

UM-SJTU Joint Institute Introduction to Algorithms VE477 Lab3 Report

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1. a. Sort and count

Code in OCaml is shown below:

```
let rec mergecount = function (*merge and count*)
      | list, [] -> list, 0
 2
 3
     | [], list -> list, 0
     | x1::xs1, x2::xs2 ->
4
 5
       if x1 \ll x2 then
         match (mergecount (xs1, x2::xs2)) with
 6
 7
           | (list, count) -> x1 :: list, count
8
         else
           match (mergecount (x1::xs1, xs2)) with
9
10
            | (list, count) -> x2 :: list, count + (List.length xs1) + 1
11
    let rec split_at n tmp list = (*split the list at specified position*)
12
     if n = 0 then List.rev tmp, list
13
14
     else match list with
15
     | [] -> List.rev tmp, []
16
     | head::tail -> split_at (n-1) (head::tmp) tail
17
    ;;
18
19
    let split list = split_at ((List.length list + 1) / 2) [] list;;
20
    (*split the list at middle*)
21
22
    let get_ele1 (a, _) = a;; (*get the first element of a two-element
    tuple*)
23
    let get_ele2 (_, a) = a;; (*get the second element*)
24
25
    let rec merge_sort = function (*sort and count*)
26
        | [] -> [], 0
        | [_] as list -> list, 0
27
        | list ->
28
29
           let 11, 12 = (split list) in
             let (s1, s2) = merge_sort l1, merge_sort l2 in
30
31
                let (result, count) = mergecount (get_ele1 s1, get_ele1 s2) in
              result, get_ele2 (s1) + get_ele2 (s2) + count
32
33
    ;;
34
    let rec print_list list = match list with (*print function*)
35
     | [] -> print_string ""
36
     | head::[] -> print_int head
37
     | head::tail -> print_int head ; print_string ", " ; print_list tail
38
39
    ;;
40
    let input = read_line();;
41
    let 11 = List.map int_of_string (Str.split (Str.regexp " ") input);;
42
43
44
    let (list, count) = merge_sort ll;;
45
    Printf.printf "%d\n" count;;
    print_string "[";;
46
47
    print_list list;;
    print_string "]\n";;
48
49
```

OCaml features a recursion method, so using OCaml to write merge sort and count can save some time and the code will be more concise.

1. b. Gale-Shapley

Code in Ocaml is shown below:

```
type man =
 2
     { m_index: int;
 3
        mutable free: bool;
 4
        women_rank: int list;
 5
       has_proposed: (int, bool) Hashtbl.t
 6
        (*Use hashtable to store proposed women*)
 7
     }
 8
    ;;
 9
    type woman =
10
     { w_index: int;
        men_rank: int list;
11
12
        mutable engaged: int option (*can be either none or an index of man*)
13
     }
14
    ;;
15
16 (*The function to read inputs and convert to a list of tuples which contains
    an index and a rank list*)
   let rec read_person i n ll =
17
     if i = n then List.rev 11
19
      else
20
       let s = read_line() in
         let l = List.map int_of_string (Str.split (Str.regexp " ") s) in
21
22
           match 1 with
            | [] -> List.rev 11
23
24
            | head::tail -> read_person (i+1) n ((i,1)::11)
25
    ;;
26
27
    (*returns true if m1 is of higher rank than m2 for w*)
28
    let prefer w m1 m2 =
29
     let rec comp = function
30
      [] -> invalid_arg "cannot find such index"
31
     | x::_ when x = m1 \rightarrow true
32
      | x:: \_ when x = m2 \rightarrow false
33
      | _::xs -> comp xs
34
      in
35
      comp w.men_rank
36
    ;;
37
38
    (*initialize the struct on men and women*)
39
    let build_structs men women n =
      List.map (fun (index, rank) -> {m_index = index; women_rank = rank; free =
    true; has_proposed = Hashtbl.create n}) men,
41
      List.map (fun (index, rank) -> {w_index = index; men_rank = rank; engaged
    = None}) women
42
    ;;
43
    let gale_helper m_list w_list n =
```

```
45 (*Create two hashtables for men and women, key is the index, value is the
    corresponding struct*)
      let men = Hashtbl.create n in
46
47
      List.iter (fun m -> Hashtbl.add men m.m_index m) m_list;
48
      let women = Hashtbl.create n in
49
      List.iter (fun w -> Hashtbl.add women w.w_index w) w_list;
50
      try
51
        while true do
         let m = List.find (fun m -> m.free) m_list in (*Find the first free
52
    man from the list of men*)
53
          let w_index = List.find (fun w -> not (Hashtbl.mem m.has_proposed w))
    m.women_rank in (*Find the woman of highest rank who has not been
    proposed by m*)
54
          Hashtbl.add m.has_proposed w_index true; (*Add the woman to the
    proposed hashtable, key is index, value is en empty unit*)
          let w = Hashtbl.find women w_index in (*The hashtable can shorter
55
    the finding process to O(1)^*)
          match w.engaged with
56
57
          | None -> w.engaged <- Some m.m_index; m.free <- false (*If the
    woman is free, let m and w be a pair*)
          | Some m'_index -> (* some pair (m', w) already exists *)
58
              if prefer w m.m_index m'_index then (*If woman "w" prefer "m"
59
    to "m'"*)
60
              begin
61
                w.engaged <- Some m.m_index;</pre>
                let m' = Hashtbl.find men m'_index in (*Quick find in
62
    hashtable*)
                m'.free <- true;</pre>
63
64
                m.free <- false</pre>
65
              end
66
        done;
67
      with Not_found -> ()
68
    ;;
69
70
    let gale men women n =
71
     let m_list, w_list = build_structs men women n in
72
      gale_helper m_list w_list n;
73
      let some = function Some v \rightarrow v | None \rightarrow -1 in
74
      List.stable_sort compare (List.map (fun w -> some w.engaged, w.w_index)
    w_list)
75
    ;;
76
77
    let get_ele1(a, _) = a;;
78
    let get_ele2 (\_, a) = a;;
79
    let rec print result =
80
     match result with
      | [] -> print_string ""
81
82
     | head::[] -> print_string "["; print_int (get_ele1 head); print_string ",
    "; print_int (get_ele2 head); print_string "]"
      | head::tail -> print_string "["; print_int (get_ele1 head); print_string
83
    ", "; print_int (get_ele2 head); print_string "], "; print tail
84
    ;;
85
86
    let () =
87
      let n = read_int() in
88
      let men= read_person 0 n [] in
      let tmp = read_line() in
89
90
      let women = read_person 0 n [] in
```

```
91  let result = gale men women n in
92  print_string tmp;
93  print_string "[";
94  print result;
95  print_string "]\n"
96 ;;
```

2. Knapsack Problem

a. Two strategies

(i). Smaller first

OCaml code:

```
let rec small_first list target result =
2
    match list with
3
    | [] -> []
4
     | head::tail ->
5
       if target = head then List.rev (head::result)
6
       else if target > head then small_first tail (target - head)
   (head::result)
7
       else []
8
   ;;
```

Counter example:

input:

```
1 | 4
2 | 2 3 4
```

Output:

```
1 | []
```

But it should output [4].

(ii). Larger first

OCaml code:

```
let rec large_first list target result =
2
   match list with
3
    | [] -> []
    | head::tail ->
4
5
       if target = head then List.rev (head::result)
6
       else if target > head then large_first tail (target - head)
   (head::result)
7
       else large_first tail target result
8
   ;;
```

Counter example:

Input:

```
1 | 7
2 | 2 3 4 5 6
```

Output:

```
1 | []
```

Should output: [2, 5] or [3, 4].

b. Proper solution

From wiki, the proper way to solve the knapsack (subset sum problem) is using dynamic programming to handle a two-dimension array. Since in OCaml lists are immutable, it is hard to perform dynamic programming. Therefore, this part is written in Python.

```
def knapsack(arr, target):
                                   # dynamic programming
 1
 2
        arr_len = len(arr)
        subset = [[False for j in range(target + 1)] for i in range(arr_len +
    1)]
        for i in range(arr_len + 1):
 4
 5
            subset[i][0] = True
 6
        for i in range(1, arr_len + 1):
 7
 8
             for j in range(1, target + 1):
 9
                 if subset[i-1][j]:
                     subset[i][j] = True
10
11
                 else:
12
                     if arr[i-1] > j:
13
                         subset[i][j] = False
14
                     else:
                         subset[i][j] = subset[i-1][j-arr[i-1]]
15
16
        return subset
17
18
19
    # To print the result, recursion can be used to judge each element.
    def print_subset(subset, arr, target, n, result):
21
        if target == 0:
             print(sorted(result))
22
23
             return
24
        if n == 0 and target != 0 and subset[1][target]:
25
             result.append(arr[0])
26
            print(sorted(result))
27
             return
28
        if subset[n][target]:
29
            print_subset(subset, arr, target, n-1, result.copy())
30
31
        if target >= arr[n] and subset[n][target - arr[n]]:
32
             result.append(arr[n])
33
            print_subset(subset, arr, target - arr[n], n-1, result)
34
35
    if __name__ == '__main__':
36
37
        target = int(input())
38
        arr = []
```

```
for i in input().split():
    arr.append(int(i))

subset = knapsack(arr, target)
print_subset(subset, arr, target, len(arr) - 1, [])
```

3. Running time of three Knapsack algorithms

a. Smaller first

By choosing different sizes of arrays, the running time is shown below:

```
1 Size: 10
              end
 2 Execution time: 0.000013
 3 Size: 20 end
 4 Execution time: 0.000017
 5 Size: 30 end
 6 Execution time: 0.000030
 7 Size: 40 end
8 Execution time: 0.000022
9 Size: 50 end
10 Execution time: 0.000011
11 Size: 100 end
12 Execution time: 0.000084
13 Size: 1000 end
14 Execution time: 0.000055
15 | Size: 10000 end
16 Execution time: 0.000452
17 | Size: 100000 end
18 Execution time: 0.003987
```

This algorithm is supposed to run in $\mathcal{O}(n)$ time, but due to extra computation like if statement and push the current element into the result list, according to the time complexity table, it is likely to have complexity $\mathcal{O}(n \log n)$

b. Larger first

Same as the above one, the output is shown below:

```
1 Size: 10 end
 2 | Execution time: 0.000006
 3 Size: 20 end
4 Execution time: 0.000011
 5 Size: 30 end
 6 Execution time: 0.000009
7 Size: 40 end
8 Execution time: 0.000009
9 Size: 50 end
10
   Execution time: 0.000059
11 | Size: 100 | end
12 Execution time: 0.000016
13 Size: 1000 end
14 | Execution time: 0.000180
15
   Size: 10000 end
16 | Execution time: 0.000509
```

```
17 | Size: 100000 end
18 | Execution time: 0.003707
```

We can see that the running time of the larger-first algorithm differs a little with the smaller-first algorithm, but it in general costs less time a little.

c. Dynamic programming algorithm

I choose the target number to be a random number in [100, 130], and the numbers in the test array are all range in [0, 25].

The test program's code is:

```
from knapsack import knapsack, print_subset
 2
    import random
 3
    import time
    import matplotlib.pyplot as plt
 4
 5
    import numpy
 6
 7
 8
    def create_list(s):
 9
        result = []
10
        for i in range(s):
11
            result.append(random.randint(0, 25))
12
        return result
13
14
15
    random.seed()
    sizes = [10, 20, 30, 40, 50, 100, 200, 300, 400, 500, 1000, 2000, 3000,
16
    4000, 5000,
             10000, 20000, 30000, 40000, 50000, 60000, 70000, 80000, 90000,
17
    100000]
    y = []
18
19
    for size in sizes:
20
        arr = create_list(size)
21
        target = random.randint(100, 130)
        start = time.time()
22
23
        subset = knapsack(arr, target)
        end = time.time()
24
25
        # print_subset(subset, arr, target, size - 1, [])
26
        print("Size:", size, "Execution time:", end-start)
27
        y.append(end-start)
28
    plt.figure()
29
30
    scatter = plt.scatter(sizes, y, color='red', s=50, alpha=0.5)
    plot ,= plt.plot(sizes, y, color="blue", linewidth=2.0, label="running"
31
32
    plt.legend(handles=[scatter, plot], labels=["scatter", "line chart"])
33
    plt.title("Running time of Knapsack --- Dynamic programming")
    plt.xlabel("Size of array")
34
    plt.ylabel("Time cost (s)")
35
36
    plt.show()
37
38
    plt.figure()
39 | p1 = numpy.polyfit(x=sizes, y=y, deg=1)
40 | yvals = numpy.polyval(p1, sizes)
    plt.scatter(sizes, y, color="red", alpha=0.5, label="scatter")
```

```
plt.plot(sizes, yvals, color="blue", linewidth=2.0, label="linear fit")
plt.legend()

plt.title("Linear fit of output")

plt.xlabel("Size of array")

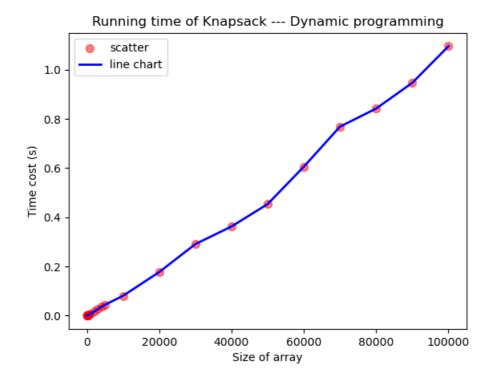
plt.ylabel("Time cost (s)")

plt.show()
```

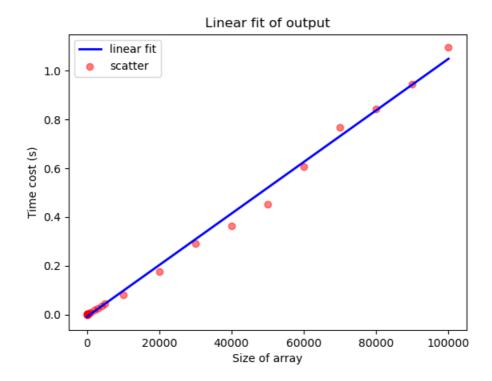
And the time output is:

```
Size: 10 Execution time: 0.0001468658447265625
    Size: 20 Execution time: 0.00020837783813476562
    Size: 30 Execution time: 0.0002999305725097656
   Size: 40 Execution time: 0.00038504600524902344
    Size: 50 Execution time: 0.00038623809814453125
 5
    Size: 100 Execution time: 0.0007736682891845703
    Size: 200 Execution time: 0.0013849735260009766
    Size: 300 Execution time: 0.0021965503692626953
8
9
    Size: 400 Execution time: 0.002763509750366211
    Size: 500 Execution time: 0.0036928653717041016
10
    Size: 1000 Execution time: 0.0071218013763427734
11
    Size: 2000 Execution time: 0.016954660415649414
    Size: 3000 Execution time: 0.02609395980834961
13
    Size: 4000 Execution time: 0.03602409362792969
14
15
    Size: 5000 Execution time: 0.04427981376647949
    Size: 10000 Execution time: 0.08069109916687012
16
    Size: 20000 Execution time: 0.1772470474243164
17
    Size: 30000 Execution time: 0.29102230072021484
18
19
    Size: 40000 Execution time: 0.3626530170440674
    Size: 50000 Execution time: 0.4538257122039795
20
    Size: 60000 Execution time: 0.6058144569396973
21
22
    Size: 70000 Execution time: 0.7685534954071045
    Size: 80000 Execution time: 0.8421218395233154
23
24
    Size: 90000 Execution time: 0.9468364715576172
    Size: 100000 Execution time: 1.0956668853759766
25
```

Using matplotlib.pyplot, we can show the scatter and line chart of the data:



Then it comes the linear fit of the obtained points:



Theoretically, the dynamic programming algorithm should have the time complexity of $\mathcal{O}(TN)$, where T is the **target number** and \mathbf{N} is the size of the input array. From the plot generated, we can obviously find out that the time increase is linear, which conforms with the theory. However, due to the target number's size, the time cost will differ.

Since the smaller-first and larger-first algorithm just traverse the test array once, they will be much faster than the dynamic programming algorithm. However, as they cannot give the right answer, only serves as a greedy algorithm, the time cost comparison will have no significance.

4. Topological sort

Topological sort of a **directed graph** is a linear ordering of its vertices such that for every directed edge uv from vertex u to vertex v, u comes before v in the ordering.

OCaml code is shown below:

```
let create_hash hash n =
 2
      for i = 0 to n - 1 do
 3
        Hashtbl.add hash i []
 4
      done
 5
    ;;
 6
 7
    let init hash e =
                         (*Initialize hashtable according to input*)
 8
     for i = 0 to e - 1 do
 9
        let (node_from, node_to) = Scanf.scanf " %d %d" (fun a b -> (a, b)) in
10
        let x = Hashtbl.find hash node_from in
        Hashtbl.remove hash node_from;
11
        if List.mem node_to x then Hashtbl.add hash node_from x
12
13
        else Hashtbl.add hash node_from (node_to :: x)
14
      done
15
    ;;
16
17
    let dfs graph visited start = (*Use dfs to find the path*)
      let rec explore path visited node =
18
        if List.mem node visited then visited
19
20
        else
21
          let new_path = node :: path in
22
          let edges = List.assoc node graph in
23
          let visited = List.fold_left (explore new_path) visited edges in
24
          node :: visited
25
     in explore [] visited start
26
    ;;
27
28
    let toposort graph =
29
     List.fold_left (fun visited (node,_) -> dfs graph visited node) [] graph
30
    ;;
31
    let rec print 1 =
32
33
     match 1 with
      | [] -> print_string ""
34
35
      | head :: [] -> print_int head
      | head :: tail -> print_int head; print_string " "; print tail
36
37
38
   let n = read_int();;
    let h = Hashtbl.create n;;
39
40
    let e = read_int();;
41
42
    create_hash h n;;
43
    init h e;;
44 let l = Hashtbl.fold (fun x y z -> (x, y) :: z) h [];;
45
   (*Transform hashtable into list*)
46 | let result = toposort 1;;
    print result;;
```