

## Example - Circle Area

$$area = \pi r^2$$

- ◆ Area of a circle:
  - `val pi = 3.14159;`
  - `val pi = 3.14159 : real`
  - `fun area (r) = pi*r*r;`
  - `val area = fn : real -> real`
  - `area 2.0;`
  - `val it = 12.56636 : real`

ML Declarations.2

## Identifiers in ML

- ◆ **val** declaration binds a name to a value.
- ◆ A name can not be used to change its value !
  - Actually a constant
- ◆ A name can be reused for another purpose

```
- val pi = "pi";  
val pi = "pi" : string
```
- ◆ If a name is declared again the new meaning is adopted afterwards

```
- pi;  
val it = "pi" : string
```

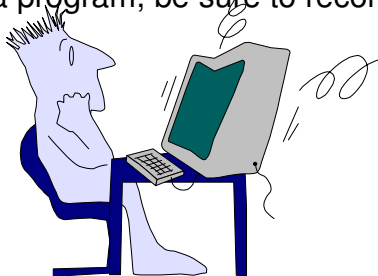
but does not affect existing uses of the name

```
- area(1.0)  
val it = 3.14159 : real
```

ML Declarations.3

## Is permanence of names a good feature?

- ◆ LUCKY: redefining a function cannot damage the system or your program.
- ◆ BUT: redefining a function called by your program may have no visible effect.
- ◆ NOTE: when modifying a program, be sure to recompile the entire file



ML Declarations.4

## 'val' and 'val rec'

- ◆ We can define function using val
  - `val sq = fn x => x*x;`
- ◆ What about recursive functions?
  - `fun f(n) = if n=0 then 1 else n * f(n-1);`
  - `val f = fn (n) => if n=0 then 1 else n * ??;`
  - `val rec f = fn (n) => if n=0 then 1  
                    else n * f(n-1);`
- ◆ 'val rec' stands for recursive definition and it is just like 'fun'

ML Declarations.5

## Pattern Matching

- ◆ Patterns can be used to simplify function definition
  - `fun factorial 0 = 1`
  - | `factorial n = n * factorial(n-1);`
  - `val factorial = fn : int -> int`
- ◆ When the function is called, the first pattern to match the actual parameter determines which expression on the right-hand-side will be evaluated.
- ◆ Patterns can consist
  - Constants - int, real, string, etc ...
  - Constructs - tuples, datatype constructs
  - Variables - all the rest
  - Underscore - a wildcard

Later ...

ML Declarations.6

## Pattern Matching

- ◆ When matching a pattern P to a value X, the matching is done recursively - "from outside to inside".
- ◆ If matching succeeded, any variable in the pattern is binded with the corresponding value in X
- ◆ There is no binding where the wildcard is used
- ◆ Example

```
- fun foo (x,1) = x
    | foo (1,_) = 0
    | foo _ = ~1;
val foo = fn : int * int -> int
- foo(3,1);
val it = 3 : int
- foo(1,3);
val it = 0 : int
- foo(2,2);
val it = ~1 : int
```

`foo(1,1) = 1`  
*Since matching is  
done in the order of  
definition*

ML Declarations.7

## Patterns in Conditional Expression

- ◆ Patterns can be used in a **case** conditional expression

● **case** E **of** P1 => E1 | ... | Pn => En

```
- case p-q of
    0 => "zero"
  | 1 => "one"
  | 2 => "two"
  | n => if n<10 then "lots"
        else "lots &lots";
```

- If P<sub>i</sub> is the first to match then the result is the value of E<sub>i</sub>
- Equivalent to an expression that defines a function by cases and applies it to E
- Scope of **case**: No symbol terminates the case expression!
  - Enclose in parentheses to eliminate ambiguity

ML Declarations.8

## Type Abbreviation

- ◆ You can give new name to existing type:

```
- type vec = real*real;  
type vec = real * real  
- infix ++;  
- fun (x1,y1) ++ (x2,y2) : vec = (x1+x2,y1+y2);  
val ++ = fn: (real * real) * (real * real) -> vec  
- (3.6,0.9) ++ (0.1,0.2) ++ (20.0,30.0);  
(23.7,31.1) : vec
```

- ◆ The new name is only an alias - it is acceptable where ever the original name is acceptable and vice versa

ML Declarations.9

## Local Declarations in Expressions

- ◆ let D in E end

```
- fun fraction(n,d)=  
    (n div gcd(n,d), d div gcd(n,d));  
val fraction = fn: int*int -> int*int  
- fun fraction(n,d)=  
    let val com = gcd(n,d)  
    in (n div com, d div com) end;  
val fraction = fn: int*int -> int*int
```

- D may be a compound declaration
  - D1;D2;...;Dn
- The semicolons are optional

- ◆ "let D in E end" Can be simulated using anonymous functions

```
- fun fraction(n,d)=  
    (fn com => (n div com, d div com)) (gcd(n,d));
```

ML Declarations.10

## Nested Local Declaration

- ◆ ML allows nested function definitions

```
- fun sqroot a =  
    let val acc=1.0e~10  
        fun findroot x =  
            let val nextx = (a/x + x)/2.0  
            in if abs (x-nextx)<acc*x  
               then nextx else findroot nextx  
            end  
        in findroot 1.0 end;  
val sqroot = fn: real -> real
```

ML Declarations.11

## Local Declarations in Declarations

- ◆ **local D1 in D2 end**
  - Behaves like the list D1;D2 in let
  - D1 is visible only within D2

- ◆ Used to hide a declaration

```
- local  
    fun itfib (n,prev,curr):int =  
        if n=1 then curr  
        else itfib (n-1,curr,prev+curr)  
    in  
        fun fib (n) = itfib(n,0,1)  
    end;  
val fib = fn : int -> int
```

- ◆ Why not simply nest the declaration of `itfib` within `fib`?

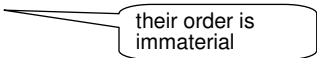
ML Declarations.12

## Comparing `let` and `local`

- ◆ - `fun fraction(n,d)=  
 let val com = gcd(n,d)  
 in (n div com, d div com) end;`
- ◆ `val fraction = fn: int*int -> real`
- ◆ - `local  
 fun itfib (p,prev,curr):int=  
 if p=1 then curr  
 else itfib (p-1,curr,prev+curr)  
 in  
 fun fib (n) = itfib(n,0,1)  
 end;  
 val fib = fn : int -> int`

ML Declarations.13

## Simultaneous Declarations (collateral)

- ◆ `val Id1 = E1 and ... and Idn = En`
  - evaluates `E1,...,En`
  - and only then declares the identifiers `Id1,...,Idn`
- ◆ Example: Swapping the values of names
  - `val x = y and y = x`
  - `val (x,y) = (y,x)`
- ◆ Note the last declaration. Actually the allowed format is
  - `val P = E;`
- ◆ So it can be used to disassemble tuples
  - `val a = (1,2,3);`
  - `val a = (1,2,3) : int * int * int`
  - `val (_,x,_) = a;`
  - `val x = 2 : int`

ML Declarations.14

## Mutually Recursive functions

◆ Example:

$$\frac{\pi}{4} = \sum_{k=0}^{\infty} \frac{1}{4k+1} - \frac{1}{4k+3} = 1 - \frac{1}{3} + \frac{1}{5} - \dots$$

```
fun pos d = neg(d-2.0) + 1.0/d
and neg d = if d>0.0 then pos(d-2.0)-1.0/d
             else 0.0;

fun sum(d,one)=
  if d>0.0 then sum(d-2.0,~one)+one/d
  else 0.0;
```

ML Declarations.15

## Translating an imperative code to mutually recursive functions

◆ Emulating `goto` statements...

```
var x:=0; y:=0; z:=0;
F: x:=x+1; goto G
G: if y<z then goto F else (y:=x+y; goto H)
H: if z>0 then (z:=z-x; goto F) else stop
```

```
- fun F(x,y,z)=G(x+1,y,z)
= and G(x,y,z)=if y<z then F(x,y,z) else H(x,x+y,z)
= and H(x,y,z)=if z>0 then F(x,y,z-x) else (x,y,z);
val F = fn : int * int * int -> int * int * int
val G = fn : int * int * int -> int * int * int
val H = fn : int * int * int -> int * int * int
- F(0,0,0);
val it = (1,1,0) : int * int * int
```

ML Declarations.16