

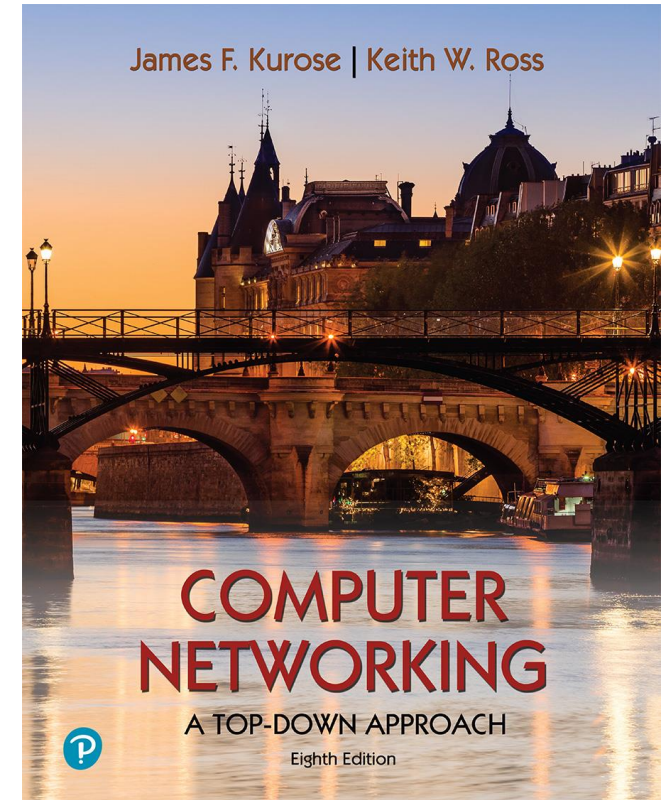
Chapter 2

Application Layer

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Adapted from the slides of the book's authors



*Computer Networking: A
Top-Down Approach*

8th edition
Jim Kurose, Keith Ross
Pearson, 2020

Application layer: overview

- Principles of network applications
- socket programming with UDP and TCP
- Web and HTTP
- E-mail, SMTP, IMAP
- The Domain Name System DNS
- P2P applications
- video streaming and content distribution networks



Application layer: overview

Our goals:

- conceptual *and* implementation aspects of application-layer protocols
 - transport-layer service models
 - client-server paradigm
 - peer-to-peer paradigm
- learn about protocols by examining popular application-layer protocols and infrastructure
 - HTTP
 - SMTP, IMAP
 - DNS
 - video streaming systems, CDNs
- programming network applications
 - socket API

Some network apps

- social networking
 - Web
 - text messaging
 - e-mail
 - multi-user network games
 - streaming stored video (YouTube, Hulu, Netflix)
 - P2P file sharing
 - voice over IP (e.g., Skype)
 - real-time video conferencing (e.g., Zoom)
 - Internet search
 - remote login
 - ...
- Q: *your* favorites?

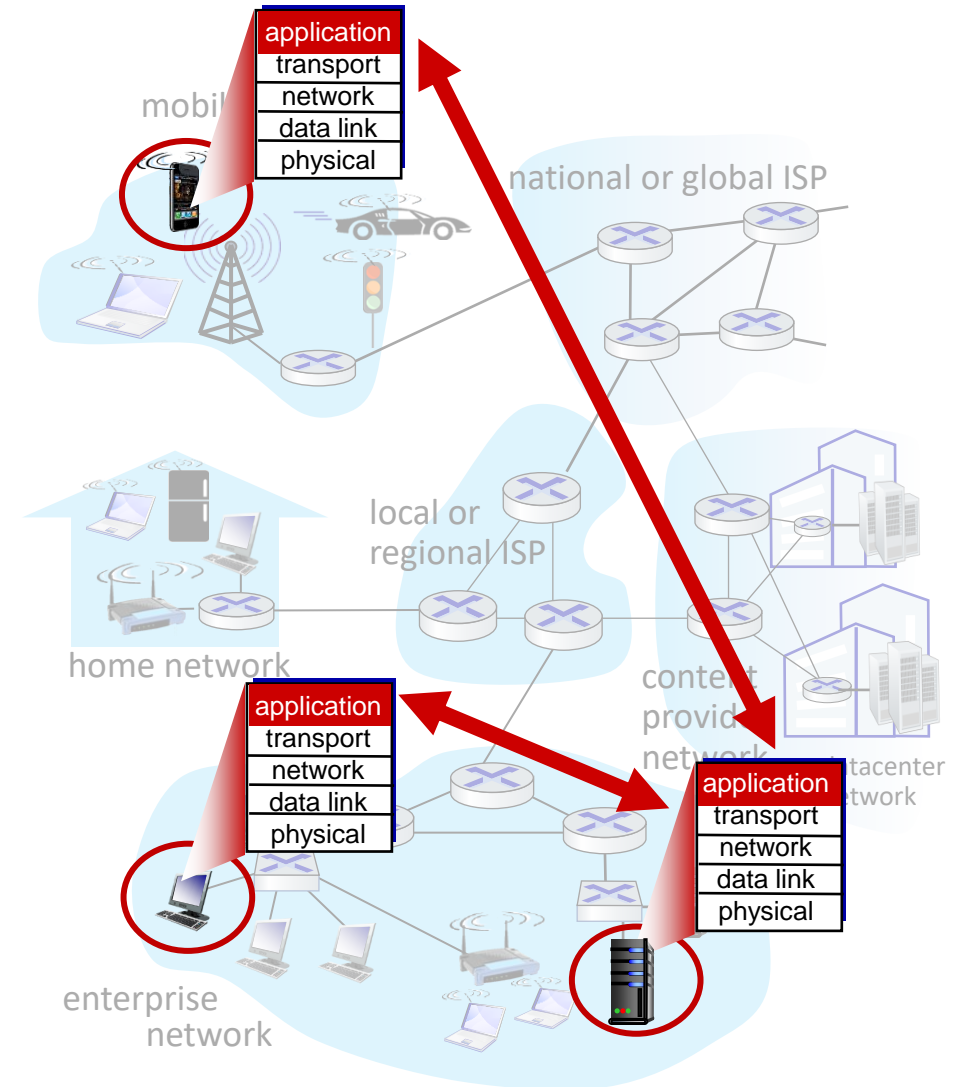
Creating a network app

write programs that:

- run on (different) end systems
- communicate over network
- e.g., web server software communicates with browser software

no need to write software for network-core devices

- network-core devices do not run user applications
- applications on end systems allow for rapid app development, propagation



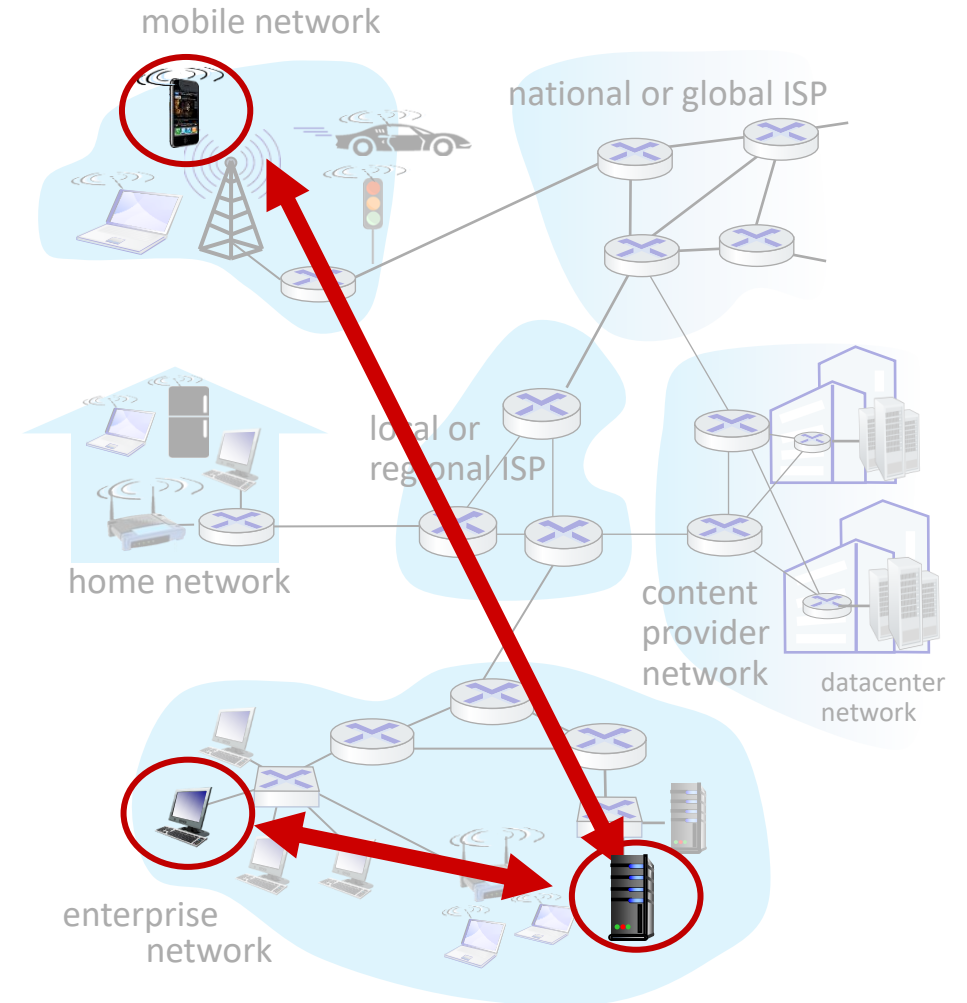
Client-server paradigm

server:

- always-on host
- permanent IP address
- often in data centers, for scaling

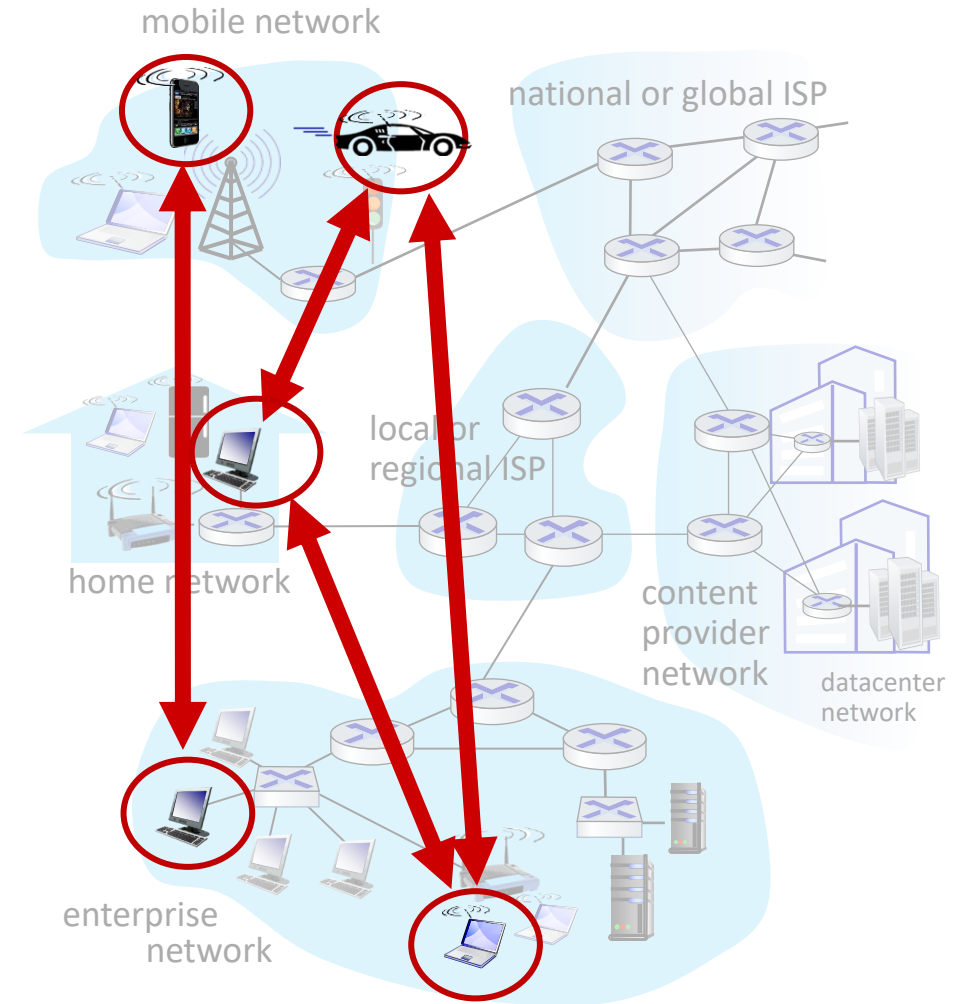
clients:

- contact, communicate with server
- may be intermittently connected
- may have dynamic IP addresses
- do *not* communicate directly with each other
- examples: HTTP, IMAP, FTP



Peer-peer architecture

- *no* always-on server
- arbitrary end systems directly communicate
- peers request service from other peers, provide service in return to other peers
 - *self scalability* – new peers bring new service capacity, as well as new service demands
- peers are intermittently connected and change IP addresses
 - complex management
- example: P2P file sharing



Processes communicating

process: program running within a host

- within same host, two processes communicate using **inter-process communication** (defined by OS)
- processes in different hosts communicate by exchanging **messages**

clients, servers

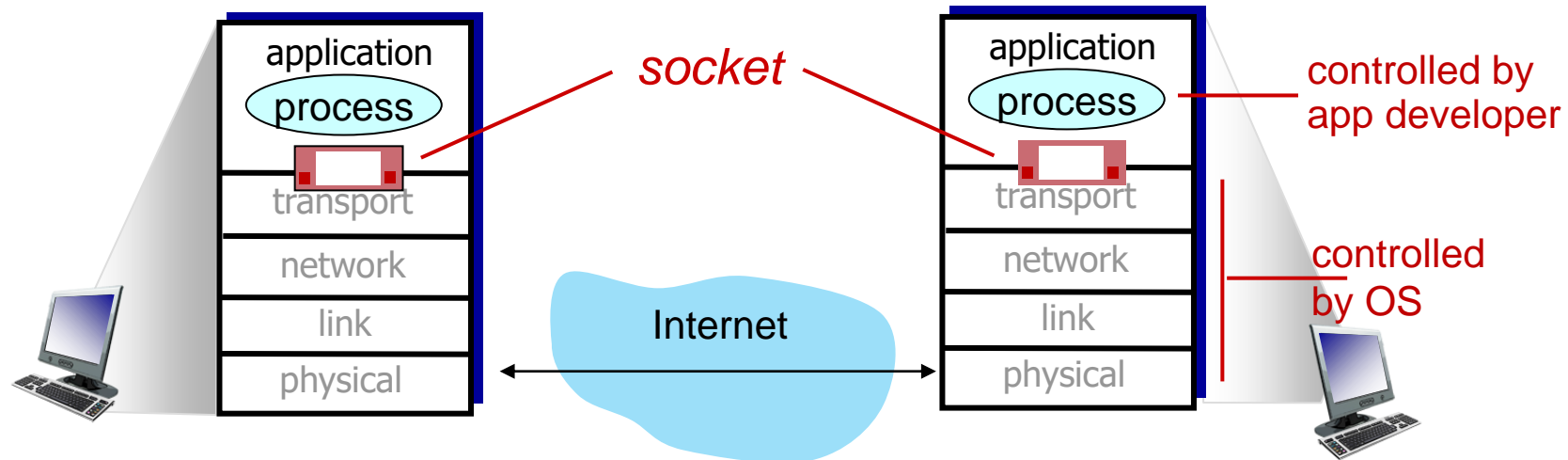
client process: process that initiates communication

server process: process that waits to be contacted

- note: applications with **P2P architectures also have client processes & server processes**

Sockets

- process sends/receives messages to/from its **socket**
- socket analogous to door
 - sending process shoves message out door
 - sending process relies on transport infrastructure on other side of door to deliver message to socket at receiving process
 - two sockets involved: one on each side



Addressing processes

- to receive messages, a process must have an *identifier*
- host device has unique 32-bit IP address
- Q: does IP address of host on which process runs suffice for identifying the process?
 - A: no, *many* processes can be running on same host
- *identifier* includes both IP address and port numbers associated with process on host.
- example port numbers:
 - HTTP server: 80
 - mail server: 25
- to send HTTP message to gaia.cs.umass.edu web server:
 - IP address: 128.119.245.12
 - port number: 80
- more shortly...

An application-layer protocol defines:

- **types of messages exchanged**,
 - e.g., request, response
- **message syntax**:
 - what fields in messages & how fields are delineated
- **message semantics**
 - meaning of information in fields
- **rules** for when and how processes send & respond to messages

open protocols:

- defined in RFCs, everyone has access to protocol definition
- allows for interoperability
- e.g., HTTP, SMTP

proprietary protocols:

- e.g., Skype, Zoom

What transport service does an app need?

data integrity

- some apps (e.g., file transfer, web transactions) require 100% reliable data transfer
- other apps (e.g., audio) can tolerate some loss

timing

- some apps (e.g., Internet telephony, interactive games) require low delay to be “effective”

throughput

- some apps (e.g., multimedia) require minimum amount of throughput to be “effective”
- other apps (“elastic apps”) make use of whatever throughput they get

security

- encryption, data integrity, ...

Transport service requirements: common apps

application	data loss	throughput	time sensitive?
file transfer/download	no loss	elastic	no
e-mail	no loss	elastic	no
Web documents	no loss	elastic	no
real-time audio/video	loss-tolerant	audio: 5Kbps-1Mbps video:10Kbps-5Mbps	yes, 10's msec
streaming audio/video	loss-tolerant	same as above	yes, few secs
interactive games	loss-tolerant	Kbps+	yes, 10's msec
text messaging	no loss	elastic	yes and no

Internet transport protocols services (Details in Chapter 3)

TCP service:

- *reliable transport* between sending and receiving process
- *flow control*: sender won't overwhelm receiver
- *congestion control*: throttle sender when network overloaded
- *connection-oriented*: setup required between client and server processes
- *does not provide*: timing, minimum throughput guarantee, security

UDP service:

- *unreliable data transfer* between sending and receiving process
- *does not provide*: reliability, flow control, congestion control, security, or connection setup.

Q: why bother? *Why* is there a UDP?

Internet applications, and transport protocols

application	application layer protocol	transport protocol
file transfer/download	FTP [RFC 959]	TCP
e-mail	SMTP [RFC 5321]	TCP
Web documents	HTTP 1.1 [RFC 7320]	TCP
Internet telephony	SIP [RFC 3261], RTP [RFC 3550], or proprietary	TCP or UDP
streaming audio/video	HTTP [RFC 7320], DASH	UDP or TCP
interactive games	WOW, FPS (proprietary)	UDP or TCP

Application Layer: Overview

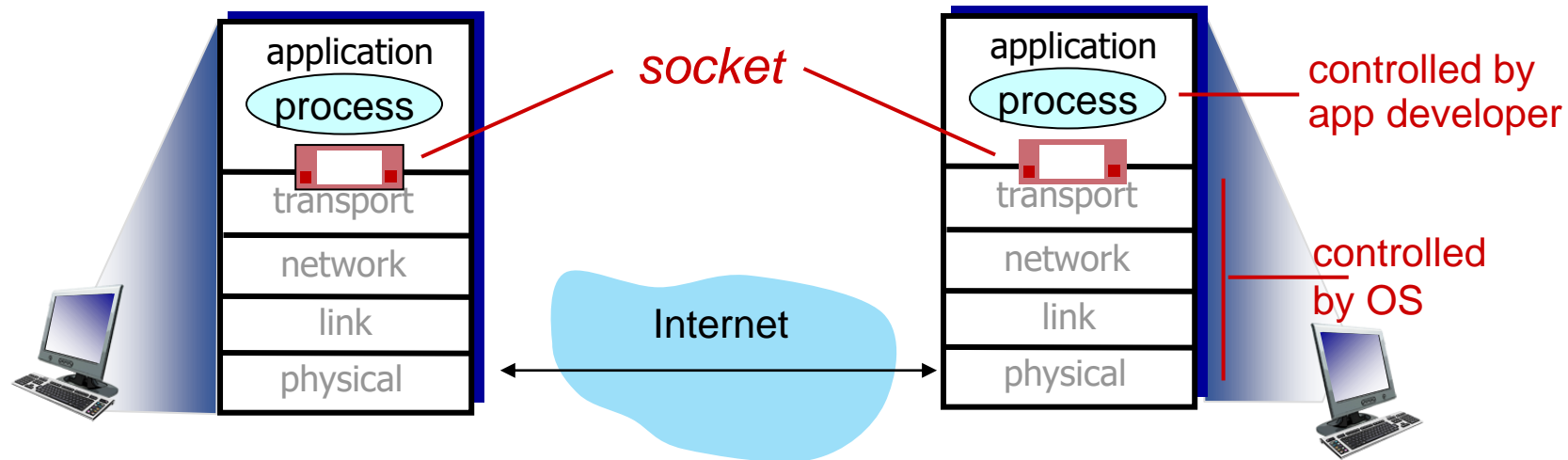
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Socket programming (also see Slide 9)

goal: learn how to build client/server applications that communicate using sockets

socket: door between application process and end-end-transport protocol



Socket programming

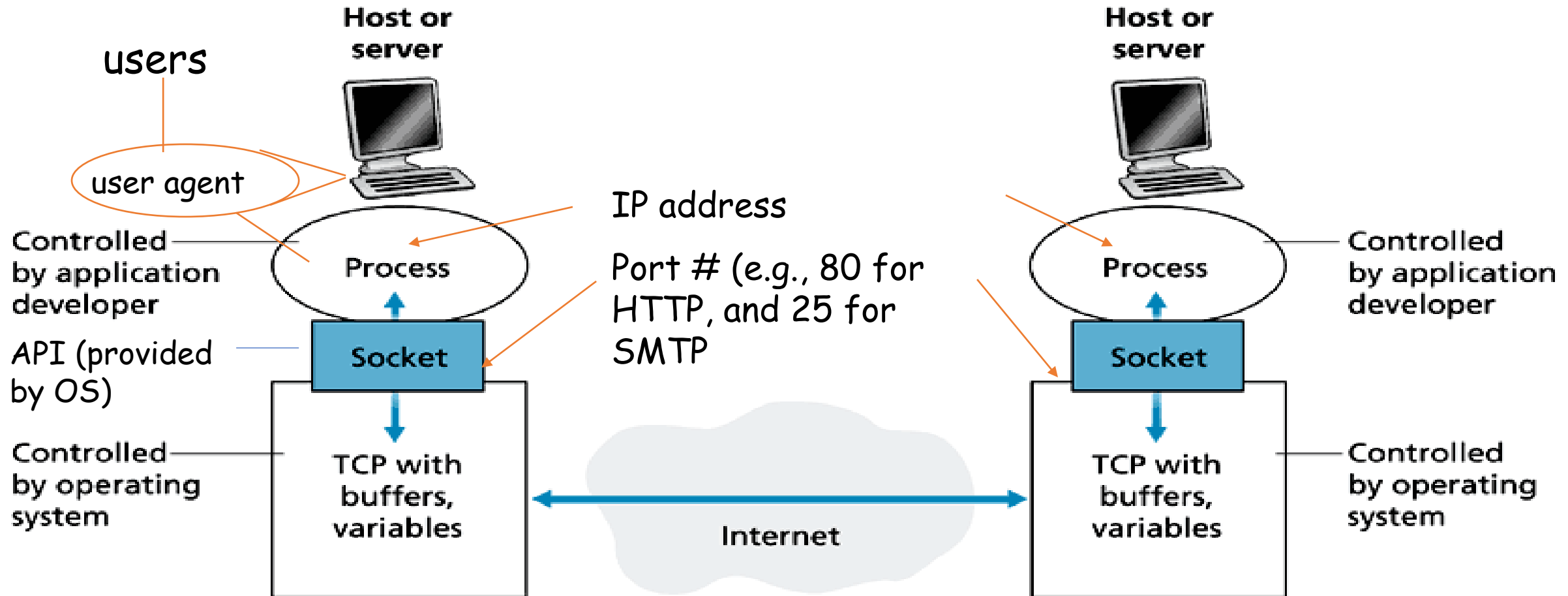
Two socket types for two transport services:

- *UDP*: unreliable datagram
- *TCP*: reliable, byte stream-oriented

More on these protocols later (in Chapter 3), but

- relevant to application programming (API)
- useful for PA1

An example: Email using HTTP or SMTP over TCP socket)



Another Application Example:

1. client reads a line of characters (data) from its keyboard and sends data to server
2. server receives the data and converts characters to uppercase
3. server sends modified data to client
4. client receives modified data and displays line on its screen

Socket programming with UDP

UDP: no “connection” between client and server:

- no handshaking before sending data
- sender (e.g. client) explicitly attaches its IP destination address and port #, in addition to the destination’s IP/port info to each packet
- receiver (e.g. server) extracts sender IP address and port# from received packet

UDP: transmitted data may be lost or received out-of-order

Application viewpoint:

- UDP provides *unreliable* transfer of groups of bytes (“datagrams”) between client and server processes

Client/server socket interaction: UDP



server (running on *serverIP*)

create socket, assigns port= x:

**serverSocket =
socket(AF_INET,SOCK_DGRAM)**

specify the type
of the UDP socket

read datagram from
serverSocket

Gets client IP address
and port y

write reply to
serverSocket
specifying
client address,
port number (y)

client



create socket (at some port y)

**clientSocket =
socket(AF_INET,SOCK_DGRAM)**

Create datagram with serverIP address
And port=x; **send** datagram via

clientSocket

Includes its own IP
address and port y by
the OS

read datagram from
clientSocket

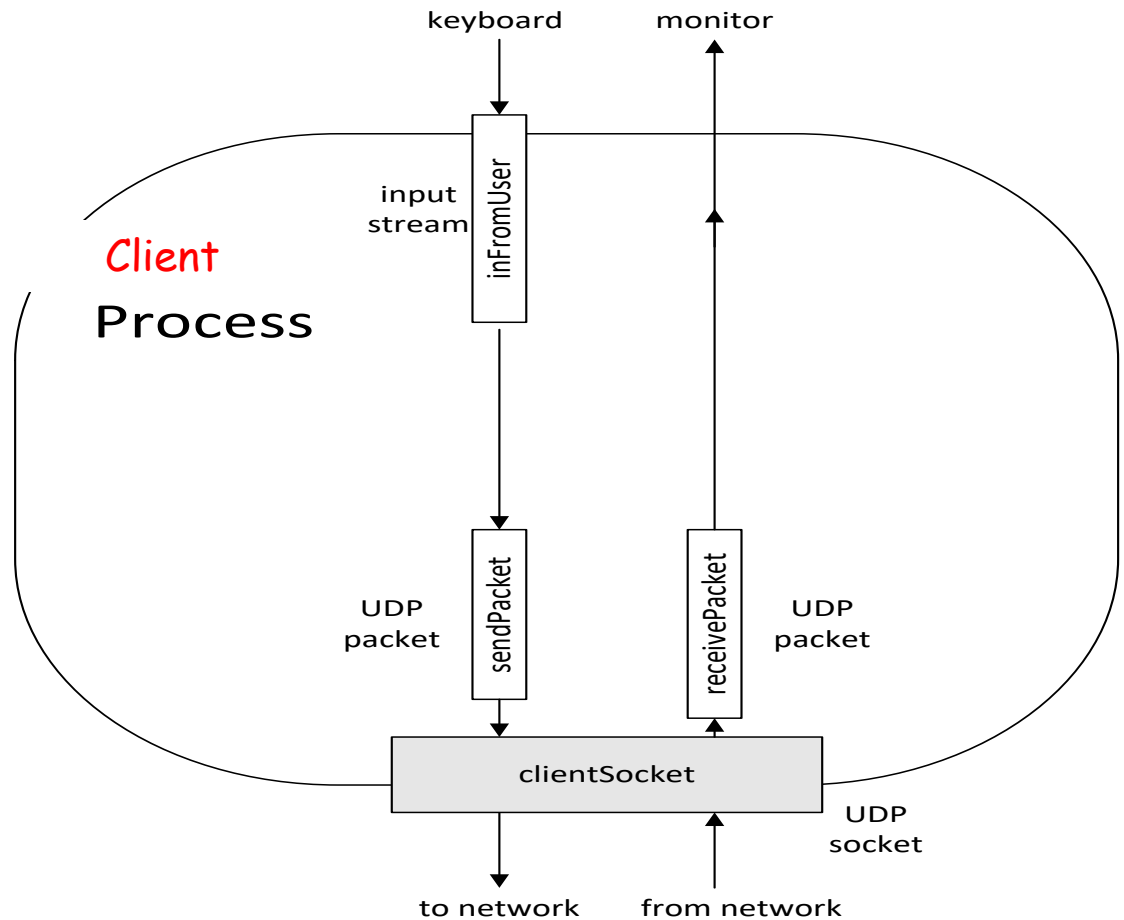
close
clientSocket

Example UDP Client (to request uppercase conversion)

Get keyboard input (in lowercase)
from users

Send to a server for uppercase
conversion

Receive the converted characters
and display them



Example app: UDP client

Python UDPClient

include Python's socket library → `from socket import *`

```
serverName = 'hostname'
```

```
serverPort = 12000
```

create UDP socket

(no need to specify port yet)

```
→ clientSocket = socket(AF_INET, SOCK_DGRAM)
```

get user keyboard input

```
→ message = raw_input('Input lowercase sentence:')

```

attach server name, port x to message; send into socket
(OS inserts this client's IP address and port y)

```
clientSocket.sendto(message.encode(),
                    (serverName, serverPort))
```

read reply characters from socket into string

```
→ modifiedMessage, serverAddress =  
    clientSocket.recvfrom(2048)
```

print out received string and close socket

```
→ print modifiedMessage.decode()
   clientSocket.close()
```


Example app: UDP server

Python UDPServer

```
from socket import *
serverPort = 12000
create UDP socket → serverSocket = socket(AF_INET, SOCK_DGRAM)
bind socket to local port number 12000 → serverSocket.bind(('', serverPort))
print ("The server is ready to receive")
loop forever → while True:
    Read from UDP socket into message, getting → message, clientAddress = serverSocket.recvfrom(2048)
    client's address (client IP and port)      modifiedMessage = message.decode().upper()
    send upper case string back to this client → serverSocket.sendto(modifiedMessage.encode(),
                                                                    clientAddress)
```