

Discrete p-values

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Why is the p-value defined as $p = \mathbb{P}_0(T \geq t)$ and not $p = \mathbb{P}_0(T > t)$?

It obviously only matters if T is discrete (or has discrete components).

Discrete test statistics are ubiquitous, and especially so in cyber-security applications (e.g. count data, presence/absence events).

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Examples

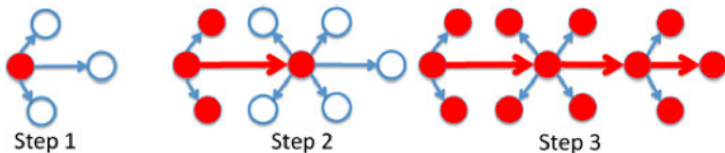
At the beginning of this module we motivated anomaly detection with the following suggestions from [the MS-ISAC guide to DDOS attacks](#) (2017):

- “Look for a large number of SYN packets, from multiple sources, over a short duration” (SYN flood)
- “Look for a large number of inbound UDP packets over irregular network ports coming from a large number of source IP addresses” (UDP flood)
- “look for a significant amount of inbound ICMP traffic from a large number of sources” (IMCP flood – e.g. SMURF attack)
- “look for a large number of inbound traffic from a significant number of source IP addresses with a destination port of 80” (HTTP GET flood)

All of these examples would produce discrete-valued anomalies.

Another example

One of the simplest and most direct approaches to intrusion detection is to monitor new connections, e.g. to detect network traversal motifs such as:



(Reproduced from “Scan statistics for the online detection of locally anomalous subgraphs”, Joshua Neil *et al.* , 2013).

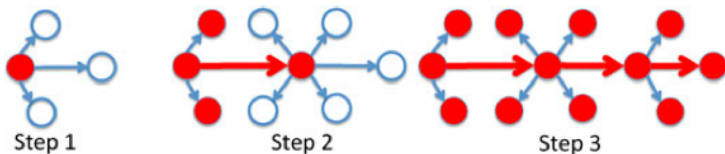
The correct way to formalise this is to consider the test statistic

$$T = \mathbb{I}(\text{edge between } i \text{ and } j),$$

which is binary-valued.

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$$T = \mathbb{I}(\text{edge between } i \text{ and } j),$$

which is binary-valued.

Partially discrete p-values occur, for example, with censored data, e.g.

$$T = \min(\text{time to next event, end of observation period/time-out}).$$

The usual stochastic order (Definition 1)

We say that a random variable X is stochastically smaller than (or stochastically dominated by) a random variable Y , denoted $X \leq_{st} Y$, in the usual stochastic order, if their cumulative distribution functions satisfy

$$F_X(a) \geq F_Y(a),$$

for all $a \in \mathbb{R}$.

There are many other stochastic orders, but this is what we mean when we say that X is stochastically smaller than Y .

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The usual stochastic order (equivalent Definition 2)

We say that a random variable X is stochastically smaller than (or stochastically dominated by) a random variable Y , denoted $X \leq_{st} Y$, in the usual stochastic order, if for any non-decreasing function f we have

$$E\{f(X)\} \leq E\{f(Y)\}.$$

It is fairly straightforward to see that Definition 2 implies Definition 1 by considering the family of functions $f_a(x) = \mathbb{1}_{[a, \infty)}$. The reverse implication can also be shown but is out of scope for this course.

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It is fairly straightforward to see that Definition 2 implies Definition 1 by considering the family of functions $f_a(x) = -\mathbb{I}(x \leq a)$. The reverse implication can also be shown but is out of scope for this course.

Discrete p-values are super-uniform

Recall the result that (week 1, Fundamentals of hypothesis testing, page 11) under the null hypothesis the p-value is uniform if the test statistic T is continuous.

Theorem (discrete p-values are super-uniform)

Under the null hypothesis, for any distribution of the test statistic T , the p-value is stochastically larger than a uniform random variable on $[0, 1]$.

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Coin-tossing example

Consider testing the hypothesis

$$H_0 : X_1, \dots, X_n \stackrel{i.i.d}{\sim} \text{Bernoulli}(1/2),$$

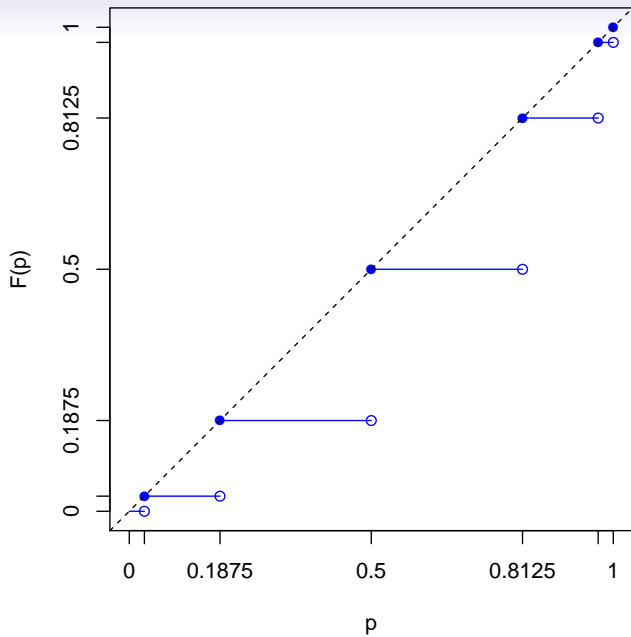
versus,

$$H_1 : X_1, \dots, X_n \stackrel{i.i.d}{\sim} \text{Bernoulli}(\mu) \quad \mu > 0.5,$$

on the basis of the test statistic $T = \sum_{i=1}^n X_i$.

Coin-tossing example ($n = 5$)

| t | $\mathbb{P}_0(T = t)$ | $p = \mathbb{P}_0(T \geq t)$ |
|-----|-----------------------|------------------------------|
| 0 | 0.03125 | 1 |
| 1 | 0.15625 | 0.96875 |
| 2 | 0.3125 | 0.8125 |
| 3 | 0.3125 | 0.5 |
| 4 | 0.15625 | 0.1875 |
| 5 | 0.03125 | 0.03125 |



Proof.

Recall that when viewed as a random variable the p-value is defined as $P = \mathbb{P}_0(T^* \geq T \mid T)$.

Let $S_0(x) = \mathbb{P}_0(T \geq x)$. Then,

$$\begin{aligned}\mathbb{P}_0(P \leq x) &= \mathbb{P}_0\{\mathbb{P}_0(T^* \geq T \mid T) \leq x\} \\ &= \mathbb{P}_0\{S_0(T) \leq x\}.\end{aligned}$$



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Let y denote the maximum value of S_0 not exceeding x , and t the supported value at which it occurs. Then,

$$\begin{aligned}\mathbb{P}_0\{S_0(T) \leq x\} &= \mathbb{P}_0\{T \geq t\} \\ &= S_0(t) = y \leq x,\end{aligned}$$

so that $F_U(x) = x \geq \mathbb{P}_0(P \leq x) = F_P(x)$, where F_U, F_P are respectively the distribution functions of a uniform random variable U and the p-value P under H_0 . □

Conservative tests

- Remember that in a classical hypothesis testing framework we reject the null hypothesis when the p-value $p \leq \alpha$.
- When the test statistic is continuous this guarantees (week 1, Fundamentals of hypothesis testing, page 12) that:
the probability of rejecting the null hypothesis if it holds is α .
- More generally, i.e. when the test is discrete or partially discrete, we have shown that
the probability of rejecting the null hypothesis if it holds does not exceed α .
- A test procedure which rejects the null hypothesis with probability less than α (resp. more than α) is called *conservative* (resp. *liberal*).
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Mid-p-values and randomised p-values

Discrete p-values computed using $\mathbb{P}_0(T \geq t)$ can turn out to be impractically conservative, and there are two proposed solutions:

1. The mid-p-value,

$$q = \frac{1}{2}\mathbb{P}_0(T \geq t) + \frac{1}{2}\mathbb{P}_0(T > t),$$

2. and the randomised p-value,

$$r = X\mathbb{P}_0(T \geq t) + (1 - X)\mathbb{P}_0(T > t),$$

where X is a randomly generated uniform random variable on $[0, 1]$.

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|-----|-----------------------|------------------------------|-----------------------|----------|-----------|
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| 1 | 0.15625 | 0.96875 | 0.8125 | 0.890625 | 0.8154765 |
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| 3 | 0.3125 | 0.5 | 0.1875 | 0.34375 | 0.2327139 |
| 4 | 0.15625 | 0.1875 | 0.03125 | 0.109375 | 0.0879319 |
| 5 | 0.03125 | 0.03125 | 0 | 0.015625 | 0.0085281 |

Consider a test statistic $T \in \{0, 1, \dots\}$ and let $v_i = \mathbb{P}_0(T \geq i)$, so that $v_0 = 1 \geq v_1 \geq v_2 \geq \dots > 0$.

If we observe $t = i$:

1. the p-value is $p = v_i$
2. the mid-p-value is $q = \frac{1}{2}v_i + \frac{1}{2}v_{i+1}$
3. the randomised p-value is $r = Xv_i + (1 - X)v_{i+1}$, where $X \sim \text{uniform}[0, 1]$

Since $X \in [0, 1]$ and $v_{i+1} < v_i$ we have:

$$v_{i+1} \leq Xv_i + (1 - X)v_{i+1} \leq v_i.$$

Furthermore, for $x \in [v_{i+1}, v_i]$,

$$\mathbb{P}_0(r \leq x) =$$

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$$\mathbb{P}_0(r \leq x) = \mathbb{P}_0 \left(X \leq \frac{x - v_{i+1}}{v_i - v_{i+1}} \right) = \frac{x - v_{i+1}}{v_i - v_{i+1}}.$$

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Differentiating gives a constant and therefore

$$r \sim \text{uniform}[v_{i+1}, v_i].$$

A quick digression

While the notation's still fresh in your minds, we'll prove this simple result.

Remember that in the coin-tossing example we observed that the distribution function of P under H_0 satisfied $F(v_i) = v_i$. In fact, more generally,

$$F(v_i) = \mathbb{P}_0(P \leq v_i) = \mathbb{P}(T \geq i) = v_i.$$

Define, as with the ordinary p-value, the random counterpart to the randomised p-value:

$$R = X\mathbb{P}_0(T^* \geq T \mid T) + (1 - X)\mathbb{P}_0(T^* > T \mid T).$$

Theorem (the randomised p-value is uniform)

Under the null hypothesis, for any distribution of the test statistic T , the randomised p-value is uniformly distributed on $[0, 1]$.

Proof

For simplicity we will prove the theorem assuming a discrete test statistic $T \in \{0, 1, \dots\}$.

For $x \in (0, 1]$, let $a = \max(y \in \{0\} \cup \{v_i\} : y < x)$ and $b = \min(y \in \{v_i\} : y \geq x)$.

Then,

$$\begin{aligned}\mathbb{P}_0(R \leq x) &= \mathbb{P}_0(R \leq x \mid R > b)\mathbb{P}_0(R > b) \\ &\quad + \mathbb{P}_0\{R \leq x \mid R \in (a, b]\}\mathbb{P}\{R \in (a, b]\}, \\ &\quad + \mathbb{P}_0(R \leq x \mid R \leq a)\mathbb{P}_0(R \leq a).\end{aligned}$$

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We have:

1. $\mathbb{P}_0(R \leq x \mid R > b)\mathbb{P}_0(R > b) = 0.$

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We have:

1. $\mathbb{P}_0(R \leq x \mid R > b)\mathbb{P}_0(R > b) = 0.$
2. $\mathbb{P}\{R \in (a, b]\} = \mathbb{P}(P = b) = F(b) - F(a).$

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3. $\mathbb{P}_0(R \leq x \mid R \leq a)\mathbb{P}_0(R \leq a) = 1 \times F(a)$.

Proof continued

Finally,

$$\mathbb{P}_0\{R \leq x \mid R \in (a, b]\} = \mathbb{P}_0(Xb + (1 - X)a \leq x)$$

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Altogether:

$$\mathbb{P}_0(R \leq x) = 0 + \frac{x - F(a)}{F(b) - F(a)}\{F(b) - F(a)\} + F(a) = x.$$

A comparison

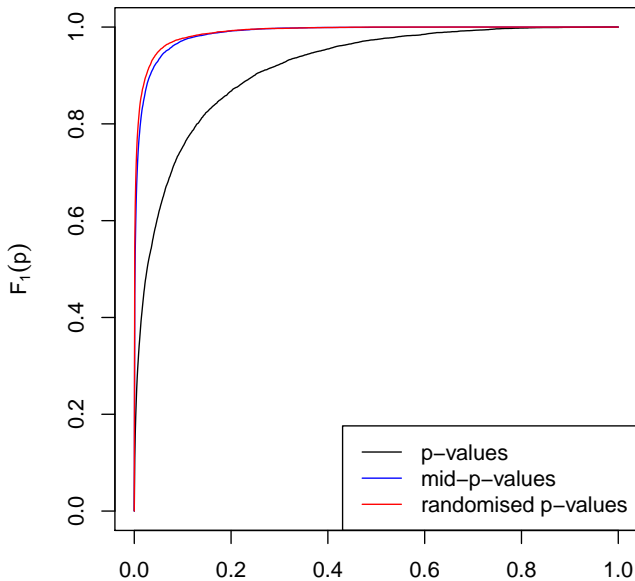
We will generate m “coin-toss” experiments with different, random numbers of tosses.

Each experiment generates m p-values, mid-p-values, and randomised p-values all computed based on a Bernoulli(0.5) null hypothesis. Each of the three sets is combined using Fisher's method.

It will be in your homework to prove that Fisher's method is conservative when ordinary discrete p-values are input. Fisher's method is only approximate with mid-p-values, but is exact with randomised p-values.

In the next example (see code), the coins are actually from a Bernoulli(0.75) distribution and we draw the power curves, $\hat{\mathbb{P}}_1(\text{combined p-value} \leq x)$, for each choice of input.

Comparison



A worrying result

“Most people will find repugnant the idea of adding yet another random element to a result which is already subject to the errors of random sampling” (Stevens, 1950).

Yet, randomised p-values often perform best in practice and in theory.

A final result about discrete p-values

Consider testing the hypothesis

$$H_0 : X \sim \text{categorical}\{\mu_1, \dots, \mu_m\},$$

versus,

$$H_1 : X \sim \text{categorical}\left\{\frac{1}{m}, \dots, \frac{1}{m}\right\},$$

often known as the “flat alternative”.

LRT

The **LRT** is (replacing probability density functions with probability mass functions)

$$\frac{1/m}{\mu_X},$$

which is equivalent to $1/\mu_X$.

Therefore, upon observing x , we compute the p-value

$$p = \mathbb{P}_0(T \geq 1/\mu_x),$$

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or “the probability of observing an outcome as unlikely as x .”