



The Network Layer

Redes de Computadores

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References

- These slides are from “Computer Networking: A Top Down Approach 5th edition. Jim Kurose, Keith Ross Addison-Wesley, April 2009”
 - With adaptations/additions by Manuel Ricardo and Pedro Brandão

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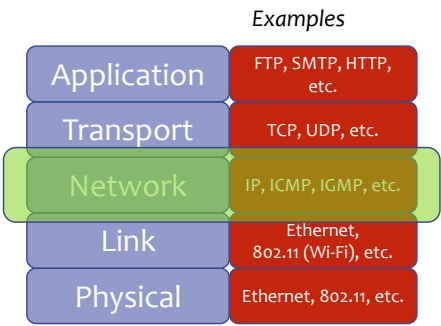
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Driving questions...

- What are the main functions of the network layer?
- What are the differences between virtual circuit and datagram networks?
- How is forwarding handled in both type of networks?
- What are the main functions of a router?
- What are the formats of IP addresses?
- How to form subnets?
- What services are provided by ARP, ICMP, DHCP and NAT? How do these protocols work?
- What are differences the between IPv4 and IPv6?

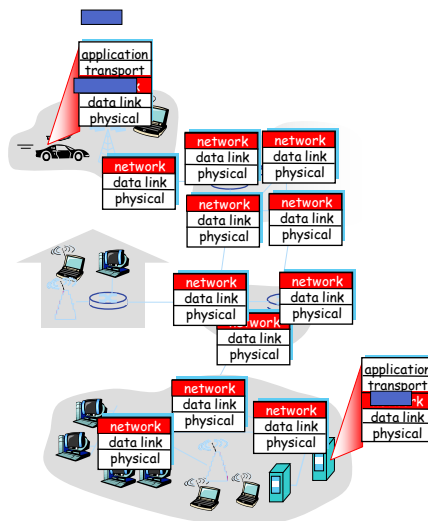
Internet protocol stack

- **Application:** network processes
- **Transport:** data transfer between processes
- **Network:** packet routing between source and destination
- **Link:** data transfer between adjacent network elements
- **Physical:** bits on the “wire”



Network Layer Overview

- Network layer
 - transports packets (datagrams)
 - from sending host to receiving host
 - functions located in every host and router
- Sender
 - encapsulates transport data into packets
 - generates packets
- Receiver
 - receives packets
 - delivers data to transport layer
- Router
 - Receives packets from input line
 - examines network layer header
 - forwards packets through adequate output line(s)

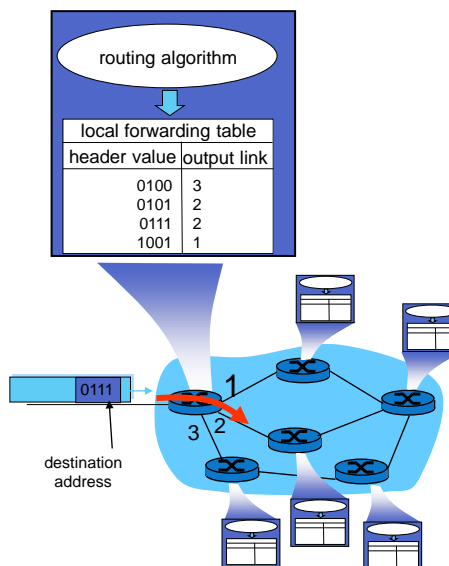


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Network Layer – Main Functions

- Routing
 - determine route taken by packets, from source to destination
 - Algorithms using cost function (usually shortest path)
 - *Analogy: process of planning trip from source to destination*
- Forwarding
 - router forwards packet **from input port to output port**
 - *Analogy: process of getting through single interchange*



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Virtual Circuits and Datagram Networks

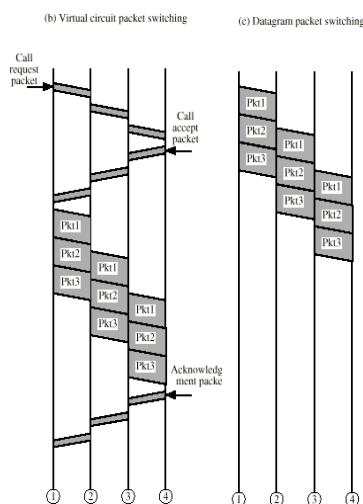
(Network layer)

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Network Layer – Connection and Connectionless Service

- Services provided by network layer
 - Virtual Circuit network
 - connection-oriented service
 - Datagram network
 - connectionless service



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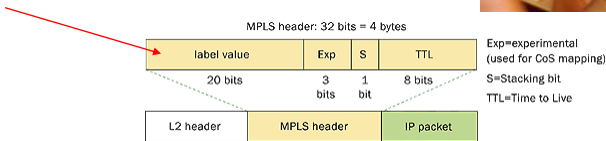
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Virtual Circuit (VC)

Image from Joseph A. Carr, in Wikipedia Telephone Exchange

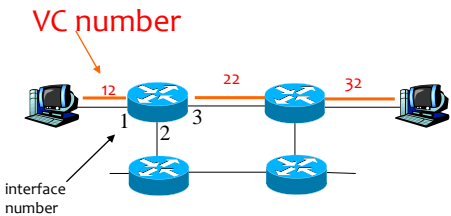


- Phases
circuit establishment → data transference → circuit termination
- Packet carries identifier of Virtual Circuit



- Path defined from source to destination
 - sequence of VC identifiers, one for each link along path
- Router
 - maintains “state” for every supported circuit
 - may allocate resources (bandwidth, buffers) per Virtual Circuit

VC - Forwarding Table



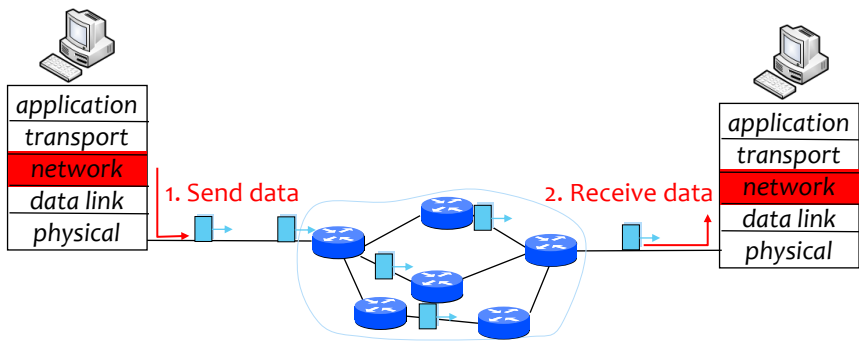
Forwarding table in northwest router:

Incoming interface	Incoming VC #	Outgoing interface	Outgoing VC #
1	12	3	22
2	63	1	18
3	7	2	17
1	97	3	87
...

Routers maintain connection state information!

Datagram Networks

- No circuit establishment; no circuit concept
- Packets
 - forwarded using destination host address
 - packets between same source-destination pair may follow different paths



Forwarding Table

- IP Address
 - 32 bits

2³² possible entries in IPv4

Destination Address Range	Output Link interface
address X through address Z	0
address W through address Y	1
address A through address K	1
address P through address R	2
Otherwise	3

- How to reduce the number of entries in the forwarding table?

Longest Prefix Matching

- Which interface?
- DA: 11001000 00010111 00010110 10100001
 - → Itf: 0
- DA: 11001000 00010111 00011000 10101010
 - Can be itf 1 or 2
 - Longest prefix 1

Prefix	Output Link interface
11001000 00010111 00010	0
11001000 00010111 00011000	1
11001000 00010111 00011	2
Otherwise	3

Virtual-Circuit versus Datagram Networks

Issue	Datagram subnet	Virtual-circuit subnet
Circuit setup	Not needed	Required
Addressing	Each packet contains the full source and destination address	Each packet contains a short VC number
State information	Routers do not hold state information about connections	Each VC requires router table space per connection
Routing	Each packet is routed independently	Route chosen when VC is set up; all packets follow it
Effect of router failures	None, except for packets lost during the crash	All VCs that passed through the failed router are terminated
Quality of service	Difficult	Easy if enough resources can be allocated in advance for each VC
Congestion control	Difficult	Easy if enough resources can be allocated in advance for each VC

Internet Protocol

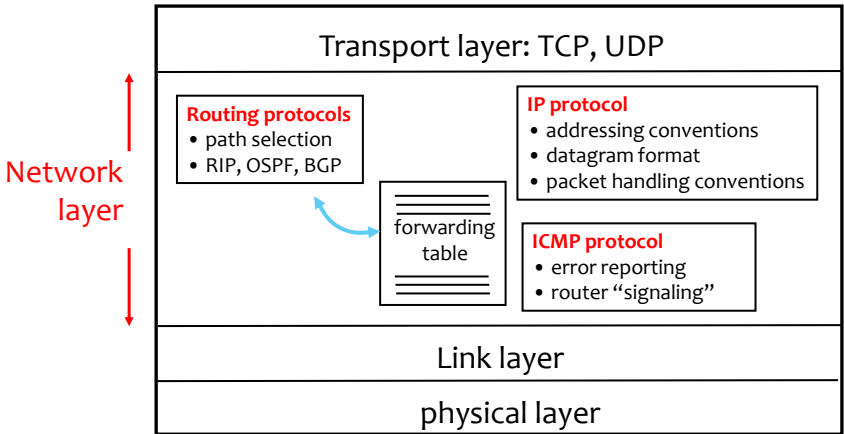
[RFC 791/STD 5](#)

[IP RFCs](#)

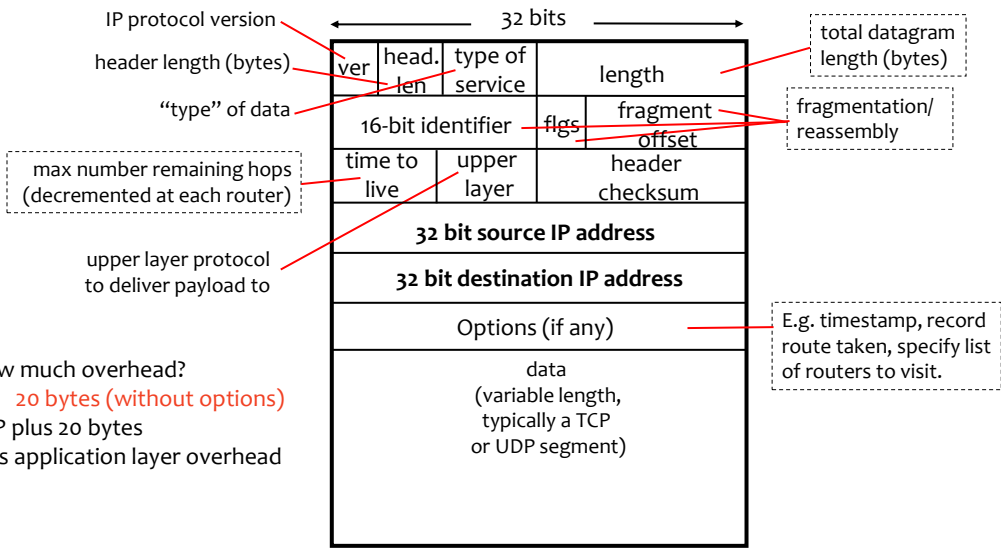
(Network layer)

The Internet Network layer

- Host, router network layer functions

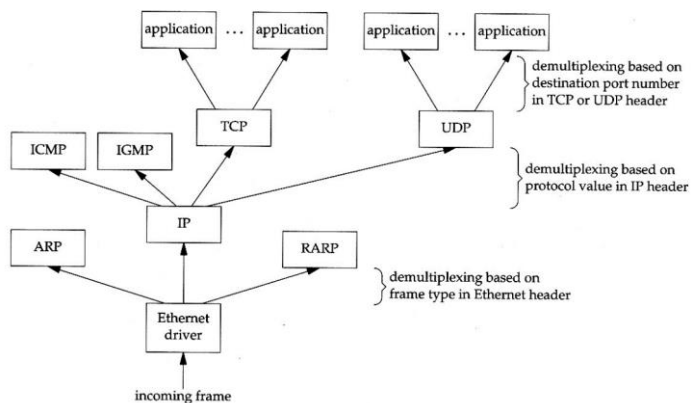


IP Datagram Format



Demultiplexing for upper layers

- Ethernet header (type)
 - IP - 0x0800
 - ARP - 0x0806
 - RARP - 0x8035
 - IPX - 0x8037
 - IPv6 - 0x86DD
 - MPLS - 0x8847
- IP header (protocol)
 - ICMP - 1
 - IGMP - 2
 - TCP - 6
 - UDP - 17
- TCP/UDP header (port)
 - FTP - 21
 - Telnet - 23
 - HTTP - 80
 - SMTP - 25



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Internet Checksum

- The Internet (**not layer 2**) uses a checksum
 - easily implementable in software
 - 1's complement sum of 16 bit words
 - Performance: d=2

```

u_short
cksum(u_short *buf, int count)
{
    register u_long sum = 0;

    while (count--)
    {
        sum += *buf++;
        if (sum & 0xFFFF0000)
        {
            /* carry occurred,
             so wrap around */
            sum &= 0xFFFF;
            sum++;
        }
    }
    return ~(sum & 0xFFFF);
}

```

- One's complement sum
 - Mod-2 addition **with carry-out**
 - Carry-out in the most-significant-bit is added to the least-significant bit
 - Get one's complement of "one's complement sum"

```

      1010011
      0110110
      carry-out ① 0001001
      Carry wrap-around 0000001
      0001010
      One's complement = 1110101

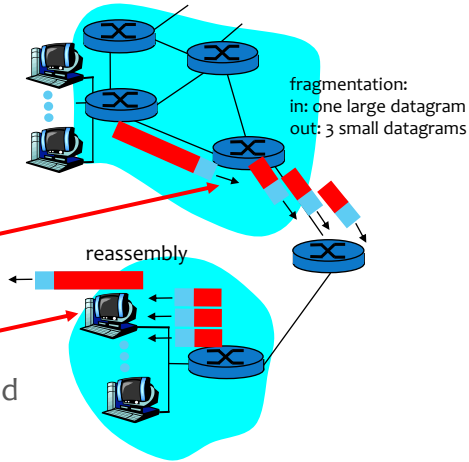
```

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IP Fragmentation and Reassembly

- Network links have MTU
 - MTU - max. transfer unit (size)
 - largest possible link-level frame
 - different link types, different MTUs
- Large IP datagram is fragmented
 - one datagram → n datagrams
 - “reassembled” at final destination
 - IP header bits used to identify, order related fragments



IP Fragmentation and Reassembly Example

- Example
- ♦ 4000 byte datagram
 - ♦ 3980 bytes data + 20 bytes IP header
 - ♦ MTU = 1500 bytes

length	ID	frag flag	offset
=4000	=x	=0	=0

One large datagram becomes several smaller datagrams

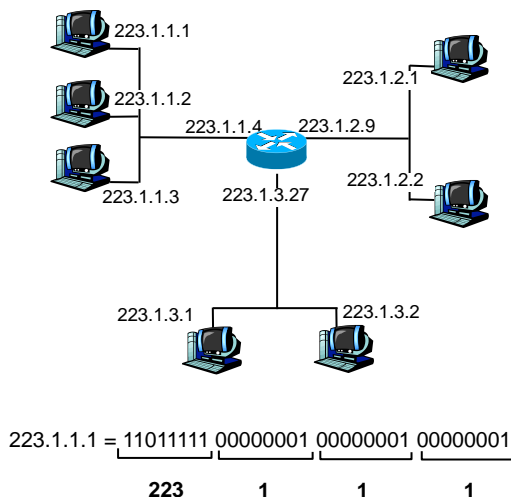
1480 bytes in data field

offset = 1480/8

length	ID	frag flag	offset
=1500	=x	=1	=0
=1500	=x	=1	=185
=1040	=x	=0	=370

IP Addressing - Introduction

- IP address
 - 32-bit identifier for host/router interface
- Interface
 - connection between host/router and physical link
 - Routers have multiple interfaces
 - IP addresses associated with interface

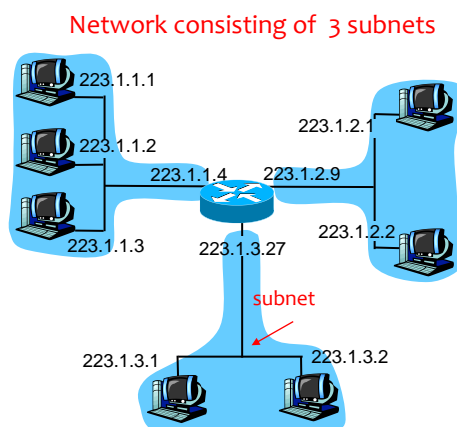


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Subnets

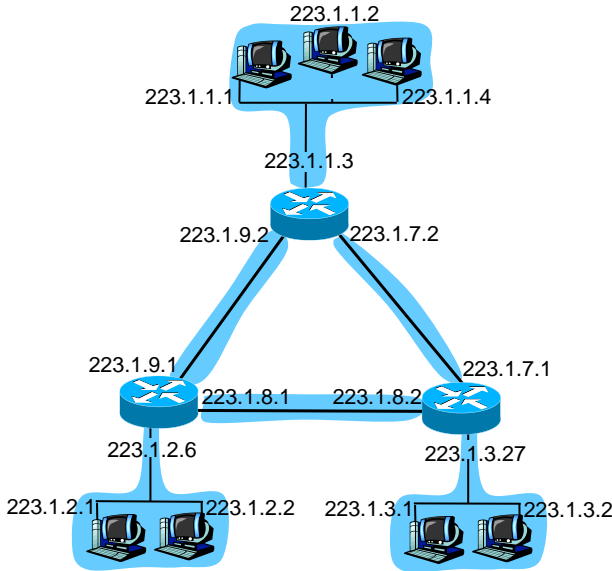
- IP address
 - subnet part → high order bits
 - host part → low order bits
- Subnet → set of interfaces
 - with same subnet part of IP address
 - can reach each other without router intervention



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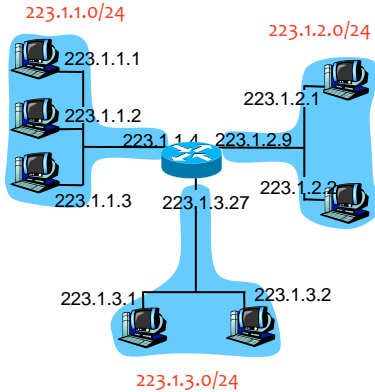
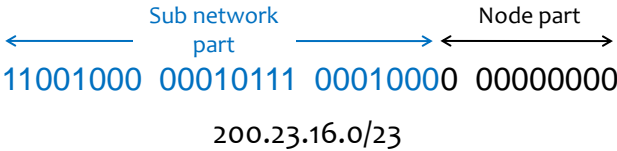
6 Subnets



IP Addressing - CIDR

CIDR: Classless InterDomain Routing ([RFC4632/BCP122*](#))

- subnet portion of address has arbitrary length
- address format → a.b.c.d/x
 - where x is # bits in subnet portion of address



* BCP – Best Current Practice

Special IP Addresses

0 0	This host
0 0 . . . 0 0 Host	A host on this network
1 1	Broadcast on the local network
Network 1 1 1 1 . . . 1 1 1 1	Broadcast on a distant network
127 (Anything)	Loopback

Forming Sub-Networks

Network **192.228.17.0/24** is divided in **8 subnetworks** → masks of 27 bits

LAN X Net ID/Subnet ID: 192.228.17.32 Subnet number: 1
R1 IP Address: 192.228.17.33 Host number: 1
A IP Address: 192.228.17.57 Host number: 25
B

LAN Y Net ID/Subnet ID: 192.228.17.64 Subnet number: 2
C IP Address: 192.228.17.65 Host number: 1
R2

LAN Z Net ID/Subnet ID: 192.228.17.96 Subnet number: 3
D IP Address: 192.228.17.97 Host number: 1

Subnetwork mask – 27 bits
subnetid – 3 bits (8 subnetworks)

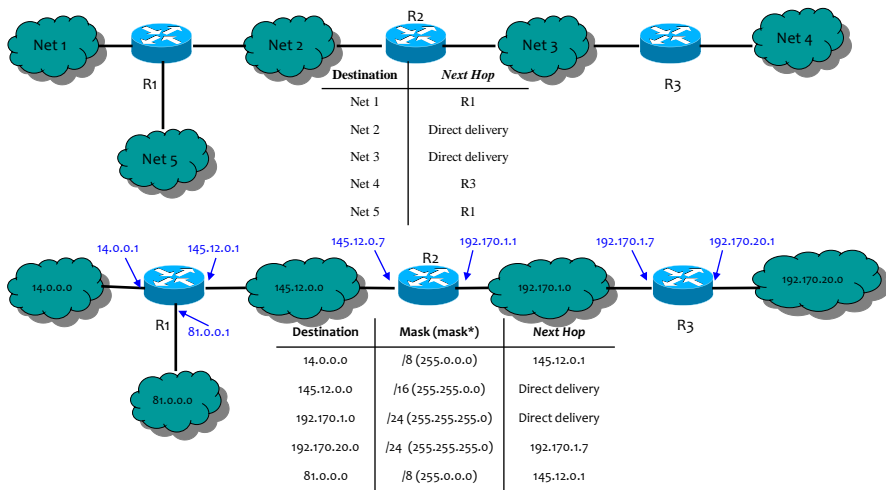
11000000 11100100 00010001 01100000
192.228.17.96/27

hostid – 5 bits
30 hosts per subnet supported
all 0 – identifies subnet
all 1 – broadcast address

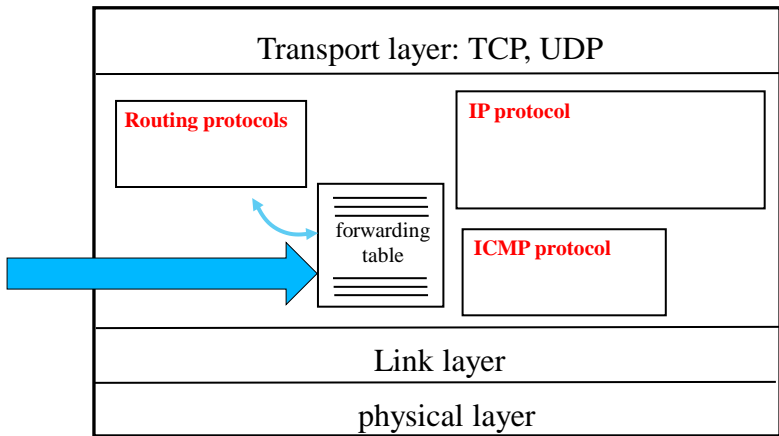
Example of subnetworks

192.228.17.0/27	(.00000000)
192.228.17.32/27	(.00100000)
192.228.17.64/27	(.01000000)
192.228.17.96/27	(.01100000)
...	
192.228.17.224/27	(.11100000)

Forwarding Table at R2



Forwarding Table use



- What is a loopback interface? What is its IP address?

IP Forwarding Function

- Forwarding table has entries in format
<networkAddress/mask, port>
- Forwarding function
 - When a datagram arrives with destination address **A**, then
 - For each entry of the forwarding table


```
val= A & mask* // e.g., mask=8, mask*=255.0.0.0 = (bin)11111111.0.0.0
```

```
if (val == networkAddress & mask*)
```

 - add corresponding output port to the set of candidate ports
 - Select the port with the largest mask → most specific route

IP Forwarding Function - Example

- $\text{frdTbl} = \{ \langle 128.32.0.0/16, 1 \rangle, \langle 128.32.192.0/18, 3 \rangle, \langle 128.0.0.0/8, 5 \rangle \}$
- Datagram with destination address $A = 128.32.195.1$
- Set of candidate output ports
 - $\rightarrow \{1, 3, 5\}$.
 - Port 1: 128.32.0.1 - 128.32.255.254
 - Port 3: 128.32.192.1 - 128.32.255.254
 - Port 5: 128.0.0.1 - 128.255.255.254
- Selected port
 - 3 largest mask, 18 bits

One online [IP Subnet Calculator](#)

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Address Resolution Protocol

[RFC 826/STD 37](#)

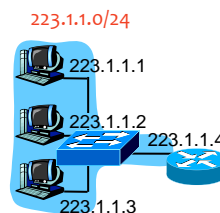
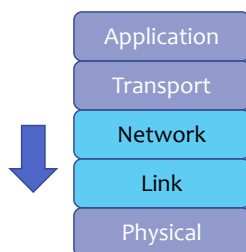
(Network layer)

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What is ARP needed for?

- On the same LAN you need to connect on the link layer
 - **Known:** IP address of destination (packet from Network Layer)
 - **Unknown:** MAC address of destination



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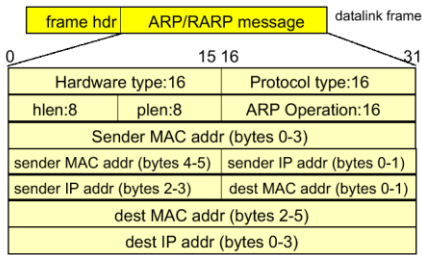
ARP – Address Resolution Protocol

- A network interface has
 - one MAC address
 - one (or more) IP address(es)
- ARP: Address Resolution Protocol
 - Protocol used to **obtain the MAC address** associated to a given IP address
- RARP – Reverse ARP
 - Protocol used to **obtain the IP address** associated to a MAC address

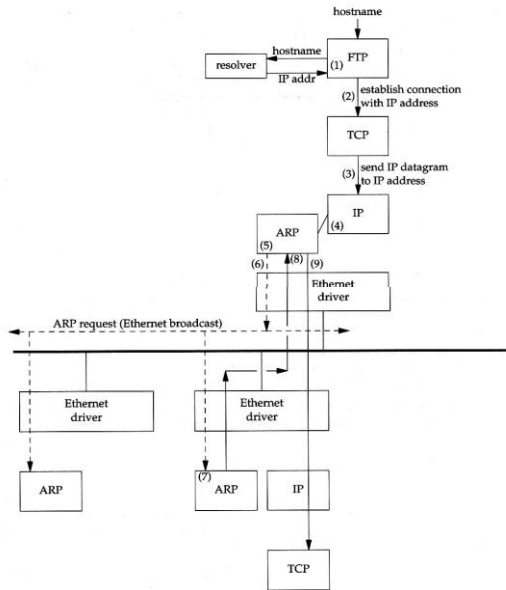
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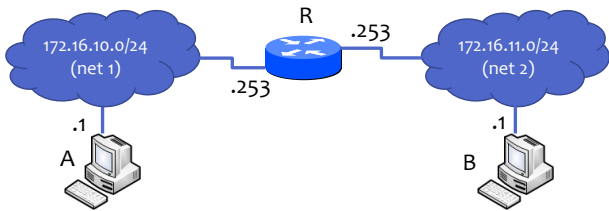
ARP Example



- **hardware type** : Ethernet=1 ARCNET=7, localtalk=11
- **protocol type** : IP=0x800
- **hlen** : length of hardware address, Ethernet=6 bytes
- **plen** : length of protocol address, IP=4 bytes
- **ARP operation** : ARP request = 1, ARP reply = 2
RARP request = 3, RARP reply = 4



ARP/IP addresses



- Assume host A sends an IP packet to host B and that this packet is forwarded by router R. What are the MAC and IP addresses (source and destination) observed?
- What roles does ARP play in this scenario?

Obtaining IP Addresses

RFC 2131 Dynamic Host Configuration Protocol

(Network layer)

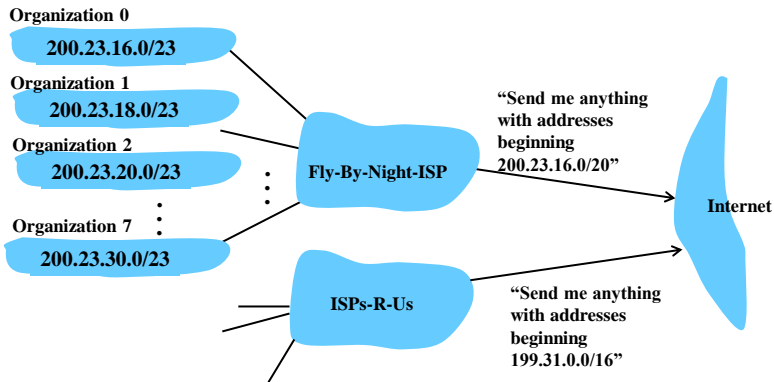
How to Obtain IP Addresses

- How does network get subnet part of IP addresss?
 - Gets allocated portion of its provider ISP’s address space

ISP's block	<u>11001000</u>	<u>00010111</u>	<u>00010000</u>	00000000	200.23.16.0/20
Organization 0	11001000	00010111	0001 <u>000</u> 0	00000000	200.23.16.0/23
Organization 1	11001000	00010111	0001 <u>001</u> 0	00000000	200.23.18.0/23
Organization 2	11001000	00010111	0001 <u>010</u> 0	00000000	200.23.20.0/23
...
Organization 7	11001000	00010111	0001 <u>111</u> 0	00000000	200.23.30.0/23

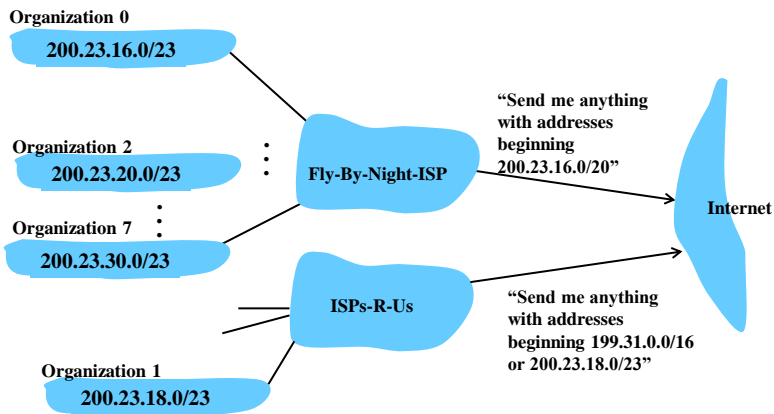
Hierarchical Addressing - Route Aggregation

- Hierarchical addressing
allows efficient advertisement of routing information



Hierarchical Addressing – More specific routes

- ISPs-R-Us has a more specific route to Organization 1



IP Addressing

- How does an ISP get block of addresses?
 - From [ICANN](#): Internet Corporation for Assigned Names and Numbers
 - ICANN
 - allocates addresses
 - manages Domain Name Service (DNS)
 - assigns domain names, resolves disputes

IP Addresses

- How does a host obtain an IP address?
 - Hard-coded by system admin in a file
 - [Windows](#): control-panel → network and sharing center → change adapter settings properties → tcp/ip → properties
 - UNIX/Linux/BSD: /etc/sysconfig/network-scripts/, /etc/network/interfaces, [ifconfig](#), [ip](#) addr
 - DHCP: Dynamic Host Configuration Protocol
 - dynamically get address from a server
 - “plug-and-play”

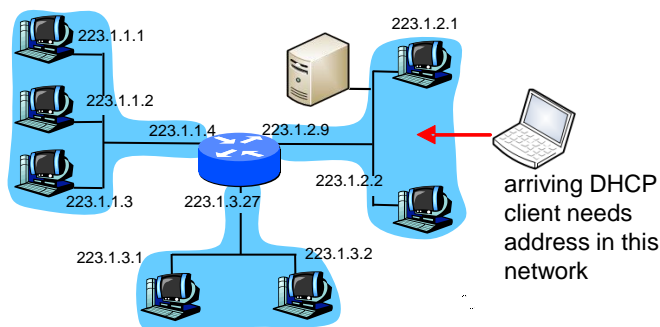
DHCP - Dynamic Host Configuration Protocol

- DHCP allows
 - host to dynamically obtain its IP address from network server when it joins network
 - It supports address reuse
- DHCP overview
 - host broadcasts “**DHCP discover**” msg
 - DHCP server responds with “**DHCP offer**” msg
 - host requests IP address “**DHCP request**” msg
 - DHCP server sends address “**DHCP ack**” msg

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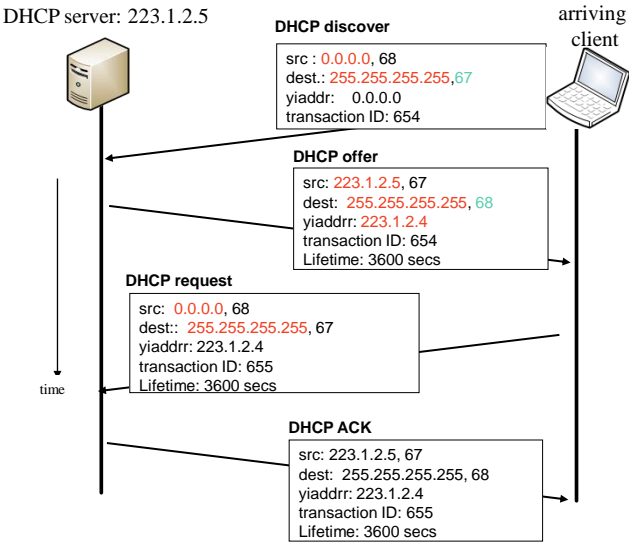
DHCP - Client-server Scenario



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DHCP Client-server



- Is it sufficient for an arriving client to acquire an IP address? What other relevant information shall this client obtain in order to start working with full functionality?

Network Address Translation

RFC 3022 Traditional IP Network Address Translator (Traditional NAT)

(Network layer)

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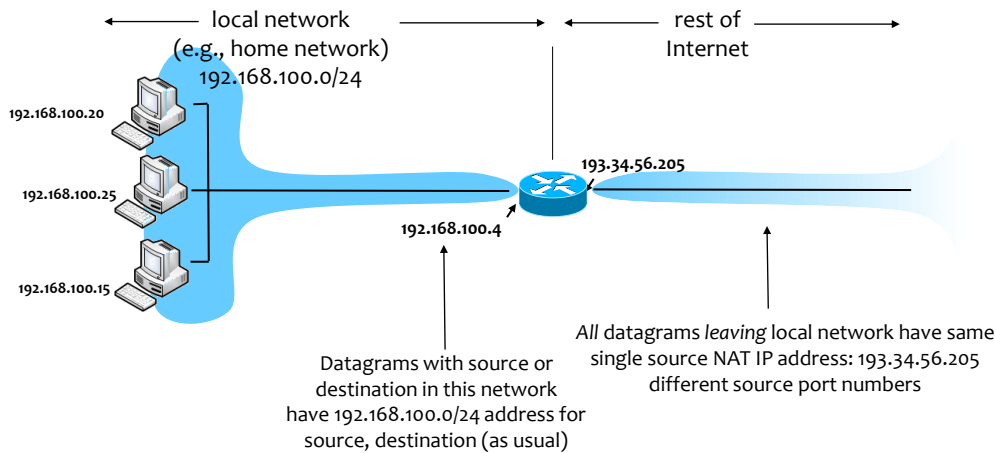
NAT: Motivation

- Shortage of IP addresses
 - Small/medium companies with ADSL connections, cable, want IPs for their machines (also domestic users).
- Uses private IP addresses
 - Not allowed in the Internet
 - 10.0.0.0/8
 - 172.16.0.0/12
 - 192.168.0.0/16

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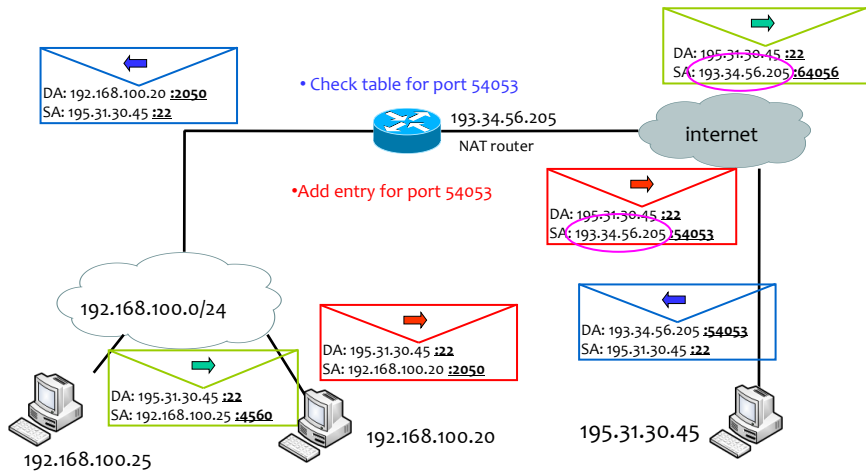
NAT - Network Address Translation



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Private Address	Private Port	Outside address	Outside port	NAT port	Protocol
192.168.100.20	2050	195.31.30.45	22	54053	tcp
192.168.100.25	4560	195.31.30.45	22	64056	tcp



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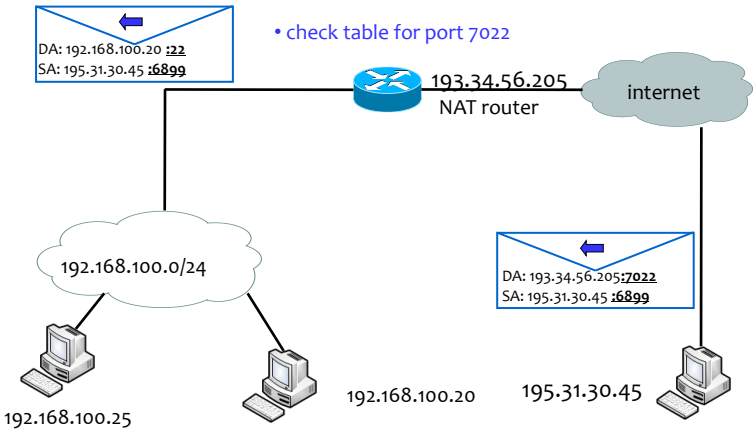
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NAT traversal problem



- As inside address are private (not seen in the internet) how can an outside host connect to an internal one?
- Possible solution statically add entry to NAT table

Private address	Private port	Outside address	Outside port	NAT port	Protocol
192.168.100.20	22			7022	tcp
192.168.100.25	80			9080	tcp



Internet Message Control Protocol

STD 5/RFC 792 Internet Control Message Protocol

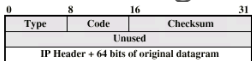
(Network layer)

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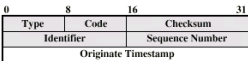
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ICMP - Internet Control Message Protocol

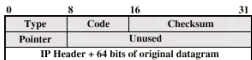
- Used by router or host
 - to send layer 3 error or control messages
 - to other hosts or routers
- Carried in IP datagrams



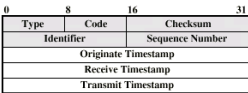
(a) Destination Unreachable; Time Exceeded; Source Quench



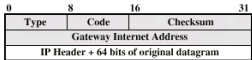
(e) Timestamp



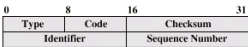
(b) Parameter Problem



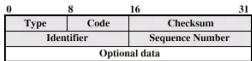
(f) Timestamp Reply



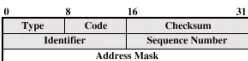
(c) Redirect



(g) Address Mask Request



(d) Echo, Echo Reply



(h) Address Mask Reply

Type	Code	Description
0	0	echo reply (ping)
3	0	dest. network unreachable
3	1	dest host unreachable
3	2	dest protocol unreachable
3	3	dest port unreachable
3	6	dest network unknown
3	7	dest host unknown
4	0	source quench (congestion control - not used)
5		Redirect
8	0	echo request (ping)
9	0	route advertisement
10	0	router discovery
11	0	TTL expired
12	0	bad IP header

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Ping – Echo Request, Echo Reply

```

machine:$ ping www.up.pt
PING www.up.pt.cdn.cloudflare.net (104.18.7.105) 56(84) bytes of data.
64 bytes from 104.18.7.105 (104.18.7.105): icmp_seq=1 ttl=56 time=8.47 ms
64 bytes from 104.18.7.105 (104.18.7.105): icmp_seq=2 ttl=56 time=7.06 ms
64 bytes from 104.18.7.105 (104.18.7.105): icmp_seq=3 ttl=56 time=7.76 ms
64 bytes from 104.18.7.105 (104.18.7.105): icmp_seq=4 ttl=56 time=7.08 ms
64 bytes from 104.18.7.105 (104.18.7.105): icmp_seq=5 ttl=56 time=7.18 ms
64 bytes from 104.18.7.105 (104.18.7.105): icmp_seq=6 ttl=56 time=6.99 ms
64 bytes from 104.18.7.105 (104.18.7.105): icmp_seq=7 ttl=56 time=7.64 ms
64 bytes from 104.18.7.105 (104.18.7.105): icmp_seq=8 ttl=56 time=6.95 ms
64 bytes from 104.18.7.105 (104.18.7.105): icmp_seq=9 ttl=56 time=7.32 ms
64 bytes from 104.18.7.105 (104.18.7.105): icmp_seq=10 ttl=56 time=6.75 ms
64 bytes from 104.18.7.105 (104.18.7.105): icmp_seq=11 ttl=56 time=6.86 ms
^C64 bytes from 104.18.7.105: icmp_seq=12 ttl=56 time=6.50 ms

--- www.up.pt.cdn.cloudflare.net ping statistics ---
12 packets transmitted, 12 received, 0% packet loss, time 11294ms
rtt min/avg/max/mdev = 6.504/7.214/8.471/0.506 ms

```

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Traceroute and ICMP

- Source sends series of UDP segments to destination
 - first segment has TTL=1
 - second segment has TTL=2, ...
 - unlikely port number
- When n^{th} datagram arrives to n^{th} router
 - router discards datagram
 - sends to source: ICMP TTL expired
 - message includes: router name & IP address
- When ICMP message arrives, source calculates RTT
- Traceroute does this 3 times for each TTL
- Stop criterion: UDP segment eventually arrives at destination host
 - Destination returns ICMP Dest port unreachable packet
 - source stops

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Traceroute and ICMP – example

```
machineFEUP$ traceroute tom.fe.up.pt
traceroute to tom.fe.up.pt (10.227.240.138), 30 hops max, 60 byte packets
 1 not.mshome.net (172.21.0.1)  0.337 ms  0.292 ms  0.270 ms
 2 172.29.0.1 (172.29.0.1)  12.832 ms  12.814 ms  12.636 ms
 3 * * *
 4 pinguim.fe.up.pt (10.227.240.138)  13.405 ms  12.966 ms  12.956 ms
```

Host did not respond with ICMP error

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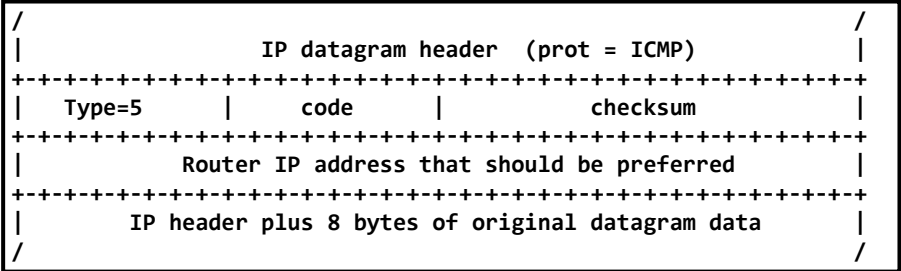
ICMP Redirect

- General routing principle of the TCP/IP architecture
 - routers have extensive knowledge of routes
 - hosts have minimal routing information → learn routes also from ICMP redirects
- ICMP redirect message
 - Sent by router R1 to source host A
 - when R1 receives a packet from A with destination = B, and R1
 - finds that the next hop is R2 and
 - A is on-link with R2
 - R1 sends ICMP redirect to A saying next hop for destination B is R2
 - A updates its forwarding table with a host route

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ICMP Redirect format



ICMP Redirect Example – Routing table in host A

```
ha$ netstat -nr
Kernel IP routing table
```

Destination	Gateway	Flags	Genmask	...	Iface
127.0.0.1	127.0.0.1	UH	255.255.255.255		lo0
193.154.156.0	193.154.156.24	U	255.255.255.0		eth0
0.0.0.0	193.154.156.1	UG	0.0.0.0		eth0

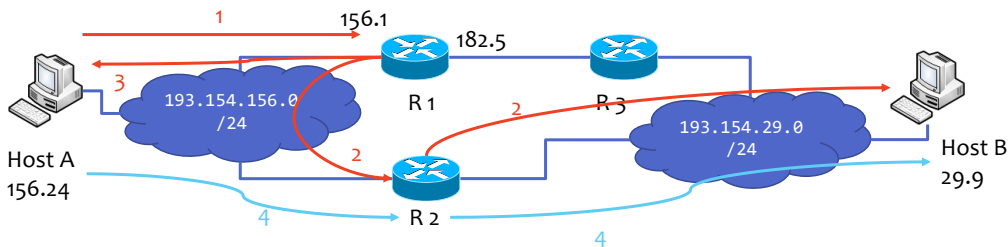
Other commands to check routing table in hosts:

- ♦ ip route
- ♦ route
- ♦ (deprecated)

Flags:

- U - route Up
- G - route to a Gateway (next hop router)
- H - route to a Host

ICMP Redirect Example



- Host A to send packet to B
 - 193.154.156.24 → 193.154.29.9
- Mesg 3:
 - ICMP of type redir
 - Src: 193.154.156.1 → Dst: 193.154.156.24
- All messages (except 3) are
 - 193.154.156.24 → 193.154.29.9

ICMP Redirect Example – Routing table in host A after mesg 3

ha\$ netstat -nr

Kernel IP routing table

Destination	Gateway	Flags	Genmask	...	Iface
127.0.0.1	127.0.0.1	UH	255.255.255.255		lo0
193.154.29.9	193.154.156.100	UGH	255.255.255.255		eth0
193.154.156.0	193.154.156.24	U	255.255.255.0		eth0
0.0.0.0	193.154.156.1	UG	0.0.0.0		eth0

Flags:
U - route Up
G - route to a Gateway (next hop router)
H - route to a Host

IPv6

STD 86/RFC 8200 Internet Protocol, Version 6 (IPv6) Specification

(Network layer)

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The Need of a New IP

- IPv4
 - Small addressing space (32 bits)
 - Non-continuous usage
 - Some solutions used to overcome these problems
 - private networks (NAT), classless networks (CDIR)
- IETF developed new IP version: IPv6
 - Same principles of IPv4
 - Many improvements
 - Header re-defined
 - First [RFC 1883, December 1995](#)

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IPv6 – Improvements

- 128 bit addresses (16 octets, 8 shorts). No classes
- Better QoS support (native flow label)
- Native security functions (peer authentication, data encryption)
- Autoconfiguration (Plug-n-play)
- Routing
- Multicast

Addresses

- Represented in hexadecimal
 - Ex. 1528:8653:294c:0000:0000:90af:0900:7654

- Zeros may be aggregated by ::
 - Ex.: 1528:8653:294c::90af:8900:7654
 - Only one :: in an address

- Masks are the same as for IPv4 CIDR
 - Loopback address ::1 /128

- Combining IPv6 and IPv4 addresses
 - ::ffff:5.6.7.8 (IPv4mapped address)
 - ::5.6.7.8 (IPv4compatible address)
 - 2002:5.6.7.8::1 (6to4 address)

128 bits →

665,570,793,348,866,943,898,599 addresses
per m² on Earth



See [RFC5952](#) for text representation

Adresses – Link-Local, Global Unicast, Anycast, Multicast

- Link-Local
 - Used for communication between hosts in the same LAN/link
 - Address built from MAC address
 - Routers do not forward packets having Link-Local destination addresses
- Global Unicast
 - Global addresses
 - Address: network prefix + computer identifier
 - Structured prefixes
 - Network aggregation; less entries in the router forwarding tables
- Anycast
 - Group address; packet is received by any (only one) member of the group
- Multicast
 - Group address; packet received by all the members of the group

Address Formats

n bits	m bits	128-n-m bits
001 global rout prefix	subnet ID	interface ID

Global Unicast Address
(2000::/3)

10 bits	54 bits	64 bits
1111111010	0	interface ID

Link-Local Unicast address
(fe80::/10)

n bits	128-n bits
subnet prefix	00000000000000

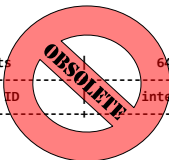
Anycast address

8	4	4	112 bits
11111111	flgs	scop	group ID

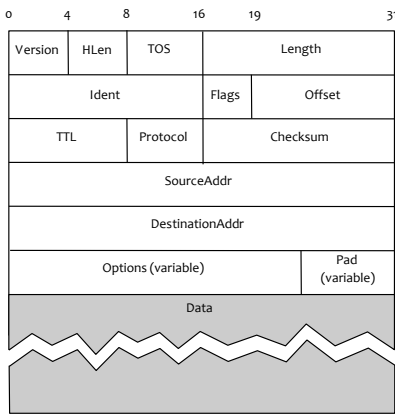
Multicast address
Scope – link, site, global, ...
(ff::/8)

10 bits	54 bits	64 bits
1111111011	subnet ID	interface ID

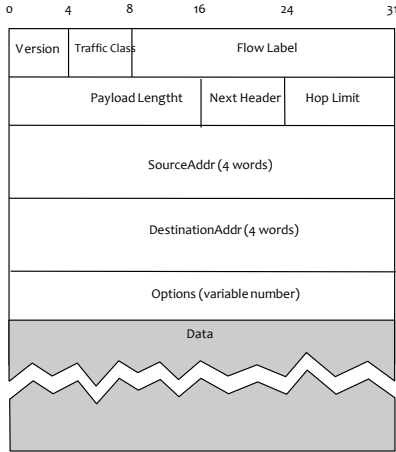
Site-Local Unicast address
(fec0::/10) (not used)



Headers IPv4 and IPv6



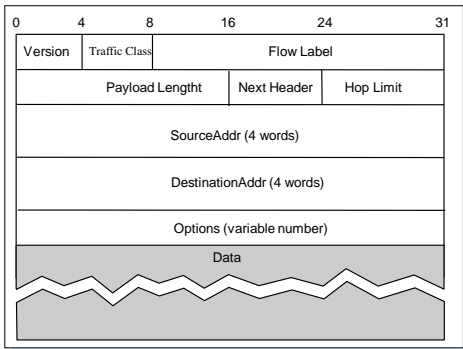
IPv4



IPv6

IPv6 Header

- Flow label → identifies packet flow
 - QoS, resource reservation
 - Packets receive same service
- Payload length
 - Header not included
 - Limited to $2^{16} - 1 = 65\,535$ bytes, but there's an option for JumboDatagrams $\leq 2^{32}$ bytes
- Next header
 - Identifies next header/extension
- Hop limit = TTL (v4)
- Options → included as extension headers
- IPv4 checksum removed, as lower layers are responsible for verification.



Extension Headers

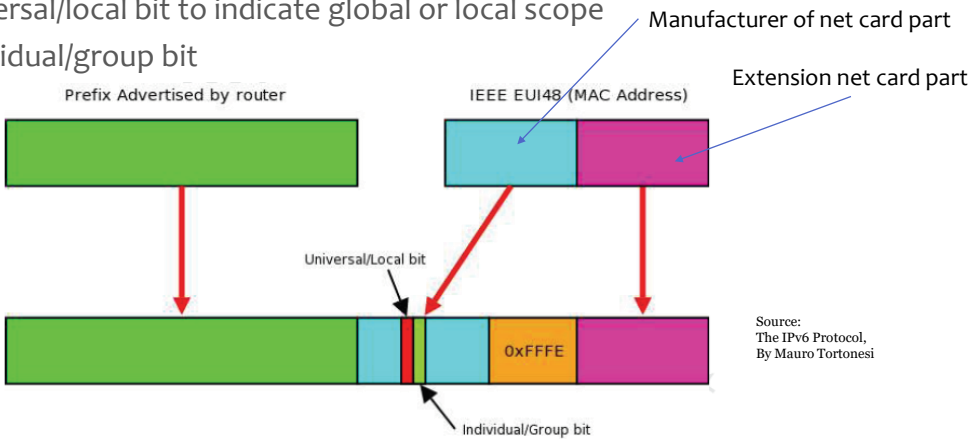
- Next Header – type of next header
 - Hop-by-Hop Options
 - the only header that is examined by intermediate nodes
 - Destination Options Header
 - Information for the destination node
 - Routing Header
 - List of nodes to be visited by the packet
 - Fragment Header
 - Authentication Header
 - Encrypted Security Payload Header
 - Transport layer headers

Example



Auto Configuration Address

- IEEE-EUI64
 - universal/local bit to indicate global or local scope
 - individual/group bit



Source:
The IPv6 Protocol,
By Mauro Tortonesi

Lookup vendors by MAC address

Protocol Neighbour Discovery (ND)

- [RFC4861](#)
- IPv6 node uses ND for
 - Find other nodes in the same link /LAN
 - Find a node MAC address
ND substitutes ARP
 - Find router(s) in its network
 - Maintaining information about neighbour nodes
- ND similar to the IPv4 functions
 - ARP IPv4
 - ICMP Router Discovery
 - ICMP Redirect

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ND Messages

- ICMP messages (over IP); using Link Local addresses
- Neighbour Solicitation
 - Sent by a host to obtain MAC address of a neighbour / to verify its presence
- Neighbour Advertisement
 - Answer to the request
- Router Advertisement
 - Information about the network prefix; periodic or under request
 - Sent by router to IP address Link Local multicast
- Router Solicitation: host solicits from router a Router Advertisement message
- Redirect: Used by a router to inform a host about the best route to a destination

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Summary

- Network layer overview
- Virtual Circuits
- Datagram Networks
- Forwarding
- ARP
- DHCP
- NAT
- ICMP
- Internet Protocol version 6

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Homework

1. Review slides
2. Read from Kurose & Ross
 - Chapter 4 – The Network Layer
(this set of slides follows mainly Kurose & Ross)
3. Or, from Tanenbaum,
 - Chapter 5 – The Network Layer
4. Answer questions at Moodle

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End of Network Layer