COMP.2030 Lab 9

1. Suppose that MIPS registers \$t0 and \$t1 are mapped to the x86 registers %eax and %ebx, respectively. Convert the MIPS instructions below left to the corresponding x86 instructions. Use the same label X in your x86 code.

```
lw $t1, X($zero)
li $t0, 1
lw $t1, X($t0)
move $t0, $t1
sw $t0, ($t1)
addi $t0, $t0, 1
```

2. Suppose that the MIPS registers \$t0, \$t1 and \$t2 registers are mapped to the x86 registers %rax, %rbx and %rcx, respectively. Convert the MIPS instruction on the left below to the corresponding x86 instructions. Use the same labels in your x86 code as shown in the MIPS code.

```
lw
          $t2, X($zero)
     li
          $t0, 1
rept:
          $t0, $t2, exit # $t0>=$t2
     bge
          $t0, $t0, 2
     sll
          $t1, X($t0)
     lw
     bgt $t2, $t1, next # $t2>$t1
     move $t2, $t1
          $t0, ($t1)
     SW
next:
         $t0, $t0, 2
     sra
     addi $t0, $t0, 1
```

1. The function fun a has the overall structure shown below:

```
long fun_a(unsigned long x){
                                      # x in %rdi
    long val = 0;
                                      fun_a:
                                                    $0, %eax
                                             movl
    while (
                             ){
                                             jmp
                                                     .L5
                                      .L6:
                                                    %rdi, %rax
                                             xorq
                                                    %rdi # shift right 1
                                             shrq
                                      .L5:
                                                    %rdi, %rdi
                                             testq
    }
                                                     .L6
                                             jne
                                             andl
                                                    $1, %eax
                                             ret
    return val;
}
```

The gcc C compiler generates the x86-64 assembly code on the right. Reverse engineer this assembly code and fill in the missing parts of the fun_a definition so the C code does the same thing.

2. The gcc C compiler generates the assembly code blow to the right. Reverse engineer the operation of this code and fill in the missing parts of the C code to the left so that it does the same thing.

```
# x in %rdi, n in %esi
Loop:
                %esi, %ecx
        movl
                $1, %edx
        movl
                $0, %eax
        mov1
                .L2
        jmp
.L3:
                %rdi, %r8
        movq
                %rdx, %r8
        andq
                %r8, %rax
        orq
                %cl, %rdx
        salq
.L2:
                %rdx, %rdx
        testq
        jne
                .L3
        ret
```