

# Concepts of PDAs and Conversion from PDAs to CFGs

\* This form will record your name, please fill your name.

1. A PDA accepts a string when:  
(1 Point)

- ☐ after consuming all the symbols of the string, the PDA is in a final state or its stack is empty
- ☐ after consuming all the symbols of the string, the PDA is in a final state and the PDA accepts by final state, or the stack of the PDA is empty and the PDA accepts by empty stack

2. A PDA is always a non-deterministic automaton  
(1 Point)

- ☐ TRUE
- ☐ FALSE

3. Any non-deterministic PDA can be converted to a deterministic PDA  
(1 Point)

☐ TRUE

☐ FALSE

4. It is possible to automatically convert any CFG to a PDA  
(1 Point)

☐ TRUE

☐ FALSE

5. The method explained in the Lecture to convert a CFG to a PDA implies:  
(1 Point)

☐ a PDA with one state and accepting by empty stack

☐ a PDA with one state and accepting by final state

6. If a CFG is ambiguous, the automatic conversion of the CFG to a PDA always results in a non-deterministic PDA:  
(1 Point)

☐ TRUE

☐ FALSE

