

# Optimization for Classical Machine Learning Problems on the GPU

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## Abstract

Constrained optimization problems arise frequently in classical machine learning. There exist frameworks addressing constrained optimization, for instance, CVXPY. However, in contrast to deep learning frameworks, GPU support is limited. Here, we provide a framework for solving constrained optimization problems on the GPU. The framework allows the user to specify constrained optimization problems in an easy-to-read modeling language. A solver is then automatically generated from this specification. When run on the GPU, the solver outperforms state-of-the-art approaches like CVXPY combined with a GPU-accelerated solver such as cuOSQP or SCS by several orders of magnitude.

## Introduction

Training classical machine learning models typically means solving an optimization problem. Hence, the design and implementation of solvers for training these models has been and still is an active research topic. While the use of GPUs is standard in training deep learning models, most solvers for classical machine learning problems still target CPUs. Easy-to-use deep learning frameworks like TensorFlow or PyTorch can be used to solve unconstrained problems on the GPU. However, many classical problems entail constrained optimization problems. So far, deep learning frameworks do not support constrained problems, not even problems with simple box constraints, that is, bounds on the variables. Optimization frameworks for classical machine learning like CVXPY (Agrawal et al. 2018; Diamond and Boyd 2016) can handle constraints but typically address CPUs. Here, we present a framework for constrained optimization that specifically targets GPUs.

Adding GPU support to an optimization framework for classical machine learning problems is not straightforward. An efficient algorithmic framework could use the limited-memory quasi-Newton method L-BFGS-B (Byrd et al. 1995) that allows to solve large-scale optimization problems with box constraints. More general constraints can then be addressed by the augmented Lagrangian approach (Hestenes 1969; Powell 1969). However, porting the L-BFGS-B method to the GPU does not provide any efficiency gain for the following reason: In each iteration the

method involves an inherently sequential Cauchy point computation that determines the variables that will be modified in the current iteration. The Cauchy point is computed by minimizing a quadratic approximation over the gradient projection path, resulting in a large number of sequential scalar computations, which means that all but one core on a multicore processor will be idle. This problem is aggravated on modern GPUs that feature a few orders of magnitude more cores than a standard CPU, e.g., 2304 instead of 18.

Let us substantiate this issue by the following example (see the appendix for details). On a non-negative least squares problem, the L-BFGS-B algorithm needs 4.9 seconds in total until convergence, where 0.6 seconds are spent in the Cauchy point subroutine. When run on a modern GPU, the same code needs 5.2 seconds in total while 4.6 seconds are spent in the Cauchy point subroutine. It can be seen that while all other parts of the L-BFGS-B algorithm can be parallelized nicely on a GPU, the inherently sequential Cauchy point computation does not and instead, dominates the computation time on the GPU; as a result, the L-BFGS-B method is not faster on a GPU than on a CPU, rendering the benefits of the GPU moot.

Here, we present a modified L-BFGS-B method that runs efficiently on the GPU, because it avoids the sequential Cauchy point computation. For instance, on the same problem as above, it needs 0.8 seconds in total on the GPU. We integrate an implementation of this method into a generic framework for constrained optimization. Experiments on several classical machine learning problems show that our implementation outperforms state-of-the-art approaches like the combination of CVXPY with GPU-accelerated solvers such as cuOSQP or SCS by several orders of magnitude.

**Contributions.** The contributions of this paper can be summarized as follows:

1. We design a provably convergent L-BFGS-B algorithm that can handle box constraints on the variables and runs efficiently on the GPU.
2. We combine our algorithm with an Augmented Lagrangian approach for solving constrained optimization problems.
3. We integrate our approach with a modeling language into a framework for constrained optimization. It outperforms comparable state-of-the-art approaches by several orders of magnitude.

## State of the art

Optimization for classical machine learning is either addressed by problem specific solvers, often wrapped in a library or toolbox, or by frameworks that combine a modeling language with a generic solver. The popular toolbox scikit-learn (Pedregosa et al. 2011) does not provide GPU support yet. Spark (Zaharia et al. 2016) has recently added GPU support and hence, enabling GPU acceleration for some algorithms from its machine learning toolbox MLlib (Meng et al. 2016). The modeling language CVXPY (Agrawal et al. 2018; Diamond and Boyd 2016) can be paired with solvers like cuOSQP (Schubiger, Banjac, and Lygeros 2020) or SCS (O’Donoghue et al. 2016, 2019) that provide GPU support.

In contrast to deep learning, classical machine learning can involve constraints, which poses extra algorithmic challenges. There are several algorithmic approaches for dealing with constraints that could be adapted for the GPU. As mentioned before, we have decided to adapt the L-BFGS-B method (Byrd et al. 1995). Its original Fortran-code (Zhu et al. 1997) is still the predominant solver in many toolboxes like scikit-learn or scipy for solving box-constrained optimization problems. In the following, we briefly describe algorithmic alternatives to the L-BFGS-B method and argue why we decided against using them for our purposes.

Proximal methods, including alternating direction method of multipliers (ADMM) have been used for solving constrained and box-constrained optimization problems. The literature on proximal methods is vast, see for instance, (Boyd et al. 2011; Parikh and Boyd 2014) for an overview. Prominent examples that relate to our work are OSQP (Stellato et al. 2020) and SCS (O’Donoghue et al. 2016, 2019). Unfortunately, both methods require a large number of iterations. One approach to mitigate this problem is to keep the penalty parameter  $\rho$ , which ties the constraints to the objective function fixed for each iteration. Then, the Cholesky-decomposition of the KKT system needs to be computed only once and can be reused in subsequent iterations, leading to a slow first iteration but very fast consecutive iterations. The OSQP solver has been shown to be the fastest general purpose solver for quadratic optimization problems (Stellato et al. 2020), beating commercial solvers like Gurobi as well as Mosek. Gurobi as well as Mosek do not provide GPU support and so far, there is no intention to do so in the near future (Glockner 2021). OSQP does not support GPU acceleration either. However, it has been ported to the GPU as cuOSQP (Schubiger, Banjac, and Lygeros 2020), where it was shown to outperform its CPU version by an order of magnitude.

The simplest way to solve box-constrained optimization problems is probably the projected gradient descent method (Nesterov 2004). However, it is as slow as gradient descent and not applicable to practical problems. Hence, there have been a number of methods proposed that try to combine the projection approach with better search directions. For instance, (van den Berg 2020) applies L-BFGS updates only to the currently active face. If faces switch between iterations, which happens in almost all iterations, it falls back to standard spectral gradient descent. A similar

approach is the non-monotone spectral projected gradient descent approach as described in (Schmidt, Kim, and Sra 2011). It also performs backtracking along the projection arc and cannot be parallelized efficiently. Another variant solves a sub-problem in each iteration that is very expensive and hence, only useful when the cost of computing the function value and gradient of the original problem is very expensive (Schmidt et al. 2009), which is typically not the case for standard machine learning problems. Another approach for solving box-constrained optimization problems has been described in (Kim, Sra, and Dhillon 2010). However, it is restricted to strongly convex problems. For small convex problems, a projected Newton method has been described in (Bertsekas 1982).

Nesterov acceleration (Nesterov 1983) has also been applied to proximal methods (Beck and Teboulle 2009; Li and Lin 2015). However, similar to Nesterov’s optimal gradient descent algorithm (Nesterov 1983), one needs several Lipschitz constants of the objective function, which are usually not known. Quasi-Newton methods do not need to know such parameters and have been shown to perform equally well or even better. Convergence rates have been obtained for quasi-Newton methods on special instances, e.g., quadratic functions with unbounded domain. Recently, improved convergence rates have been proved in (Rodomanov and Nesterov 2021) for the unbounded case.

## Algorithm

Here, we present our extension of the L-BFGS algorithm for minimizing a differentiable function  $f: \mathbb{R}^n \rightarrow \mathbb{R}$ , which can additionally handle box constraints on the variables, i.e.,  $l \leq x \leq u$ ,  $l, u \in \mathbb{R}^n \cup \{\pm\infty\}$ , and which runs efficiently on the GPU.

**Notation.** A sequence of scalars or vectors is denoted by upper indices, e.g.,  $x^1, \dots, x^k$ . The projection of a vector  $x \in \mathbb{R}^n$  to the coordinate set  $S \subseteq \{1, \dots, n\} =: [n]$  is denoted as  $x[S]$ . If  $S$  is a singleton  $\{i\}$ , then  $x[S]$  is just the  $i$ -th coordinate  $x_i$  of  $x$ . Similarly, the projection of a square matrix  $B$  onto the rows and columns in an index set  $S$  is denoted by  $B[S, S]$ . Finally, the Euclidean norm of  $x$  is denoted by  $\|x\|$ , and the corresponding scalar product between vectors  $u$  and  $v$  is denoted as  $\langle u, v \rangle$ .

Like the original L-BFGS-B algorithm, our extension runs in iterations until a convergence criterion is met. Likewise, it distinguishes between fixed and free variables in each iteration, i.e., variables that are fixed at their boundaries and variables that are optimized in the current iteration. In contrast to the original L-BFGS-B algorithm, we can avoid the inherently sequential Cauchy point computation by determining the fixed and free variables directly. Given  $\varepsilon > 0$ , we compute in iteration  $k$  the index set (working set)

$$S^k = [n] \setminus \left( \{i \mid (x^k)_i \leq l_i + \varepsilon \text{ and } \nabla f(x^k)_i \geq 0\} \cup \{i \mid (x^k)_i \geq u_i - \varepsilon \text{ and } \nabla f(x^k)_i \leq 0\} \right) \quad (1)$$

of free variables. Here,  $\{i \mid (x^k)_i \leq l_i + \varepsilon \text{ and } \nabla f(x^k)_i \geq 0\}$  holds the indices of optimization variables that, at iteration  $k$ , are within an  $\varepsilon$ -interval of the lower bound. Analogously,  $\{i \mid (x^k)_i \geq u_i - \varepsilon \text{ and } \nabla f(x^k)_i \leq 0\}$  holds all

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**Algorithm 1: GPU-efficient L-BFGS-B Method**


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**Input:** initial iterate  $x^0$  with  $l \leq x^0 \leq u$

- 1: **Initialization:** set  $k \leftarrow 0$
  - 2: **repeat**
  - 3:   compute  $\nabla f(x^k)$  and set  $S^k$  of free variables (Eq. 1)
  - 4:   solve  $-\nabla f(x^k)[S^k] = B^k[S^k, S^k]d^k[S^k]$  using L-BFGS modified two-loop algorithm, see appendix
  - 5:   set  $d^k[\bar{S}^k] = 0$
  - 6:    $p^k = \text{projectDirection}(x^k, \nabla f(x^k), d^k)$
  - 7:    $x^{k+1} = x^k + \alpha^k p^k$  using line search with appropriate upper bound on  $\alpha^k$
  - 8:    $y^k = \nabla f(x^{k+1}) - \nabla f(x^k)$  and  $s^k = x^{k+1} - x^k$
  - 9:   store new curvature pair  $(y^k, s^k)$
  - 10:   set  $k \leftarrow k + 1$
  - 11: **until** converged
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**Algorithm 2: projectDirection( $x^k, \nabla f(x^k), d^k$ )**


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- 1: compute  $z^k = x^k + d^k$  and project  $z^k$  onto feasible region
  - 2: compute  $p^k = z^k - x^k$
  - 3: **if**  $\langle p^k, \nabla f(x^k) \rangle \leq -\varepsilon \|p^k\|^2$  and  $\|p^k\|^2 \geq \varepsilon$  **then**
  - 4:   **return**  $p^k$
  - 5: **else**
  - 6:   set  $p^k = d^k$
  - 7:   set  $(p^k)_i = 0 \ \forall i$  with  $(d^k)_i < 0$  and  $(x^k)_i \leq l_i + \varepsilon$
  - 8:   set  $(p^k)_i = 0 \ \forall i$  with  $(d^k)_i > 0$  and  $(x^k)_i \geq u_i - \varepsilon$
  - 9:   **return**  $p^k$
  - 10: **end if**
- 

indices, where the optimization variable is close to the upper bound. The complement of  $S^k$ , i.e., the index set of all non-free (fixed) variables is denoted by  $\bar{S}^k$ . Algorithm 1 computes the quasi-Newton search direction  $d^k[S^k]$  only on the free variables (Line 4). It then projects this direction onto the feasible set using Algorithm 2. If it is a feasible descent direction, it takes a step into this direction. Otherwise, it takes a step into the original quasi-Newton direction until it hits the boundary of the feasible region. While the original L-BFGS-B algorithm uses a line search with quadratic and cubic interpolation to satisfy the strong Wolfe conditions, we observed that this does not provide any benefit for optimization problems from machine learning. Hence, our implementation uses a simple backtracking line search to satisfy the Armijo condition (Nocedal and Wright 1999). Even when the function is convex and satisfies the curvature condition for all variables, it does not necessarily satisfy the curvature condition for the set of free variables with indices in  $S^k$ . Hence, satisfying the strong Wolfe conditions is not necessary and instead the curvature condition is checked for the current set of free variables in the modified two-loop Algorithm 3 (see the appendix). The following theorem asserts that our algorithm converges to a stationary point.

**Theorem 1.** *Let  $f$  be a differentiable function with an  $L$ -Lipschitz continuous gradient. If  $f$  is bounded from below, then Algorithm 1 converges to a feasible stationary point.*

*Proof.* We have for any differentiable function with  $L$ -Lipschitz continuous gradient that

$$f(x + \alpha p) \leq f(x) + \langle \nabla f(x), \alpha p \rangle + \frac{L}{2} \|\alpha p\|^2$$

see, e.g., (Nesterov 2004). For computing the search direction  $p^k$ , we distinguish two cases in Algorithm 2. In the first case, we have

$$\langle p^k, \nabla f(x^k) \rangle \leq -\varepsilon \|p^k\|^2 \text{ and } \|p^k\|^2 \geq \varepsilon$$

(Algorithm 2, Line 3). Hence, we have

$$\begin{aligned} f(x^{k+1}) &= f(x^k + \alpha^k p^k) \\ &\leq f(x^k) + \alpha^k \langle \nabla f(x^k), p^k \rangle + \frac{L}{2} \|\alpha^k p^k\|^2 \\ &\leq f(x^k) - \alpha^k \varepsilon \|p^k\|^2 + (\alpha^k)^2 \frac{L}{2} \|p^k\|^2 \end{aligned}$$

If we set  $\alpha^k = \frac{\varepsilon}{L}$ , we get

$$f(x^{k+1}) \leq f(x^k) - \frac{\varepsilon^2}{2L} \|p^k\|^2 \leq f(x^k) - \frac{\varepsilon^3}{2L}.$$

Hence, the objective function reduces at least by a positive constant that is bounded away from 0 in each iteration.

In the second case, we have the following. For a function  $f$  with  $L$ -Lipschitz continuous gradient and curvature pairs that satisfy the curvature condition  $\langle y^k, s^k \rangle \geq \varepsilon \|y^k\|^2$ , the smallest and largest eigenvalue of the Hessian approximation  $B^k$  can, in general, be lower and upper bounded by two constants  $c$  and  $C$ , see e.g., (Mokhtari and Ribeiro 2015). Since we require this curvature condition to hold only for the index set  $S^k$  that is active in iteration  $k$  (see Algorithm 3 in the appendix), we can lower and upper bound the eigenvalues of the submatrix  $B^k[S^k, S^k]$  of the Hessian approximation by  $0 < c$  and  $C < \infty$ . The quasi-Newton direction  $d^k$  is computed by solving the equation

$$-\nabla f(x^k)[S^k] = B^k[S^k, S^k]d^k[S^k].$$

Multiplying both sides of this equation by  $(d^k[S^k])^\top$  gives

$$\begin{aligned} -\langle \nabla f(x^k)[S^k], d^k[S^k] \rangle &= (d^k[S^k])^\top B^k[S^k, S^k]d^k[S^k] \\ &\geq c \|d^k[S^k]\|^2. \end{aligned}$$

Since  $d^k[\bar{S}^k] = 0$ , this inequality can further be simplified to

$$\langle \nabla f(x^k), d^k \rangle \leq -c \|d^k\|^2. \quad (2)$$

Since we are in the second case (Algorithm 2, Lines 6–8) we know that for all  $i \in S^k$ , if  $d_i^k < 0$  and  $x_i \leq l_i + \varepsilon$  it must hold that  $\nabla f(x^k)_i < 0$ , otherwise  $i \notin S^k$ . In this case we have  $d_i^k \cdot \nabla f(x^k)_i > 0$  and at the same time we set  $p_i^k = 0$ . Hence, we have  $d_i^k \cdot \nabla f(x^k)_i > p_i^k \cdot \nabla f(x^k)_i$ . The case with  $d_i^k > 0$  and  $x_i \geq u_i + \varepsilon$  follows analogously. Hence, summing over all indices  $i$ , we can conclude  $\langle \nabla f(x^k), d^k \rangle \geq \langle \nabla f(x^k), p^k \rangle$ . Combining this inequality with Equation (2), we get  $-c \|d^k\|^2 \geq \langle \nabla f(x^k), p^k \rangle$ . Hence, we finally get

$$\begin{aligned} f(x^{k+1}) &= f(x^k + \alpha^k p^k) \\ &\leq f(x^k) + \alpha^k \langle \nabla f(x^k), p^k \rangle + \frac{L}{2} \|\alpha^k p^k\|^2 \\ &\leq f(x^k) - \alpha^k c \|d^k\|^2 + (\alpha^k)^2 \frac{L}{2} \|d^k\|^2, \end{aligned}$$

because  $\|p^k\| \leq \|d^k\|$ . Since all  $x_i^k$  with  $d_i^k \neq 0$  are at least  $\varepsilon$  away from the boundary, we can pick  $\alpha^k$  at least  $\min_i \frac{\varepsilon}{|d_i^k|} = \frac{\varepsilon}{\|d^k\|_\infty}$ . If  $\frac{c}{L} \leq \frac{\varepsilon}{\|d^k\|_\infty}$ , we can set  $\alpha^k = \frac{c}{L}$  and obtain the same result as in the first case. Otherwise, we set  $\alpha^k = \frac{\varepsilon}{\|d^k\|_\infty}$  and obtain

$$\begin{aligned} f(x^{k+1}) &= f(x^k + \alpha^k p^k) \\ &\leq f(x^k) - \frac{\varepsilon}{\|d^k\|_\infty} c \|d^k\|^2 + \left( \frac{\varepsilon}{\|d^k\|_\infty} \right)^2 \frac{L}{2} \|d^k\|^2 \\ &\leq f(x^k) - \frac{\varepsilon}{\|d^k\|_\infty} c \|d^k\|^2 + \frac{\varepsilon}{\|d^k\|_\infty} \frac{c}{L} \frac{L}{2} \|d^k\|^2 \\ &= f(x^k) - \frac{\varepsilon}{\|d^k\|_\infty} \frac{c}{2} \|d^k\|^2 \\ &\leq f(x^k) - \frac{\varepsilon c}{2} \|d^k\|, \end{aligned}$$

where the last line follows from  $\|d^k\|_\infty \leq \|d^k\|$ . It remains to lower bound the Euclidean norm of  $d^k$ . We have  $-\nabla f(x^k)[S^k] = B^k[S^k, S^k]d^k[S^k]$ . Taking the squared norm on both sides, we have

$$\begin{aligned} \|\nabla f(x^k)[S^k]\|^2 &= \|B^k[S^k, S^k]d^k[S^k]\|^2 \\ &= (d^k[S^k])^\top (B^k[S^k, S^k])^\top B^k[S^k, S^k]d^k[S^k] \\ &\leq C^2 \|d^k[S^k]\|^2 = C^2 \|d^k\|^2 \end{aligned}$$

since the eigenvalue of the submatrix  $B^k[S^k, S^k]$  can be upper bounded by the constant  $C$ . As long as Algorithm 1 has not converged, we know for the norm of the projected gradient that  $\|\nabla f(x^k)[S^k]\| \geq \varepsilon$ . Thus, we finally have

$$f(x^{k+1}) \leq f(x^k) - \frac{\varepsilon c}{2} \|d^k\| \leq f(x^k) - \frac{\varepsilon^2 c}{2C}.$$

Hence, also in the second case, the function value decreases by a positive constant in each iteration.

Thus, we make progress in each iteration by at least a small positive constant. Since  $f$  is bounded from below, the algorithm will converge to a stationary point. Finally, note that  $x^0$  is feasible. By construction and induction over  $k$  it follows that  $x^k$  is feasible for all  $k$ .  $\square$

Algorithm 1 uses the `projectDirection` subroutine (Algorithm 2). It becomes apparent from the proof that one can skip the projection branch (Line 3) and instead always follow the modified quasi-Newton direction and still obtain convergence guarantees. However, here, we use a projection as it was similarly suggested in (Morales and Nocedal 2011) which often reduces the number of iterations in the original L-BFGS-B algorithm (Byrd et al. 1995).

## Complete Framework

In the previous section we have described our approach for solving optimization problems with box constraints that can be efficiently run on a GPU. We extend this approach to also handle arbitrary constraints by using an augmented Lagrangian approach. This extension allows to solve constrained optimization problems of the form

$$\min_x f(x) \text{ s.t. } h(x) = 0, g(x) \leq 0, \text{ and } l \leq x \leq u, \quad (3)$$

where  $x \in \mathbb{R}^n$ ,  $f: \mathbb{R}^n \rightarrow \mathbb{R}$ ,  $h: \mathbb{R}^n \rightarrow \mathbb{R}^m$ ,  $g: \mathbb{R}^n \rightarrow \mathbb{R}^p$  are differentiable functions, and the equality and inequality constraints are understood component-wise.

The augmented Lagrangian of Problem (3) is the following function

$$\begin{aligned} L(x, \lambda, \mu, \rho) &= f(x) + \frac{\rho}{2} \|h(x) + \lambda/\rho\|^2 \\ &\quad + \frac{\rho}{2} \|(g(x) + \mu/\rho)_+\|^2, \end{aligned} \quad (4)$$

where  $\lambda \in \mathbb{R}^m$  and  $\mu \in \mathbb{R}_{\geq 0}^p$  are Lagrange multipliers,  $\rho > 0$  is a constant, and  $(v)_+$  denotes  $\max\{v, 0\}$ . The Lagrange multipliers are also referred to as dual variables.

The augmented Lagrangian Algorithm (see Algorithm 4 in the appendix) runs in iterations. In each iteration it minimizes the augmented Lagrangian function, Eq. (4), subject to the box constraints using Algorithm 1 and updates the Lagrange multipliers  $\lambda$  and  $\mu$ . If Problem (3) is convex, the augmented Lagrangian algorithm returns a global optimal solution. Otherwise, it returns a local optimum (Bertsekas 1999).

We integrated our solver with the modeling framework presented in (Laue, Mitterreiter, and Giesen 2019) that allows to specify the optimization problem in a natural, easy-to-read modeling language, see for instance the example to the right. Based on the matrix and tensor calculus methods presented in (Laue, Mitterreiter, and Giesen 2018), the framework

```
parameters
  Matrix A
  Vector b
variables
  Vector x
min
  norm2(A*x-b)
st
  sum(x) == 1
  x >= 0
```

then generates Python code that computes function values and gradients of the objective function and the constraints. The code maps all linear expressions to NumPy statements. Since any NumPy-compatible library can be used within the generated code this allows us to replace NumPy by CuPy to run the solvers on the GPU. We extended the modeling language and the Python code generator to our needs here. An anonymous interface can be found at <http://aaai2022.pythonanywhere.com>.

Our framework and solvers are solely written in Python, which makes them easily portable as long as NumPy-compatible interfaces are available. Here, we use the CuPy library (Okuta et al. 2017) in order to run the generated solvers on the GPU. An anonymized version of the code can be found in the supplemental material. Later, we will make the code available via github under the GNU LGPL license.

## Experiments

The purpose of the following experiments is to show the efficiency of our approach on a set of different classical, that is, non-deep, machine learning problems. For that purpose, we selected a number of well-known classical problems that are given as constrained optimization problems, where the optimization variables are either vectors or matrices. All these problems can also be solved on CPUs. In the supplemental material, we provide results that show that the

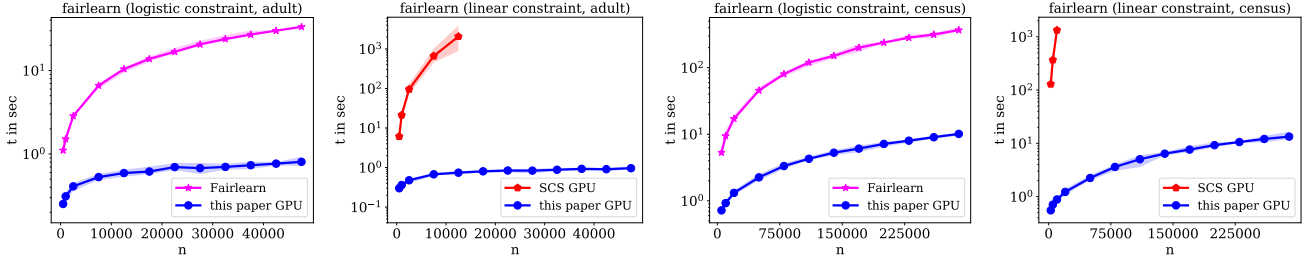


Figure 1: Running times for the logistic regression problem with fairness constraints. The two plots on the left show the running times for the adult data set and the two plots on the right for the census data set. For each data set, one plot shows the running times when the logistic loss is used in the fairness constraint and one plot when the linear loss is used.

GPU version of our framework significantly outperforms the efficient multi-core CPU framework GENO (Laue, Mitterreiter, and Giesen 2019). Here, we compare our framework on the GPU to CVXPY paired with the cuOSQP and SCS solvers. CVXPY has a similar easy-to-use interface and also allows to solve general constrained optimization problems. Note however, that CVXPY is restricted to convex problems and cuOSQP to convex quadratic problems. It was shown that cuOSQP outperforms its CPU version OSQP by about a factor of ten (Schubiger, Banjac, and Lygeros 2020). Our experiments confirm this observation.

To the best of our knowledge, pairing CVXPY with cuOSQP or SCS *are the only two approaches comparable to ours*. Another framework that has been released recently that can solve convex, constrained optimization problems on the GPU is cvxpylayers (Agrawal et al. 2019). However, its focus is on making the solution of the optimization problem differentiable with respect to the input parameters for which it needs to solve a slightly larger problem. Internally, it uses CVXPY combined with the SCS solver in GPU-mode. Hence, this framework is slower than the original combination of CVXPY and SCS.

In all our experiments we made sure that the solvers that were generated by our framework always computed a solution that was *better* than the solution computed by the competitors in terms of objective function value and constraint satisfaction. All experiments were run on a machine equipped with an Intel i9-10980XE 18-core processor running Ubuntu 20.04.1 LTS with 128 GB of RAM, and a Quadro RTX 4000 GPU that has 8 GB of GDDR6 SDRAM and 2304 CUDA cores. Our framework took always less than 10 milliseconds for generating a solver from its mathematical description. The code and all examples can be found in the supplement which also includes a Docker image for running all experiments.

## Fairness in Machine Learning

In classical machine learning approaches, one usually minimizes a regularized empirical risk in order to make correct predictions on unseen data. However, due to various causes, e.g., bias in the data, it can happen that some group of the input data is favored over another group. Such favors can be mitigated by the introduction of constraints that explicitly prevent them. This is the goal of fairness in machine learn-

ing which has gained considerable attention over the last few years. Here, we consider fairness for binary classification.

There are a number of different fairness definitions (Agarwal et al. 2018; Barocas, Hardt, and Narayanan 2019; Donini et al. 2018), see the Fairlearn project (Bird et al. 2020, 2021) for an introduction and overview. Here, we follow the exposition and formulation in (Donini et al. 2018). Let  $D = \{(x^1, y^1), \dots, (x^m, y^m)\}$  be a labeled data set and  $A$  and  $B$  be two groups, i.e., subsets of the data set. Then, one seeks to find a classifier that is statistically independent of the group membership  $A$  and  $B$ . Depending on the type of groups  $A$  and  $B$ , respectively, different types of fairness constraints are obtained. Since statistical independence is defined with respect to the true data distribution, which is typically unknown, one replaces the expectation over the true distribution by the empirical risk. Hence, one solves the following constrained optimization problem

$$\min_f \hat{L}_D(f) + \lambda \cdot r(f) \quad \text{s.t.} \quad \hat{L}_A(f) = \hat{L}_B(f),$$

where  $f: X \rightarrow \mathbb{R}$  is a function or model,  $l: \mathbb{R} \times Y \rightarrow \mathbb{R}$  is a loss function,  $\hat{L}_D(f) = \frac{1}{|D|} \sum_{(x^i, y^i) \in D} l(f(x^i), y^i)$  is the empirical risk of  $f$  over the data set  $D$ , and  $r(\cdot)$  is the regularizer.

Ideally, one would like to use the same loss function for the risk minimization  $\hat{L}_D(f)$  as in the fairness constraint  $\hat{L}_A(f) = \hat{L}_B(f)$ . The logistic loss is often used for classification. However, when the logistic loss is used in the fairness constraint, the problem becomes non-convex, even for a linear classifier. Using our framework, we can still solve this problem. However, we cannot compare its performance to cuOSQP or SCS since they only allow to solve convex problems. Thus, we compare it to the exponentiated gradient approach (Agarwal et al. 2018) paired with the Liblinear solver (Fan et al. 2008). Note, that this approach does not run on the GPU. However, to provide a better global picture, we still include it here. Only when the loss function in the fairness constraint is linear as in (Donini et al. 2018; Zemel et al. 2013), the problem becomes convex. We also consider this case and compare it to SCS. Note, the problem cannot be solved by cuOSQP since it contains exponential cones.

We used the same setup, the same data sets, and the same preprocessing as described in the Fairlearn package (Bird et al. 2020). We used the adult data set (48,842 data points

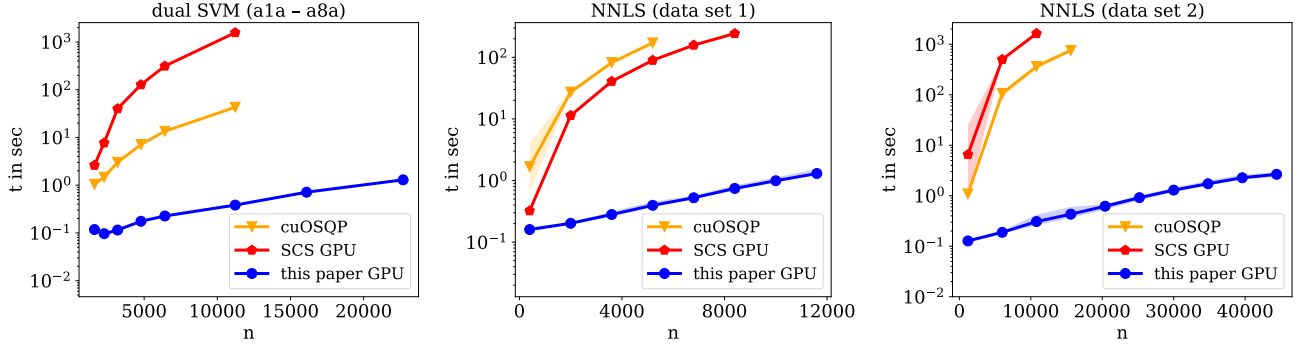


Figure 2: The plot on the left shows the running times for the SVM problem on the adult data set. The plot in the middle and the plot on the right show the running times for the non-negative least squares problem when run on the first and second data set, respectively.

Solver	Data sets						
	cod-rna	covtype	ijcnn1	mushrooms	phishing	a9a	w8a
this paper GPU	0.6	0.1	1.8	0.1	1.6	0.3	1.4
cuOSQP	55.7	failed	206.5	22.0	163.5	32.8	36.1
SCS GPU	2342.0	31.3	N/A	7994.9	N/A	1094.0	1227.1

Table 1: Running times in seconds for the dual SVM problem. All data sets were subsampled to 10,000 data points due to time and memory requirements of the cuOSQP and SCS solver. N/A indicates that the solver did not finish within 10,000 seconds.

with 120 features) and the census-income data set (299,285 data points with 400 features) each with ‘female’ and ‘male’ as the two subgroups. For each experiment, we sampled  $m$  data points from the full data set. Figure 1 shows the running times. Our framework provides similar results in terms of quality as the exponentiated gradient approach, when the logistic loss is used in the fairness constraint, and it is orders of magnitude faster than SCS on the GPU, when the linear loss is used in the fairness constraint.

## Dual SVM

Support vector machines (SVMs) are a classical yet still relevant classification method. When combined with a kernel, they are usually solved in the dual problem formulation, which reads as

$$\min_a \frac{1}{2} (a \odot y)^\top K (a \odot y) - \|a\|_1 \text{ s.t. } y^\top a = 0, 0 \leq a \leq c,$$

where  $K \in \mathbb{R}^{n \times n}$  is a positive semidefinite kernel matrix,  $y \in \{-1, +1\}^n$  are the corresponding binary labels,  $a \in \mathbb{R}^n$  are the dual variables,  $\odot$  is the element-wise multiplication, and  $c \in \mathbb{R}_+$  is the regularization parameter.

We used all data sets from the LibSVM data sets website (Lin and Fan 2021) that had more than 8000 data points with fewer than 1000 features such that a kernel approach is reasonable. We applied a standard Gaussian kernel with bandwidth parameter  $\gamma = 1$  and regularization parameter  $c = 1$ . Table 1 shows the running times for the data sets when subsampled to 10,000 data points. Figure 2 shows the running times for an increasing number of data points based on the original subsampled adult data set (Lin and

Fan 2021). It can be seen that our approach outperforms cuOSQP as well as SCS by several orders of magnitude. The cuOSQP solver ran out of memory for problems with more than 10,000 data points. While there is a specialized solver for solving these SVM problems on the GPU (Wen et al. 2018), the focus here is on general purpose frameworks.

## Non-negative Least Squares

Non-negative least squares is an extension of the least squares regression problem that requires the output to be non-negative. See (Slawski 2013) for an overview on the non-negative least squares problem. It is given as the following optimization problem

$$\min_x \|Ax - b\|_2^2 \text{ s.t. } x \geq 0,$$

where  $A \in \mathbb{R}^{m \times n}$  is the design matrix and  $b \in \mathbb{R}^m$  the response vector. We ran two sets of experiments, similarly to the comparisons in (Slawski 2013), where it was shown that different algorithms behave quite differently on these problems. For experiment (i), we generated a random data matrix  $A \in \mathbb{R}^{2000 \times 6000}$ , where the entries of  $A$  were sampled uniformly at random from the unit interval and a sparse vector  $x \in \mathbb{R}^{6000}$  with non-zero entries sampled from the standard Gaussian distribution and a sparsity of 0.01. The response variables were then generated as  $y = \sqrt{0.003} \cdot Ax + 0.003 \cdot z$ , where  $z \sim \mathcal{N}(0, 1)$ . For experiment (ii),  $A \in \mathbb{R}^{6000 \times 3000}$  was drawn from a Gaussian distribution and  $x$  had a sparsity of 0.1. The response variable was generated as  $y = \sqrt{1/6000} \cdot Ax + 0.003 \cdot z$ , where  $z \sim \mathcal{N}(0, 1)$ . The differences between the two experiments are: (1) The Gram



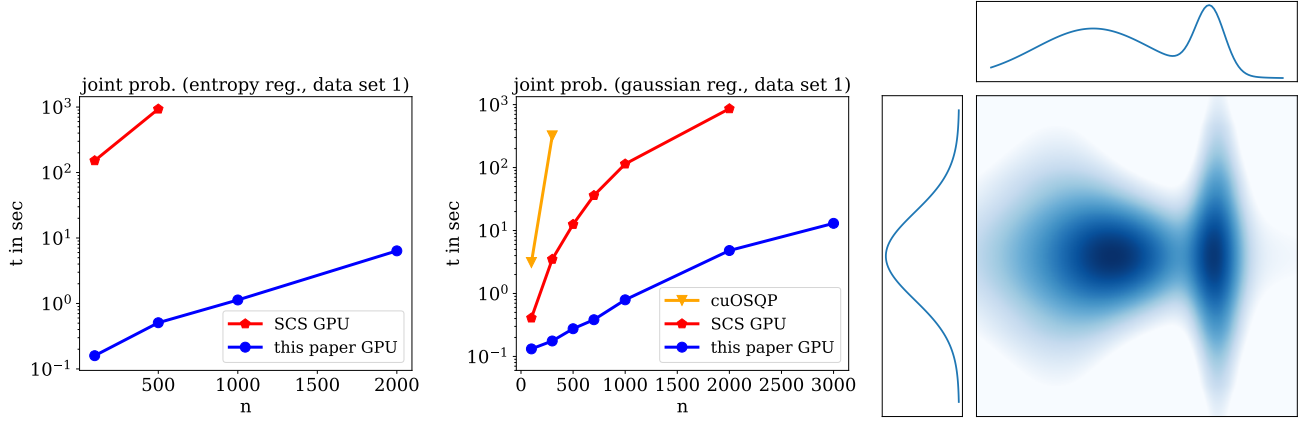


Figure 3: Running times for computing the joint probability distribution from two marginal distributions. The left plot shows the running times when the entropy prior is used and the plot in the middle when a Gaussian prior is used. The right figure visualizes the probabilities.

matrix  $A^\top A$  is singular in experiment (i) and regular in experiment (ii), (2) the design matrix  $A$  has isotropic rows in experiment (ii) but not in experiment (i), and (3)  $x$  is significantly sparser in (i) than in (ii). To evaluate the runtime behavior for increasing problem size, we scaled the problem sizes to  $A \in \mathbb{R}^{2000t \times 6000t}$  in the first experiment and to  $A \in \mathbb{R}^{6000t \times 3000t}$  in the second experiment for a parameter  $t \in [0, 6]$ . For each problem instance we performed ten runs and report the average running time along with the standard deviation in Figure 2. We stopped including SCS and cuOSQP into the experiments once their running time exceeded 1000 seconds. It can be seen that SCS is faster than cuOSQP on the first set of experiments and slower than cuOSQP on the second set. However, our approach outperforms SCS and cuOSQP by several orders of magnitude in both sets of experiments.

### Joint Probability Distribution

Given two discrete probability distributions  $u \in \mathbb{R}^m$  and  $v \in \mathbb{R}^n$ , we are interested in their joint probability distribution  $P \in \mathbb{R}^{m \times n}$ . This problem has been studied intensively before, see, e.g., (Cuturi 2013; Frogner and Poggio 2019; Muzellec et al. 2017). With additional knowledge, it can be reduced to a regularized optimal transport problem. Many different regularizers have been used, for instance, an entropy, a Gaussian, or more generally, a Tsallis regularizer. The corresponding optimization problem is the following constrained optimization problem over positive matrices

$$\min_P \langle M, P \rangle + \lambda \cdot r(P) \quad \text{s. t.} \quad P \cdot \mathbf{1} = u, P^\top \cdot \mathbf{1} = v, 0 \leq P,$$

where  $M \in \mathbb{R}^{m \times n}$  is the cost matrix,  $r(\cdot)$  is the regularizer,  $\mathbf{1}$  is the all-ones vector, and  $\lambda \in \mathbb{R}_+$  is the regularization parameter. In our experiments we used the entropy and the Gaussian regularizer that are both special cases of the Tsallis regularizer. In the special case that the regularizer is the entropy,  $m = n$ , and the cost matrix  $M$  is a metric, (Cuturi 2013) showed that the problem can be solved using Sinkhorn’s algorithm (Knopp and Sinkhorn 1967). Similar

results are known for other special cases (Janati et al. 2020). However, in general Sinkhorn’s algorithm cannot be used as it is the case in the present experiments, since the cost matrix is not a metric.

Here, we used synthetic data sets since the real-world data sets that are usually used for this task are very small, typically  $m, n \leq 20$ . We created two sets of synthetic data sets. For the first set of data sets, we let a Gaussian and a mixture of two Gaussians be the marginals, see Figure 3. Then, we discretized both distributions to obtain the marginal vectors  $u$  and  $v$ . In this case, we set  $m = n$ . Hence, when  $n = 1000$ , the corresponding optimization problem has  $10^6$  optimization variables with lower bound constraints and 2000 equality constraints. The cost matrix  $M$  was fixed to be the discretized version  $uu^\top$  of a two-dimensional Gaussian, and the regularization parameter was set as  $\lambda = \frac{1}{2}$ . We ran two sets of experiments on this data set, one where  $r(\cdot)$  is the entropy regularizer and another one with the Gaussian regularizer. Figure 3 shows the running times for both experiments for varying problem sizes. It can be seen that our approach outperforms cuOSQP and SCS by several orders of magnitude. The cuOSQP solver ran out of memory already on very small problems. The second data set was created as in (Frogner and Poggio 2019). On this data set, the speedup over cuOSQP and SCS is even more pronounced. Detailed results can be found in the appendix.

### Conclusion

We presented a framework for solving constrained optimization problems on the GPU efficiently. The framework allows to specify a constrained optimization problem in an easy-to-read modeling language and then generates solvers in portable Python code that outperform competing state-of-the-art approaches on the GPU by several orders of magnitude. Using GPUs also for classical, that is, non-deep, machine learning becomes increasingly important as hardware vendors started to combine CPUs and GPUs on a single chip like Apple’s M1 chip.

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