Algorithmics I

Tutorials – Sheet 2

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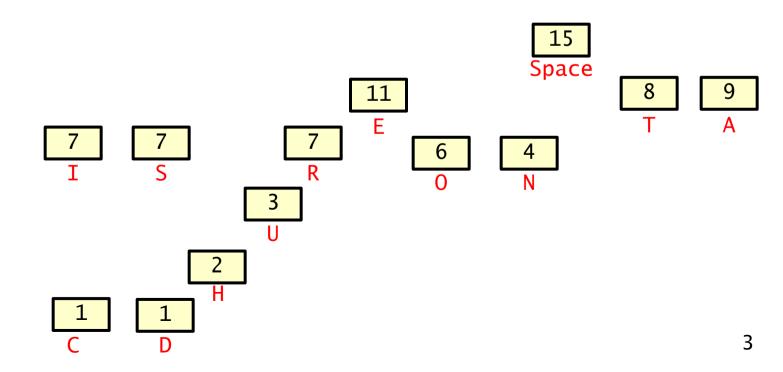
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Explain how the Huffman-tree-building algorithm can be extended so that it also determines the weighted path length of the tree that it constructs

Character frequencies:

First add leaves of Huffman tree

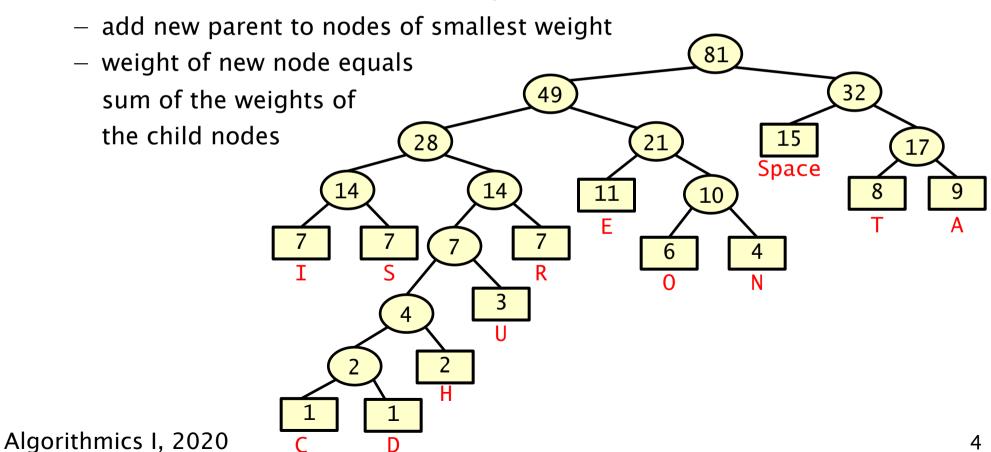
characters with their frequencies label the leaf nodes



Character frequencies:

Space	Е	Α	Т	Ι	S	R	0	N	U	Н	С	D
15	11	9	8	7	7	7	6	4	3	2	1	1

Next, while there is more than one parentless node



Weighted path length (WPL) of a tree T

 $-\Sigma$ (weight)×(distance from root) where sum is over all leaf nodes

Algorithm for calculating the weight:

- initialise a variable w to zero
- each time a new internal node is created add weight of this node to w
- upon termination return w

So why is w equal to the weighted path length?

- for each leaf node, its weight is added to w every time an ancestor of the leaf is created
- the number of ancestors of a leaf equals its distance from the root

Apply LZW algorithm to:

peter piper picked a peck of pickled pepper

Assume the dictionary initially contain the 128 characters of the basic ascii character set, represented by the code words 0,1,...,127

Next available code word is 128

Will leave this for you to do - the result will be included in the notes

LZW compression - Pseudo code

```
set current text position i to 0;
initialise codeword length k (say to 8);
initialise the dictionary d;
while (the text t is not exhausted) {
 identify the longest string s, starting at position i of text t
 that is represented in the dictionary d;
// there is such string, as all strings of length 1 are in d
output codeword for the string s; // using k bits
// move to the next position in t
i += s.length(); // move forward by the length of string just encoded
 c = character at position i in t; // character in next position
 add string s+c to dictionary d, paired with next available codeword;
// this involves adding a new leaf node if d is represented by a trie
// may have to increment the codeword length k to make this possible
```

text = peter piper picked a peck of pickled pepper

dictionary: "space":32 a:97 c:99 d:100 e:101 f:102 i:105 k:107 l:108 o:111 p:112

r:114 t:116

ASCII encoding of the characters appearing in the text to be compressed (given in the question)

step	position in string	longest string in dictionary	b	add to dictionary	code
1	1				

text = peter piper picked a peck of pickled pepper

dictionary: "space":32 a:97 c:99 d:100 e:101 f:102 i:105 k:107 l:108 o:111 p:112

r:114 t:116

I will use the integer representation for the bit sequence to simplify the presentation

step	position in string	longest string in dictionary	b	add to dictionary	code
1	1	р	112		

text = peter piper picked a peck of pickled pepper

dictionary: "space":32 a:97 c:99 d:100 e:101 f:102 i:105 k:107 l:108 o:111 p:112 r:114 t:116 pe:128

step	position in string	longest string in dictionary	b	add to dictionary	code
1	1	р	112	pe	128

text = peter piper picked a peck of pickled pepper

dictionary: "space":32 a:97 c:99 d:100 e:101 f:102 i:105 k:107 l:108 o:111 p:112 r:114 t:116 pe:128

step	position in string	longest string in dictionary	b	add to dictionary	code
1	1	р	112	pe	128
2	2	е	101		

text = peter piper picked a peck of pickled pepper

dictionary: "space":32 a:97 c:99 d:100 e:101 f:102 i:105 k:107 l:108 o:111 p:112 r:114 t:116 pe:128 et:129

step	position in string	longest string in dictionary	b	add to dictionary	code
1	1	р	112	pe	128
2	2	е	101	et	129

text = peter piper picked a peck of pickled pepper

dictionary: "space":32 a:97 c:99 d:100 e:101 f:102 i:105 k:107 l:108 o:111 p:112 r:114 t:116 pe:128 et:129

step	position in string	longest string in dictionary	b	add to dictionary	code
1	1	р	112	pe	128
2	2	е	101	et	129
3	3	t	116		

text = peter piper picked a peck of pickled pepper

dictionary: "space":32 a:97 c:99 d:100 e:101 f:102 i:105 k:107 l:108 o:111 p:112 r:114 t:116 pe:128 et:129 te:130

step	position in string	longest string in dictionary	b	add to dictionary	code
1	1	р	112	pe	128
2	2	е	101	et	129
3	3	t	116	te	130

text = peter piper picked a peck of pickled pepper

dictionary: "space":32 a:97 c:99 d:100 e:101 f:102 i:105 k:107 l:108 o:111 p:112 r:114 t:116 pe:128 et:129 te:130 er:131

step	position in string	longest string in dictionary	b	add to dictionary	code
1	1	р	112	pe	128
2	2	е	101	et	129
3	3	t	116	te	130
4	4	е	101	er	131

text = peter piper picked a peck of pickled pepper

dictionary: "space":32 a:97 c:99 d:100 e:101 f:102 i:105 k:107 l:108 o:111 p:112 r:114 t:116 pe:128 et:129 te:130 er:131 "r ":132

step	position in string	longest string in dictionary	b	add to dictionary	code
1	1	р	112	pe	128
2	2	е	101	et	129
3	3	t	116	te	130
4	4	е	101	er	131
5	5	r	114	"r "	132

text = peter piper picked a peck of pickled pepper

dictionary: "space":32 a:97 c:99 d:100 e:101 f:102 i:105 k:107 l:108 o:111 p:112 r:114 t:116 pe:128 et:129 te:130 er:131 "r ":132 " p":133

step	position in string	longest string in dictionary	b	add to dictionary	code
1	1	р	112	pe	128
2	2	е	101	et	129
3	3	t	116	te	130
4	4	е	101	er	131
5	5	r	114	"r "	132
6	6	" "	32	" p"	133

text = peter piper picked a peck of pickled pepper

dictionary: "space":32 a:97 c:99 d:100 e:101 f:102 i:105 k:107 l:108 o:111 p:112 r:114 t:116 pe:128 et:129 te:130 er:131 "r ":132 " p":133 pi:134

step	position in string	longest string in dictionary	b	add to dictionary	code
1	1	р	112	pe	128
2	2	е	101	et	129
3	3	t	116	te	130
4	4	е	101	er	131
5	5	r	114	"r "	132
6	6	""	32	" p"	133
7	7	р	112	pi	134

text = peter piper picked a peck of pickled pepper

dictionary: "space":32 a:97 c:99 d:100 e:101 f:102 i:105 k:107 l:108 o:111 p:112 r:114 t:116 pe:128 et:129 te:130 er:131 "r ":132 " p":133 pi:134 ip:135

step	position in string	longest string in dictionary	b	add to dictionary	code
1	1	р	112	pe	128
2	2	е	101	et	129
3	3	t	116	te	130
4	4	е	101	er	131
5	5	r	114	"r "	132
6	6	""	32	" p"	133
7	7	р	112	pi	134
8	8	i	105	ip	135

text = peter piper picked a peck of pickled pepper

dictionary: "space":32 a:97 c:99 d:100 e:101 f:102 i:105 k:107 l:108 o:111 p:112 r:114 t:116 pe:128 et:129 te:130 er:131 "r ":132 " p":133 pi:134 ip:135 per:136

step	position in string	longest string in dictionary	b	add to dictionary	code
9	1	pe	128	per	136

text = peter piper picked a peck of pickled pepper

dictionary: "space":32 a:97 c:99 d:100 e:101 f:102 i:105 k:107 l:108 o:111 p:112
r:114 t:116 pe:128 et:129 te:130 er:131 "r ":132 " p":133 pi:134 ip:135 per:136 r
p:137

step	position in string	longest string in dictionary	b	add to dictionary	code
9	9	pe	128	per	136
10	11	"r "	132	"r p"	137

text = peter piper picked a peck of pickled pepper

dictionary: "space":32 a:97 c:99 d:100 e:101 f:102 i:105 k:107 l:108 o:111 p:112
r:114 t:116 pe:128 et:129 te:130 er:131 "r ":132 " p":133 pi:134 ip:135 per:136 r
p:137 pic:138

step	position in string	longest string in dictionary	b	add to dictionary	code
9	9	pe	128	per	136
10	11	"r "	132	"r p"	137
11	13	pi	133	pic	138

text = peter piper picked a peck of pickled pepper

dictionary: "space":32 a:97 c:99 d:100 e:101 f:102 i:105 k:107 l:108 o:111 p:112
r:114 t:116 pe:128 et:129 te:130 er:131 "r ":132 " p":133 pi:134 ip:135 per:136 r
p:137 pic:138 ck:139

step	position in string	longest string in dictionary	b	add to dictionary	code
9	9	pe	128	per	136
10	11	"r "	132	"r p"	137
11	13	pi	133	pic	138
12	15	С	99	ck	139

text = peter piper picked a peck of pickled pepper

dictionary: "space":32 a:97 c:99 d:100 e:101 f:102 i:105 k:107 l:108 o:111 p:112
r:114 t:116 pe:128 et:129 te:130 er:131 "r ":132 " p":133 pi:134 ip:135 per:136 r
p:137 pic:138 ck:139 ke:140

step	position in string	longest string in dictionary	b	add to dictionary	code
9	9	pe	128	per	136
10	11	"r "	132	"r p"	137
11	13	pi	133	pic	138
12	15	С	99	ck	139
13	16	16 k	107	ke	140
					_

text = peter piper picked a peck of pickled pepper

dictionary: "space":32 a:97 c:99 d:100 e:101 f:102 i:105 k:107 l:108 o:111 p:112
r:114 t:116 pe:128 et:129 te:130 er:131 "r ":132 " p":133 pi:134 ip:135 per:136 r
p:137 pic:138 ck:139 ke:140 ed:141

step	position in string	longest string in dictionary	b	add to dictionary	code
9	9	pe	128	per	136
10	11	"r "	132	"r p"	137
11	13	pi	133	pic	138
12	15	С	99	ck	139
13	16	k	107	ke	140
14	17	е	101	ed	141

text = peter piper picked a peck of pickled pepper

dictionary: "space":32 a:97 c:99 d:100 e:101 f:102 i:105 k:107 l:108 o:111 p:112
r:114 t:116 pe:128 et:129 te:130 er:131 "r ":132 " p":133 pi:134 ip:135 per:136 r
p:137 pic:138 ck:139 ke:140 ed:141 "d ":142

step	position in string	longest string in dictionary	b	add to dictionary	code
9	9	pe	128	per	136
10	11	"r "	132	"r p"	137
11	13	pi	133	pic	138
12	15	С	99	ck	139
13	16	k	107	ke	140
14	17	е	101	ed	141
15	18	d	100	"d "	142

text = peter piper picked a peck of pickled pepper

dictionary: "space":32 a:97 c:99 d:100 e:101 f:102 i:105 k:107 l:108 o:111 p:112
r:114 t:116 pe:128 et:129 te:130 er:131 "r ":132 " p":133 pi:134 ip:135 per:136 r
p:137 pic:138 ck:139 ke:140 ed:141 "d ":142 " a":143

step	position in string	longest string in dictionary	р	add to dictionary	code
9	9	pe	128	per	136
10	11	"r "	132	"r p"	137
11	13	pi	133	pic	138
12	15	С	99	ck	139
13	16	k	107	ke	140
14	17	е	101	ed	141
15	18	d	100	"d "	142
16	19	66 33	97	" a"	143

text = peter piper picked a peck of pickled pepper

dictionary: "space":32 a:97 c:99 d:100 e:101 f:102 i:105 k:107 l:108 o:111 p:112
r:114 t:116 pe:128 et:129 te:130 er:131 "r ":132 " p":133 pi:134 ip:135 per:136 r
p:137 pic:138 ck:139 ke:140 ed:141 "d ":142 " a":143 "a ":144

step	position in string	longest string in dictionary	b	add to dictionary	code
17	20	a	97	"a "	144

text = peter piper picked a peck of pickled pepper

dictionary: "space":32 a:97 c:99 d:100 e:101 f:102 i:105 k:107 l:108 o:111 p:112
r:114 t:116 pe:128 et:129 te:130 er:131 "r ":132 " p":133 pi:134 ip:135 per:136 r
p:137 pic:138 ck:139 ke:140 ed:141 "d ":142 " a":143 "a ":144 " pe":145

step	position in string	longest string in dictionary	b	add to dictionary	code
17	20	a	97	"a "	144
18	21	" p"	133	" pe"	145

text = peter piper picked a peck of pickled pepper

dictionary: "space":32 a:97 c:99 d:100 e:101 f:102 i:105 k:107 l:108 o:111 p:112 r:114 t:116 pe:128 et:129 te:130 er:131 "r ":132 " p":133 pi:134 ip:135 per:136 r p:137 pic:138 ck:139 ke:140 ed:141 "d ":142 " a":143 "a ":144 " pe":145 ec:146

step	position in string	longest string in dictionary	b	add to dictionary	code
17	20	a	97	"a "	144
18	21	" p"	133	" pe"	145
19	23	е	101	ec	146

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dictionary: "space":32 a:97 c:99 d:100 e:101 f:102 i:105 k:107 l:108 o:111 p:112
r:114 t:116 pe:128 et:129 te:130 er:131 "r ":132 " p":133 pi:134 ip:135 per:136 r
p:137 pic:138 ck:139 ke:140 ed:141 "d ":142 " a":143 "a ":144 " pe":145 ec:146
"ck ":147

step	position in string	longest string in dictionary	b	add to dictionary	code
17	20	a	97	"a "	144
18	21	" p"	133	" pe"	145
19	23	е	101	ec	146
20	24	ck	139	"ck "	147

text = peter piper picked a peck of pickled pepper

dictionary: "space":32 a:97 c:99 d:100 e:101 f:102 i:105 k:107 l:108 o:111 p:112 r:114 t:116 pe:128 et:129 te:130 er:131 "r ":132 " p":133 pi:134 ip:135 per:136 r p:137 pic:138 ck:139 ke:140 ed:141 "d ":142 " a":143 "a ":144 " pe":145 ec:146 "ck ":147 " o":148

step	position in string	longest string in dictionary	b	add to dictionary	code
17	20	a	97	"a "	144
18	21	" p"	133	" pe"	145
19	23	е	101	ec	146
20	24	ck	139	"ck "	147
21	25	<i>""</i>	32	" o"	148

text = peter piper picked a peck of pickled pepper

dictionary: "space":32 a:97 c:99 d:100 e:101 f:102 i:105 k:107 l:108 o:111 p:112 r:114 t:116 pe:128 et:129 te:130 er:131 "r ":132 " p":133 pi:134 ip:135 per:136 r p:137 pic:138 ck:139 ke:140 ed:141 "d ":142 " a":143 "a ":144 " pe":145 ec:146 "ck ":147 " o":148 of:149

step	position in string	longest string in dictionary	b	add to dictionary	code
17	20	a	97	"a "	144
18	21	" p"	133	" pe"	145
19	23	е	101	ec	146
20	24	ck	139	"ck "	147
21	25	""	32	" o"	148
22	26	0	111	of	149

text = peter piper picked a peck of pickled pepper

dictionary: "space":32 a:97 c:99 d:100 e:101 f:102 i:105 k:107 l:108 o:111 p:112 r:114 t:116 pe:128 et:129 te:130 er:131 "r ":132 " p":133 pi:134 ip:135 per:136 r p:137 pic:138 ck:139 ke:140 ed:141 "d ":142 " a":143 "a ":144 " pe":145 ec:146 "ck ":147 " o":148 of:149 "f ":150

step	position in string	longest string in dictionary	b	add to dictionary	code
17	20	a	97	"a "	144
18	21	" p"	133	" pe"	145
19	23	е	101	ec	146
20	24	ck	139	"ck "	147
21	25	""	32	" o"	148
22	26	0	111	of	149
23	27	f	102	"f "	150

text = peter piper picked a peck of pickled pepper

dictionary: "space":32 a:97 c:99 d:100 e:101 f:102 i:105 k:107 l:108 o:111 p:112 r:114 t:116 pe:128 et:129 te:130 er:131 "r ":132 " p":133 pi:134 ip:135 per:136 r p:137 pic:138 ck:139 ke:140 ed:141 "d ":142 " a":143 "a ":144 " pe":145 ec:146 "ck ":147 " o":148 of:149 "f ":150 " pi":151

step	position in string	longest string in dictionary	b	add to dictionary	code
17	20	a	97	"a "	144
18	21	" p"	133	" pe"	145
19	23	е	101	ec	146
20	24	ck	139	"ck "	147
21	25	<i>""</i>	32	" o"	148
22	26	0	111	of	149
23	27	f	102	"f "	150
24	28	" p"	133	" pi"	151

text = peter piper picked a peck of pickled pepper

dictionary: "space":32 a:97 c:99 d:100 e:101 f:102 i:105 k:107 l:108 o:111 p:112 r:114 t:116 pe:128 et:129 te:130 er:131 "r ":132 " p":133 pi:134 ip:135 per:136 r p:137 pic:138 ck:139 ke:140 ed:141 "d ":142 " a":143 "a ":144 " pe":145 ec:146 "ck ":147 " o":148 of:149 "f ":150 " pi":151 ic:152

step	position in string	longest string in dictionary	b	add to dictionary	code
25	30	i	105	ic	152

text = peter piper picked a peck of pickled pepper

dictionary: "space":32 a:97 c:99 d:100 e:101 f:102 i:105 k:107 l:108 o:111 p:112 r:114 t:116 pe:128 et:129 te:130 er:131 "r ":132 " p":133 pi:134 ip:135 per:136 r p:137 pic:138 ck:139 ke:140 ed:141 "d ":142 " a":143 "a ":144 " pe":145 ec:146 "ck ":147 " o":148 of:149 "f ":150 " pi":151 ic:152 ckl:153

step	position in string	longest string in dictionary	b	add to dictionary	code
25	30	i	105	ic	152
26	31	ck	139	ck1	153

text = peter piper picked a peck of pickled pepper

dictionary: "space":32 a:97 c:99 d:100 e:101 f:102 i:105 k:107 l:108 o:111 p:112 r:114 t:116 pe:128 et:129 te:130 er:131 "r ":132 " p":133 pi:134 ip:135 per:136 r p:137 pic:138 ck:139 ke:140 ed:141 "d ":142 " a":143 "a ":144 " pe":145 ec:146 "ck ":147 " o":148 of:149 "f ":150 " pi":151 ic:152 ckl:153 le:154

step	position in string	longest string in dictionary	b	add to dictionary	code
25	30	i	105	ic	152
26	31	ck	139	ck1	153
27	33	1	108	1e	154

text = peter piper picked a peck of pickled pepper

dictionary: "space":32 a:97 c:99 d:100 e:101 f:102 i:105 k:107 l:108 o:111 p:112 r:114 t:116 pe:128 et:129 te:130 er:131 "r ":132 " p":133 pi:134 ip:135 per:136 r p:137 pic:138 ck:139 ke:140 ed:141 "d ":142 " a":143 "a ":144 " pe":145 ec:146 "ck ":147 " o":148 of:149 "f ":150 " pi":151 ic:152 ckl:153 le:154 "ed ":155

step	position in string	longest string in dictionary	b	add to dictionary	code
25	30	i	105	ic	152
26	31	ck	139	ck1	153
27	33	1	108	1e	154
28	34	ed	141	"ed "	155
					_

text = peter piper picked a peck of pickled pepper

dictionary: "space":32 a:97 c:99 d:100 e:101 f:102 i:105 k:107 l:108 o:111 p:112 r:114 t:116 pe:128 et:129 te:130 er:131 "r ":132 " p":133 pi:134 ip:135 per:136 r p:137 pic:138 ck:139 ke:140 ed:141 "d ":142 " a":143 "a ":144 " pe":145 ec:146 "ck ":147 " o":148 of:149 "f ":150 " pi":151 ic:152 ckl:153 le:154 "ed ":155 " pep":156

step	position in string	longest string in dictionary	b	add to dictionary	code
25	30	i	105	ic	152
26	31	ck	139	ck1	153
27	33	7	108	1e	154
28	34	ed	141	"ed "	155
29	36	" pe"	145	" pep"	156

text = peter piper picked a peck of pickled pepper

dictionary: "space":32 a:97 c:99 d:100 e:101 f:102 i:105 k:107 l:108 o:111 p:112 r:114 t:116 pe:128 et:129 te:130 er:131 "r ":132 " p":133 pi:134 ip:135 per:136 r p:137 pic:138 ck:139 ke:140 ed:141 "d ":142 " a":143 "a ":144 " pe":145 ec:146 "ck ":147 " o":148 of:149 "f ":150 " pi":151 ic:152 ckl:153 le:154 "ed ":155 " pep":156 pp:157

step	position in string	longest string in dictionary	b	add to dictionary	code
25	30	i	105	ic	152
26	31	ck	139	ck1	153
27	33	1	108	1e	154
28	34	ed	141	"ed "	155
29	36	" pe"	145	" pep"	156
30	39	р	112	рр	157

text = peter piper picked a peck of pickled pepper

dictionary: "space":32 a:97 c:99 d:100 e:101 f:102 i:105 k:107 l:108 o:111 p:112 r:114 t:116 pe:128 et:129 te:130 er:131 "r ":132 " p":133 pi:134 ip:135 per:136 r p:137 pic:138 ck:139 ke:140 ed:141 "d ":142 " a":143 "a ":144 " pe":145 ec:146 "ck ":147 " o":148 of:149 "f ":150 " pi":151 ic:152 ckl:153 le:154 "ed ":155 " pep":156 pp:157

step	position in string	longest string in dictionary	b	add to dictionary	code
25	30	i	105	ic	152
26	31	ck	139	ckl	153
27	33	1	108	le	154
28	34	ed	141	"ed "	155
29	36	" pe"	145	" pep"	156
30	39	р	112	рр	157
31	40	per	136		

Suppose that the LZW compression algorithm is applied to the text:

ababababab...

of length 100 (i.e. alternating sequence of a's and b's and 100 chars)

Determine the compression ratio and space saved

- assume that each character in the source file occupies 8 bits
- the initial dictionary size is 128
- the initial codeword size is 8

Hint:

- apply the LZW algorithm until a pattern emerges
- use this to compute the steps required to reach the 100th (and last) char
- from this you can find the number of code words used to encode the text

LZW compression - Pseudo code

```
set current text position i to 0;
initialise codeword length k (say to 8);
initialise the dictionary d;
while (the text t is not exhausted) {
 identify the longest string s, starting at position i of text t
 that is represented in the dictionary d;
 // there is such string, as all strings of length 1 are in d
 output codeword for the string s; // using k bits
// move to the next position in t
i += s.length(); // move forward by the length of string just encoded
 c = character at position i in t; // character in next position
 add string s+c to dictionary d, paired with next available codeword;
 // may have to increment the codeword length k to make this possible
```

text = ababababababababababababab...

dictionary: a:0 b:1

text = abababababababababababababab...

dictionary: a:0 b:1

step	position in string	longest string in dictionary	add to dictionary	code
1	1	a		
2				
3				
4				
5				
6				
7				
8				
9				
10				

text = abababababababababababab...

dictionary: a:0 b:1 ab:2

step	position in string	longest string in dictionary	add to dictionary	code
1	1	a	ab	2
2				
3				
4				
5				
6				
7				
8				
9				
10				

text = ababababababababababababab...

dictionary: a:0 b:1 ab:2

step	position in string	longest string in dictionary	add to dictionary	code
1	1	a	ab	2
2	2	b		
3				
4				
5				
6				
7				
8				
9				
10				

text = abababababababababababab...

dictionary: a:0 b:1 ab:2 ba:3

step	position in string	longest string in dictionary	add to dictionary	code
1	1	a	ab	2
2	2	b	ba	3
3				
4				
5				
6				
7				
8				
9				
10				

text = ababababababababababababab...

dictionary: a:0 b:1 ab:2 ba:3

step	position in string	longest string in dictionary	add to dictionary	code
1	1	a	ab	2
2	2	b	ba	3
3	3	ab		
4				
5				
6				
7				
8				
9				
10				

text = abababababababababababab...

dictionary: a:0 b:1 ab:2 ba:3 aba:4

step	position in string	longest string in dictionary	add to dictionary	code
1	1	a	ab	2
2	2	b	ba	3
3	3	ab	aba	4
4				
5				
6				
7				
8				
9				
10				

text = ababababababababababababab...

dictionary: a:0 b:1 ab:2 ba:3 aba:4

step	position in string	longest string in dictionary	add to dictionary	code
1	1	a	ab	2
2	2	b	ba	3
3	3	ab	aba	4
4	5	aba		
5				
6				
7				
8				
9				
10				

text = abababababababababababab...

dictionary: a:0 b:1 ab:2 ba:3 aba:4 abab:5

step	position in string	longest string in dictionary	add to dictionary	code
1	1	a	ab	2
2	2	b	ba	3
3	3	ab	aba	4
4	5	aba	abab	5
5				
6				
7				
8				
9				
10				

text = ababababababababababababab...

dictionary: a:0 b:1 ab:2 ba:3 aba:4 abab:5

step	position in string	longest string in dictionary	add to dictionary	code
1	1	a	ab	2
2	2	b	ba	3
3	3	ab	aba	4
4	5	aba	abab	5
5	8	ba		
6				
7				
8				
9				
10				

text = abababababababababababab...

dictionary: a:0 b:1 ab:2 ba:3 aba:4 abab:5 bab:6

step	position in string	longest string in dictionary	add to dictionary	code
1	1	a	ab	2
2	2	b	ba	3
3	3	ab	aba	4
4	5	aba	abab	5
5	8	ba	bab	6
6				
7				
8				
9				
10				

text = ababababababababababababab...

dictionary: a:0 b:1 ab:2 ba:3 aba:4 abab:5 bab:6

step	position in string	longest string in dictionary	add to dictionary	code
1	1	a	ab	2
2	2	b	ba	3
3	3	ab	aba	4
4	5	aba	abab	5
5	8	ba	bab	6
6	10	bab		
7				
8				
9				
10				

dictionary: a:0 b:1 ab:2 ba:3 aba:4 abab:5 bab:6 baba:7

step	position in string	longest string in dictionary	add to dictionary	code
1	1	a	ab	2
2	2	b	ba	3
3	3	ab	aba	4
4	5	aba	abab	5
5	8	ba	bab	6
6	10	bab	baba	7
7				
8				
9				
10				

dictionary: a:0 b:1 ab:2 ba:3 aba:4 abab:5 bab:6 baba:7

step	position in string	longest string in dictionary	add to dictionary	code
1	1	a	ab	2
2	2	b	ba	3
3	3	ab	aba	4
4	5	aba	abab	5
5	8	ba	bab	6
6	10	bab	baba	7
7	13	abab		
8				
9				
10				

dictionary: a:0 b:1 ab:2 ba:3 aba:4 abab:5 bab:6 baba:7 ababa:8

step	position in string	longest string in dictionary	add to dictionary	code
1	1	a	ab	2
2	2	b	ba	3
3	3	ab	aba	4
4	5	aba	abab	5
5	8	ba	bab	6
6	10	bab	baba	7
7	13	abab	ababa	8
8				
9				
10				

text = abababababababababababab...

dictionary: a:0 b:1 ab:2 ba:3 aba:4 abab:5 bab:6 baba:7 ababa:8

step	position in string	longest string in dictionary	add to dictionary	code
1	1	a	ab	2
2	2	b	ba	3
3	3	ab	aba	4
4	5	aba	abab	5
5	8	ba	bab	6
6	10	bab	baba	7
7	13	abab	ababa	8
8	17	ababa		
9				
10				

text = ababababababababababababab...

dictionary: a:0 b:1 ab:2 ba:3 aba:4 abab:5 bab:6 baba:7 ababa:8 ababab:9

step	position in string	longest string in dictionary	add to dictionary	code
1	1	a	ab	2
2	2	b	ba	3
3	3	ab	aba	4
4	5	aba	abab	5
5	8	ba	bab	6
6	10	bab	baba	7
7	13	abab	ababa	8
8	17	ababa	ababab	9
9				
10				

text = abababababababababababab...

dictionary: a:0 b:1 ab:2 ba:3 aba:4 abab:5 bab:6 baba:7 ababa:8 ababab:9

step	position in string	longest string in dictionary	add to dictionary	code
1	1	a	ab	2
2	2	b	ba	3
3	3	ab	aba	4
4	5	aba	abab	5
5	8	ba	bab	6
6	10	bab	baba	7
7	13	abab	ababa	8
8	17	ababa	ababab	9
9	22	baba		
10				

text = ababababababababababababa...

dictionary: a:0 b:1 ab:2 ba:3 aba:4 abab:5 bab:6 baba:7 ababa:8 ababab:9
 babab:10

step	position in string	longest string in dictionary	add to dictionary	code
1	1	a	ab	2
2	2	b	ba	3
3	3	ab	aba	4
4	5	aba	abab	5
5	8	ba	bab	6
6	10	bab	baba	7
7	13	abab	ababa	8
8	17	ababa	ababab	9
9	22	baba	babab	10
10				

text = ababababababababababababab...

dictionary: a:0 b:1 ab:2 ba:3 aba:4 abab:5 bab:6 baba:7 ababa:8 ababab:9
 babab:10 bababa:11

step	position in string	longest string in dictionary	add to dictionary	code
1	1	a	ab	2
2	2	b	ba	3
3	3	ab	aba	4
4	5	aba	abab	5
5	8	ba	bab	6
6	10	bab	baba	7
7	13	abab	ababa	8
8	17	ababa	ababab	9
9	22	baba	babab	10
10	26	babab	bababa	11

Looking at how much of text is used up in each step we find the

pattern: 1,1,2,3,2,3,4,5,4,5,6,7,6,7,8,9,8,9,10,11, ...

step	position in string	longest string in dictionary	add to dictionary	code
1	1	a	ab	2
2	2	b	ba	3
3	3	ab	aba	4
4	5	aba	abab	5
5	8	ba	bab	6
6	10	bab	baba	7
7	13	abab	ababa	8
8	17	ababa	ababab	9
9	22	baba	babab	10
10	26	babab	bababa	11

Looking at how much of text is used up in each step we find the

pattern: 1,1,2,3,2,3,4,5,4,5,6,7,6,7,8,9,8,9,10,11,...

- therefore the number of steps required to encode the text we need to
 know how many of these terms are need for the summation to equal 100
- after a bit of adding up we find the first 19 terms sum to 100

End up with 21 code words so never have to increase the dictionary

- 2 code words to start with and then add 1 code word in each step
- 8 bits to encode the text so can fit 28 code words

19 steps means 19 (8 bit) code words in the compressed file

- input size 100 · 8 and output size 19 · 8
- compression ratio: 152/800=0.19 and space saved: $(1-0.19) \times 100=81\%$

Find the distance between the strings

s = agcgatc and t = ctacgaccg

and derive an optimum alignment

Find the distance between the strings

and derive an optimum alignment

Recall the recurrence relation is given by:

$$d(i,j) = \begin{cases} d(i-1,j-1) & \text{if } s[i-1]=t[j-1] \\ \\ 1 + min\{d(i,j-1), \ d(i-1,j), \ d(i-1,j-1)\} & \text{otherwise} \end{cases}$$

subject to d(i,0)=i and d(0,j)=j for all $i \le n-1$ and $j \le m-1$

s\t		0	1	2	3	4	5	6	7	8	9
			С	t	a	С	g	a	С	С	g
0											
1	a										
2	g										
3	C										
4	g										
5	a										
6	t		_	_	_				-	-	_
7	С										

s\t		0	1	2	3	4	5	6	7	8	9
			C	t	a	C	g	a	С	С	g
0		0	1	2	3	4	5	6	7	8	9
1	a	1									
2	g	2									
3	C	3									
4	g	4									
5	a	5									
6	t	6			_	_		_	-	_	
7	С	7									

table is initialised by filling in row 0 and column 0 using the initial conditions of the recurrence relation

s\t		0	1	2	3	4	5	6	7	8	9
			С	t	a	С	g	a	С	С	g
0		0	1	2	3	4	5	6	7	8	9
1	a	1	1								
2	g	2									
3	С	3									
4	g	4									
5	a	5									
6	t	6									
7	С	7									

The entries are calculated one by one by application of the formula

-d(1,1)=1+0=1 (since $s[0]\neq t[0]$ and min(d(1,0),d(0,1),d(1,1))=0)

s\t		0	1	2	3	4	5	6	7	8	9
			С	t	a	С	g	a	С	С	g
0		0	1	2	3	4	5	6	7	8	9
1	a	1	1	2							
2	g	2									
3	С	3									
4	g	4									
5	a	5									
6	t	6									
7	С	7									

The entries are calculated one by one by application of the formula

-d(1,2)=1+1=2 (since $s[0]\neq t[1]$ and min(d(1,0),d(0,2),d(0,1))=1)

s\t		0	1	2	3	4	5	6	7	8	9
			С	t	a	С	g	a	С	С	g
0		0	1	2	3	4	5	6	7	8	9
1	a	1	1	2	2						
2	g	2									
3	C	3									
4	g	4									
5	a	5									
6	t	6									
7	С	7									

The entries are calculated one by one by application of the formula

$$- d(1,3)=2$$
 (since $s[0]=t[2]$ and $d(0,2)=2$)

s\t		0	1	2	3	4	5	6	7	8	9
			С	t	a	С	g	a	С	С	g
0		0	1	2	3	4	5	6	7	8	9
1	a	1	1	2	2	3	4	5	6	7	8
2	g	2	2	2	3	3	3	4	5	6	7
3	O	3	2	3	3	3	4	4	4	5	6
4	g	4	3								
5	a										
6	t										
7	С										

The entries are calculated one by one by application of the formula

-d(4,1)=1+2=3 (since s[3] \neq t[0] and min(d(4,0),d(3,1),d(3,0))=2)

s\t		0	1	2	3	4	5	6	7	8	9
			C	t	a	С	g	a	С	С	g
0		0	1	2	3	4	5	6	7	8	9
1	a	1	1	2	2	3	4	5	6	7	8
2	g	2	2	2	3	3	3	4	5	6	7
3	С	3	2	3	3	3	4	4	4	5	6
4	g	4	3	3	4	4	3	4	5	5	5
5	a	5	4	4	3	4	4	3	4	5	6
6	t	6	5	4	4	4	5	4	4	5	6
7	C	7	6	5	5	4	5	5	4	4	5

The entries are calculated one by one by application of the formula

- the final table: d(7,9)=5 so the string distance is 5

s\t		0	1	2	3	4	5	6	7	8	9
			C	t	a	С	g	a	С	С	g
0		0	1	2	3	4	5	6	7	8	9
1	a	1	1	2	2	3	4	5	6	7	8
2	g	2	2	2	3	3	3	4	5	6	7
3	U	3	2	3	3	3	4	4	4	5	6
4	g	4	3	3	4	4	3	4	5	5	5
5	a	5	4	4	3	4	4	3	4	5	6
6	t	6	5	4	4	4	5	4	4	5	6
7	С	7	6	5	5	4	5	5	4	4	- 5

traceback: d(7,9)=1+d(7,8)

s\t		0	1	2	3	4	5	6	7	8	9
			C	t	a	С	g	a	С	С	g
0		0	1	2	3	4	5	6	7	8	9
1	a	1	1	2	2	3	4	5	6	7	8
2	g	2	2	2	3	3	3	4	5	6	7
3	C	3	2	3	3	3	4	4	4	5	6
4	g	4	3	3	4	4	3	4	5	5	5
5	a	5	4	4	3	4	4	3	4	5	6
6	t	6	5	4	4	4	5	4	4_	5	6
7	С	7	6	5	5	4	5	5	4	4	- 5

traceback: d(7,8)=d(6,7)

s\t		0	1	2	3	4	5	6	7	8	9
			C	t	a	С	g	a	С	С	g
0		0	1	2	3	4	5	6	7	8	9
1	a	1	1	2	2	3	4	5	6	7	8
2	g	2	2	2	3	3	3	4	5	6	7
3	C	3	2	3	3	3	4	4	4	5	6
4	g	4	3	3	4	4	3	4	5	5	5
5	a	5	4	4	3	4	4	3	4	5	6
6	t	6	5	4	4	4	5	4	4	5	6
7	С	7	6	5	5	4	5	5	4	4	- 5

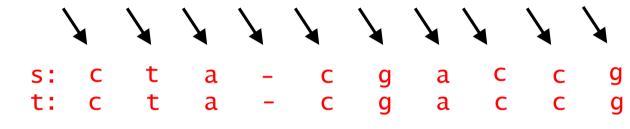
traceback: d(6,7)=1+d(5,6)

s\t		0	1	2	3	4	5	6	7	8	9
			C	t	a	С	g	a	С	С	g
0		0	-1	- 2	3	4	5	6	7	8	9
1	a	1	1	2	2	3	4	5	6	7	8
2	g	2	2	2	3	3	3	4	5	6	7
3	C	3	2	3	3	3	4	4	4	5	6
4	g	4	3	3	4	4	3	4	5	5	5
5	a	5	4	4	3	4	4	3	4	5	6
6	t	6	5	4	4	4	5	4	4	5	6
7	С	7	6	5	5	4	5	5	4	4	- 5

traceback:

Corresponding alignment:

insertions ertino at childeletion at children at children



```
step: h \leftarrow h \leftarrow d \leftarrow v \leftarrow d \leftarrow d \leftarrow d \leftarrow d \leftarrow h
```

Interpretation

- v (vertical) steps deletions
- h (horizontal) steps insertions
- d (diagonal) steps as substitutions or matches

A string u is a subsequence of a string s if u can be obtained from s by deleting zero or more characters. A string u is a common subsequence of s and t if it is a subsequence of both s and t

Design a dynamic programming algorithm to determine the length of the longest common subsequence (LCS) of two strings s and t

What are the time and space complexities of your algorithm?

Hint: base your algorithm on evaluating 1(i,j), the length of the LCS of the ith prefix of s and the jth prefix of t

Design a dynamic programming algorithm to determine the length of the longest common subsequence (LCS) of two strings s and t

$$1(i,j) = \begin{cases} if s[i-1]=t[j-1] \\ otherwise \end{cases}$$

We have found one element in common

Design a dynamic programming algorithm to determine the length of the longest common subsequence (LCS) of two strings s and t

$$1(i,j) = \begin{cases} 1(i-1,j-1) + 1 & \text{if } s[i-1]=t[j-1] \\ & \text{otherwise} \end{cases}$$

We have found one element in common

LCS is 1+ LCS of what remains

Design a dynamic programming algorithm to determine the length of the longest common subsequence (LCS) of two strings s and t

$$1(i,j) = \begin{cases} 1(i-1,j-1) + 1 & \text{if } s[i-1]=t[j-1] \\ & \text{otherwise} \end{cases}$$

Elements do not match, so need to delete one of them

 no point deleting both as the other may match and want longest common subsequence

Design a dynamic programming algorithm to determine the length of the longest common subsequence (LCS) of two strings s and t

$$1(i,j) = \begin{cases} 1(i-1,j-1) + 1 & \text{if } s[i-1]=t[j-1] \\ \max\{1(i,j-1), 1(i-1,j)\} & \text{otherwise} \end{cases}$$

Elements do not match, so need to delete one of them

- no point deleting both as the other may match
- want longest common subsequence

Design a dynamic programming algorithm to determine the length of the longest common subsequence (LCS) of two strings s and t

$$1(i,j) = \begin{cases} 1(i-1,j-1) + 1 & \text{if } s[i-1]=t[j-1] \\ \max\{1(i,j-1), 1(i-1,j)\} & \text{otherwise} \end{cases}$$

Base case what is the longest common subsequence with any string and the empty string?

subject to 1(i,0)=0 and 1(0,j)=0 for all $0 \le i \le n-1$ and $0 \le j \le m-1$

Design a dynamic programming algorithm to determine the length of the longest common subsequence (LCS) of two strings s and t

$$1(i,j) = \begin{cases} 1(i-1,j-1) + 1 & \text{if } s[i-1]=t[j-1] \\ \max\{1(i,j-1), 1(i-1,j)\} & \text{otherwise} \end{cases}$$

subject to l(i,0)=0 and l(0,j)=0 for all $0 \le i \le n-1$ and $0 \le j \le m-1$

The time and space complexities are both O(mn) for strings of lengths m and n, same reasoning as for the string distance algorithm

s\t		0	1	2	3	4	5	6	7	8	9
			С	t	a	С	g	a	С	С	g
0											
1	a										
2	g										
3	U										
4	g										
5	a										
6	t										
7	U				_					_	

s\t		0	1	2	3	4	5	6	7	8	9
			C	t	a	С	g	a	С	С	g
0		0	0	0	0	0	0	0	0	0	0
1	a	0									
2	g	0									
3	С	0									
4	g	0									
5	a	0									
6	t	0									
7	С	0									

1(i,0)=0 and 1(0,j)=0 for all $0 \le i \le n-1$ and $0 \le j \le m-1$

s\t		0	1	2	3	4	5	6	7	8	9
			C	t	a	С	g	a	С	С	g
0		0	0	0	0	0	0	0	0	0	0
1	a	0	0								
2	g	0									
3	O	0									
4	g	0									
5	a	0									
6	t	0									
7	С	0									

1(1,1)=0 (since $s[0] \neq t[0]$ and max(1(1,0),1(0,1))=0)

s\t		0	1	2	3	4	5	6	7	8	9
			C	t	a	С	g	a	C	С	g
0		0	0	0	0	0	0	0	0	0	0
1	a	0	0	0	1						
2	g	0	0								
3	O	0	1								
4	g	0	1								
5	a	0	1								
6	t	0	1		-	_	_	_	_		-
7	С	0	1								

1(1,3)=1 (since s[0]=t[2] and 1(0,2)+1=1)

s\t		0	1	2	3	4	5	6	7	8	9
			U	t	a	С	g	a	С	С	g
0		0	0	0	0	0	0	0	0	0	0
1	a	0	0	0	1	1	1	1	1	1	1
2	g	0	0	0	1						
3	C	0	1								
4	g	0	1								
5	a	0	1								
6	t	0	1		_	-	_	_	_		
7	C	0	1								

1(2,3)=1 (since $s[2]\neq t[3]$ and max(1(2,1),1(1,3))=0)

s\t		0	1	2	3	4	5	6	7	8	9
			С	t	a	С	g	a	С	С	g
0		0	0	0	0	0	0	0	0	0	0
1	a	0	0	0	1	1	1	1	1	1	1
2	g	0	0	0	1	1	2	2	2	2	2
3	С	0	1	1							
4	g	0	1	1							
5	a	0	1	1							
6	t	0	1	2							
7	C	0	1								

1(6,2)=1 (since s[5]=t[1] and 1(5,1)+1=2)

s\t		0	1	2	3	4	5	6	7	8	9
			C	t	a	С	g	a	С	С	g
0		0	0	0	0	0	0	0	0	0	0
1	a	0	0	0	1	1	1	1	1	1	1
2	g	0	0	0	1	1	2	2	2	2	2
3	U	0	1	1	1	2	2	2	3	3	3
4	g	0	1	1	1	2	3	3	3	3	4
5	a	0	1	1	2	2	3	4	4	4	4
6	t	0	1	2	2	2	3	4	4	4	4
7	U	0	1	2	2	3	3	4	5	5	5

1(7,9)=0 (since s[6] \neq t[8] and max(1(6,9),1(7,8))={4,5}=5)

s\t		0	1	2	3	4	5	6	7	8	9
			U	t	a	С	g	a	C	С	g
0		0	0	- 0	0	0	0	0	0	0	0
1	a	0	0	0	1	1	1	1	1	1	1
2	g	0	0	0	1	1	2	2	2	2	2
3	C	0	1	1	1	2	2	2	3	3	3
4	g	0	1	1	1	2	3	3	3	3	4
5	a	0	1	1	2	2	3	4	4	4	4
6	t	0	1	2	2	2	3	4	4	4	4
7	С	0	1	2	2	3	3	4	5	5	5

Substring: a c g a c

```
s: - - a g c g a - t c - t: c t a - c g a c - c g

step: h \leftarrow h \leftarrow d \leftarrow v \leftarrow d \leftarrow d \leftarrow h \leftarrow v \leftarrow d \leftarrow h
```

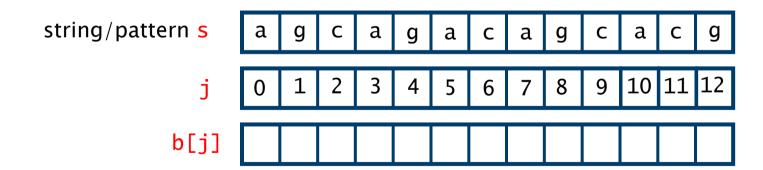
Interpretation

- v (vertical) steps delete current element from s
- h (horizontal) steps delete current element from t
- d (diagonal) steps skip (keep) current element in both s and t

A border of a string s is a substring that is both a prefix and a suffix and must be a strict subsequence

Border table b: array which has the same size as the string

```
- b[j] = the length of the longest border of s[0..j-1]
= max \{ k \mid s[0..k-1] = s[j-k..j-1] \land k < j \}
```

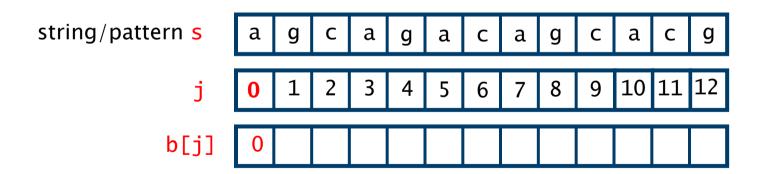


Construct the above border table

A border of a string s is a substring that is both a prefix and a suffix and must be a strict subsequence

Border table b: array which has the same size as the string

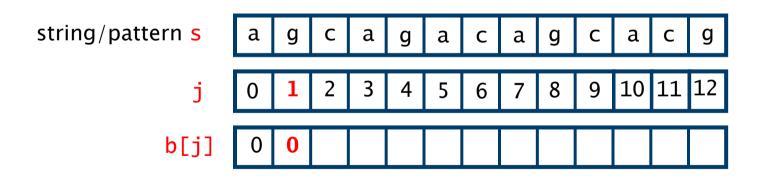
```
- b[j] = the length of the longest border of s[0..j-1]
= max \{ k \mid s[0..k-1] = s[j-k..j-1] \land k < j \}
```



 $s[0..-1] = empty_string$

A border of a string s is a substring that is both a prefix and a suffix and must be a strict subsequence

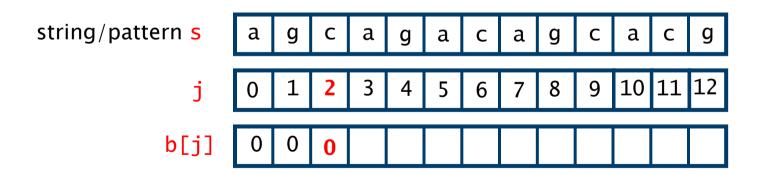
```
- b[j] = the length of the longest border of s[0..j-1]
= max \{ k \mid s[0..k-1] = s[j-k..j-1] \land k < j \}
```



$$s[0..0] = a$$

A border of a string s is a substring that is both a prefix and a suffix and must be a strict subsequence

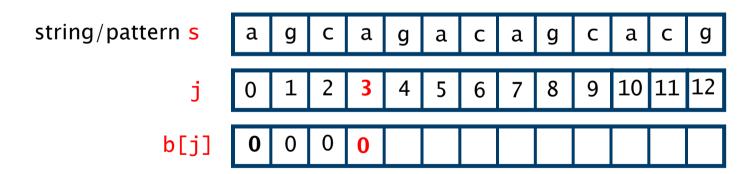
```
- b[j] = the length of the longest border of s[0..j-1]
= max \{ k \mid s[0..k-1] = s[j-k..j-1] \land k < j \}
```



$$s[0..1] = ag$$

A border of a string s is a substring that is both a prefix and a suffix and must be a strict subsequence

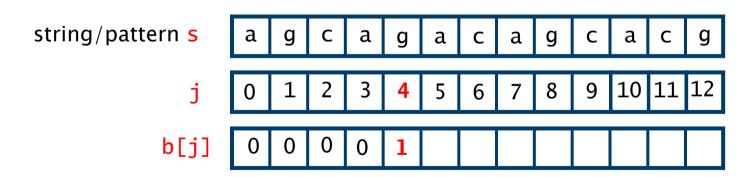
```
- b[j] = the length of the longest border of s[0..j-1]
= max \{ k \mid s[0..k-1] = s[j-k..j-1] \land k < j \}
```



$$s[0..2] = agc$$

A border of a string s is a substring that is both a prefix and a suffix and must be a strict subsequence

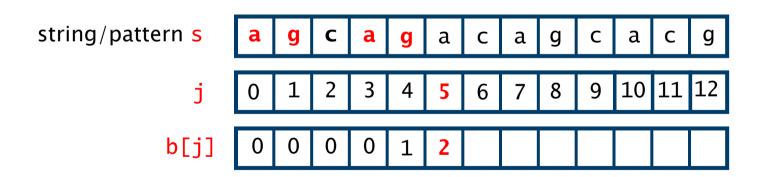
```
- b[j] = the length of the longest border of s[0..j-1]
= max \{ k \mid s[0..k-1] = s[j-k..j-1] \land k < j \}
```



```
s[0..3] = agca and s[0..1-1] = s[4-1..3] = a
```

A border of a string s is a substring that is both a prefix and a suffix and must be a strict subsequence

```
- b[j] = the length of the longest border of s[0..j-1]
= max \{ k \mid s[0..k-1] = s[j-k..j-1] \land k < j \}
```

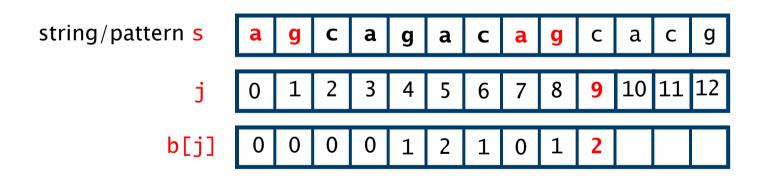


```
s[0..4] = agcag and s[0..2-1] = s[5-2..4] = ag
```

A border of a string s is a substring that is both a prefix and a suffix and must be a strict subsequence

Border table b: array which has the same size as the string

```
- b[j] = the length of the longest border of s[0..j-1]
= max \{ k \mid s[0..k-1] = s[j-k..j-1] \land k < j \}
```

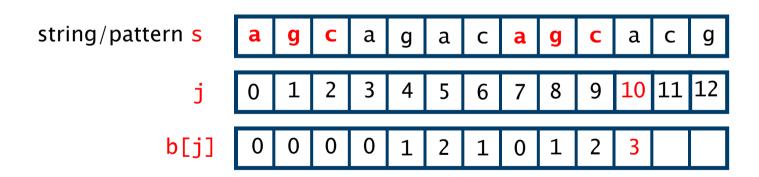


s[0..8] = agcagacag and s[0..2-1] = s[9-2..8] = ag

A border of a string s is a substring that is both a prefix and a suffix and must be a strict subsequence

Border table b: array which has the same size as the string

```
- b[j] = the length of the longest border of s[0..j-1]
= max \{ k \mid s[0..k-1] = s[j-k..j-1] \land k < j \}
```

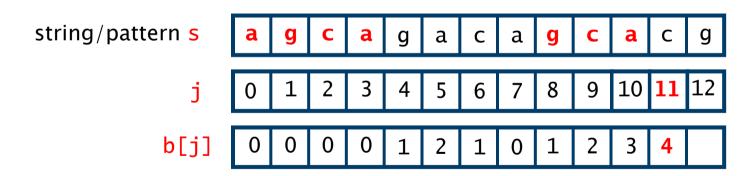


s[0..9] = agcagacagc and s[0..3-1] = s[9-3..9] = agc

A border of a string s is a substring that is both a prefix and a suffix and must be a strict subsequence

Border table b: array which has the same size as the string

```
- b[j] = the length of the longest border of s[0..j-1]
= max \{ k \mid s[0..k-1] = s[j-k..j-1] \land k < j \}
```

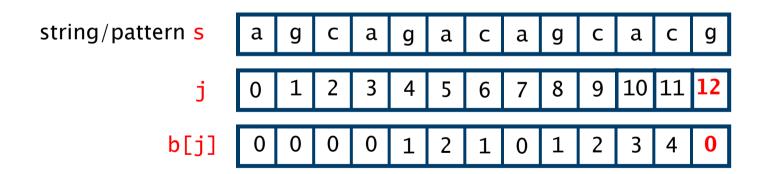


s[0..10] = agcagacagca and s[0..4-1] = s[10-4..10] = agca

A border of a string s is a substring that is both a prefix and a suffix and must be a strict subsequence

Border table b: array which has the same size as the string

```
- b[j] = the length of the longest border of s[0..j-1]
= max \{ k \mid s[0..k-1] = s[j-k..j-1] \land k < j \}
```



s[0..11] = agcagacagcac

Boyer-Moore Algorithm: the string is scanned right-to-left

- text character involved in a mismatch is used to decide next comparison
- involves pre-processing the string to record the position of the last occurrence of each character c in the alphabet
- therefore the alphabet must be fixed in advance of the search

Last occurrence array p

- p[c] position of character c in the string s
- i.e. equals $p[c]=max\{k \mid s[k]=c\}$ if such a k exists and -1 otherwise

On finding a mismatch there is a jump step in the algorithm

- if the mismatch is between s[j] and t[i]
- 'slide s along' so that position p[t[i]] of s aligns with t[i]
 - · i.e. align last position in s of character t[i] with position i of t
- if this moves s in the 'wrong direction', instead move s one position right
- if t[i] does not appear in string, 'slide string' passed t[i]
 - · i.e. align position -1 of s with position i of t

```
agcgcctgatagcgacagt
agcga
```

- 1 comparison (c and a differ)
- c appears in string so move along so last c of string line up with c in text

```
agcgcctgatagcgacagt
agcga
```

- 1 comparison (c and a differ)
- c appears in string so move along so last c of string line up with c in text

```
agcgcctgatagcgacagt
agcga
```

- 1 comparison (c and a differ)
- 1 comparison (t and a differ)
- t does not appear in string so move past t

```
agcgcctgatagcgacagt
agcga
```

- 1 comparison (c and a differ)
- 1 comparison (t and a differ)
- t does not appear in string so move past t

```
agcgcctgatagcgacagt
agcga
```

- 1 comparison (c and a differ)
- 1 comparison (t and a differ)
- 1 comparison (g and a differ)
- g appears in string so move along so last g of string line up with g in text

```
agcgcctgatagcgacagt
agcga
```

- 1 comparison (c and a differ)
- 1 comparison (t and a differ)
- 1 comparison (g and a differ)
- g appears in string so move along so last g of string line up with g in text

```
agcgcctgatagcgacagt
agcga
```

- 1 comparison (c and a differ)
- 1 comparison (t and a differ)
- 1 comparison (g and a differ)
- 1 comparison (c and a differ)
- c appears in string so move along so last c of string line up with c in text

```
agcgcctgatagcgacagt
agcga
```

- 1 comparison (c and a differ)
- 1 comparison (t and a differ)
- 1 comparison (g and a differ)
- 1 comparison (c and a differ)
- c appears in string so move along so last c of string line up with c in text

```
agcgcctgatagcgacagt
agcga
```

- 1 comparison (c and a differ)
- 1 comparison (t and a differ)
- 1 comparison (g and a differ)
- 1 comparison (c and a differ)
- 5 comparisons and found string
- 9 comparisons in total