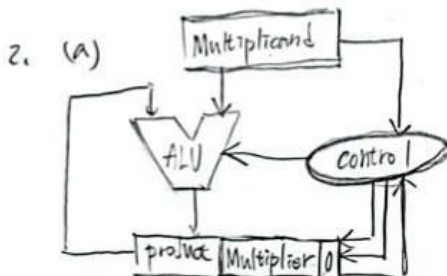
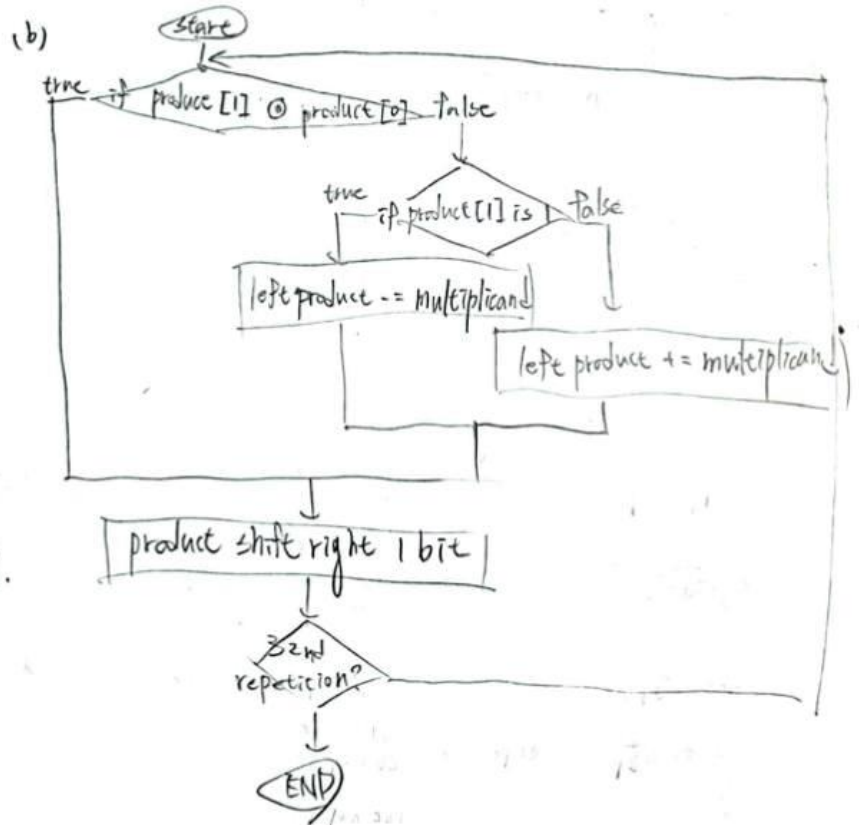


1. $147 = 10010011_{10}$ $28 = 00011100$

Iteration	step	Remainder	Divisor
0	shift left Remainder 1 bit	0000001 00100110	00011100
1	left Remainder \geq Divisor	11100101 00100110	00011100
2	if left Remainder < 0		
3	left Remainder \neq Divisor	00000001 00100110	
4	shift left Remainder 1 bit	00000010 01001100	
5	if not repetition 9th		
...			
19	shift left Remainder 1 bit	00100100 11000000	
20	if not repetition 19th		
21	left Remainder \geq Divisor	00001000 11000000	
22	if left Remainder > 0		
23	shift left Remainder 1 bit and set 1	00010001 10000001	
24	if not repetition 9th		
25	left Remainder \geq Divisor	11110101 10000001	
26	if left Remainder < 0		
27	left Remainder \neq Divisor	00010001 10000001	
28	shift left Remainder 1 bit	00100011 00000010	
29	if not repetition 9th		
30	left Remainder \geq Divisor	00000111 00000010	
31	if left Remainder > 0		
32	shift left Remainder 1 bit and set 1	00001110 00000010	
33	if not repetition 9th		
34	left Remainder shift left 1 bit	00000111 00000010	



(c) in some cases,
Booth's algorithm can achieve
the goal with less computation
than the original multiplication method.



8.

1) $23.25 = 10111.01_{(2)} = 1.011101 \times 2^4$

Exponent = 4 + bias = 4 + 127 = 131 = 10000011₍₂₎

Fraction = 011101₍₂₎

S = 1 (because -23.25 is negative number \Rightarrow (-1)')

$\frac{1}{3} \frac{10000011}{\text{Exponent}} \frac{011101000000000000000000}{\text{Fraction}}$

(b) $\frac{1}{3} \frac{0111001}{\text{Exponent}} \frac{011010000000000000000000}{\text{Fraction}}$
 Negative $-127 - 5 = -132$
 $-1.01101_2 \times 2^{-6} = -1.40625 / 64 = -0.02197265625$
 $0.25 - 0.125 + 0.03125 = 0.15625$

(c)

$\frac{0}{\text{Sign}} \frac{11111110}{\text{Exponent}} \frac{111111111111111111111111}{\text{Fraction}}$
 positive $\text{Max } 127 + 127 = 254$

(d) $\frac{0}{\text{Sign}} \frac{00000000}{\text{Exponent}} \frac{000000000000000000000000}{\text{Fraction}}$
 positive $\text{min } -127 + 127 = 0$

	sign	Exponent	Fraction	object represented
e)	0	1111111	000000000000000000000000	$+\infty$
	1	1111111	000000000000000000000000	$-\infty$
	1/0	1111111	$\neq 000000000000000000000000$	NaN

(f) Negative 130

$1 \ 100.0001 \ 0000 \ 0010 \ 1011 \ 0001 \ 1111 \ 0010$

$-1.0001_2 \times 2^{(130-127)} = 1.0210554599161963 \times 8 = 8.16844367980957$

4. (a) $0.3125 = 0.0101_2 \Rightarrow 1.0101 \times 2^0$
 $1 \ 0111 \ 01010 \ 00000 \#$

(b) $10011000.01 + 0.0101010101$
 $= 1.00110000 \times 10^7 + 1.01010101 \times 10^{-2}$
 $= 1.00110000 \times 10^7 + 0.0000000101010101 \times 10^7$
 $= 1.0011000101010101 \times 10^7$
 $\begin{array}{l} 0 \ 0011 \ 0011000101010101 \\ \hline \text{GRS} \end{array} \left\{ \begin{array}{l} \text{Guard} = 0 \text{ 捨去} \\ \text{Round} = 1 \\ \text{Sticky} = 0(0\dots 0) \end{array} \right.$

$\Rightarrow 0 \ 0011 \ 00110001010101 \#$

(c) $110.0111001 \times 1000.01101$
 $= 1.100111001 \times 2^2 \times 1.00001101 \times 2^3$
 $= 110110.00101100101 \times 2^{(2+3)}$
 $= 1.1011000101100101 \times 2^{(5+5)}$
 $\begin{array}{l} 0 \ 01010 \ 10110001011001010 \\ \hline \text{GRS} \end{array} \left\{ \begin{array}{l} \text{Guard} = 0 \text{ 捨去} \\ \text{Round} = 1 \\ \text{Sticky} = 0(0\dots 0) \end{array} \right.$
 $\Rightarrow 0 \ 01010 \ 10110001011001 \#$