

1. Pseudo code:

for  $i$  in range  $(1, n-1)$ :current  $\leftarrow$  s.headwhile current.next  $\neq$  nullif current.value  $>$  current.next.value

swap(current, current.next)

current  $\leftarrow$  current.nextTime Complexity =  $O(n^2)$ 2. in-order traversal  $\Rightarrow$  finding the  $k$ -th largest keyreverse traversal  $\Rightarrow$  finding the  $k$ -th smallest key, by binary search tree's orderliness.Time Complexity =  $O(n)$ 

3.

Pseudo code:

function merge-heaps( $T_1, T_2$ ):if not  $T_1$  then return  $T_2$ if not  $T_2$  then return  $T_1$ if  $T_1.value > T_2.value$  then swap( $T_1, T_2$ ) $T_1.right = \text{merge-heaps}(T_1.right, T_2)$ swap( $T_1.right, T_1.left$ )return  $T_1$  $T = \text{merge-heaps}(T_1, T_2)$ Time Complexity =  $O(k_1 + k_2)$ 

4. (a) For each node, we can choose whether to swap its left and right subtrees.  
 Therefore, if  $T$  has  $n$  nodes, the total number of binary trees isomorphic to  $T$  is  $2^n$ . two choice

(2) pseudo code:

isomorph - distance ( $T_1, T_2$ ):

if not  $T_1$  and not  $T_2$  then return 0

if not  $T_1$  or not  $T_2$  then return

if  $T_1.value = T_2.value$  then return  $\text{float}('inf')$  Trees are not isomorphic

dist 1 = isomorph - distance ( $T_1.left, T_2.left$ ) + isomorph - distance ( $T_1.right, T_2.right$ )

dist 2 = isomorph - distance ( $T_1.left, T_2.right$ ) + isomorph - distance ( $T_1.right, T_2.left$ )

return  $\min(\text{dist}1, \text{dist}2) + 1$  if  $\text{dist}1 \neq \text{dist}2$  else 0

time complexity =  $O(n)$