Embedded Systems - Real-Time Scheduling

part 1

Timing constraints: origin and characterization

Basic concepts of *real-time*Real-Time Systems requirements

Real-Time Computing

The results of the computations must be

- Logicaly correct
- Produced in time

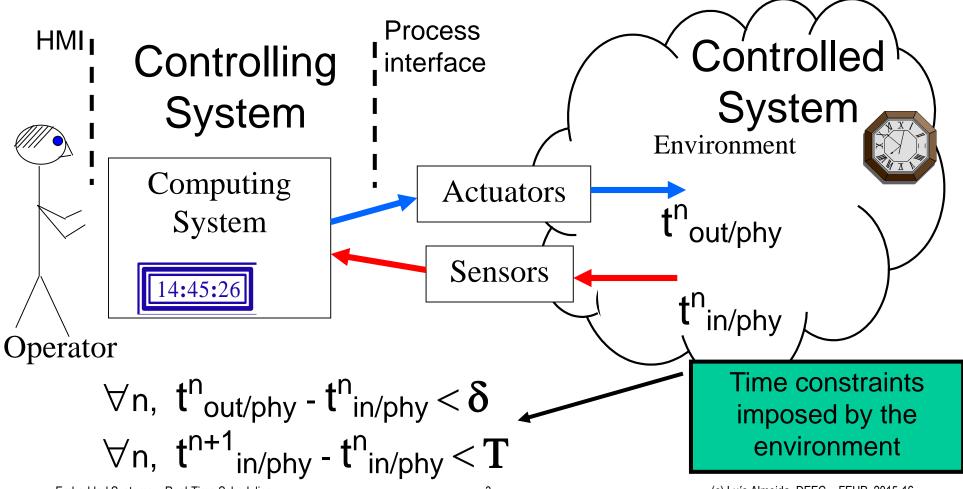
(Stankovic, 1988)

Timeliness



Logical correctness

Real-Time System



Purpose of studying Real-Time Systems

- The main objective of a Real-Time Systems (RTS) discipline is to develop appropriate techniques for the
 - Design / synthesis,
 - Analysis and
 - Verification

of such systems that allow obtaining **guarantees** that their **temporal behavior is adequate to the dynamics** of the process they must interact with.

Special features of Real-Time Systems

Given their **dependence with time**, some of the most important aspects that characterize a real-time system are:

- Execution time of its computations
- Response time to events
- Regularity of generating periodic events

Special features of Real-Time Systems

These aspects are influenced by several internal parameters:

- Execution time
 - Code structure (language, conditionals, cycles)
 - Operating system or kernel (system calls)
 - DMA, caches, instructions pipeline
- Response time and regularity
 - Multi-tasking
 - Access to shared resources (buses, disks, comm ports,...)
 - HW Interrupts

Requirements of Real-Time Systems

Typical **requirements** imposed on real-time systems:

- Functional
- Temporal
- Dependability

• ...

Non-Functional

Functional requirements

Example

Acquiring environment data

 Sample process variables (real-time entities) either continuous or discrete

Direct Digital Control (DDC)

Realize low-level feedback control (sensor/actuators loops)

Human-Machine Interface (HMI)

 Provide information on the system state, carry out logging, perform configuration

8

Functional requirements

Acquiring environment data

The **process real-time entities**, upon acquisition, are internally accessible to the computing system through local images (internal variables).

Each **local image** of a real-time entitiy has a **limited temporal validity** due to the process dynamics.

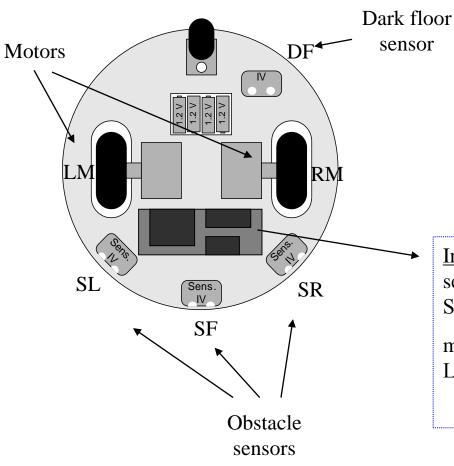
The **set of local images** of the real-time entities forms the **real-time database**.

The **real-time database** needs to be **updated** whenever there is a **change** in the state/value of a real-time entity.

Functional requirements

Example:

A simple mobile robot



RT entities:

sensors:

SL, SF, SR, DF

Motors speeds:

LM, RM

<u>Internal images:</u>

sensors:

SL', SF', SR', DF'

motors speeds: LM', RM'

RT database

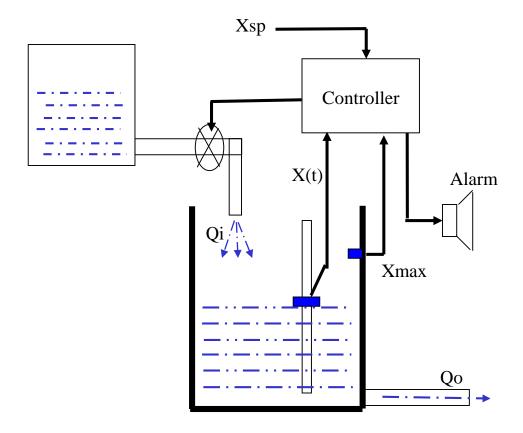
Arise from the dynamics of the process to be controlled/monitored

Sets constraints:

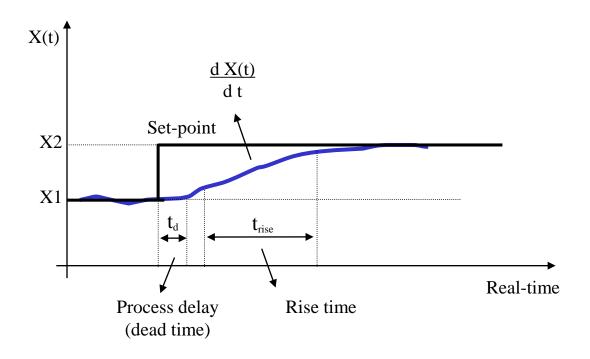
- To the observation delays of the system state
- To the computing delays of the control/actuation values
- To the delay variations of the previous (jitter)

In some cases, such constraints must be met in all instances (including the worst-case) and not only on average terms.

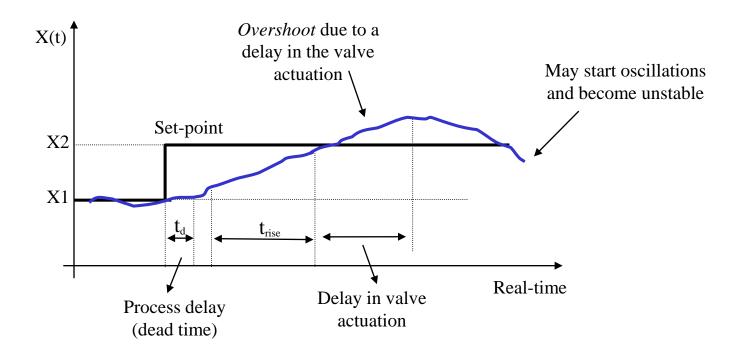
Control of the level of a liquid in a tank



Control of the level of a liquid in a tank



Delayed actuation –control degradation



(Control systems)

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Sampling period − T<sub>s</sub> (~< 1/10 t<sub>rise</sub> - quasi-continuous control)
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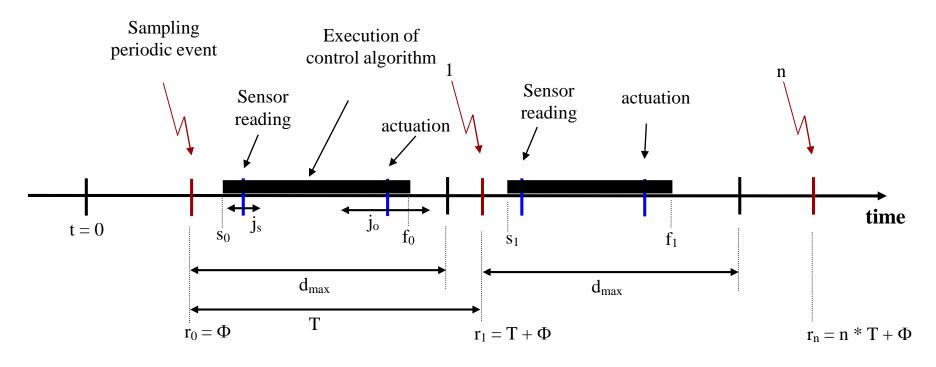
Maximum delay in the valve response - d_{max} (< T_s) (control delay - can be compensated in many cases with a model)

Variations in the level reading (jitter) – $j_{s,max}$ ($<< d_{max}$)

Variations in the valve actuation delay (jitter) - $j_{o,max}$ ($<< d_{max}$)

(harder to compensate – degradation of the quality of the control)

Maximum delay in alarm signaling - dal,max



 r_n – activation

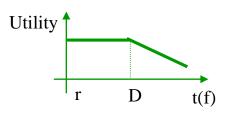
 s_n – start of execution

 f_n – end of execution

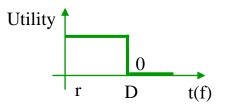
Classification of real-time constraints:

(according to the utility of the result to the application)

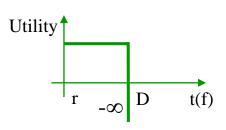
 Soft – Time constraint according to which the associated value keeps some utility to the application even after a deadline D, despite a possible degradation of the quality of service (videophone)



 Firm — Time constraint according to which the associated value loses any utility to the application after a deadline D (pre-paid services validation)



 Hard – Time constraint that, when not met, may generate a catastrophic failure (critical control)



Classification of Real-Time Systems:

(according to the type of time constraints they have to cope with)

- Soft Real-Time There are firm or soft time constraints, only (e.g., simulators, multimedia systems)
- Hard Real-Time There is at least one hard time constraint.
 These are safety-critical systems (e.g. transportation systems, weapons control, process control of dangerous materials...)

Dependability Requirements

Real-Time Systems are frequently used in **critical applications** either in terms of **safety** and **economy** (e.g., power plants, traffic control (air, rail, road), process industry...).

This leads to a specific requirement for:

High reliability – In hard real-time systems one typically finds ultra high reliability requirements (λ <10⁻⁹ failures/hour – cannot be verified experimentally! 1 failure in 114155 years!).

Dependability Requirements

<u>Important aspects concerning critical safety:</u>

- Stable interfaces between the critical and non-critical subsystems in order to prevent error propagation.
- Well defined worst-case scenarios. The system must possess the adequate resources to withstand a worst-case operational scenario without need for probabilistic arguments, i.e., it should provide service guarantees even in such cases.
- Architecture composed by autonomous subsystems, whose properties can be verified independently from the other subsystems (composability).

Wrapping up

- Notion of real-time and real-time system
- Purpose of studying RTS obtain guarantees of adequate temporal behavior
- Aspects to consider: execution time, response time, regularity of generating/handling periodic events
- •RTS requirements: functional, temporal and dependability
- Notion of real-time database
- Soft, firm and hard time constraints
- Hard real time vs soft real time
- Importance of the worst-case scenario

Embedded Systems - Real-Time Scheduling

part 2

Computing models

Models of tasks with explicit time constraints

Logic and temporal control (by events -ET and by time -TT)

Computing models

Transformational model

A program starts and finishes, transforming input data in output results.

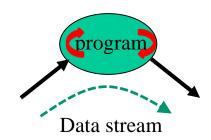
Reactive model

program data A program executes indefinitely a sequence of interactions, for example operating over a continuous input data stream.

Input data

Real-time model

 Reactive model in which the output stream has to be synchronized with the input data stream, thus imposing time constraints on the program execution.



Output

Determinism versus Predictability

(within reactive systems)

Determinism

 Feeding a program with the same sequence of data inputs originates the same sequence of data outputs.
 (logical property)

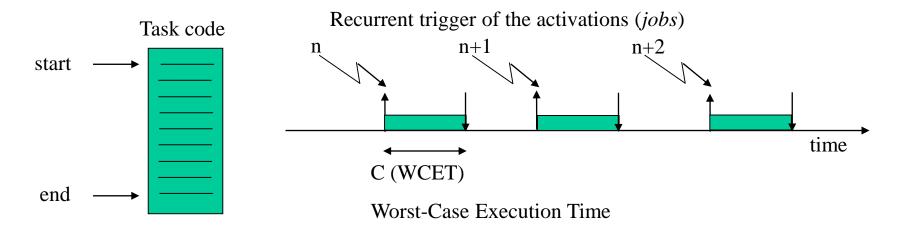
Predictability

 feeding a program with the same input sequence will generate the same output sequence with a known delay or within a know time window.

(temporal determinism – logical and temporal property)

Definition of task (recurrent)

Infinite sequence of activations (instances or jobs), each executing a set of operations (a **given function**).

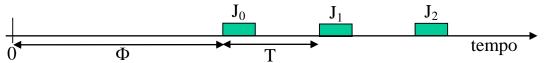


Concerning recurrency of arrivals, a task can be

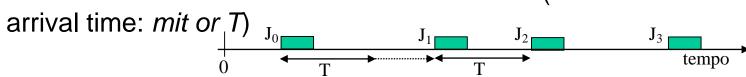
periodic

Instance *n* activated at $a_n = n^*T + \Phi$

sporadic

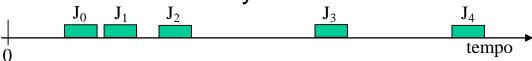


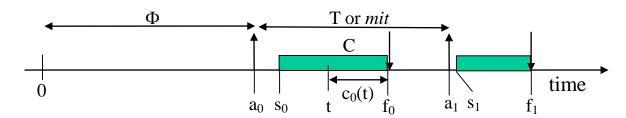
Minimal time between consecutive activations (minimum inter-



aperiodic

Can only be characterized stochastically





Characterization of tasks

C – Worst-case execution time (WCET)

T – period (periodic)

 Φ -relative phase/offset = instant of 1st activation (periodic)

mit – minimum inter-arrival time (sporadic)

a_n – activation instant of the nth instance

s_n – start of execution of the nth instance

f_n – finishing time of the nth instance

c_n(t) – maximum remaining execution time at instant t of the nth instance

Tasks requirements:

- Temporal time constraints on the finishing instants or for generating certain output events.
- Precedence establish a certain order of execution among tasks.
- Resource usage need for using shared resources (e.g. communication ports, buffers in shared memory, global variables, system peripherals...). May imply the use of atomic operations (which execution sequence cannot be interrupted)

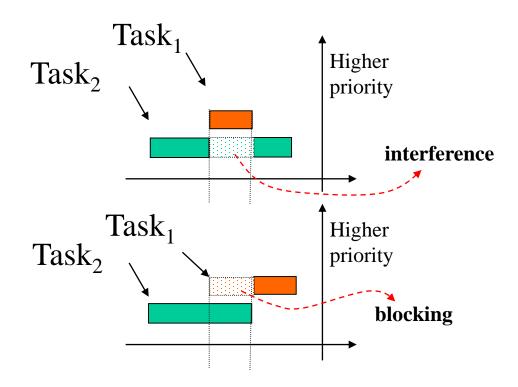
Preemption

- When a task can be temporarily suspended for the execution of another one with higher priority, it is preemptable.
- When a system makes use of the preemptability of tasks is it called preemptive.
- A set of tasks is said to admit full preemption when all its tasks can be preempted in any point of their execution (independent tasks)

Note: the access to **shared resources** (thus tasks with **dependencies**) may impose restrictions on the level of preemptability of a task.

Execution with preemption

Execution without preemption



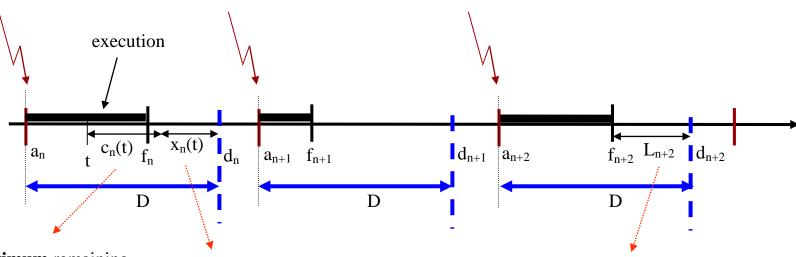
The temporal requirements can be of diverse types:

- Deadline Upper bound to the finish instant of a task.
- Window Lower and upper bounds to the finish instant of a task.
- Synchronism Upper bound to the interval between two output events (there are types, too).
- Distance Upper bound to the interval between the finishing time of two consecutive jobs

The *deadline* type is the most common!

Activation event Deadline

 \forall n, f_n - a_n < D (relative) \forall n, f_n < d_n (absolute)



 $c_n(t) = maximum$ remaining execution time at instant t

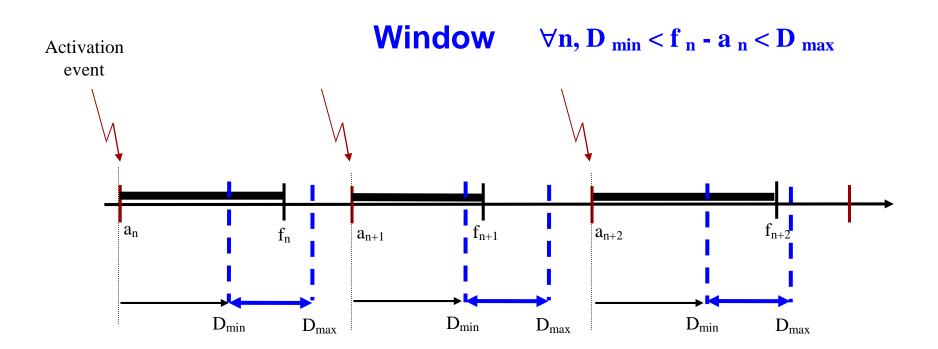
 $\mathbf{x}_{\mathbf{n}}(\mathbf{t}) = \mathbf{minimum}$ task slack

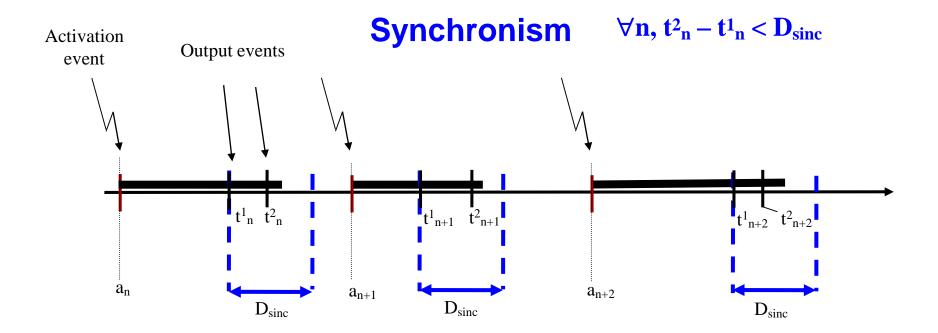
$$x_n(t) = d_n - t - c_n(t)$$

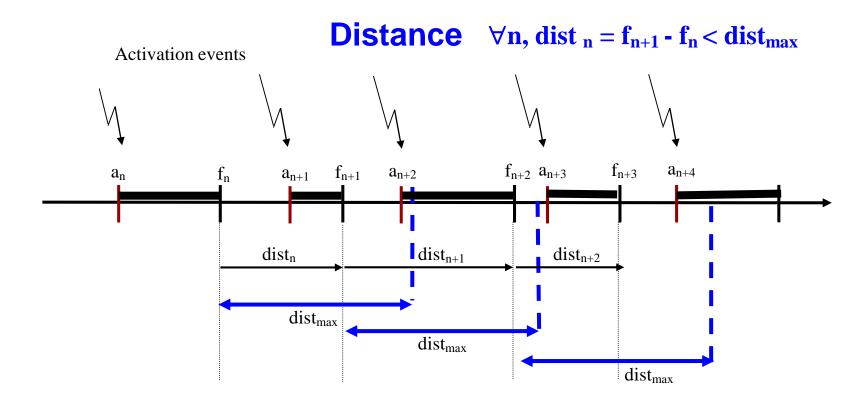
 $L_n = f_n - d_n = delay of nth$ activation

 $L_n \le 0 \rightarrow \text{timely termination}$

 $L_n > 0 \rightarrow late termination$







Examples of tasks characterization:

• Periodic:
$$\tau_i = \tau_i \ (C_i, \Phi_i, T_i, D_i)$$

$$\tau_1 = \tau_1 \ (2,5,10,10) \qquad \tau_2 = \tau_2 \ (3,10,20,20)$$

• Sporadic: Similar to periodic but with mit_i in the place of T_i and Φ_i is not typically used (could mean a minimal interval until the first activation).

$$\tau_i = \tau_i (C_i, mit_i, D_i)$$

$$\tau_1 = \tau_1 (2,5,5) \qquad \tau_2 = \tau_2 (3,10,7)$$

Logic and temporal control

Logic control

Program control flow, i.e., effective sequence of operations to be executed – **fundamental to determine C (WCET)**

Temporal control

Control of the **execution instants** of the program operations (e.g., activations, verification of time constraints,...)

Triggering tasks (or functions)

By time (time-triggered)

The execution of tasks (or functions) is triggered by a control signal based on the progression of time (e.g., a periodic interrupt).

By events (event-triggered)

The execution of tasks (or functions) is triggered by an asynchronous control signal generated upon a change in the system state (e.g., by an external interrupt).

Activating tasks by Time

time-triggered (TT) systems

- Typically used in automatic control (sampling of continuous variables).
- There is a common time base (allows establishing offset relationships or relative phasing)
- CPU utilization is constant even when there are no variations in the system state.
 - Well defined worst-case situation

Activating tasks by Events

event-triggered (ET) systems

- Typically used in monitoring sporadic conditions in the system state (e.g., alarms or asynchronous service requests).
- CPU utilization (and not only CPU!) is variable according to the frequency of event occurrence.
 - Poorly defined worst-case situation
 either probabilistic arguments are used
 or the maximum event rate must be bounded

Example: Consider the following task sets and determine the worst-case response time of each task

• TT
$$\{\tau_i = \tau_i \ (C_i=1, \Phi_i=i, T_i=5, D_i=T_i \ i=1..5)\}$$

• ET
$$\{\tau_i = \tau_i \ (C_i=1, (\Phi_i=0), mit_i=5, D_i=mit_i \ i=1..5)\}$$

Moreover, determine the average and maximum **CPU utilization**** utilization = sum_{i=1..N}(execution time/activation period) **
for both cases, considering that on average the ET tasks are activated once every 100 time units.

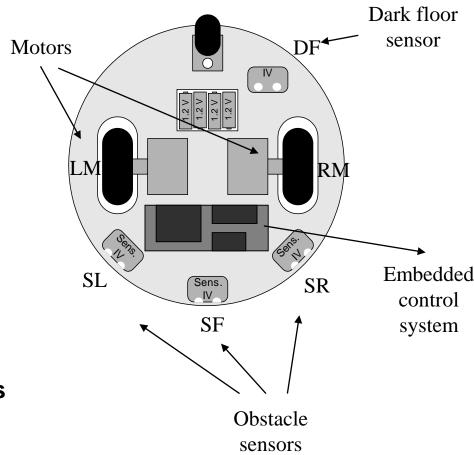
Modeling a simple mobile robot

Functional requirements:

- Differential motion control
- If dark floor, then stop and finish
- Else if wall, then follow wall
- Else if obstacle, then go around or move back
- Else move straight

Non-functional aspects:

- Max speed 50 cm/s; time to stop 0.3s
- Obst. sensors range 5-25 cm
- Dark floor sensor within 1s



A procedure-based SW modeling approach

Tasks and dependencies

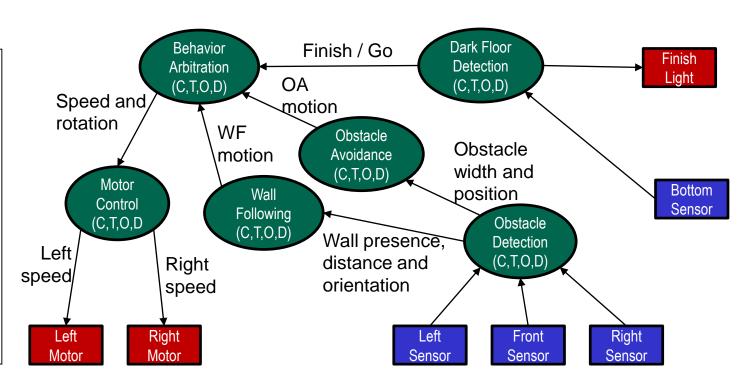
RT entities: sensors, motors, light

Inter-task communication:

Global variables (good to represent them explicitly in the model as shared data items)

Periodic (TT) approach:

- C execution time
- T period
- O offset
- D deadline

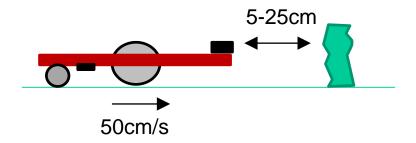


Outputs / Actuators

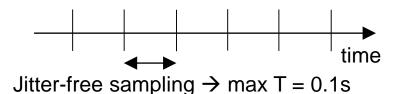
inputs / Sensors

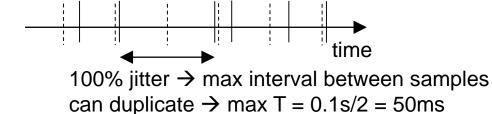
Time-constraints

Obstacle detection



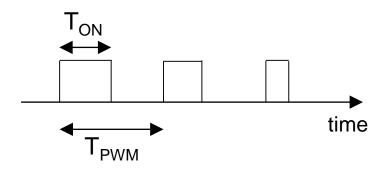
20cm range, @ 50cm/s → 0.4s window to detect the obstacle 0.3s for stopping → actual time for detecting the obstacle 0.1s Obstacle sampling must guarantee no more than 0.1s between samples





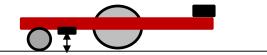
Time-constraints

Motor control (PWM)



Max $T_{PWM} \rightarrow 1/10$ trise (~1s) \rightarrow max T_{PWM} =100ms $T_{ON} = f(v), v \rightarrow$ desired speed

Dark floor detection



Reliable detection → several consecutive detections + hysteresis (say 5 samples)
1s / 5samples → max T = 200ms

Behavior tasks

Depends on desired reactivity

Normally equal to Obstacle detection

Wrapping up

- Computing models (real-time model)
- •Real-time tasks: periodic, sporadic and aperiodic
- •Temporal constraints: **deadline**, window, synchronism and distance
- Logic control and temporal control
- Event-triggered versus time-triggered tasks
- •Example model and time constraints from a simple case study