

	0	1	2	3	4	5	6	7	8
0				X		X			
1			X				X		
2									
3			X				X		
4				X		X			

48  
24

	0	1	2	3	4	5	6	7	8
0	0	0	0	X	3	X	5	5	1
1	0	0	X	3	3	3	X	4	1
2	0	0	0	0	X	3	3	3	1
3	0	0	X	0	0	0	X	2	1
4	0	0	0	X	0	X	1	1	1

	0	1	2	3	4	5	6	7	8
0	1	1	1	X	0	X	0	0	0
1	1	2	X	0	0	0	X	0	0
2	1	3	3	3	X	0	0	0	0
3	1	4	X	3	3	3	X	0	0
4	1	5	5	X	3	X	0	0	0

60  
33

	0	1	2	3	4	5	6
0	1	1	1	1	1	1	1
1	1	2	X	1	X	1	2
2	1	X	0	1	1	X	2
3	1	1	1	X	1	1	3
4	1	X	1	1	1	X	2

$dp[i][j]$   
 $= dp[i-1][j]$  ←  
 $+ dp[i][j-1]$  ←  
 $dp[i][j]$

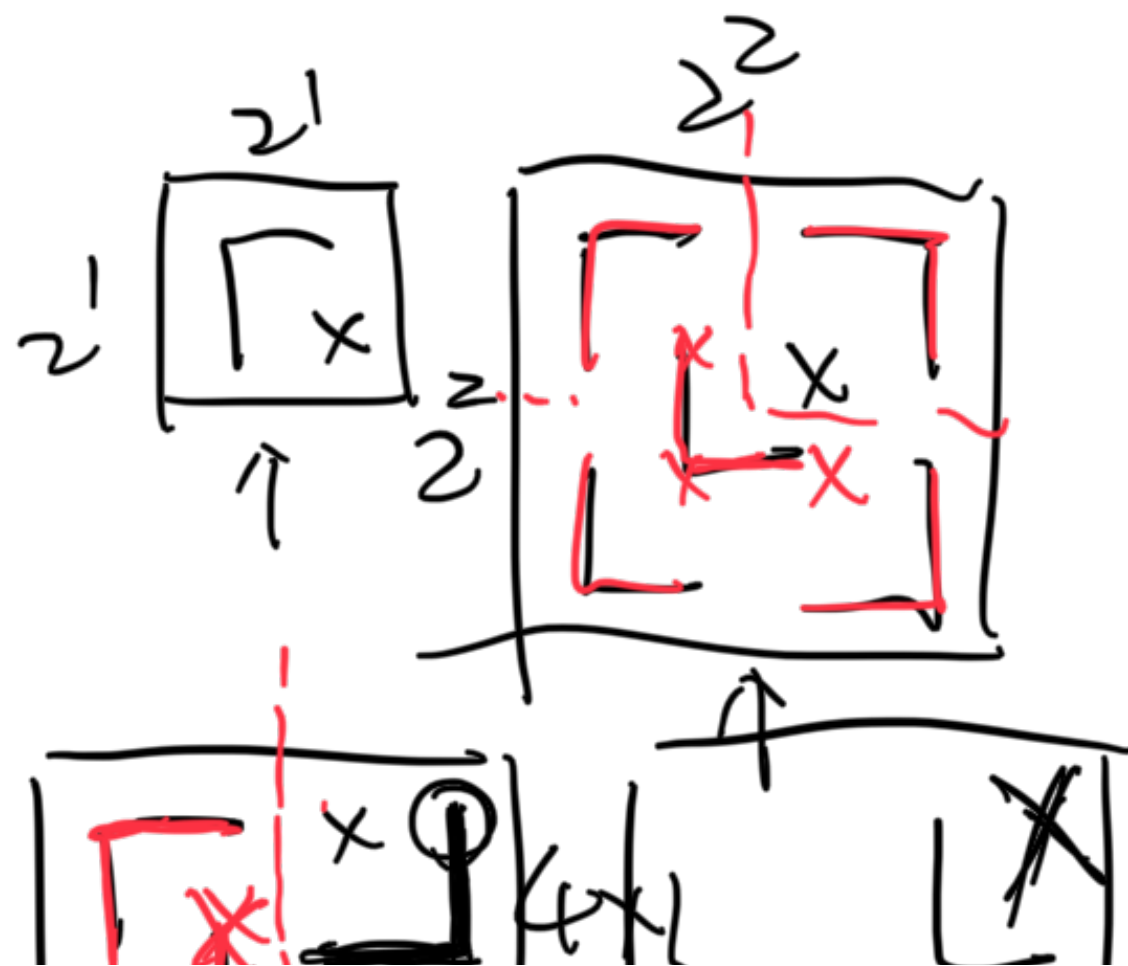
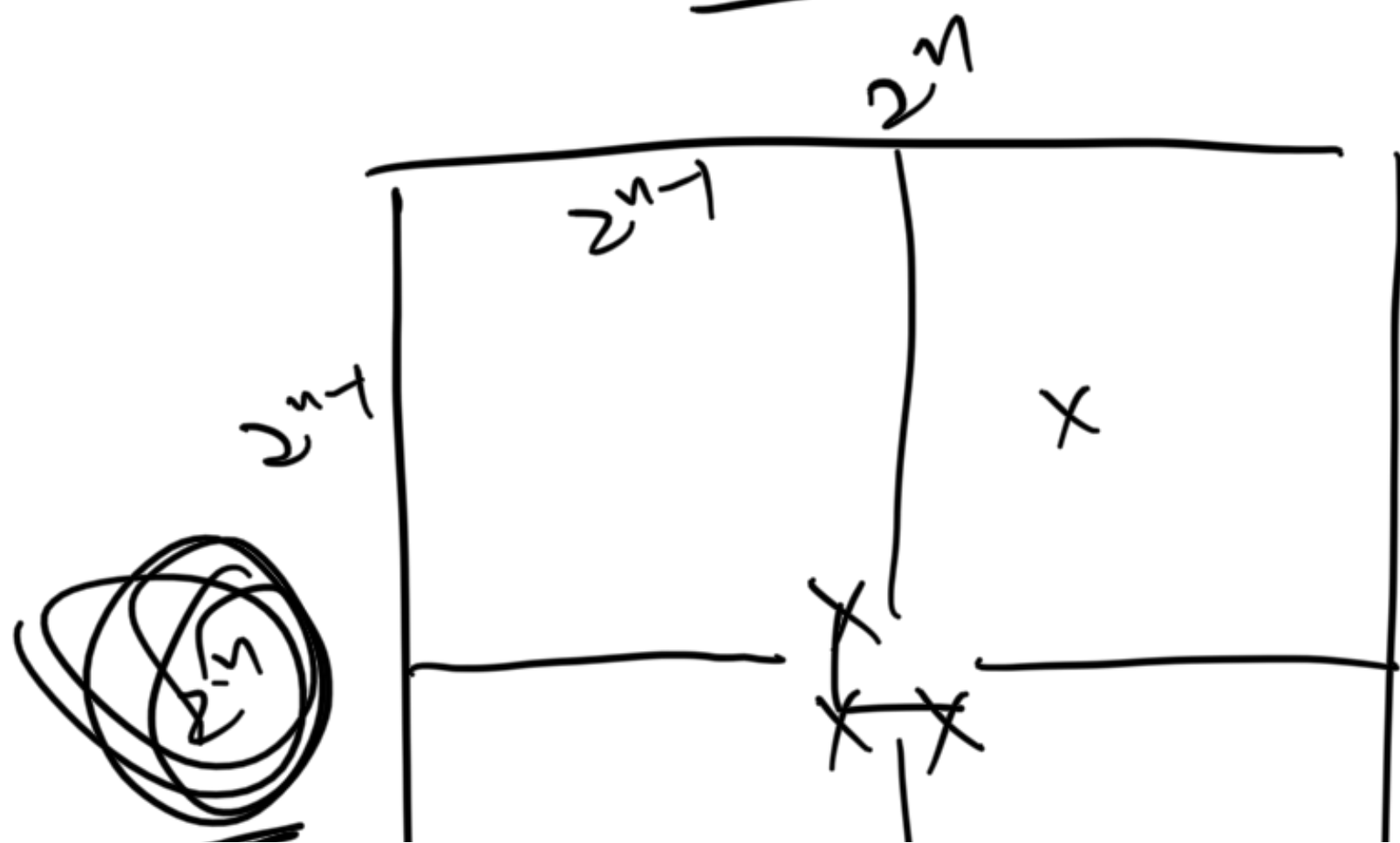
$$\begin{array}{c|cccccc} 4 & 1 & \times & 1 & 1 & 1 & \times & 3 \\ 5 & 1 & 1 & \times & 1 & \times & 0 & 3 \\ 6 & 1 & 2 & 2 & 3 & 3 & 3 & 6 \end{array} \leftarrow$$

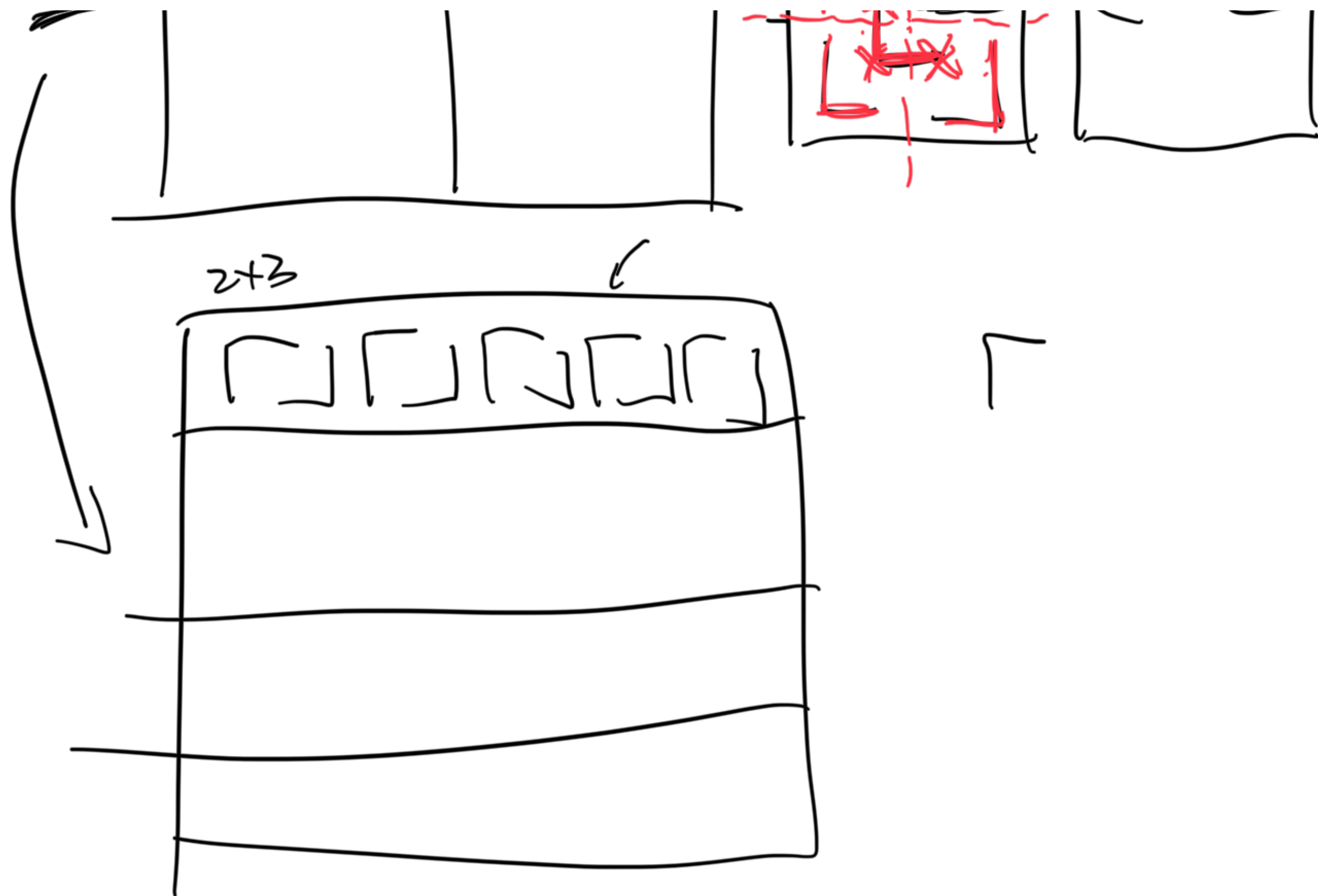
$M(0,0)$  走到  $(i,j)$  的方法数

$dp[a][b][c][d]$

$\underline{a, b, c, d} \rightarrow \underline{a-1, b, c, d}$   
 $\searrow$   
 $\underline{a, b-1, c, d}$

状态转移方程





$dp[i][j]$

用前  $i$  种花, 共  $j$  盆摆放的方法数



$n=2$   $m=4$

3 2

A B  
A B  
A

AAB B  
AAH B

B

4

dp	0	1	2	3	4
0	0	0	0	0	0
1	1	1	1	1	0
2	1	2	3	3	2

BB →

A B  
AA AB BB  
AAH AAB AB B

AAAB  
AABB

dp[1][2]  
前1种花放2个位置  
BB

前2种花放到4个位置

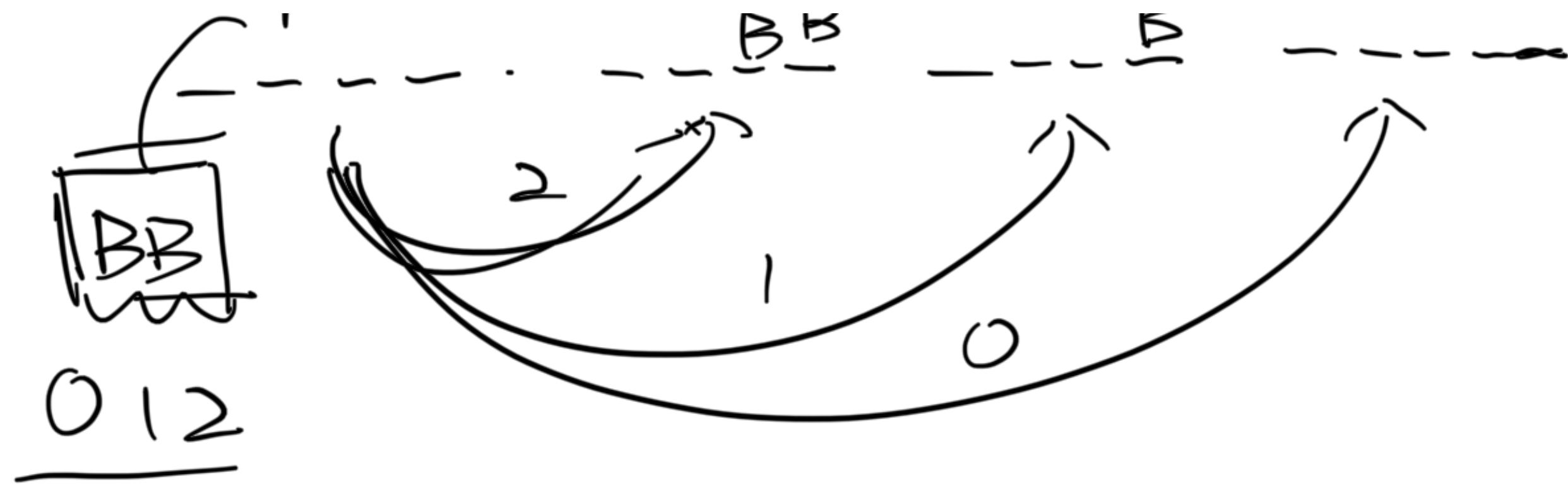
2种花  
4个位置

1种花  
2个位置

只能用前1种花放2个位置

1种花 2个      1种 3个      1种 4个

$$\rightarrow dp[2][4] = dp[1][2] + dp[1][3] + dp[1][4] + \dots$$



$$\rightarrow dp[2][3] = dp[1][3] + dp[1][2] + dp[1][1]$$

$$\rightarrow dp[2][2] = dp[1][1] + dp[1][2] + \underbrace{dp[1][0]}$$

$$\rightarrow dp[2][1] = \underbrace{dp[1][1] + dp[1][0]}$$

$dp[i]$ .

$$\underbrace{dp[2][j]} = dp[1][j] + dp[1][j-1] + \dots + dp[1][j-a[2]]$$

$$= dp[i][j] + \dots + dp[i][\max(0, j - a[i])]$$

$$dp[i][j] = dp[i-1][j] + dp[i-1][j-1] + \dots + dp[i-1][j - a[i]]$$

i: 2 ~ n  
j: 0 ~ m

②

$$dp[i][j] = 0$$

for (int k=0; k <= a[i]; k++)  
if (j-k < 0) break;

$$dp[i][j] += dp[i-1][j-k]$$

①

for (int i=0; i <= a[1]; i++)

$$dp[1][i] = 1$$

on t[i]

dp[n][m]

for (i ~ 3) t[i]

for (j 0 ~ 10)

t[i]-1



$$dp[3][70] = \max(dp[2][70], dp[2][70-1] + 2)$$

$$dp[2][69] = \max(dp[1][69], dp[0][69-69] + 1)$$

		0	1	2	3	4	68	69	70
70	100	0	0	0	0	0	0	0	0
69	1	0	0	0	0	0	0	0	0
1	2	0	0	0	0	0	0	0	0
2	3	0	2	2	2	2	2	2	3

↑ t    ↑ v

$$dp[3][1] = \max(dp[2][1], dp[2][1-1] + 2)$$

for (int i = 1; i <= n; i++)  
 for (int j = t[i]; j <= m; j++)



$$dp[i-1][j - a[i]] + v[i] \quad dp[i-1][j]$$

$$dp[i][j] = \max(dp[i-1][j])$$

$j - t[i] \sim 70$

$$dp[i][j - t[i]] + v[i]$$

$$j - t[i] < 0$$

		0	1	2	3	4	...	...	...	...	68	69	70
71	1	0	0	0	0	0					0	0	0
69	1	2	0	0	0	0					0	1	0
1	2	3	0	2	0	0					0	1	3

$$dp[i][j]$$

$$t = j$$

$$70$$