

1.)

$$OPT(v) = \max_{u:v \in \text{adj}(u)} \{OPT(u) + C(u,v)\}$$

$v$  in topological order

Base cases :

$$OPT(s) = 0$$

$$OPT(v) = -\infty$$

Solution : Find largest of  $OPT(v)$ , backtrack and find the largest path.

$\Rightarrow$  Complexity :  $O(m+n)$

2.)

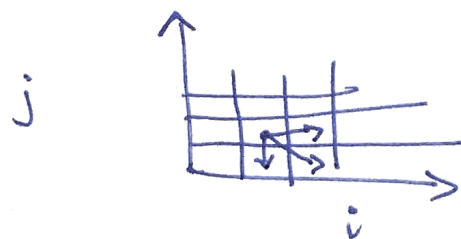
Sol. 1 : Find LCS of  $S$  &  $S^R$

$$LCS[i,j] = 1 + LCS[i-1, j-1], \text{ if } S[i] = T[j]$$

$$LCS[i,j] = \max(LCS[i-1, j], LCS[i, j-1]), \text{ if } S[i] \neq T[j]$$

Sol 2 :

$$OPT(i,j) = \begin{cases} \text{if } S_i \neq S_j, & OPT(i+1, j-1) + 2 \\ \text{else,} & \max(OPT(i+1, j), OPT(i, j-1)) \end{cases}$$



Move to the cell according to the opt directions, to set the ~~ans~~ answer, through backtracking.

$\Rightarrow$  Complexity =  $\Theta(n^2)$

3.)

Base cases :

$$\text{maxCoin}[i][i] = \text{coins}[i]$$

$$\text{left} = \begin{cases} \text{if } (i+2 \leq j) \\ \text{maxCoin}[i+2][j] \\ \text{else} \end{cases}$$

$$\text{bottom} = \begin{cases} \text{if } (j-2 \geq i) \\ \text{maxCoin}[i][j-2] \\ \text{else} \\ 0 \end{cases}$$

$$\text{maxCoin}[i][j] = \max \left( \text{coins}[i] + \min(\text{left}, \text{maxCoin}[i+1][j-1]), \text{coins}[j] + \min(\text{maxCoin}[i-1][j-1], \text{bottom}) \right)$$

$$\Rightarrow \text{Complexity: } \Theta(n^2)$$

4.)

$$0 \leq i \leq n-1$$

$$\text{cutRod}(n) = \max(\text{price}[i] + \text{cutRod}(n-i-1))$$

$$\Rightarrow \text{Complexity} = \Theta(n)$$