \mid \mathbf{h}(\mathbf{x}) = \mathbf{0}, \mathbf{g}(\mathbf{x}) \leq \mathbf{0}, \mathbf{x}^L \leq \mathbf{x} \leq \mathbf{x}^U\}$

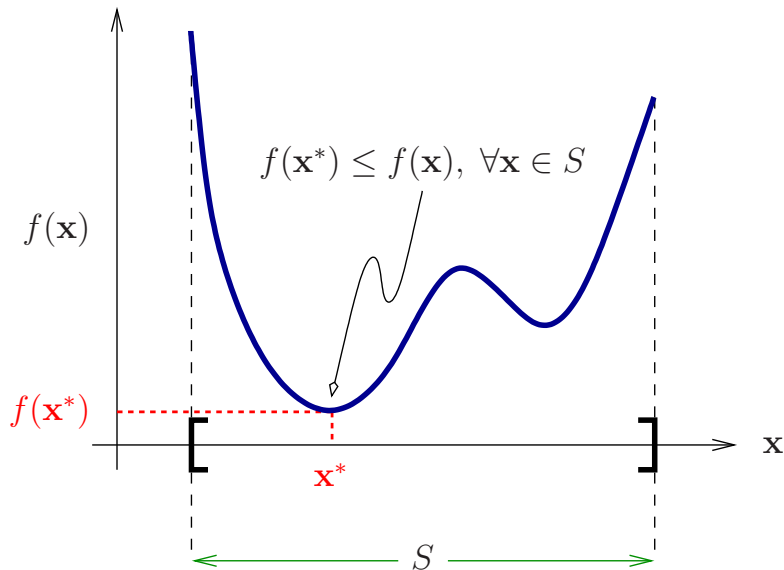
## (Global) Optimum

An **optimal solution**,  $\mathbf{x}^*$ , is a **feasible** choice for decision variables with objective function value **superior** to any other feasible point. For a

minimisation problem:  $f(\mathbf{x}^*) \leq f(\mathbf{x}), \quad \forall \mathbf{x} \in S$

- ❶ The **optimal value**  $f^*$  in an optimisation model is the objective function value of any optimal solutions:  $f^* = f(\mathbf{x}^*)$  — It is **unique**!
- ❷ But, an optimisation model may have:
  - ▶ a **unique** optimal solution
  - ▶ several **alternative** optimal solutions
  - ▶ **no** optimal solutions (unbounded or infeasible models)

## Illustration of a (Global) Minimum, $x^*$



# Defining Local Optimality

## Neighbourhood

The **neighbourhood**  $N_\delta(\mathbf{x})$  of a point  $\mathbf{x}$  consists of all nearby points; that is, all points within a small distance  $\delta > 0$  of  $\mathbf{x}$ :

$$N_\delta(\mathbf{x}) \triangleq \{\mathbf{y} \mid \|\mathbf{y} - \mathbf{x}\|_2 < \delta\}$$

## Local Optimum

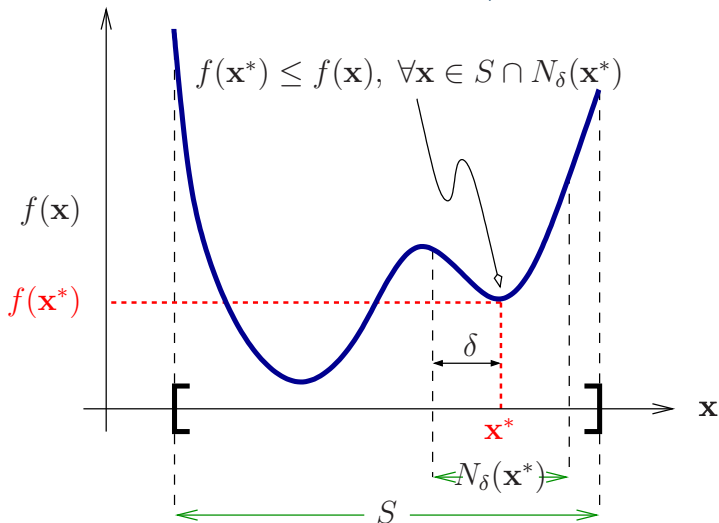
A point  $\mathbf{x}^*$  is a **local optimum** for the function  $f : \mathbb{R}^n \rightarrow \mathbb{R}$  on the set  $S$  if it is feasible ( $\mathbf{x}^* \in S$ ) and if sufficiently small neighbourhoods surrounding it contain no points that are both feasible and superior in objective value:

$$\exists \delta > 0 : \quad f(\mathbf{x}^*) \leq f(\mathbf{x}), \quad \forall \mathbf{x} \in S \cap N_\delta(\mathbf{x}^*)$$

## Remarks:

- 1 Global optima are **always** local optima
- 2 Local optima **may not be** global optima

## Illustration of a Local Minimum, $x^*$

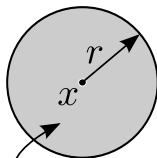
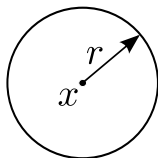


What is a **neighbourhood**  $N_\delta(x^*)$ ? Backing up for definitions ...

A neighbourhood of a point  $x \in \mathbb{R}^n$

## Neighbourhood

A **neighbourhood** of a point  $x \in \mathbb{R}^n$  is the set,  $\{y \mid \|x - y\|_2 < r\}$ , for some  $r > 0$ .



Neighbourhood of  $x \in \mathbb{R}^n$

## Formal definition

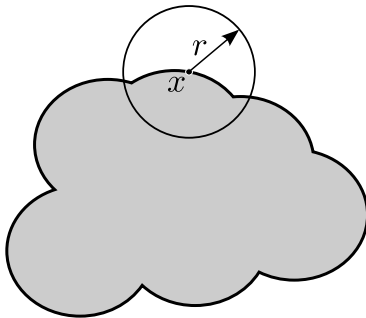
A **neighbourhood** of  $p$  in metric space  $X$  is a set  $N_r(p)$  consisting of all  $q$  such that  $d(p, q) < r$  for distance metric  $d$  and some  $r > 0$ . The number  $r$  is called the **radius** of  $N_r(p)$ .



# Boundary points

## Boundary point

A point  $x$  is called a **boundary point** if **every** neighbourhood of  $x$  contains a point in the set and a point outside the set.



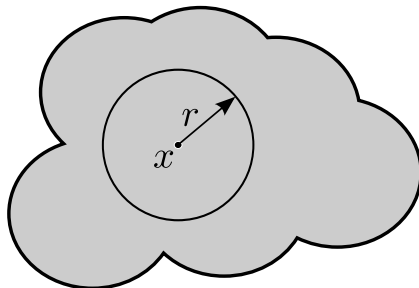
## Sanity Check

Give an example of a boundary point?

# Interior points

## Interior point

A point  $x$  is called an **interior point** if all points within **some** neighbourhood of  $x$  are contained in the set.



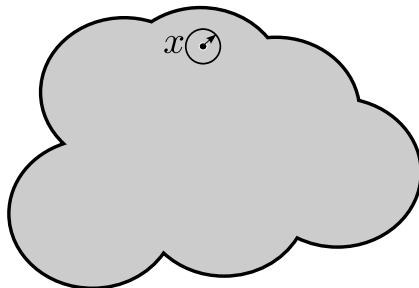
## Sanity Check

Give an example of an interior point?

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## Sanity Check

Give an example of an interior point?

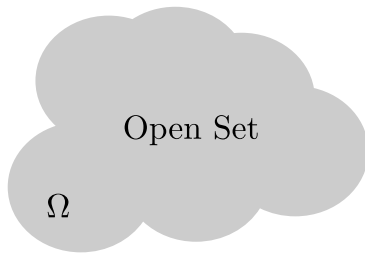
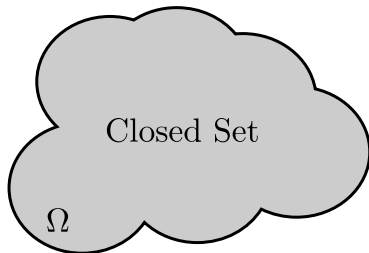
# Closed/Open Sets

## Closed set

A set is **closed** if it contains its boundary.

## Open set

A set is **open** if it contains no boundary points.



## Sanity Check

Examples of closed and open sets?

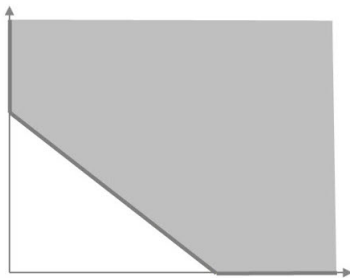
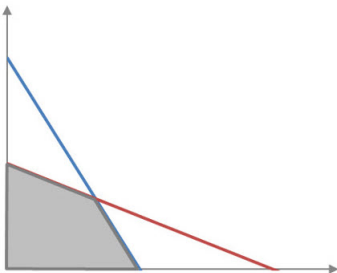
# Bounded & Compact Sets

## Bounded Set

A set that is contained in a ball of finite radius is **bounded**.

## Compact set

A **compact** set is both **closed** and **bounded**.



## Sanity Check

Which set is bounded?

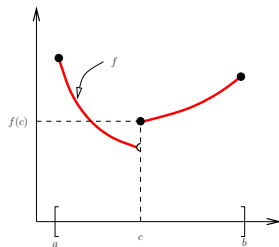
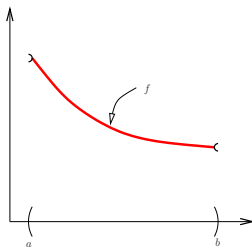
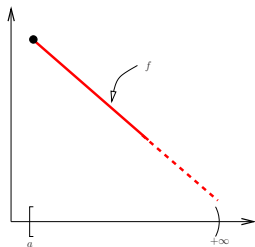
# Asserting Existence of Optima

How do we know an optimal solution exists?

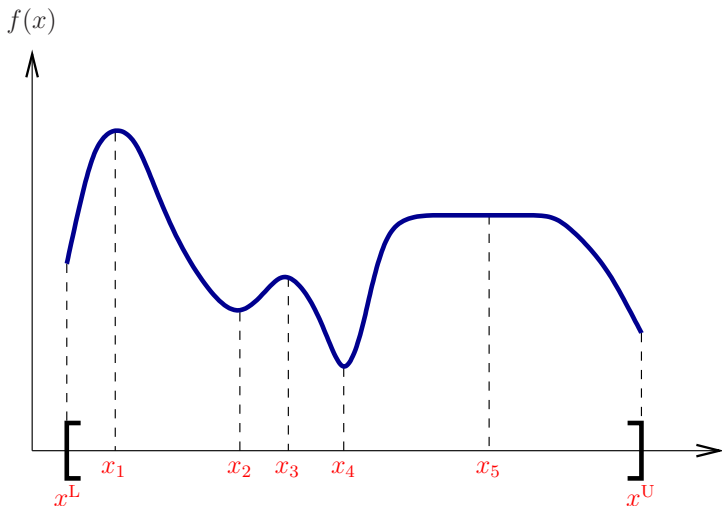
## Weierstrass Theorem

Let  $S \subset \mathbb{R}^n$  be a nonempty, compact set, and let  $f : S \rightarrow \mathbb{R}$  be continuous on  $S$ . Then, the problems  $\min / \max_{x \in S} f(x)$  attain their optimal values; that is, there exist optimal solution points for either problems.

Why **boundedness** of  $S$ ?    Why **closedness** of  $S$ ?    Why **continuity** of  $f$ ?



# Global vs Local Optima



## Sanity Check

Identify the minima & maxima types for  $f$  on  $S := [x^L, x^U]$

# Min vs Argmin

## Terminology

If  $\mathbf{x}^*$  is a global minimiser we write

$$f^* = \min_{\mathbf{x} \in \Omega} f(\mathbf{x})$$

$$\mathbf{x}^* = \arg \min_{\mathbf{x} \in \Omega} f(\mathbf{x})$$

**Examples** Assume  $x \in \mathbb{R}$ . What is  $\arg \min f(x)$ ? What is  $\min f(x)$ ?

- $f(x) = (x + 1)^2 + 3$ 
  - ▶  $\arg \min f(x) = -1$ ;  $\min f(x) = 3$
- $f(x) = (x + 1)^2 (x - 1)^2 + 3$ 
  - ▶  $\arg \min f(x) = \{-1, 1\}$ ;  $\min f(x) = 3$

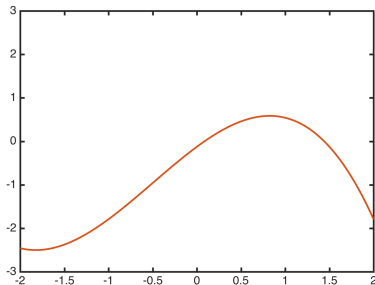
## Sanity Check

What is  $\arg \min_{\mathbf{x} \in \Omega} f(\mathbf{x})$  on the previous slide.



# Minimisation & Maximisation are Equivalent

$$\min_{x \in \Omega} f(x) = -\max_{x \in \Omega} -f(x).$$



$$\min f(x) = -\frac{1}{2}(x+0.5)^2 + 2x - \frac{1}{3}x^3$$

$$x \in [-2, 2]$$

$$\rightarrow x^* \approx -1.823$$

$$f^* \approx -2.502$$

$$-\max -f(x) = \frac{1}{2}(x+0.5)^2 - 2x + \frac{1}{3}x^3$$

$$x \in [-2, 2]$$

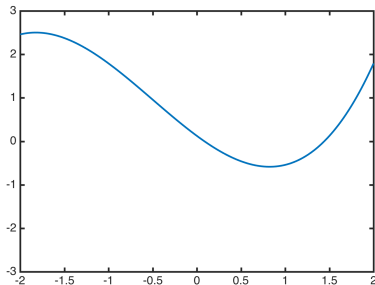
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Most optimisation problems in this class will be **minimisation** problems.

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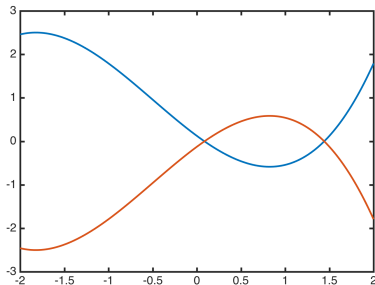
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# Minimisation & Maximisation are Equivalent (proof)

Assume that both the minimum and maximum of  $f$  are attained within set  $\Omega$  and show that:

$$\min_{\mathbf{x} \in \Omega} f(\mathbf{x}) = -\max_{\mathbf{x} \in \Omega} -f(\mathbf{x}).$$

Proof.

- The minimum of  $f$  satisfies:

$$\min_{\mathbf{x} \in \Omega} f(\mathbf{x}) = f^* \leq f(\mathbf{x}), \quad \forall \mathbf{x} \in \Omega$$

- The maximum of  $-f$  satisfies:

$$\max_{\mathbf{x} \in \Omega} -f(\mathbf{x}) = F^* \geq -f(\mathbf{x}), \quad \forall \mathbf{x} \in \Omega$$

- If  $-F^* < f^*$  then there exists an  $\bar{\mathbf{x}} \in \Omega$  such that  $f(\bar{\mathbf{x}}) < f^*$  contradicting the optimality of  $f^*$ .
- If  $-f^* > F^*$  then there exists an  $\bar{\mathbf{x}} \in \Omega$  such that  $-f(\bar{\mathbf{x}}) > F^*$  contradicting the optimality of  $F^*$ .

Therefore  $f^* = -F^*$ .



# Definitions: Types of Optimisation Models (1/4)

## Unconstrained Optimisation

$$\min_{\mathbf{x} \in \mathbb{R}^n} f(\mathbf{x})$$

## Example Soft-margin Support Vector Machine

$$\min_{\mathbf{w}, b} f(\mathbf{w}, b) = \left[ \frac{1}{n} \sum_{i=1}^n \max \left( 0, 1 - y_i (\mathbf{w}^\top \mathbf{x}_i + b) \right) \right] + \lambda \|\mathbf{w}\|_2^2$$

## Constrained Optimisation

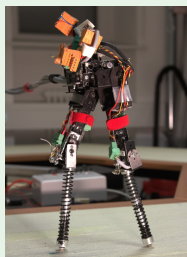
$$\min_{x \in \Omega} f(x)$$

# Definitions: Types of Optimisation Models (2/4)

## Box-Constrained Optimisation

$$\min_{\mathbf{x} \in [\mathbf{x}^L, \mathbf{x}^U]} f(\mathbf{x})$$

### Example Make a robot walk



Objective: optimise performance, e.g., as measured by velocity, within the range of reasonable gait parameters. But there is no closed-form mathematical formulation for the velocity as a function of the input parameters. To solve this problem, the authors use a type of **black-box optimisation** called **Bayesian optimisation**.

<http://wp.doc.ic.ac.uk/sml/project/bayesian-optimization/>

# Definitions: Types of Optimisation Models (3/4)

## Linear Programming (LP)

$$\begin{aligned}\min \quad & \mathbf{c}^\top \mathbf{x} \\ & \mathbf{A} \mathbf{x} \leq \mathbf{b} \\ & \mathbf{x} \in \mathbb{R}^n\end{aligned}$$

## Quadratic Programming (QP)

$$\begin{aligned}\min \quad & \frac{1}{2} \mathbf{x}^\top \mathbf{Q} \mathbf{x} + \mathbf{c}^\top \mathbf{x} \\ & \mathbf{A} \mathbf{x} \leq \mathbf{b} \\ & \mathbf{x} \in \mathbb{R}^n\end{aligned}$$

Portfolio optimisation is often formulated as a QP

Balance expected value of the portfolio's return versus financial risk measures.

# Definitions: Types of Optimisation Models (4/4)

## Nonlinear Optimisation (NLP)

$$\begin{array}{ll} \min_{\mathbf{x} \in X} f(\mathbf{x}) & \leftarrow \text{Objective function} \\ \text{s.t. } \mathbf{h}(\mathbf{x}) = \mathbf{0} & \leftarrow \text{Equality constraints} \\ \mathbf{g}(\mathbf{x}) \leq \mathbf{0} & \leftarrow \text{Inequality constraints} \\ \mathbf{x} \in [\mathbf{x}^L, \mathbf{x}^U] = X \subset \mathbb{R}^n & \leftarrow \text{Continuous variable bounds} \end{array}$$

$$f : \mathbb{R}^n \rightarrow \mathbb{R}, \mathbf{h} : \mathbb{R}^n \rightarrow \mathbb{R}^m, \mathbf{g} : \mathbb{R}^n \rightarrow \mathbb{R}^p$$

## Historical Note

I keep saying *program* when I mean *optimisation problem* because the first well-defined optimisation problems solved *training and logistics schedules*. George Dantzig referred to the solution, or proposed plan, as a *program*.



# Mixed-Integer Nonlinear Optimisation (C343 covers mixed-integer optimisation)

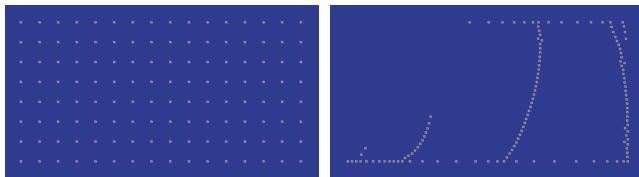
## Mixed-Integer Nonlinear Optimisation (MINLP)

$$\begin{array}{ll} \min_{\mathbf{x} \in X, \mathbf{y} \in Y} f(\mathbf{x}, \mathbf{y}) & \leftarrow \text{Objective function} \\ \text{s.t. } \mathbf{h}(\mathbf{x}, \mathbf{y}) = \mathbf{0} & \leftarrow \text{Equality constraints} \\ \mathbf{g}(\mathbf{x}, \mathbf{y}) \leq \mathbf{0} & \leftarrow \text{Inequality constraints} \\ \mathbf{x} \in [\mathbf{x}^L, \mathbf{x}^U] = X \subset \mathbb{R}^n & \leftarrow \text{Continuous variable bounds} \\ \mathbf{y} \in \{0, 1\}^{n_y} & \leftarrow \text{Binary variables} \end{array}$$

- Widely used in industry MIP
- Commercial [IBM CPLEX, 2014] MIQP
  - ▶ Passed academic codes from Imperial & CMU in 2015 [ISMP]
- Strong commercial goal MIQCQP
- Wait-&-See MINLP

# Other types of optimisation

- Dynamic optimisation is useful for optimisation-based control;
- Partial differential equation-constrained optimisation has loads of engineering applications, e.g. in the *Grantham Institute for Climate Change* at Imperial: <http://arxiv.org/pdf/1304.1768.pdf>
  - ▶ Oceanic tide stream generators could harvest renewable energy;
  - ▶ *Optimisation objective*: How to place the turbines within the site and individually tune them for maximal power output?
  - ▶ *Optimisation constraints*: Need to accurately consider tidal flow, turbine wakes, and the resulting power output.



(a) Initial turbine positions (128 turbines)    (b) Optimised turbine positions (128 turbines)

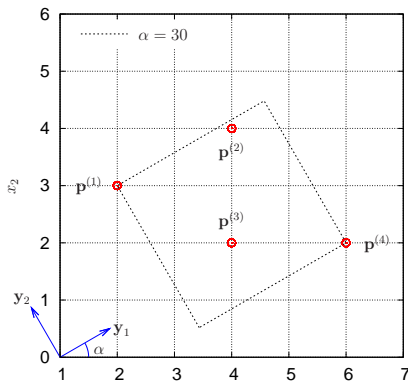
Figure: <http://arxiv.org/pdf/1304.1768.pdf>

# A First Example Problem – Statement

**Workshop.** Find a rectangle with minimum area enclosing the set of points  $\{\mathbf{p}^{(1)} := (2, 3), \mathbf{p}^{(2)} := (4, 4), \mathbf{p}^{(3)} := (4, 2), \mathbf{p}^{(4)} := (6, 2)\}$

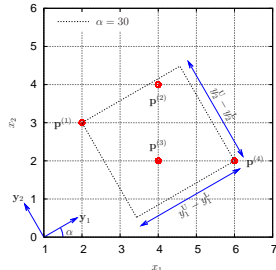
➡ Formulate this problem as an NLP

Hint: Consider the coordinate system  $(y_1, y_2)$  aligned with the axis of the rectangle, as obtained by rotating  $(x_1, x_2)$  by the angle  $\alpha$



# A First Example Problem – NLP Formulation

## ● Formulation using Projections



$$\min_{\mathbf{y}_1, \mathbf{y}_2, y_1^U, y_1^L, y_2^U, y_2^L} (y_1^U - y_1^L)(y_2^U - y_2^L)$$

$$\text{s.t. } y_1^U := \max\{\mathbf{y}_1^\top \mathbf{p}^{(k)}, k = 1, \dots, 4\}$$

$$y_1^L := \min\{\mathbf{y}_1^\top \mathbf{p}^{(k)}, k = 1, \dots, 4\}$$

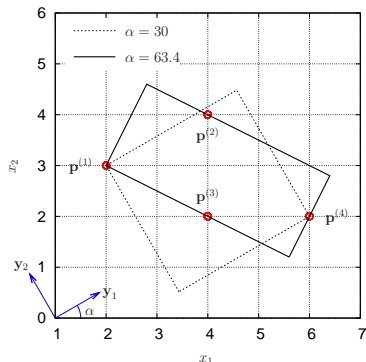
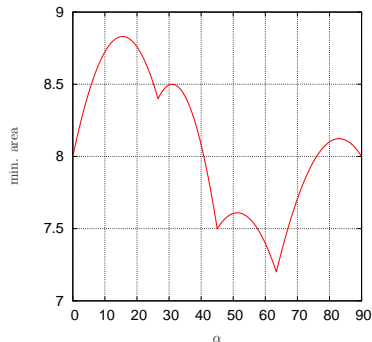
$$y_2^U := \max\{\mathbf{y}_2^\top \mathbf{p}^{(k)}, k = 1, \dots, 4\}$$

$$y_2^L := \min\{\mathbf{y}_2^\top \mathbf{p}^{(k)}, k = 1, \dots, 4\}$$

$$\mathbf{y}_1^\top \mathbf{y}_2 = 0$$

$$\|\mathbf{y}_1\|_2 = 1; \|\mathbf{y}_2\|_2 = 1$$

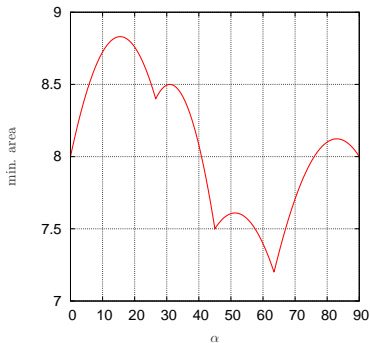
# A First Example Problem – Results



## Multi-Start Gradient Descent: 20,000 Random Initial points

Freq	Area	$\alpha$	$y_1^U$	$y_1^L$	$y_2^U$	$y_2^L$
32.6%	8.00	0.0	6.000	2.000	4.000	2.000
<b>30.5%</b>	<b>7.20</b>	<b>63.4</b>	<b>5.367</b>	<b>3.578</b>	<b>-0.447</b>	<b>-4.472</b>
15.2%	8.00	90.0	4.000	2.000	-2.000	-6.000
12.1%	7.50	45.0	5.657	3.536	0.707	-2.828
9.7%	8.40	26.6	6.261	3.130	1.789	-0.894

# Motivation for next time



What is going on? We are getting different answers depending on where we initialise the algorithm. When is a local optimum also the global solution?

Key concept: **convexity**. If a problem is convex, then a local optimum solution must be a global solution. (More on this Wednesday!!)