

Week 1 Quiz

What is the complexity of heap sort?

- a. $O(n)$
- b. $O(n^2)$
- c. $O(n \log n)$
- d. $O(n^2 \log n)$

Week 1 Quiz

What is the complexity of heap sort?

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- c. $O(n \log n)$
- d. $O(n^2 \log n)$

Heap can be built in $O(n)$ and then $O(n \log n)$ to remove maximal elements from the heap and preserve the heap structure n times

Week 1 Quiz

If you perform Radix sort on a string of length m using the factor b of m .

Which of the following statements is true (there is only one correct statement)?

- a. The number of iterations is m/b and number of buckets is $2b$
- b. The number of iterations is m/b and number of buckets is b^2
- c. The number of iterations is b/m and number of buckets is $2b$
- d. The number of iterations is b/m and number of buckets is b^2

Week 1 Quiz

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Which of the following statements is true (there is only one correct statement)?

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- b. The number of iterations is m/b and number of buckets is 2^b
- c. The number of iterations is b/m and number of buckets is $2b$
- d. The number of iterations is b/m and number of buckets is b^2

Split string of length m into m/b blocks of size b
and need a bucket for each bit string of length b

Week 1 Quiz

As we have seen, every comparison based sorting algorithm must use at least $O(n \log n)$ comparisons.

This is because the decision tree for the algorithm must have at least:

- a. n^2 leaf nodes
- b. $2n$ branch nodes
- c. $n!$ leaf nodes
- d. $n!$ branch nodes

Week 1 Quiz

As we have seen, every comparison based sorting algorithm must use at least $O(n \log n)$ comparisons.

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Leaf nodes are the different possible orderings and $n!$ possibilities

Week 2 Quiz

What is the complexity of breath first search when using an adjacency list representation for an undirected graph with n vertices and m edges?

- a. $O(n+m)$
- b. $O(\log n + \log m)$
- c. $O(n^2)$
- d. $O(nm)$

Week 2 Quiz

What is the complexity of breath first search when using an adjacency list representation for an undirected graph with **n** vertices and **m** edges?

- a. **$O(n+m)$**
- b. $O(\log n + \log m)$
- c. $O(n^2)$
- d. $O(nm)$

Go through the adjacency list of each vertex once and each edge appears in each edge twice.

Week 2 Quiz

If all edges have the same weight in an undirected graph, which algorithm will find the shortest path between two nodes more efficiently?

- a. Dijkstra
- b. Bellman–Ford
- c. Depth–First Search
- d. Breadth–First Search

Week 2 Quiz

If all edges have the same weight in an undirected graph, which algorithm will find the shortest path between two nodes more efficiently?

- a. Dijkstra
- b. Bellman–Ford
- c. Depth–First Search
- d. Breadth–First Search

a and **b** are less efficient and **c** cannot be used for finding shortest paths

Week 2 Quiz

In a weighted, directed graph if we apply Dijkstra's algorithm to find the shortest path between two nodes. If we add **1** to all the edge weights, does the shortest path always remain the same?

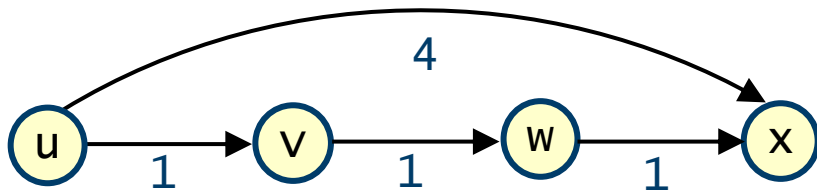
- a. **true**
- b. **false**

Week 2 Quiz

In a weighted, directed graph if we apply Dijkstra's algorithm to find the shortest path between two nodes. If we add **1** to all the edge weights, does the shortest path always remain the same?

- a. true
- b. **false**

A path of length **k** will increase its distance by **k** (**1** is added to the weight of each edge in the path), so longer paths are penalised more

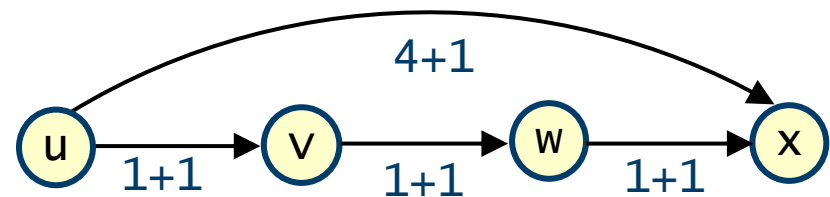
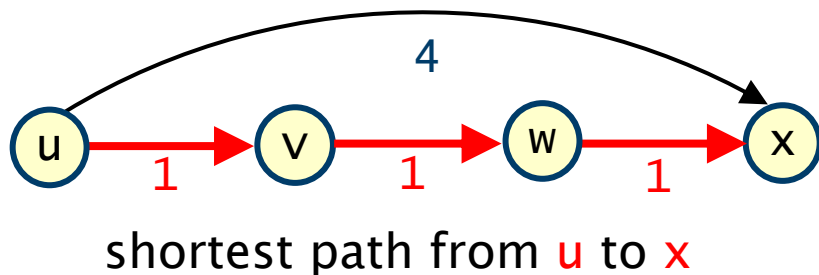


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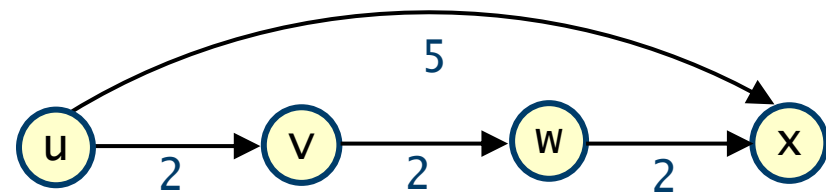
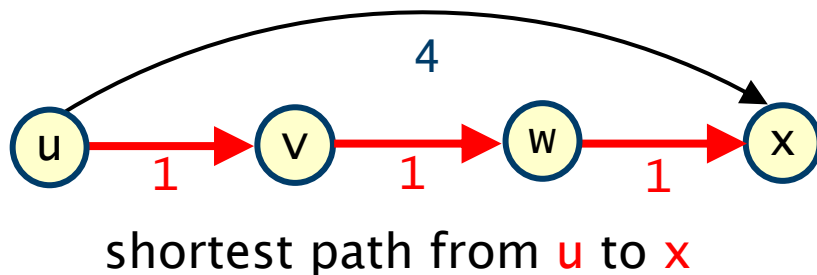


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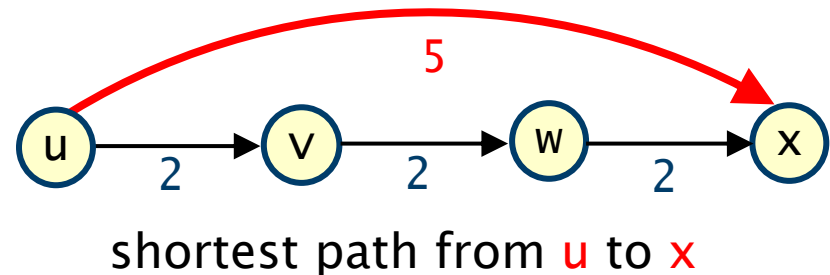
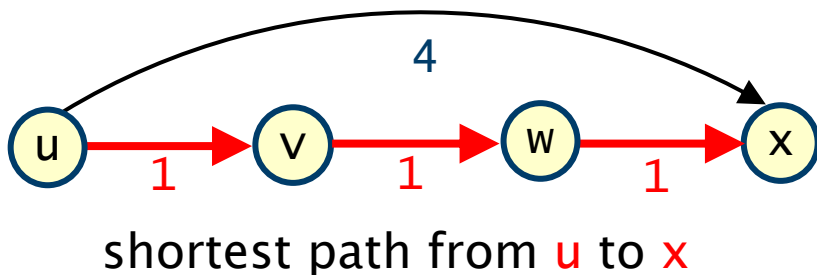


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Week 3 Quiz

In a weighted, undirected graph if we apply Prim Jarnak's algorithm to find the shortest path between two nodes. If we add **1** to all the edge weights, does the minimum weight spanning tree always remain the same?

- a. **true**
- b. **false**

Week 3 Quiz

In a weighted, undirected graph if we apply Prim Jarnak's algorithm to find the shortest path between two nodes. If we add **1** to all the edge weights, does the minimum weight spanning tree always remain the same?

- a. **true**
- b. false

Edges are chosen based on their weights and all increase by the same amount

- also all spanning trees have the same number of edges
 - if **n** vertices then **n-1** edges in a spanning tree

So weight of each spanning tree increases by the same amount

Week 3 Quiz

Let G be an undirected connected graph with distinct edge weights, e_{\max} be the edge with maximum weight and e_{\min} the edge with minimum weight.

Which of the following statements is false?

- a. Every minimum spanning tree of G must contain the edge with weight e_{\min}
- b. If the edge with weight e_{\max} is in a minimum spanning tree, then its removal must disconnect G
- c. No minimum spanning tree contains the edge with weight e_{\max}
- d. G has a unique minimum spanning tree

Week 3 Quiz

Let G be an undirected connected graph with distinct edge weights, e_{\max} be the edge with maximum weight and e_{\min} the edge with minimum weight.

Which of the following statements is **false**?

- a. Every minimum spanning tree of G must contain the edge with weight e_{\min}

true: this edge always added when one vertex of the edge is a tv and the other vertex a ntv

Week 3 Quiz

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- b. If the edge with weight e_{\max} is in a minimum spanning tree, then its removal must disconnect G

true: to include e_{\max} must be no other option

- can make local optimal choices (as the Prim Jarnik algorithm does)

Week 3 Quiz

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- c. No minimum spanning tree contains the edge with weight e_{\max}

false: see case b which to a case when this edge is included

Week 3 Quiz

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false: can be more than one spanning tree

– as we have seen in the lectures

Week 3 Quiz

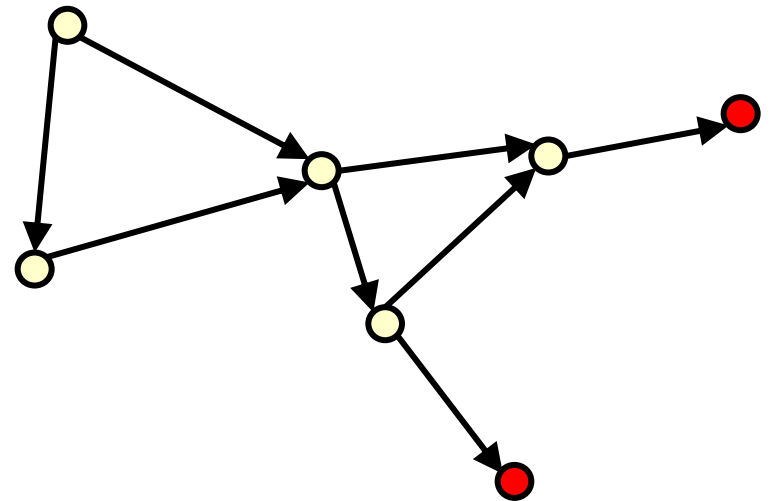
In a topological ordering, if v is a sink, it is true that $lab(u) < lab(v)$ for all non-sink vertices u

- a. true
- b. false

Week 3 Quiz

In a topological ordering, if **v** is a sink, it is true that **$Tab(u) < Tab(v)$** for all non-sink vertices **u**

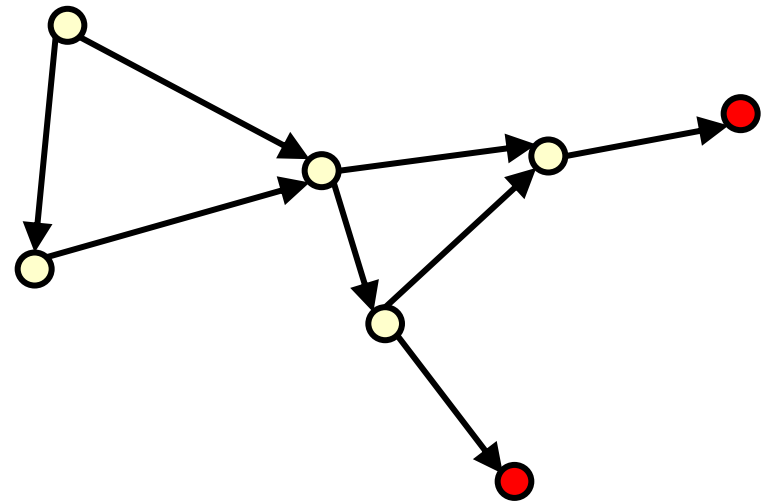
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Week 3 Quiz

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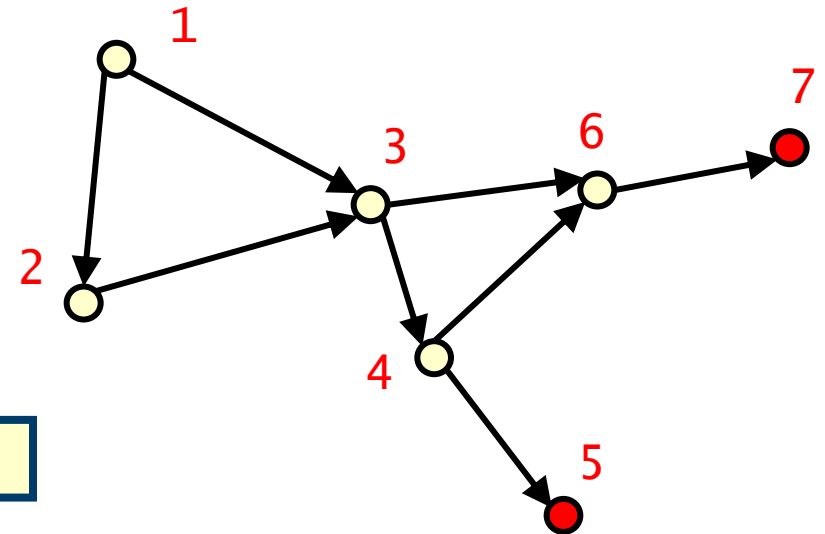
- a. true
- b. false



Week 3 Quiz

In a topological ordering, if **v** is a sink, it is true that **$lab(u) < lab(v)$** for all non-sink vertices **u**

- a. true
- b. false



would be true if there was only one sink

Week 4 Quiz

Assuming the character frequencies have already been found, what is the complexity of building a Huffman tree if there are m characters?

- a. $O(m)$
- b. $O(m \log m)$
- c. $O(m^2)$
- d. $O(n^3)$

Week 4 Quiz

Assuming the character frequencies have already been found, what is the complexity of building a Huffman tree if there are m characters?

- a. $O(m)$
- b. $O(m \log m)$
- c. $O(m^2)$
- d. $O(n^3)$

Add leaf nodes with weights equal to character frequencies $O(m)$
Find two parentless nodes with smallest weights and add parent with weight equal to the sum of weights

- m times we remove and add elements to a heap
- with each operation $O(\log m)$

Week 4 Quiz

What property of the dictionary holds throughout the running of the LZW algorithm?

- a. If string **s** is in the dictionary, then all prefixes of **s** are in the dictionary
- b. If string **s** is in the dictionary, then no prefixes of **s** are in the dictionary
- c. All strings of length **k** are in the dictionary when the codeword length is **k**
- d. No code word is a prefix of another code word

Week 4 Quiz

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- c. All strings of length **k** are in the dictionary when the codeword length is **k**
- d. No code word is a prefix of another code word

identify the longest string **s**, starting at position **i** of text **t**
that is represented in the dictionary **d**;

...
add string **s+c** to dictionary **d**, paired with next available codeword;

Week 4 Quiz

In the recurrence relation for finding the optimal string distance when transforming the string **x** into the string **y**. When the characters do not match the distance **$d(i, j)$** equals:

$$1 + \min\{ d(i, j-1), d(i-1, j), d(i-1, j-1) \}$$

What does the case **$d(i-1, j-1)$** correspond to?

- a. Deleting an element from x
- b. Inserting an element in x
- c. **Substituting an element in **x****
- d. Inserting an element in y

After substitution matched element of **x** with that of **y**

Week 5 Quiz

In the table is required by the KMP algorithm what is included in entry i when searching for the pattern/string s ?

- a. The shortest border of the substring $s[0..i-1]$
- b. The longest border of the substring $s[0..i-1]$
- c. The last occurrence of the character $s[i]$
- d. The first occurrence of the character $s[i]$

Week 5 Quiz

In the table is required by the KMP algorithm what is included in entry **i** when searching for the pattern/string **s**?

- a. The shortest border of the substring $s[0..i-1]$
- b. **The longest border of the substring $s[0..i-1]$**
- c. The last occurrence of the character $s[i]$
- d. The first occurrence of the character $s[i]$

Look at how the algorithm works

Week 5 Quiz

What is the length of the longest border of **abcabca**?

- a. 2
- b. 4
- c. 5

Week 5 Quiz

What is the length of the longest border of **abcabca**?

- a. 2 not a border: **ab**cabca abcab**ca**
- b. 4 longest border: **abcab**ca ab**cab**ca
- c. 5 not a border: **abcab**ca ab**cab**ca

Week 5 Quiz

In the table is required by the Boyer Moore algorithm what is included in entry i when searching for the pattern/string s in the text t ?

- a. The shortest border of the substring $t[0..i-1]$
- b. The longest border of the substring $t[0..i-1]$
- c. The last occurrence of the character $t[i]$
- d. The first occurrence of the character $t[i]$

Week 5 Quiz

In the table is required by the Boyer Moore algorithm what is included in entry i when searching for the pattern/string s ?

- a. The shortest border of the substring $t[0..i-1]$
- b. The longest border of the substring $t[0..i-1]$
- c. The last occurrence of the character $t[i]$
- d. The first occurrence of the character $t[i]$

Look back at the slides on how the algorithm works

Week 6 Quiz

Suppose problem **X** is in P, problem **Y** is in class NP, and there is a polynomial reduction from **X** to **Y**. Which of the following is true?

- a. Problem **Y** is in class P
- b. Problem **Y** is NP-complete
- c. Both of the above
- d. Neither of the above

Week 6 Quiz

Suppose problem **X** is in P, problem **Y** is in class NP, and there is a polynomial reduction from **X** to **Y**. Which of the following is true?

- a. Problem Y is in class P
- b. Problem Y is NP-complete
- c. Both of the above
- d. Neither of the above

All this shows is that **Y** is “at least as hard” as **X**, but might map problems of **X** to easy problems of **Y**

Week 6 Quiz

Suppose problem **X** is NP-complete, problem **Y** is in class NP, and there is a polynomial reduction from **X** to **Y**. Which of the following is true?

- a. Problem **Y** is in class P
- b. Problem **Y** is NP-complete
- c. Both of the above
- d. None of the above

Week 6 Quiz

Suppose problem **X** is NP-complete, problem **Y** is in class NP, and there is a polynomial reduction from **X** to **Y**. Which of the following is true?

- a. Problem Y is in class P
- b. **Problem Y is NP-complete**
- c. Both of the above
- d. None of the above
 - if **X** in NP-complete, then we can reduce any NP problem to **X** in polynomial time
 - by the hypothesis we can reduce **X** to **Y** in polynomial time
 - therefore we can reduce any NP problem to **Y** in polynomial time
 - first reduce it to **X** and then from **X** reduce to **Y**
 - and since **Y** is in NP, it follows **Y** is NP-complete

Week 6 Quiz

Suppose problem **X** is NP-complete, problem **Y** is in class **P**, and there is a polynomial reduction from **X** to **Y**. If **Z** is NP-complete, which of the following is true?

- a. Problem **Z** is in class **P**
- b. Problem **Z** is in the class **NP**
- c. Both of the above
- d. None of the above

Week 6 Quiz

Suppose problem **X** is NP-complete, problem **Y** is in class P, and there is a polynomial reduction from **X** to **Y**. If **Z** is NP-complete, which of the following is true?

- a. Problem **Z** is in class **P**
 - b. Problem **Z** is in the class **NP**
 - c. Both of the above
 - d. None of the above
-
- we can solve **X** in polynomial time by reducing to **Y** and solving **Y**
 - since **X** is NP-complete this means all NP-complete problems can be solved in polynomial time including **Z**
 - since **Z** is NP-complete, by definition it is in NP

Week 7 Quiz

Which of the following strings is a member of the language over $\{a, b\}$ defined by the regular expression $(aa|ba)^*(bb)^*?$

- a. aabbba
- b. aaaabb
- c. babbaa
- d. bbaa

Week 7 Quiz

Which of the following strings is a member of the language over $\{a, b\}$ defined by the regular expression $(aa|ba)^*(bb)^*?$

- a. aabbba
 - ends **ba** and these is only achievable when **aa** or **ba** precedes the **ba**
- b. aaaabb**
 - **aaaa** comes from **aa** twice followed by **bb** which comes from **bb** once
- c. babbaa
 - ends **aa** and these is only achievable when **aa** or **ba** precedes the **aa**
- d. bbbaa
 - ends **aa** and these is only achievable when **aa** or **ba** precedes the **aa**

Week 7 Quiz

Over the alphabet $\{a, b\}$, which of the following the regular expressions represents the language which consists of strings that start and end with different symbols.

- a. $a(a|b)^*b$
- b. $b(a|b)^*a$
- c. $(a(a|b)^*b) \mid (b(a|b)^*a)$
- d. All of the above

Week 7 Quiz

Over the alphabet $\{a, b\}$, which of the following the regular expressions represents the language which consists of all strings that start and end with different symbols.

- a. $a(a|b)^*b$
- b. $b(a|b)^*a$
- c. $(a(a|b)^*b) \mid (b(a|b)^*a)$
- d. All of the above

a (b) are strings that start with a (b) and end with b (a)

Week 7 Quiz

Over the alphabet $\{a, b\}$, which of the following the regular expressions represents the language which consists of strings that starts with **ab** and ends with **ba**.

- a. **aba*b*ba**
- b. **ab(ab)*ba**
- c. **ab (a|b)*ba**
- d. All of the above

Week 7 Quiz

Over the alphabet $\{a, b\}$, which of the following the regular expressions represents the language which consists of strings that starts with ab and ends with ba .

- a. aba^*b^*ba starts with ab , then ≥ 0 a s then ≥ 0 b s, ends with ba
- b. $ab(ab)^*ba$ starts with ab , then ≥ 0 ab s, ends with ba
- c. $ab(a|b)^*ba$ all strings that starts with ab and end with ba
- d. All of the above

question should have said
consists of all strings that

will mark both c and d as correct
(ignore moodle I will export the
individual marks for quiz
questions and fix)

Week 8 Quiz

Which of the following can be computed by a deterministic finite state automaton? (Select all that are computable).

- a. Strings which have more **a**s than **b**s
- b. Strings which have no single **a**s (i.e. any **a** is either preceded or followed by another **a**)
- c. Strings that start and end with the same character
- d. Strings that start with a certain number of **a**s end with the same number of **a**s

Week 8 Quiz

Which of the following can be computed by a deterministic finite state automaton? (Select all that are computable).

- a. Strings which have more as than bs
 - b. Strings which have no single **a**s (i.e. any **a** is either preceded or followed by another **a**)
 - c. Strings that start and end with the same character
 - d. Strings that start with a certain number of as end with the same number of as
-
- b and c can be expressed by the regular expressions $(b^*aaa^*b^*)^*$ and $(a(a|b)^*a) | (b(a|b)^*b)$
 - a and d both require counting the number of a certain character and keeping track of this, but we only have finitely many states

Week 8 Quiz

Does nondeterminism increase the languages that can be expressed by finite state automata?

- a. yes
- b. no

Week 8 Quiz

Does nondeterminism increase the languages that can be expressed by finite state automata?

- a. yes
- b. no**

For any NFA one can build (using the subset construction) a DFA that recognises the same language

- this can cause an exponential blow up in the state space

Week 8 Quiz

Does nondeterminism increase the languages that can be expressed by pushdown automata?

- a. yes
- b. no

Week 8 Quiz

Does nondeterminism increase the languages that can be expressed by pushdown automata?

- a. **yes**
- b. no

For example for the language of palindromes, we need to guess when we have reached the "middle" of the input